

# Lecture 3: Multicast Routing

7COM1030 – Multicast and Multimedia Networking

Dr. Xianhui Che (Cherry)

[x.che@herts.ac.uk](mailto:x.che@herts.ac.uk)

School of Computer Science  
University of Hertfordshire, UK



# Topics

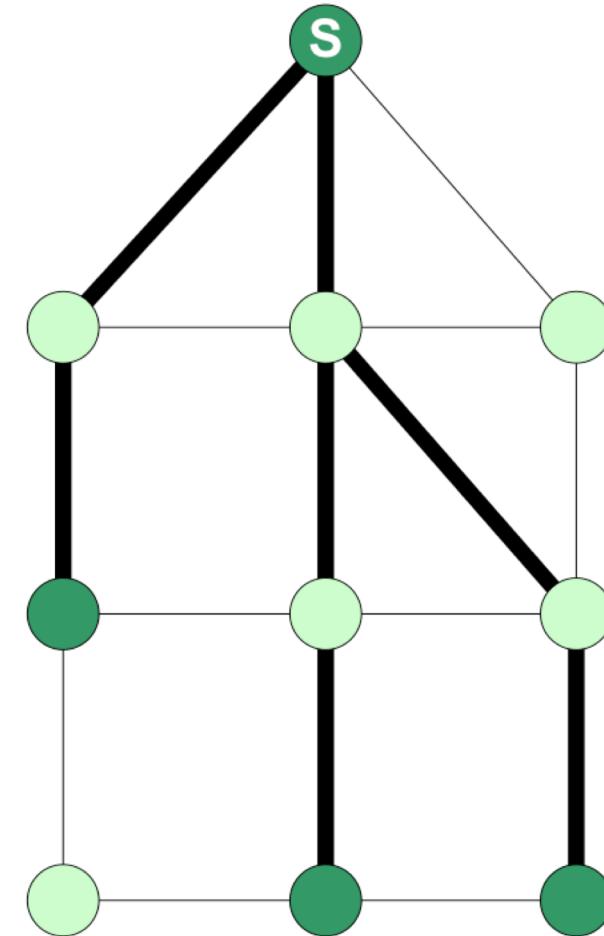
- ▶ Multicast Routing Solutions
- ▶ Source Based Trees
- ▶ Core Based Trees

# Multicast Routing Protocols in the Internet

- ▶ **Distance Vector Multicast Routing Protocol (DVMRP):**
  - First multicast routing protocol
  - Implements flood-and-prune
- ▶ **Multicast Open Shortest Path First (MOSPF):**
  - Multicast extensions to OSPF. Each router calculates a shortest-path tree based on link state database
  - Not widely used
- ▶ **Core Based Tree (CBT):**
  - First core-based tree routing protocol
- ▶ **Protocol Independent Multicast (PIM):**
  - Runs in two modes: PIM Dense Mode (PIM-DM) and PIM Sparse Mode (PIM-SM).
  - PIM-DM builds source-based trees using flood-and-prune
  - PIM-SM builds core-based trees as well as source-based trees with explicit joins.

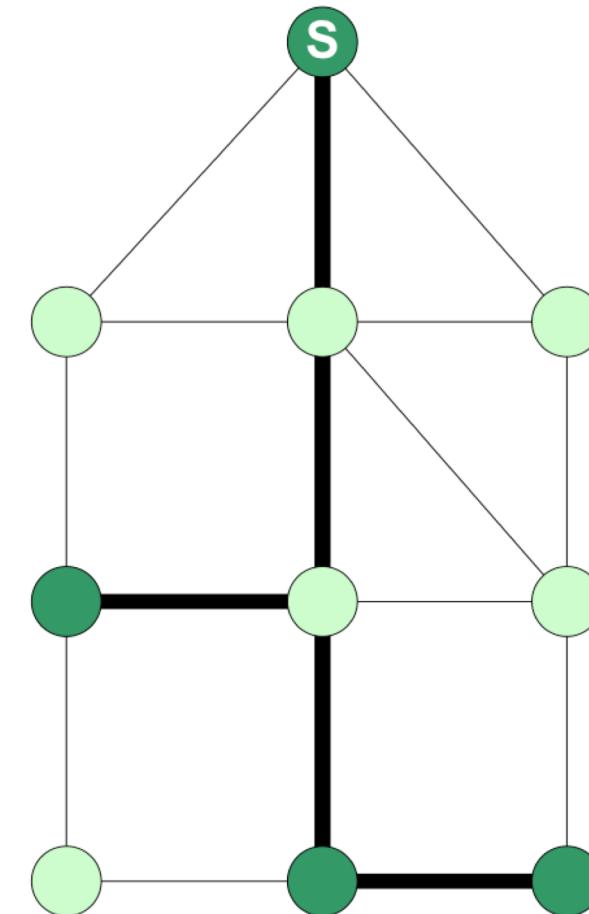
# Multicast Routing Solution 1

- ▶ **Problem:** Embed a tree such that all multicast group members are connected by the tree
- ▶ **Solution 1: Shortest Path Tree or source-based tree**
  - Build a tree that minimizes the path cost from the source to each receiver
    - *Good tree if there is a single sender*
    - *If there are multiple senders, need one tree per sender*
    - *Easy to compute*



# Multicast Routing Solution 2

- ▶ **Problem:** Embed a tree such that all multicast group members are connected by the tree
- ▶ **Solution 2: Minimum-Cost Tree**  
Build a tree that minimizes the total cost of the edges
  - *Good solution if there are multiple senders*
  - *Very expensive to compute (not practical for more than 30 nodes)*



# Multicast Routing in Practice

► Routing Protocols implement one of two approaches:

