

Given a positive integer  $n$ , the *complexity* of  $n$  is the minimum number of ones that can be used to represent  $n$ , using only the operations of addition and multiplication, as well as parenthesisation.

For example, we have the following representations:

$$\begin{aligned} 6 &= (1 + 1 + 1) \times (1 + 1). \\ 8 &= (1 + 1) \times (1 + 1) \times (1 + 1). \\ 9 &= (1 + 1 + 1) \times (1 + 1 + 1). \\ 12 &= (1 + 1 + 1 + 1) \times (1 + 1 + 1). \\ 19 &= (1 + 1 + 1) \times (1 + 1 + 1) \times (1 + 1) + 1. \end{aligned}$$

The first twenty entries are given for you.

$n$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
	1	2	3	4	5	5	6	6	6	7	8	7	8	8	8	8	9	8	9	9

- Show that every positive integer can be represented by a string of ones, along with addition, multiplication, and parenthesisation operations; that is, the complexity of  $n$  is always finite.
- Given a positive integer  $n$ , describe an  $O(n^2)$  algorithm to compute the minimum number of one's (1's) using only the operations of addition and multiplication, as well as parentheses, whose expression equals  $n$ .

**Note.** This is also known as the Mahler-Popken complexity. Here is the [OEIS](#) entry.

### Solution.

- Any integer can be represented as a sum of ones, so we have an upper bound as simply  $n = \sum_{i=1}^n 1$ .
- Begin by initialising a Dynamic Programming table of size  $n + 1$ ,  $DP$ , with  $DP[i] = i$  for each element. Now, we must explore all pairs  $(a, b)$  such that  $a \times b = n$  or  $a + b = n$  to ensure that we cover all different ways to arrive at that number.
  - For each value  $a$  from 1 up until  $\lceil \sqrt{n} \rceil$ , we check to see if  $i = 0 \bmod a$ . If so, we take  $b = i \div a$  and update  $DP[i] = \min\{DP[i], DP[a] + DP[b]\}$ .
  - For each value  $a$  from 1 up until  $\lceil \frac{n}{2} \rceil$ , we take  $b = i - a$  and again update  $DP[i] = \min\{DP[i], DP[a] + DP[b]\}$ .

Finally, we return the result  $DP[n]$  as the ideal value. Our multiplication and addition checks take  $O\left(\sqrt{n} + \frac{n}{2}\right) = O(n)$ , and since we do this for all numbers up until  $n$ , we have an overall time complexity of  $O(n^2)$ .