

Precise static modeling of Ethereum "memory"

Ethereum: a programmable blockchain

• Decentralized register of transactions that also serves as the execution environment for smart contracts:

Programs stored on a blockchain that run when predetermined conditions are met

Written mainly in a Turing-complete programming language - **Solidity**The code is first compiled to bytecode and then is deployed to the blockchain.

Smart contracts are mostly only available as byte code and not as source code!

A gas cost is paid for performing transactions.

Proportional to the complexity of the computation

Ethereum: a programmable blockchain

An Ethereum user (sender) can submit a transaction proposal to the Ethereum network, containing:

- (i) sender and receiver fields
- (ii) a message
- (iii) value
- (iv) a gas budget

i	sender: 0xBEEFBABE; receiver: 0xC0CAC01A
ii	deliverCans(amount: 10, to: 0xCAFED00D)
iii	2 ETH (~ \$400)
iv	10000 gas units @ 2 * 10-7ETH each

EVM - Ethereum Virtual Machine

- Abstract machine description on which smart contracts are executed by EVM emulators.
- Computes the state of the Ethereum network after a new block is added to the chain.
- Stack based
- The EVM memory model distinguishes **storage** and **memory**:

Storage: a persistent data store, kept on the blockchain's state

Memory: used as a temporary store for all sorts of data (zero-set when any transaction is started)

Smart contract example

```
Variable stored
contract Example{
                       on persistent
 string onStorage;
                          storage
 function setlt(string memory newStr) public {
   onStorage = newStr;
 function getHash() public view returns (bytes32){
   return keccak256(onStorage);
         Hashes memory
             contents
```

"Memory" hides implicit computation!

```
mapping ( string => string ) mTokens ; ...
function getToken ( string pDocumentHash ) view public returns ( string )
{ return mTokens [ pDocumentHash ]; }
```

```
var var0 = 0x053b; var0 = func_06C6(); var var1 = 0x02;
var temp0 = arg0; var var2 = temp0;
var var3 = memorv[0x40:0x60]: var var4 = var3:
var var5 = var2 + 0x20: var var6 = memory[var2:var2 + 0x20]:
var var7 = var6: var var8 = var4: var var9 = var5:
if (var7 < 0x20) (
label_0573:
  var temp1 = 0x0100 ** (0x20 - var7) - 0x01: var temp2 = var8:
  memory[temp2:temp2 + 0x20] = (memory[var9:var9 + 0x20] & ~temp1) | (memory[temp2:temp2 + 0x20] & temp1);
  var temp3 = var6 + var4;
  memory[temp3:temp3 + 0x20] = var1:
  var temp4 = memory[0x40:0x60];
  var temp5 = keccak256(memory[temp4:temp4+(temp3+0x20)-temp4]);
  var temp6 = storage[temp5];
  var temp7 = (!(temp6 & 0x01) * 0x0100 - 0x01 & temp6) / 0x02:
  var temp8 = memory[0x40:0x60];
  memory[0x40:0x60] = temp8 + (temp7+0x1f) / 0x20 + 0x20 + 0x20:
  var1 = temp8: var2 = temp5: var3 = temp7:
  memory[var1:var1 + 0x20] = var3:
  var4 = var1 + 0x20: var5 = var2:
  var temp9 = storage[var5];
  var6 = (!(temp9 & 0x01) * 0x0100 - 0x01 & temp9) / 0x02:
  if (!var6) {
  label_063A:
   return var1;
  } else if (0x1f < var6) {
   var temp10=var4: var temp11 = temp10 + var6: var4=temp11;
   memory[0x00:0x20] = var5:
   var temp12 = keccak256(memory[0x00:0x20]);
   memory[temp10:temp10 + 0x20] = storage[temp12]:
   var5 = temp12 + 0x01; var6 = temp10 + 0x20;
   if (var4 <= var6) { goto label_0631; }
  label 061D:
   var temp13 = var5; var temp14 = var6;
   memory[temp14:temp14 + 0x20] = storage[temp13];
   var5 = temp13 + 0x01; var6 = temp14 + 0x20;
   if (var4 > var6) { goto label_061D; }
  label_0631:
   var temp15 = var4; var temp16 = temp15+(var6 - temp15&0x1f);
   var6 = temp15; var4 = temp16;
   goto label 063A:
  } else {
   var temp17 = var4;
   memory[temp17:temp17+0x20] = storage[var5]/0x0100 * 0x0100:
   var4 = temp17 + 0x20: var6 = var6:
   goto label_063A:
} else {
label_0559:
 var temp18 = var9; var temp19 = var8;
 memory[temp19:temp19 + 0x20] = memory[temp18:temp18 + 0x20];
  var8 = temp19 + 0x20: var9 = temp18 + 0x20: var7 = var7-0x20:
  if (var7 < 0x20) { goto label_0573; }
  else { goto label_0559; }
```

function getToken(var arg0) returns (var r0) {

Smart contract example (decompiled)