- **Source Based Tree:**

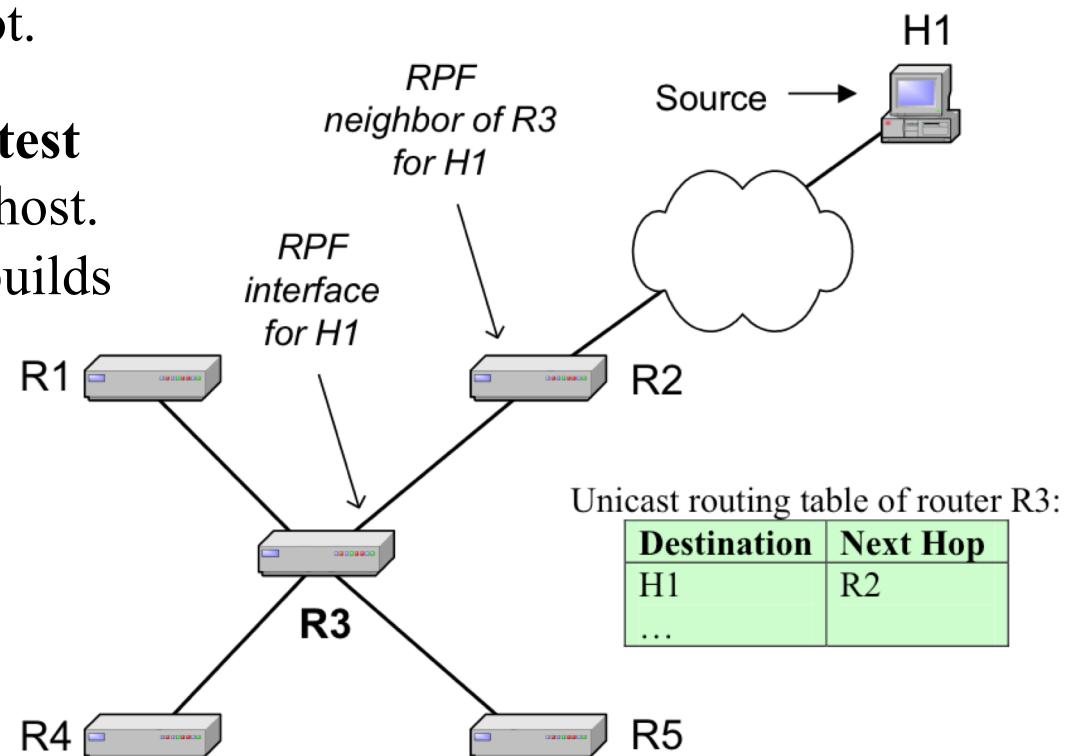
- Essentially implements Solution 1.
  - Builds one shortest path tree for each sender
  - Tree is built from receiver to the sender
  - **reverse shortest path / reverse path forwarding**

- **Core-based Tree:**

- Build a single distribution tree that is shared by all senders
  - Does not use Solution 2 (because it is too expensive)
  - Selects one router as a “core” (also called “rendezvous point” )
  - All receivers build a shortest path to the core
  - **reverse shortest path /reverse path forwarding**

# Reverse Path Forwarding (RPF)

- ▶ RPF builds a shortest path tree in a distributed fashion by taking advantage of the unicast routing tables.
- ▶ **Main concept:** Given the address of the root of the tree (e.g., the sending host), a router selects as its upstream neighbor in the tree, which is the next-hop neighbor for forwarding unicast packets to the root.
- ▶ This concept leads to a **reverse shortest path** from any router to the sending host. The union of reverse shortest paths builds a **reverse shortest path tree**.



# Multicast Routing table

- ▶ Routing table entries for source-based trees and for core-based trees are different
  - **Source-based tree:** (Source, Group) or (S, G) entry.
  - **Core-based tree:** (\*, G) entry.

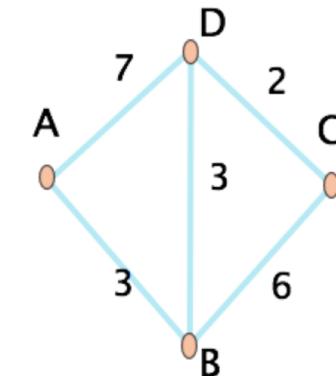
| Source IP address | Multicast group | Incoming interface (RPF interface) | Outgoing interface list |
|-------------------|-----------------|------------------------------------|-------------------------|
| S1                | G1              | I1                                 | I2, I3                  |
| *                 | G2              | I2                                 | I1, I3                  |

# Recap of Unicast Routing Protocols

- ▶ Depending on the routing protocol, different factors may be taken into account, e.g. number of hops, latency, bandwidth, etc.
- ▶ Shortest path routing aims to find a path that has the least amount of cost.
- ▶ There are two types algorithms to update the routing tables.

| Category  | Link State | Distance Vector |
|-----------|------------|-----------------|
| Algorithm | Dijkstra   | Bellman–Ford    |
| Example   | OSPF       | RIP             |

# Distance Vector Example



(1) First exchange for all nodes to find best 1-hop routes

| To   | A        | B        | C        | D        |
|------|----------|----------|----------|----------|
| says | says     | says     | says     | says     |
| A    | 0        | $\infty$ | $\infty$ | $\infty$ |
| B    | $\infty$ | 0        | $\infty$ | $\infty$ |
| C    | $\infty$ | $\infty$ | 0        | $\infty$ |
| D    | $\infty$ | $\infty$ | $\infty$ | 0        |

| A           | B      | C           | D      |
|-------------|--------|-------------|--------|
| learns      | learns | learns      | learns |
| Cost        | Cost   | Cost        | Cost   |
| Next        | Next   | Next        | Next   |
| 0 --        | 3 A    | $\infty$ -- | 7 A    |
| 3 B         | 0 --   | 6 B         | 3 B    |
| $\infty$ -- | 6 C    | 0 --        | 2 C    |
| 7 D         | 3 D    | 2 D         | 0 --   |

(2) Second exchange for all nodes to find best 2-hop routes

| To   | A        | B    | C        | D    |
|------|----------|------|----------|------|
| says | says     | says | says     | says |
| A    | 0        | 3    | $\infty$ | 7    |
| B    | 3        | 0    | 6        | 3    |
| C    | $\infty$ | 6    | 0        | 2    |
| D    | 7        | 3    | 2        | 0    |

| A      | B      | C      | D      |
|--------|--------|--------|--------|
| learns | learns | learns | learns |
| Cost   | Cost   | Cost   | Cost   |
| Next   | Next   | Next   | Next   |
| 0 --   | 3 A    | 9 B    | 6 B    |
| 3 B    | 0 --   | 5 D    | 3 B    |
| 9 D    | 5 D    | 0 --   | 2 C    |
| 6 B    | 3 D    | 2 D    | 0 --   |

(3) Third exchange for all nodes to find best 3-hop routes

| To   | A    | B    | C    | D    |
|------|------|------|------|------|
| says | says | says | says | says |
| A    | 0    | 3    | 9    | 6    |
| B    | 3    | 0    | 5    | 3    |
| C    | 9    | 5    | 0    | 2    |
| D    | 6    | 3    | 2    | 0    |

