Introduction to the Boolean Satisfiability Problem

Spring 2017

CSCE 235H Introduction to Discrete Structures

URL: cse.unl.edu/~cse235h

All questions: Piazza

Satisfiability Study

- 7 weeks
- 25 min lectures in recitation
- ~2 hours of homework per week
- Goals:
 - Exposure to fundamental research in CS
 - Understand how to model problems
 - Learn to use SAT solver, MiniSAT

• Given:

A Boolean formula

Question:

— Is there an assignment of truth values to the Boolean variables such that the formula holds true?

$$a \vee (\neg a \wedge b)$$

$$(a \lor \neg a) \to (b \land \neg b)$$

$$a \vee (\neg a \wedge b)$$

SATISFIABLE a=true, b=true

$$(a \lor \neg a) \to (b \land \neg b)$$

$$a \vee (\neg a \wedge b)$$

SATISFIABLE a=true, b=true

$$(a \lor \neg a) \to (b \land \neg b)$$

UNSATISFIABLE

Left side of implication is a tautology.

Right side of implication is a contradiction.

True cannot imply false.

Applications of SAT

- Scheduling
- Resource allocation
- Hardware/software verification
- Planning
- Cryptography

Conjunctive Normal Form

- Variable
- Literal
- Clause
- Formula

$$a, b, p, q, x_1, x_2$$

$$a, \neg a, q, \neg q, x_1, \neg x_1$$

$$(a \lor \neg b \lor c)$$

$$(a \lor \neg b \lor c)$$

$$\wedge (b \vee c)$$

$$\wedge (\neg a \vee \neg c)$$

- All Boolean formulas can be converted to CNF
- The \rightarrow , \leftrightarrow , \oplus operators can be rewritten in terms of \neg , \vee , \wedge
- \neg , \vee , \wedge can be rearranged using
 - De Morgan's Laws
 - Distributive Laws
 - Double Negative
- May result in exponential size increase of the formula

$$(a \lor \neg a) \to (b \land \neg b) \equiv$$

$$(a \lor \neg a) \to (b \land \neg b) \equiv$$

$$\neg(a \lor \neg a) \lor (b \land \neg b) \equiv$$

$$(a \lor \neg a) \to (b \land \neg b) \equiv$$
$$\neg (a \lor \neg a) \lor (b \land \neg b) \equiv$$
$$(\neg a \land a) \lor (b \land \neg b) \equiv$$

$$(a \lor \neg a) \to (b \land \neg b) \equiv$$

$$\neg (a \lor \neg a) \lor (b \land \neg b) \equiv$$

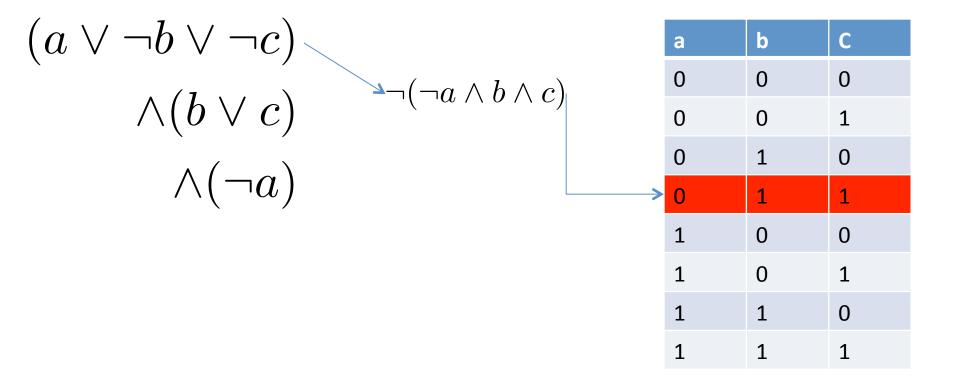
$$(\neg a \land a) \lor (b \land \neg b) \equiv$$