0x116: PUSH1 0x1 0x116: v116 = 0x10x116: v116 = 0x10x1c5: PUSH10x0 0x1c5: v1c5 = 0x00x1c7: MSTORE v1c5(0x0) = v116(0x1)0x1c7: MSTORE 0x1cc: v1cc = SHA3 [v116(0x1)]0x1c8: PUSH10x20 0x1c8: v1c8 = 0x200x1ca: PUSH10x0 0x1ca: v1ca = 0x0 (\ldots) 0x1cc: SHA3 0x1cc: v1cc = SHA3 v1ca(0x0) v1c8(0x20)(...) (\ldots) **EVM** bytecode Gigahorse **Register based representation** memory modeling complete

Precise Memory Modeling

The analysis tracks symbolic values based on the **free-memory pointer**, offset by a constant.

constant memory address (0x40/decimal 64)

Use constant and symbolic indexes to infer high-level properties:

- Arguments passed through memory, to statements that read from it
- Array allocations, reads, writes
- Access to the data returned by external calls

Analysis - Input

• The analysis is expressed as a set of declarative rules, using datalog.

Concrete instructions

(s: S) :[r: V := ADD/SUB/MUL(a: V U C, b: V U C)]

(s: S) :[MSTORE(addr : V, from : V)]

(s: S) :[to : V := MLOAD(addr : V)]

Arithmetic operations

Memory related operations

Generic instructions / Syntactic patterns
STATEMENTUSESMEMORY(s: S, start: V, len: V)

Memory consuming operations

V is a set of program variables C is a set of constants. $C \subseteq N256$ S is a set of statement identifiers N256 is the set of 256-bit unsigned integers M is a set of symbolic values FreePtr is the free-memory pointer

Analysis - Interfacing with other analysis modules

Variable_Value(v : V, c : C)

Constant propagation/folding

Flows(from: V, to: V)

Alias(x : V, y : V)

MatchingMSTORE(ms: S, s: S)

MSTORE statement 'ms' writes for memory

UnchangedFreePointer(s1: S, s2: S)

Analysis - Symbolic Value creation

```
FreePointerBasedValue(val, mload, 0):-
    mload: [to := MLOAD(FreePtr)], val = mload ++ "0x0"

Variable_Value+(to, val),
FreePointerBasedValue(val, mload, res):-
[to := ADD(numVar, freePtrBasedVar)],
Variable_Value(numVar, numVal1),
Variable_Value+(freePtrBasedVar, freePtrBasedVal),
FreePointerBasedValue(freePtrBasedVal, mload, numVal2),
res = numVal1 + numVal2,
val = mload ++ numVal1 + numVal2.
```

Variable_Value+(to, val),

Symbolic value creation on MLOADs on the free-memory pointer

Constant folding of the numeric part of the symbolic values

Analysis - Sample of final outputs

```
StatementUsesMemoryAtIndex(stmt, relativeIndex, actual):-
StatementUsesMemory(stmt, startVar, lenVar),
Variable_Value+(startVar, startVal),
Variable_Value(lenVar, lenVal),
MatchingMSTORE(mstore, stmt),
mstore: [MSTORE(indexVar, actual)],
Variable_Value+(indexVar, indexVal),
FreePointerDiff(indexVal, startVal, relativeIndex),
lenVal > relativeIndex >= 0
```

StatementUsesMemory_ActualMemoryArg(stmt, i, actual) :- order(StatementUsesMemoryAtIndex(stmt, _, actual)) = i

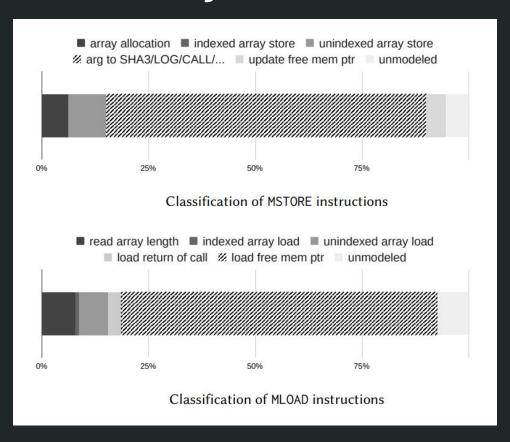
Matches MSTOREs that write for a memory consuming statement to the relative offset they write to

Client Analyses

The next 2 analyses require **memory modeling**, as they uses operations that are done through non-trivial uses of memory.