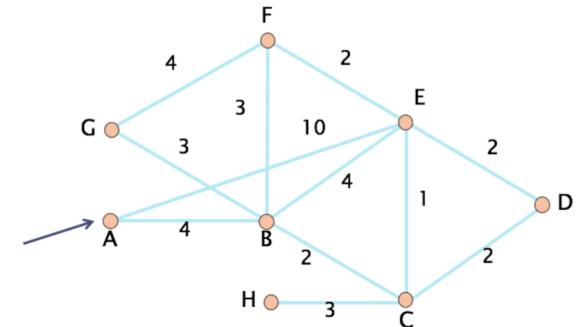
| A      | B      | C      | D      |
|--------|--------|--------|--------|
| learns | learns | learns | learns |
| Cost   | Cost   | Cost   | Cost   |
| Next   | Next   | Next   | Next   |
| 0 --   | 3 A    | 8 D    | 6 B    |
| 3 B    | 0 --   | 5 D    | 3 B    |
| 8 B    | 5 D    | 0 --   | 2 C    |
| 6 B    | 3 D    | 2 D    | 0 --   |

(4) Fourth and subsequent exchanges; converged

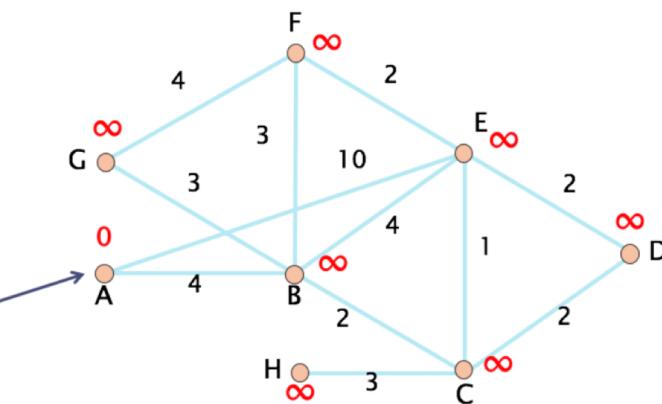
| To   | A    | B    | C    | D    |
|------|------|------|------|------|
| says | says | says | says | says |
| A    | 0    | 3    | 8    | 6    |
| B    | 3    | 0    | 5    | 3    |
| C    | 8    | 5    | 0    | 2    |
| D    | 6    | 3    | 2    | 0    |

| A      | B      | C      | D      |
|--------|--------|--------|--------|
| learns | learns | learns | learns |
| Cost   | Cost   | Cost   | Cost   |
| Next   | Next   | Next   | Next   |
| 0 --   | 3 A    | 8 D    | 6 B    |
| 3 B    | 0 --   | 5 D    | 3 B    |
| 8 B    | 5 D    | 0 --   | 2 C    |
| 6 B    | 3 D    | 2 D    | 0 --   |

