Code Documentation

1. NFA to DFA conversion

I. Epsilon NFA to DFA

```
Step 1 – Consider M={Q, \Sigma, \delta,q0,F) is NFA with \epsilon. We have to convert this NFA with \epsilon to equivalent DFA denoted by M0=(Q0,\Sigma,\delta0,q0,F0) Then obtain, \epsilon\text{-closure}(q0)=\{p1,p2,p3,.....pn\} then [p1,p2,p2,....pn] becomes a start state of DFA now[p1,p2,p3,....pn] \in Q0
```

```
epsilon_closures = {}
   # Get all epsilon closure of all states
   for state in self.transitions:
        epsilon closures[state] = self.epsilon closure(state)
    print("Epsilon Closures:",epsilon closures)
   states = []
8 states closures = []
9 ts = \{\}
10 newAlphabet = self.alphabet
   newAlphabet.remove('eps')
11
   i=0
12
# Add the epsilon closures of the first state
states closures.append(list(epsilon closures.values())[0])
15 j = len(states closures)
   newAcceptingStates = []
```

```
Step 2 – We will obtain \delta transition on [p1,p2,p3,...pn] for each input. \delta \ O([p1,p2,p3,...pn],a) = \epsilon \text{-closure}(\delta(p1,a) \ U \ \delta(p2,a2) \ U..... \ \delta(pn,a)) = U \ (i=1 \ to \ n) \ \epsilon \text{-closure} \ d(pi,a) Where a is input \in \Sigma
```

```
• • •
   while j>0:
        closures = states_closures[i]
        1 = list(closures)
        if len(set(self.acceptingStates).intersection(1)):
           newAcceptingStates.append('q'+str(i))
        ts['q'+str(i)] = {}
        states.append('q'+str(i))
        print('q'+str(i))
        print(closures)
        for alpha in newAlphabet:
           nextState = set()
            # Get all next states based on the input alphabet
            for st in 1:
               nextState |= set(self.transitions[st][alpha])
           ep = set()
            for n in nextState:
                if len(n)>0:
                   ep |= epsilon_closures[n]
           print(alpha,ep)
            if not(ep in states_closures) and len(ep)>0:
                states_closures.append(ep)
                j+=1
            if len(ep)>0:
                ts['q'+str(i)][alpha] = 'q'+str(states_closures.index(ep))
                ts['q'+str(i)][alpha] = []
        print(ts)
        j-=1
        i+=1
```

Step 3 – The state obtained $[p1,p2,p3,...pn] \in Q0$.

The states containing final state in pi is a final state in DFA

II. NFA to DFA without epsilon

Steps for converting NFA to DFA:

Step 1: Initially $Q' = \phi$

Step 2: Add q0 of NFA to Q'. Then find the transitions from this start state.

```
1
2 states = [self.states[0]]
3 ts = {}
4 # Add first state
5 ts[self.states[0]] = getTransitionsOf(self.states[0], self.transitions)
6 i = len(states)
```

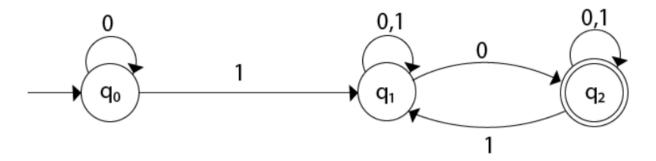
Step 3: In Q', find the possible set of states for each input symbol. If this set of states is not in Q', then add it to Q'.

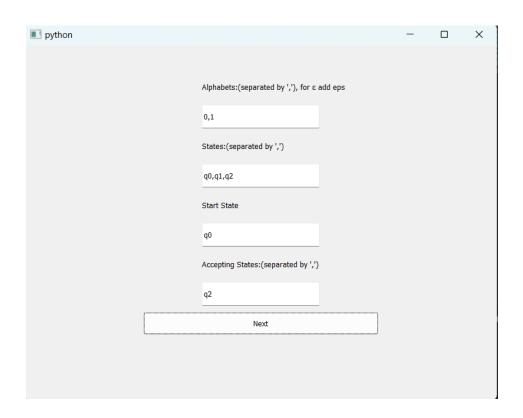
```
1
2 # Loop whenever a state is added
3 while i>=0:
4  # Loop to get (To states) from the transitions dictionary
5  for a, toStates in ts[states[i-len(states)]].items():
6  # If the same input alphabet goes to more than one state then name this state with '-' separated between the names of state
7  temp = '-'.join(toStates)
8  # If this state is not already added then add it
9  if not(temp in states) and temp != '':
10  newTransitions = getTransitionsof(temp, self.transitions)
11  states.append(temp)
12  ts[temp] = newTransitions
13  i+=1
14 i-=1
```