- Passing the arguments to an external call
- Getting its return value back

Client Analyses - Evaluation



94,59%

92,72%

Client Analysis - Taint Analysis

Tainted ERC20 Token transfer

• The transfer(address **recipient**, uint256 **amount**) function(part of the ERC20 interface) when called transfers **amount** tokens. The tainted ERC20 token transfer is **vulnerable!**

```
allows anyone to take
contract Victim {
                                 over the contract's
 address owner;
                               ownership and bypass
 function init () public {
                                      the guard
    owner = msg . sender ; ... }
 function withdrawTokens (address _tokenContract) public returns (bool) {
    require ( msg . sender == owner );
    Token token = Token ( _tokenContract );
    uint256 amount = token . balanceOf ( address (this));
    return token . transfer ( owner , amount );}
```

Client Analysis - Taint Analysis

Introduction of new Guard Conditions

• Are introduced new guard conditions that use an external contract as a source of authority, calling it to approve or reject an attempt to call a sensitive operation.

```
contract Guarded {
  address auth = ...;
  function sensitiveOperation () public {
    require ( msg . sender == auth . owner () );
  function otherSensitiveOperation () public {
    require ( auth . isOwner ( msg . sender ));
```

Client Analysis - Taint Analysis

Manual inspection for the tainted ERC20 token transfer vulnerability.

MD5	LOC	TP/FP	Comment
7eacf	1441	0 / 1	requires sender to destroy token
c09fb	146	0 / 1	can only be used to send to untainted investors
cee49	244	1/0	composite
486df	490	0 / 1	unrecognized guard
92a49	511	0 / 1	unrecognized guard
17c8f	64	1/0	by design, airdrop
b1092	201	1/0	by design, airdrop
a4f0e	52	1/0	by design, airdrop
8bfbd	479	0 / 1	unrecognized guard
af93f	264	1/0	composite
f02a4	204	1/0	by design, airdrop
b30d4	1069	1/0	composite
78fcb	64	1/0	by design, airdrop
82815	652	1/0	composite
6ecdb	864	0 / 1	complex logic, tokens sent will be compensated
394d2	394	0 / 1	unrecognized guard
e3129	429	0 / 1	unrecognized guard
4f9ac	56	0 / 1	requires caller to transfer tokens first
29976	224	1/0	composite
95b19	237	1/0	composite
503f3	698	0/3	unrecognized guard
0f8ab	268	1/0	composite
040e6	599	0 / 1	unrecognized guard
fb0ae	227	1/0	composite
33350	74	1/0	unguarded transfer
Total:	9951	14 / 13	

Client Analysis - Precise Gas Consumption

• EIP-1884, part of the Istanbul hard fork, repriced many resource intensive operations which led to the increasing cost of SLOADs (200 to 800 gas), causing **fallback functions** (2300 gas limit) to fail.

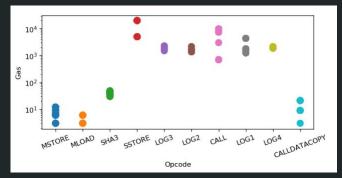
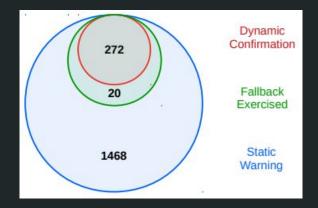


Image 1: gas costs per opcode vary by orders of magnitude

unnamed external function without any input or output parameters

• Smart contracts can't have code whose execution cost can be easily manipulated by user input.(instructions that use memory have a variable gas price!)

Client Analysis - Precise Gas Consumption



• 93% precision of the gas analysis, meant to find fallback functions that would fail after EIP-1884 repricing

Client Analysis - Repeated Calls

Pattern: Two identical external calls, with one preceding the other.

same callee contract, target method and arguments

Why should it be avoided?

- Same external call could return different values
- Call cost gas

Evaluation:

Precise Static modeling analysis	Securify
85.96%	16.09%



Client Analysis - Repeated Calls (example)

```
interface Untrusted(
 function getBenificiary() external returns (address payable);
contract Victim{
 function isFriend (address addr) private returns (bool){(...)}
 function givEth(Untrusted untrustedAddress) public
                                                                          Might be
                                                                          exploited
    if (isFriend (untrustedAddress.getBenificiary())){
                                                                         returning 2
      untrustedAddress.getBenificiary().transfer(1000 ether);
                                                                       different values
```