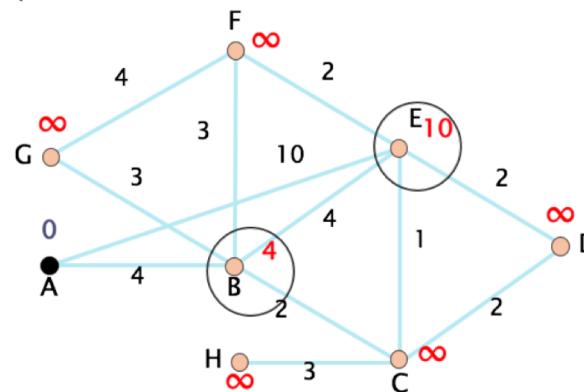
# Link State Example



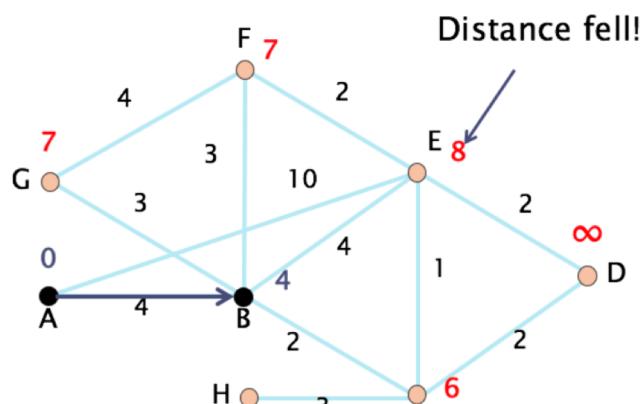
(1) Initialisation



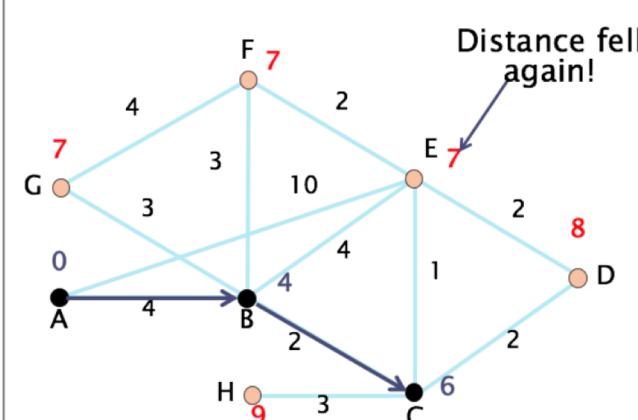
(2) Relax around A



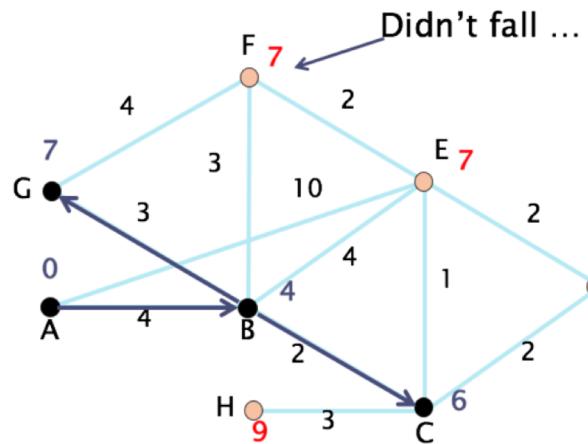
(3) Relax around B



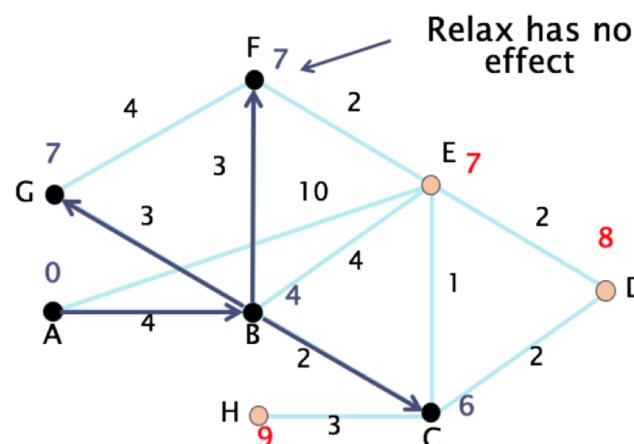
(4) Relax around C



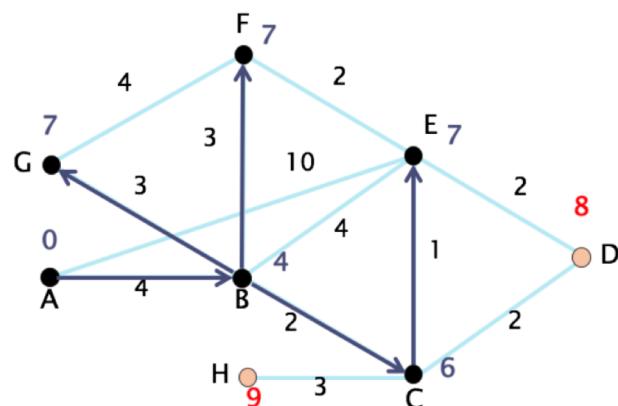
(5) Relax around G (say)



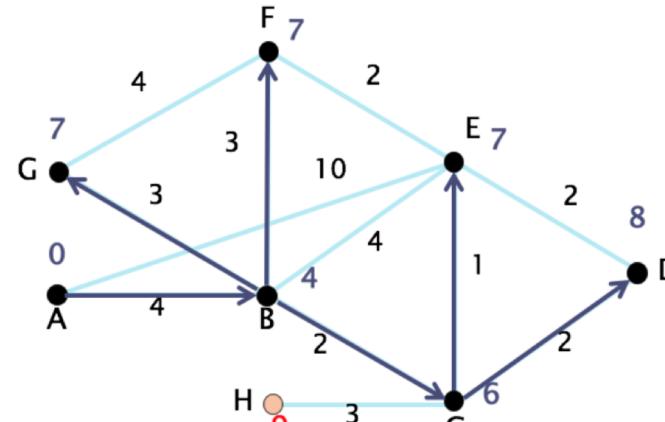
(6) Relax around F (say)