Step 4: In DFA, the final state will be all the states which contain F(final states of NFA)

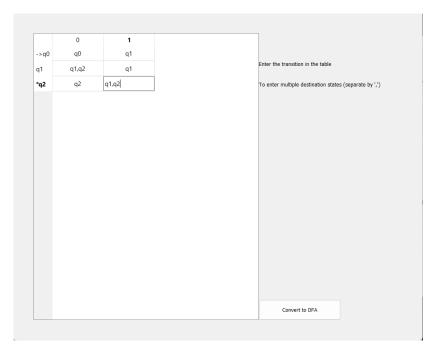
Screenshots:

For example:

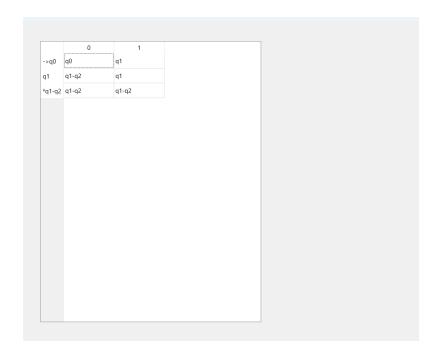




Input transition table:



Output:



2. CFG to PDA conversion

Step 1: The code defines a PDA (Pushdown Automaton) class, which simulates a context-free grammar (CFG) using a pushdown automaton.

Initialize the PDA instance: The __init__ method initializes a new PDA instance with the following parameters:

variables: A set of variables (non-terminal symbols) in the CFG.

terminals: A set of terminal symbols in the CFG.

start variable: The start variable for the CFG.

production rules: A dictionary containing the production rules for each variable.

```
class PDA:
    def __init__(self, variables, terminals, start_variable, production_rules):
        self.states = ['q_start', 'q_push_start_var', 'q_loop', 'q_final']
        self.variables = variables
        self.terminals = terminals
        self.start_variable = start_variable
        self.production_rules = production_rules
        self.transitions = self.create_transitions()
```

Step 2: The create_transitions method generates the transition rules for the PDA, based on the given CFG. The transitions are created for the following scenarios:

- Transition from q_start to q_push_start_var, pushing the stack symbol \$.
- Transition from q_push_start_var to q_loop, pushing the start variable.
- Transitions for production rules, in the q_loop state. If a variable A produces a string w, the PDA pushes the symbols of w in reverse order onto the stack.
- Transitions for terminal symbols, in the q_loop state, reading the terminal symbol from the input and popping it from the stack.
- Transition from q loop to q final, popping the stack symbol \$.

```
def create_transitions(self):
  transitions = {}
  # Transition from q_start to q_push_start_var, pushing the stack symbol $
  transitions[('q_start', '', '')] = [('q_push_start_var', '$')]
  # Transition from q_push_start_var to q_loop, pushing the start variable
  transitions[('q_push_start_var', '', '')] = [('q_loop', self.start_variable)]
  # Transitions for production rules, in q_loop state
  state_counter = 1
   # Loop through the production rules for each variable
   for variable in self.variables:
       for prod_rule in self.production_rules[variable]:
          previous_state = 'q_loop'
           for symbol in enumerate(prod_rule[::-1]):
              new_state = f'q_alt_{state_counter}
              state_counter += 1
              transitions[(previous_state, '', variable)] = [(new_state, symbol)]
              previous_state = new_state
              variable = '' # Clear the variable after the first iteration
          # Transition back to q_loop after pushing all symbols
          transitions[(previous_state, '', '')] = [('q_loop', '')]
```

```
# Transitions for terminal symbols, in q_loop state
for terminal in self.terminals:
    # For each terminal symbol a, with input a and pop a, push ɛ
    transitions[('q_loop', terminal, terminal)] = [('q_loop', '')]

# Transition from q_loop to q_final, popping the stack symbol $
transitions[('q_loop', '', '$')] = [('q_final', '')]

return transitions
```