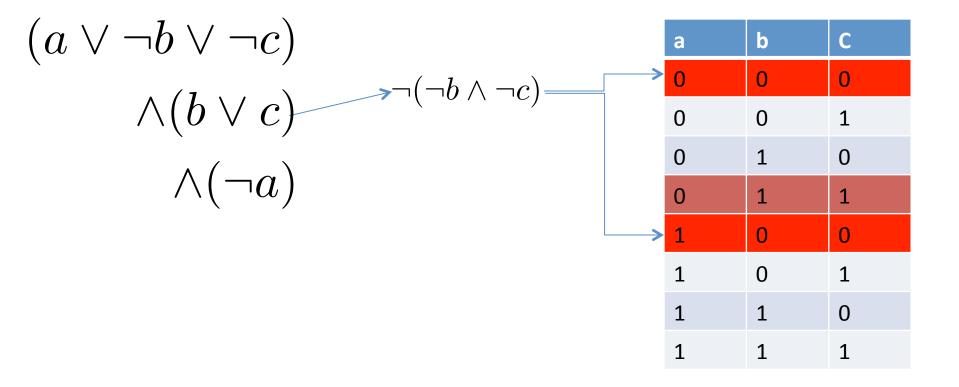
$$(\neg a \lor b) \land (\neg a \lor \neg b) \land (a \lor b) \land (a \lor \neg b)$$

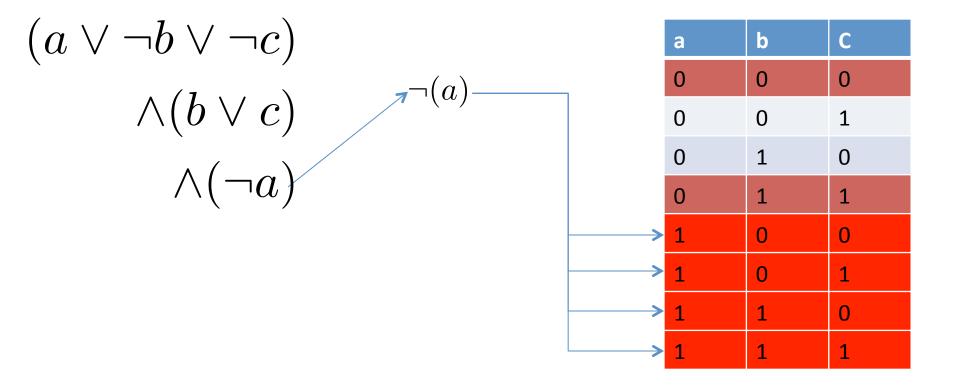
- Every clause must be satisfied by at least one true literal
- Total possible number of solutions increases as number of variables increases
- Clauses constrain the possible solutions
- Smaller clauses are more constraining

$$(a \lor \neg b \lor \neg c)$$
$$\land (b \lor c)$$
$$\land (\neg a)$$

a	b	C
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1







$$(a \lor \neg b \lor \neg c)$$
$$\land (b \lor c)$$
$$\land (\neg a)$$

а	b	C
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

Determining SAT/UNSAT

- All that is required to show satisfiability is to find a valid solution
- Many techniques available:
 - Guessing and checking
 - Systematic search
 - Inference

- Construct a binary tree of all combinations
- Proceeds in a depth first manner
- Each level corresponds to a variable
- Each branch corresponds to a truth assignment
- Branches of the tree are 'pruned' when the assignment cannot be extended in a satisfiable manner

$$(a \lor b \lor c)$$

$$\wedge (\underline{\neg a} \vee \neg b)$$

$$\wedge (\neg b \vee \neg c)$$

$$\wedge (\neg c \vee \neg a)$$

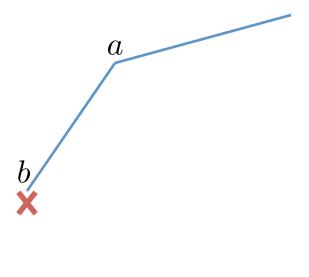


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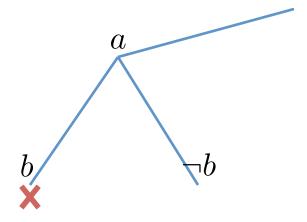


$$(\underline{a} \vee \underline{b} \vee c)$$

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$$\wedge (\underline{\neg b} \vee \neg c)$$

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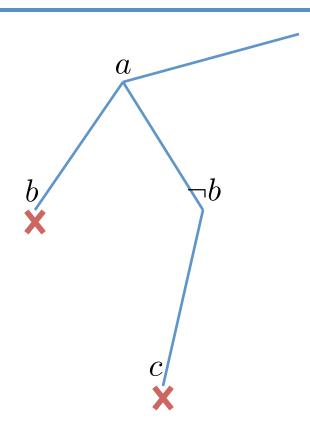