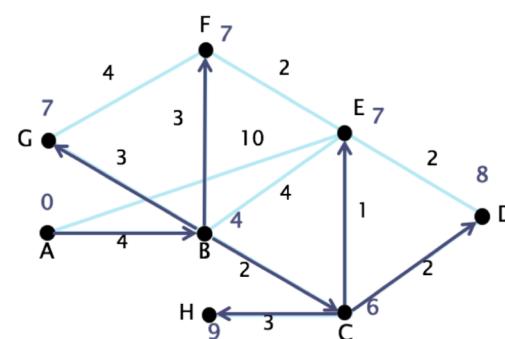
(7) Relax around E



(8) Relax around D



(9) Finally, H... done

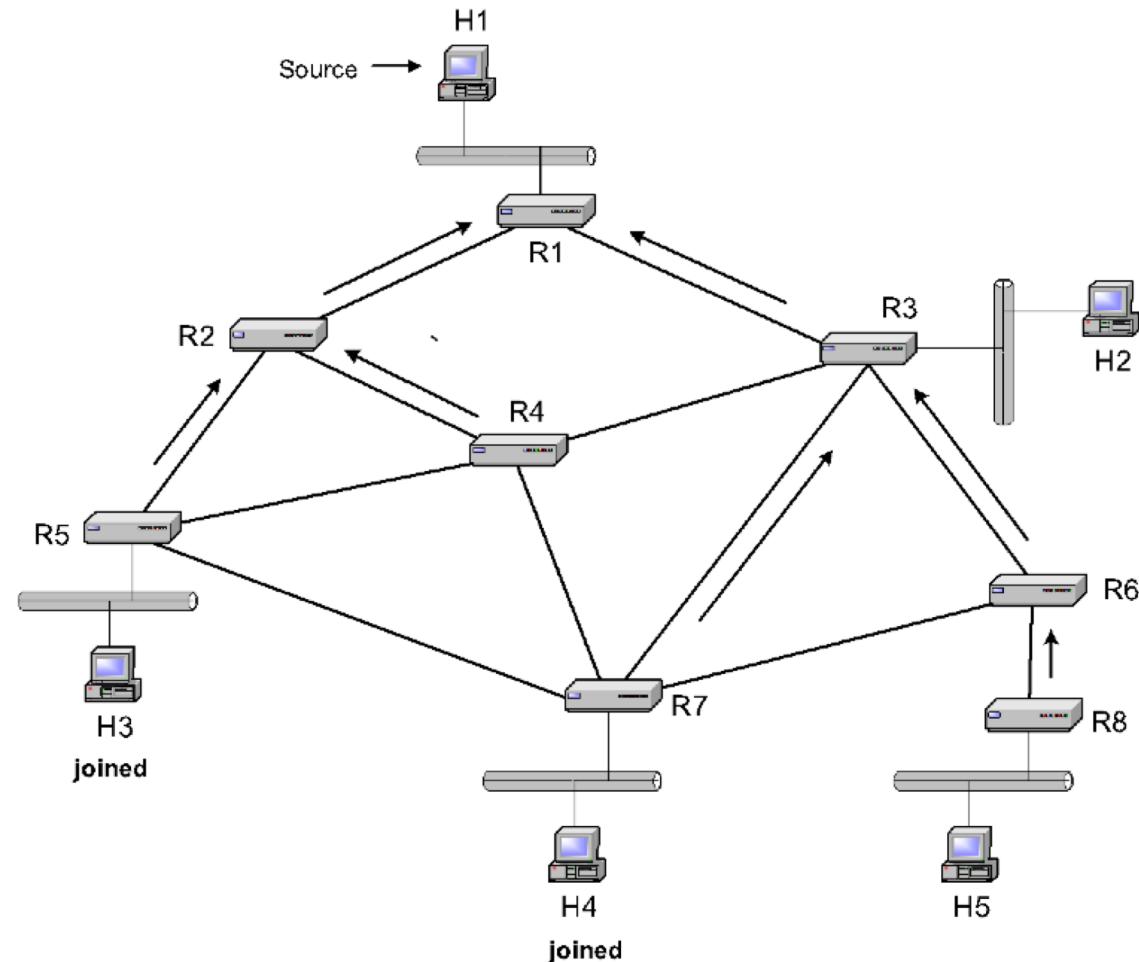


# Topics

- ▶ Multicast Routing Solutions
- ▶ Source Based Trees
- ▶ Core Based Trees

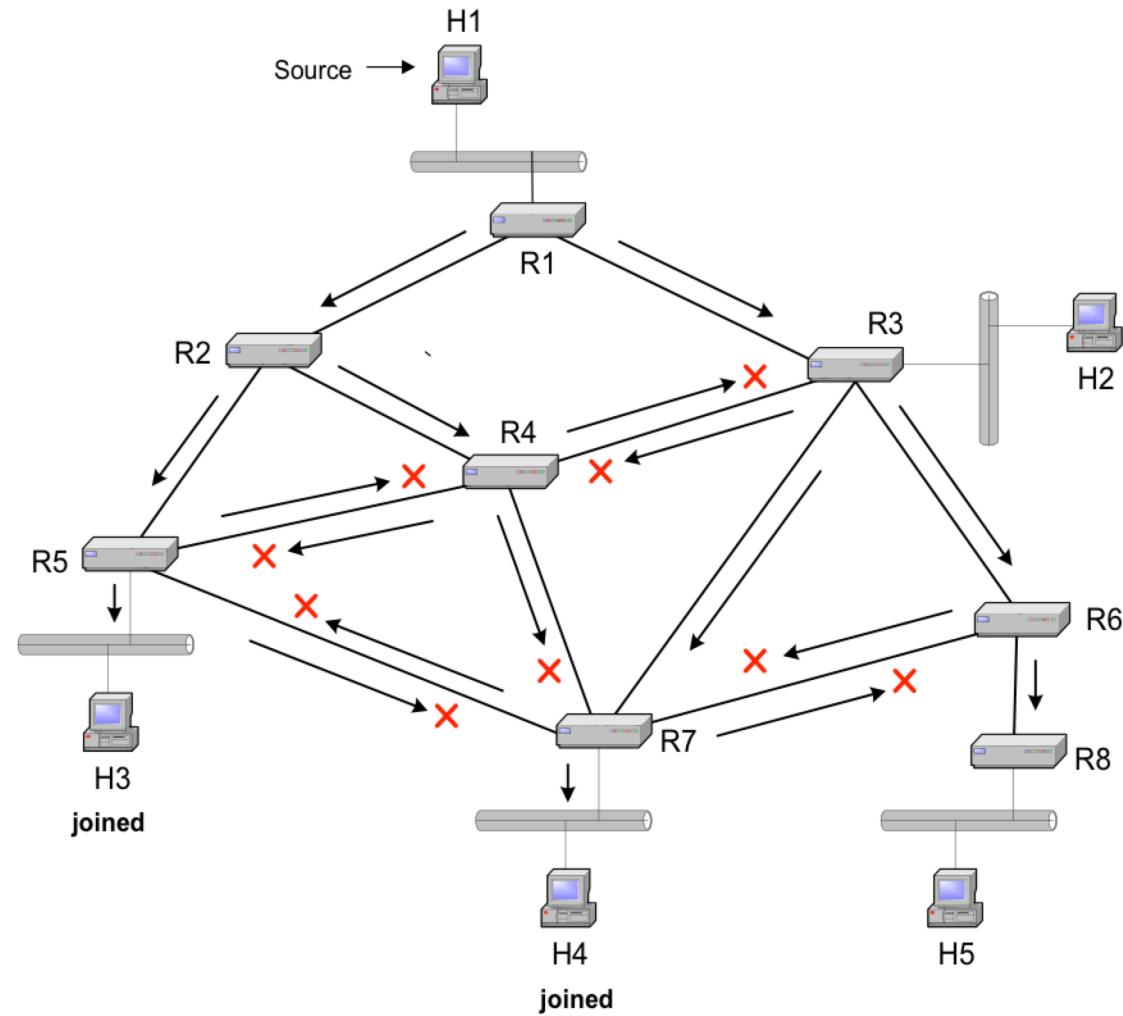
# Source-Based Tree

- ▶ Set routing tables according to RPF
- ▶ Flood-and-Prune



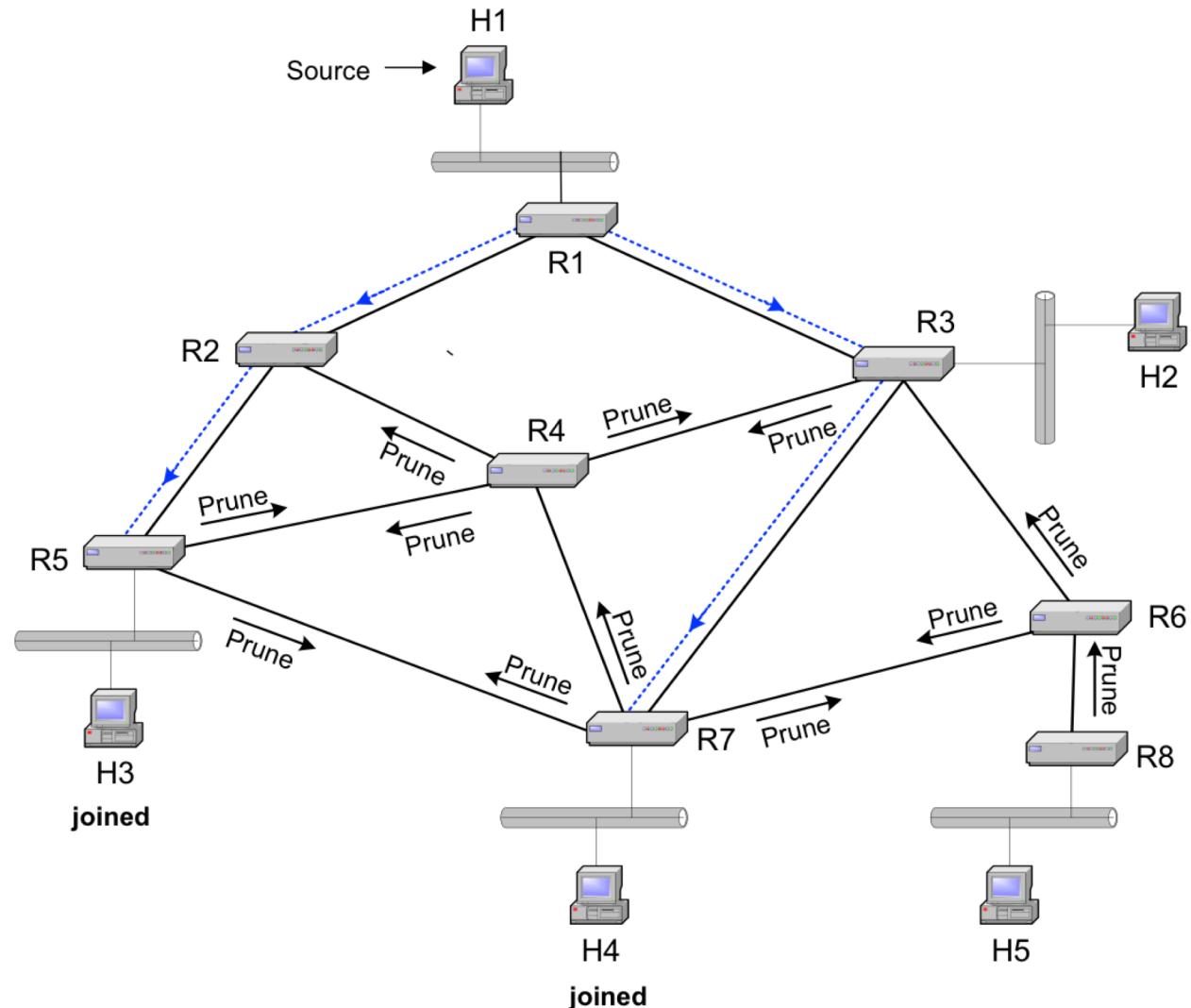