Step 3: The main function processes the given CFG by creating a PDA instance and printing the generated transition rules.

```
# Main function for testing the PDA class
def main():
    # Define variables, terminals, start variable, and production rules
variables = {'S','B'}
terminals = {'a', 'b','c'}
start_variable = 'S'
production_rules = {
    'S': ['aBc', 'ab'],
    'B': ['SB', '']

# Create a PDA object with the given parameters
pda = PDA(variables, terminals, start_variable, production_rules)

# Print the transitions in a human-readable format
for key, values in pda.transitions.items():
    for value in values:
        print(f"From {key[0]} with input '{key[1]}' and pop '{key[2]}': to {value[0]} and push '{value[1]}'")

print(f"")
print(f"m)
print(f"another format\n")

# Print the transitions in another format
for key, values i (variable) values: Any):
    for value in values:
        print(f"{key[0]}, '{key[1]}', '{key[2]}': --> {value[0]}, '{value[1]}'")
```

past cfg

```
variables = {'S','B'}
terminals = {'a', 'b','c'}
start_variable = 'S'
production_rules = {
    'S': ['aBc', 'ab'],
    'B': ['SB', '']
}

# Create a PDA object with the given parameters
pda = PDA(variables, terminals, start_variable, production_rules)
```

```
PS C:\Users\ziada> & C:\Users\ziada\AppData\Local\Programs\Python\Python311\python.exe "c:\Users\ziada\OneDrive\Prom q_start with input '' and pop '': to q_push_start_var and push '$'
From q_push_start_var with input '' and pop '': to q_loop and push 'S'
From q_loop with input '' and pop 'S: to q_alt_1 and push '(0, 'c')'
From q_alt_1 with input '' and pop '': to q_alt_2 and push '(1, 'B')'
From q_alt_2 with input '' and pop '': to q_alt_3 and push '(2, 'a')'
From q_alt_3 with input '' and pop '': to q_loop and push ''
From q_loop with input '' and pop '': to q_loop and push ''
From q_alt_4 with input '' and pop '': to q_loop and push ''
From q_alt_5 with input '' and pop '': to q_loop and push ''
From q_loop with input '' and pop 'B: to q_alt_6 and push '(0, 'B')'
From q_alt_6 with input '' and pop '': to q_loop and push ''
From q_alt_6 with input '' and pop '': to q_loop and push ''
From q_loop with input '' and pop 'c: to q_loop and push ''
From q_loop with input 'c' and pop 'c: to q_loop and push ''
From q_loop with input 'c' and pop 'b: to q_loop and push ''
From q_loop with input 'a' and pop 's': to q_loop and push ''
From q_loop with input 'a' and pop 's': to q_loop and push ''
From q_loop with input 'a' and pop 's': to q_loop and push ''
From q_loop with input 'a' and pop 's': to q_loop and push ''
From q_loop with input '' and pop 's': to q_loop and push ''
From q_loop with input '' and pop 's': to q_loop and push ''
```

Another Format

```
PROBLEMS
                           DEBUG CONSOLE
                                                TERMINAL
From q_loop with input '' and pop '$': to q_final and push
another format
q_start, '', '': --> q_push_start_var, '$'
q_push_start_var, '', '': --> q_loop, 'S'
q_loop, '', 'S': --> q_alt_1, '(0, 'c')'
q_alt_1, '', '': --> q_alt_2, '(1, 'B')'
q_alt_2, '', '': --> q_alt_3, '(2, 'a')'
q_alt_3, '', '': --> q_loop, ''
q_loop, '', '': --> q_loop,
q_alt_4, '', '': --> q_alt_5, '(1, 'a')'
q_alt_5, '', '': --> q_loop, ''
q_loop, '', 'B': --> q_alt_6, '(0, 'B')'
q_alt_6, '', '': --> q_alt_7, '(1, 'S')'
q_alt_7, '', '': --> q_loop,
q_loop, 'c', 'c': --> q_loop, ''
q_loop, 'b', 'b': --> q_loop, ''
q_loop, 'a', 'a': --> q_loop, ''
q_loop, '', '$': --> q_final, ''
```