# Flooding

- ▶ Forward packets on all non-RPF interfaces
- ▶ Receiver drops packets not received on RPF interface



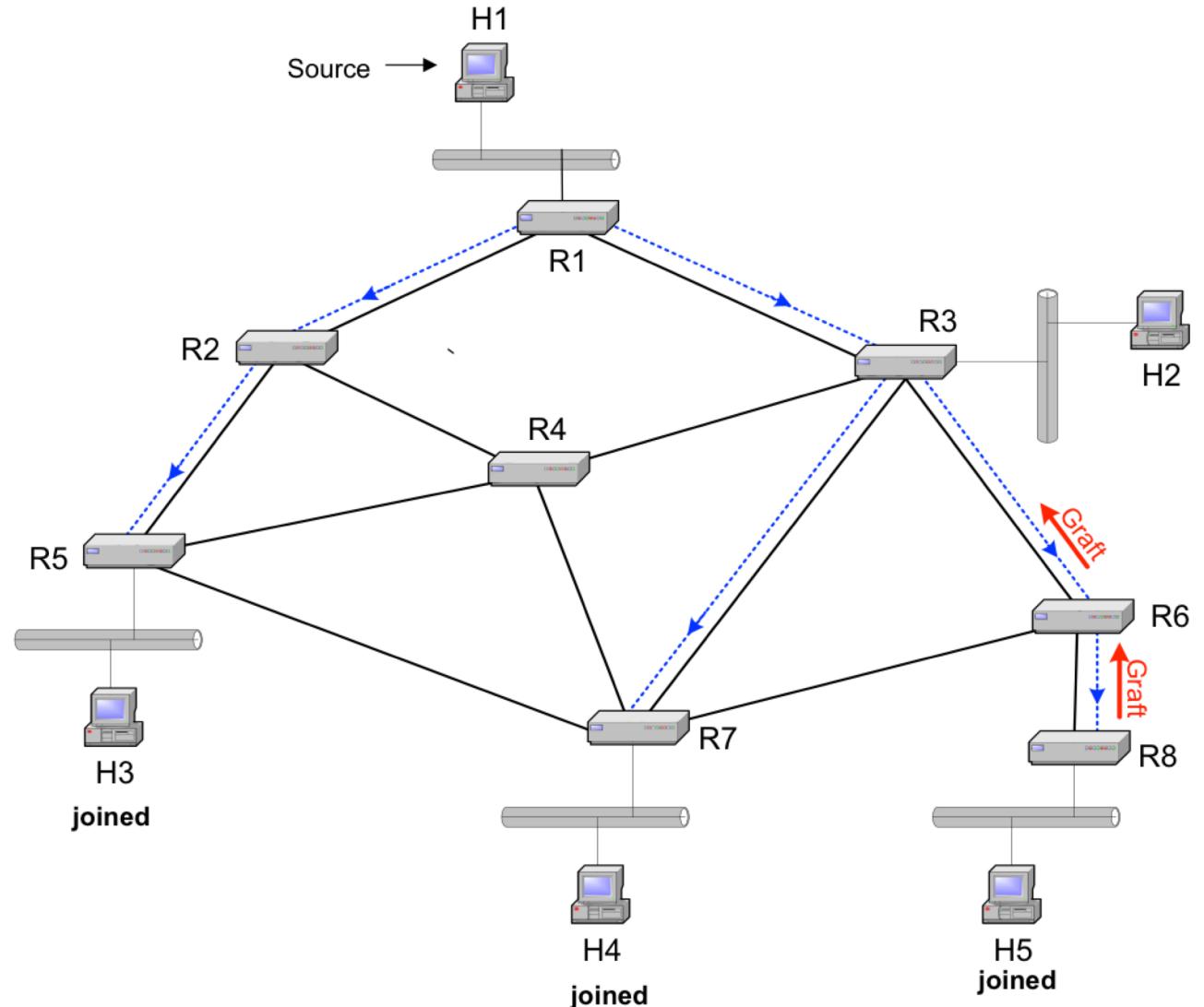
# Pruning

- Prune message temporarily disables a routing table entry
- **Effect:** Removes a link from the multicast tree  
No multicast messages are sent on a pruned link
- Prune message is sent in response to a multicast packet
- *Question: Why is routing table only temporarily disabled?*



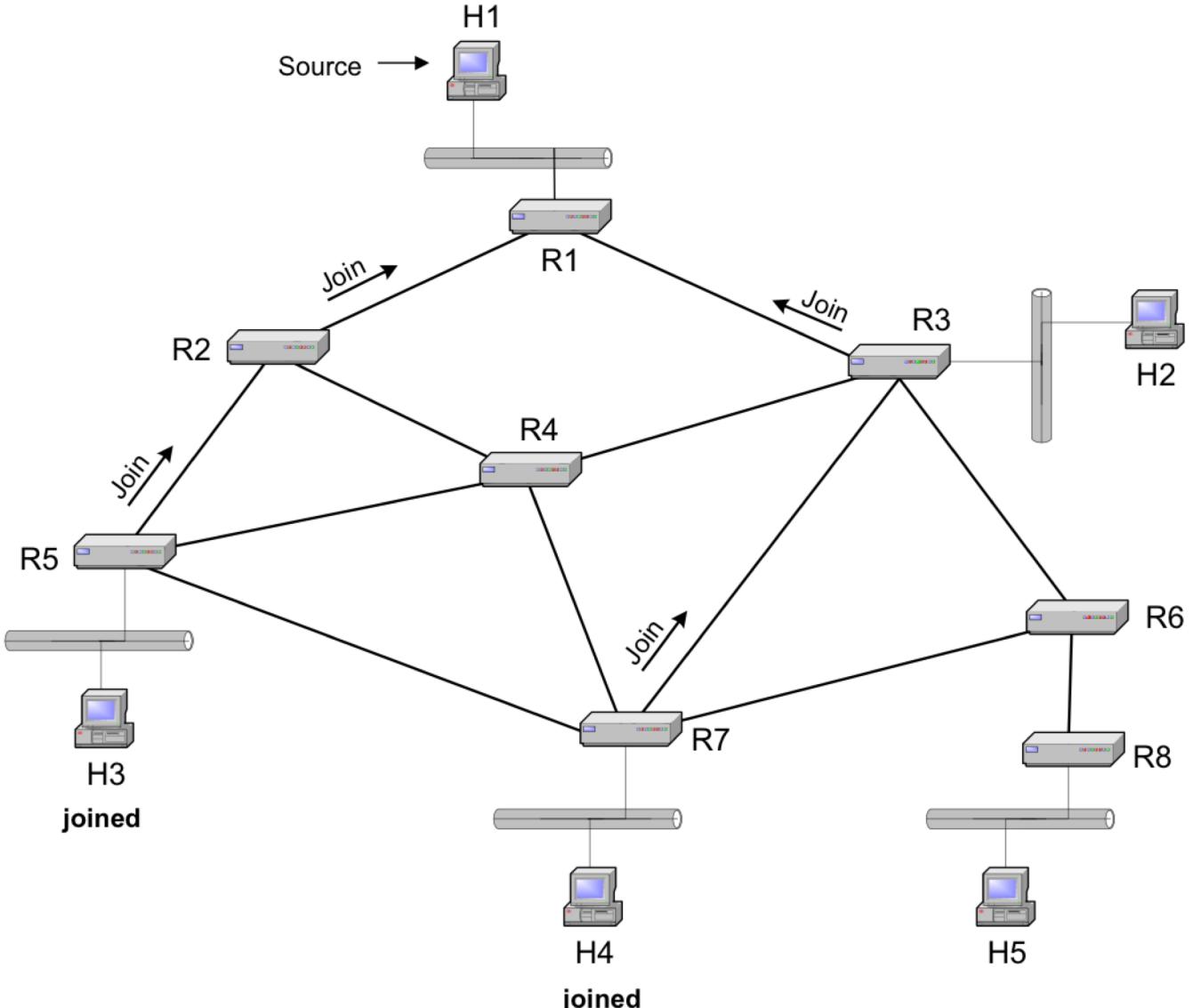
# Grafting

- When a receiver joins, one needs to re-activate a pruned routing table entry
- Sending a *Graft* message disables prune, and re-activates routing table entry.



## Alternative: Explicit Join

- ▶ This only works if the receiver knows the source
- ▶ Receiver sends a *Join* message to RPF neighbor
- ▶ Join message creates  $(S,G)$  routing table entry
- ▶ *Join* message is passed on

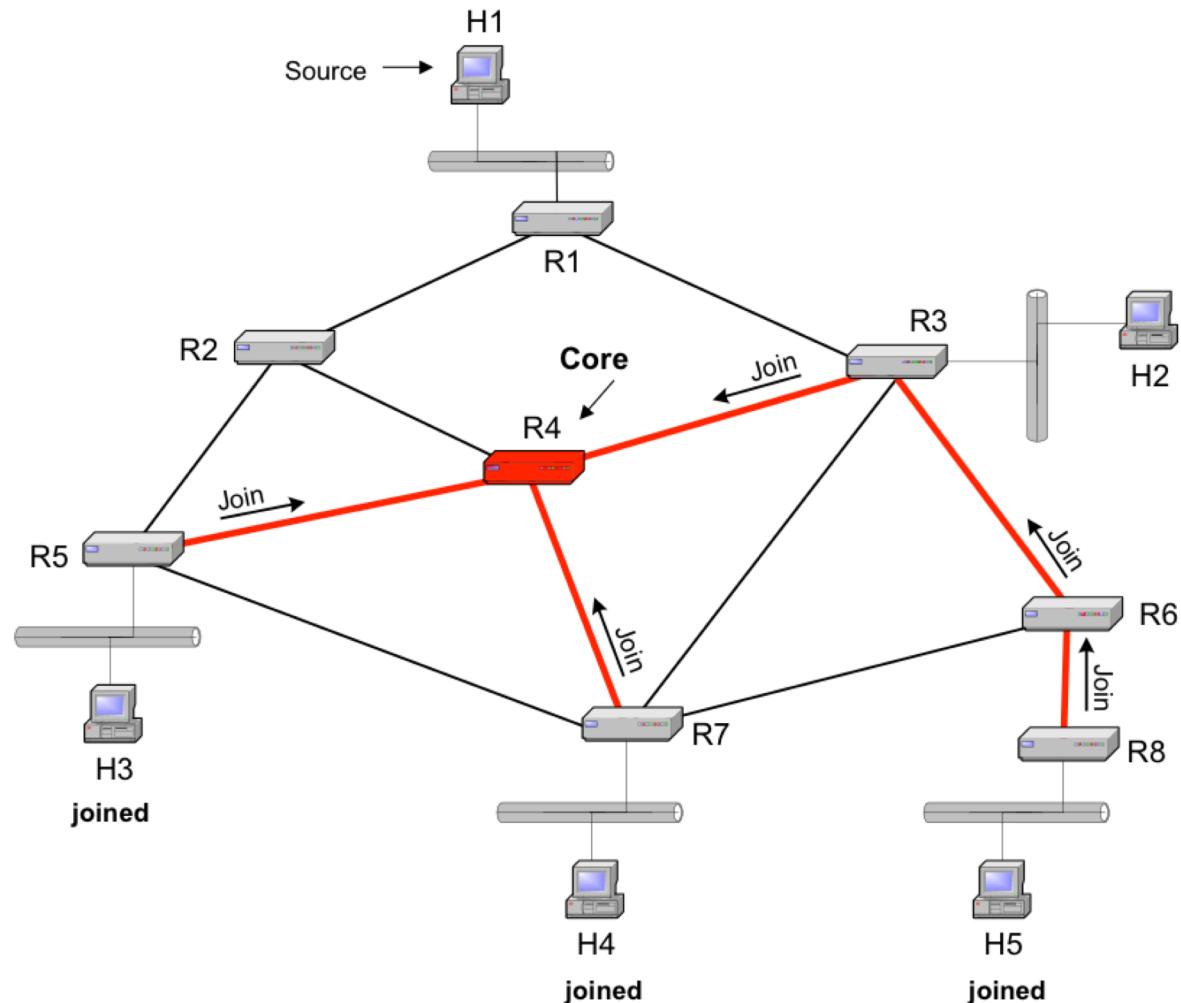


# Topics

- ▶ Multicast Routing Solutions
- ▶ Source Based Trees
- ▶ Core Based Trees

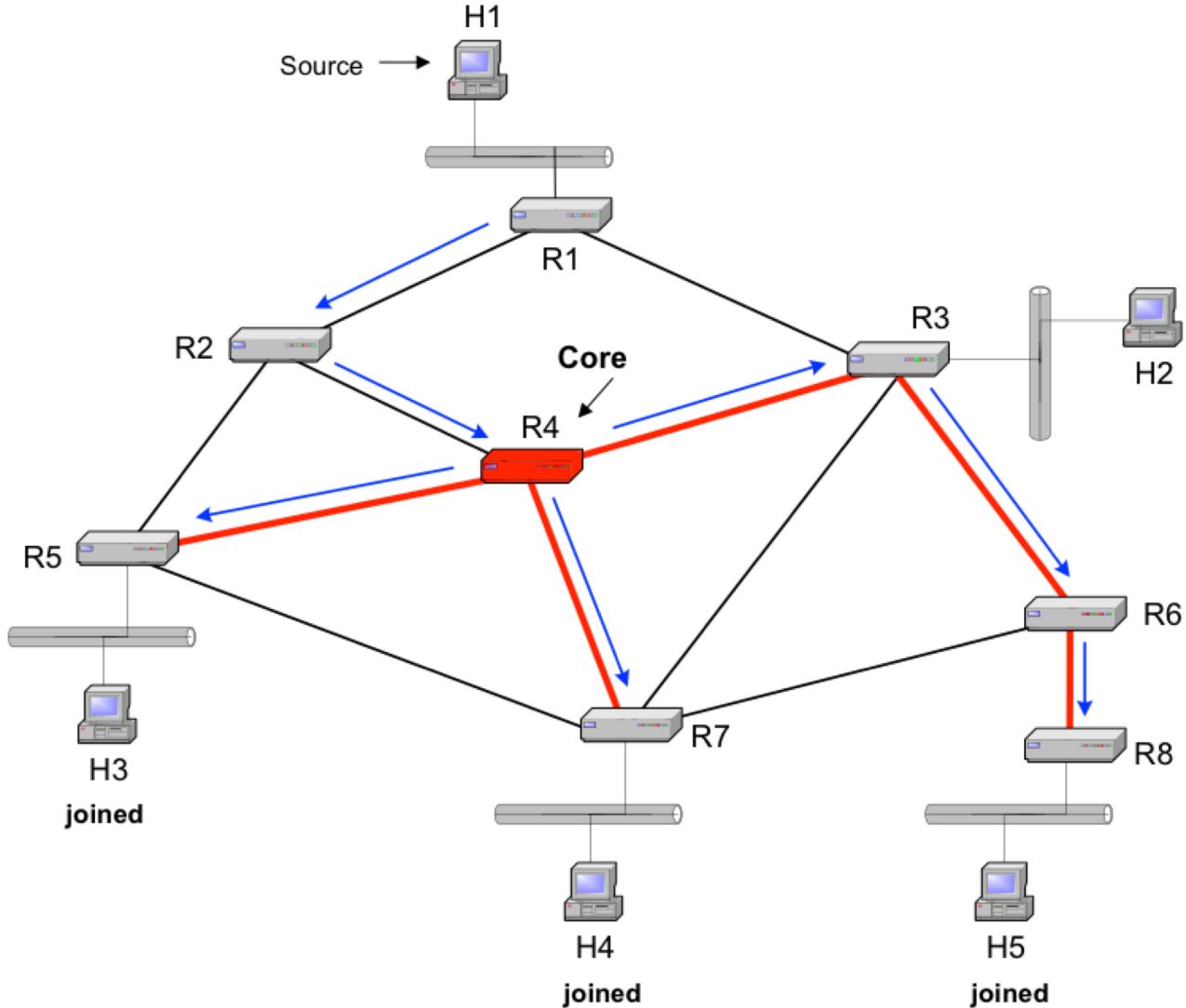
# Core-Based Tree

- ▶ Core is called rendezvous-point (**RP**)
- ▶ Receivers know RP (statically configured or dynamically elected)
- ▶ Receiver sends a *Join* message to RPF neighbour with respect to core
- ▶ *Join* message creates  $(*, G)$  routing table entry



# Core-Based Tree

- ▶ Source sends data to the core
- ▶ Core forwards data according to routing table entry



# Questions?

- ▶ **Email:** x.che@herts.ac.uk
- ▶ **Office:** LB218
- ▶ **Tel:** 01707 286206