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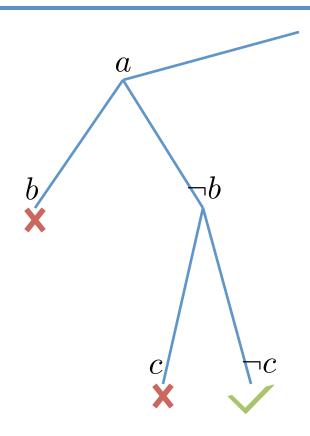


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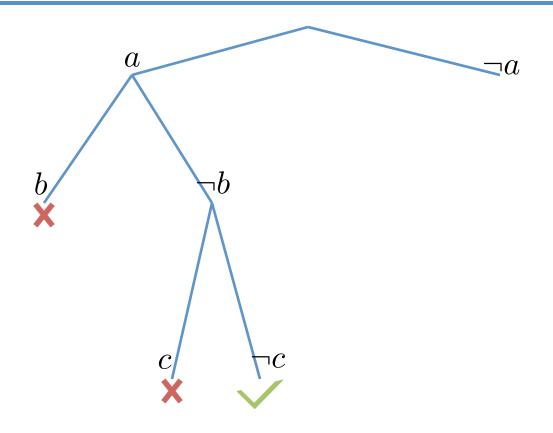


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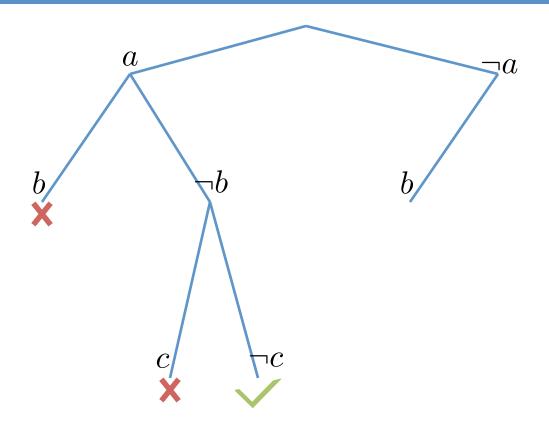


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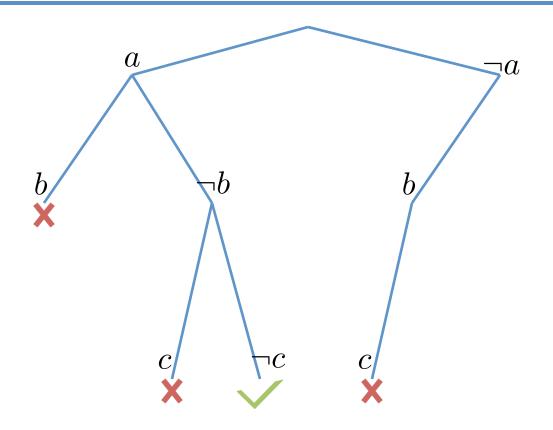


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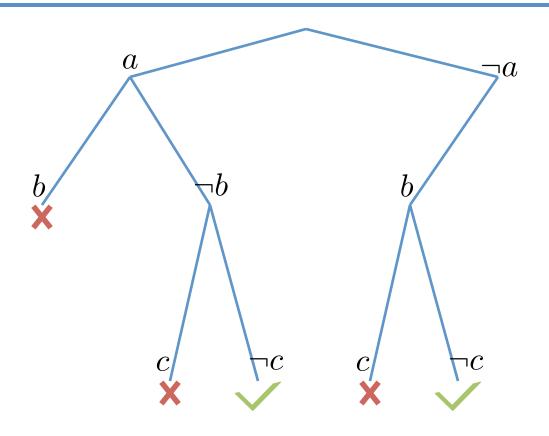


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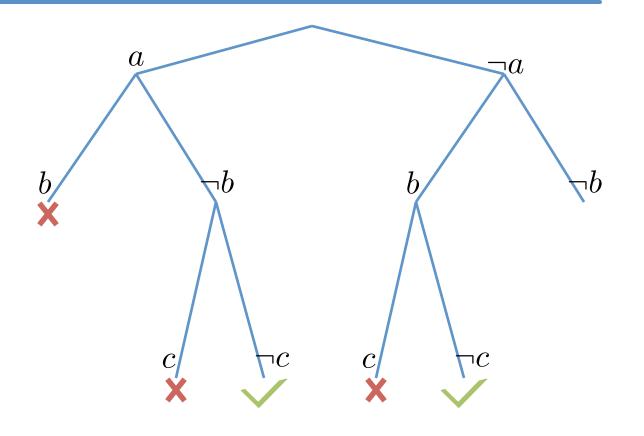


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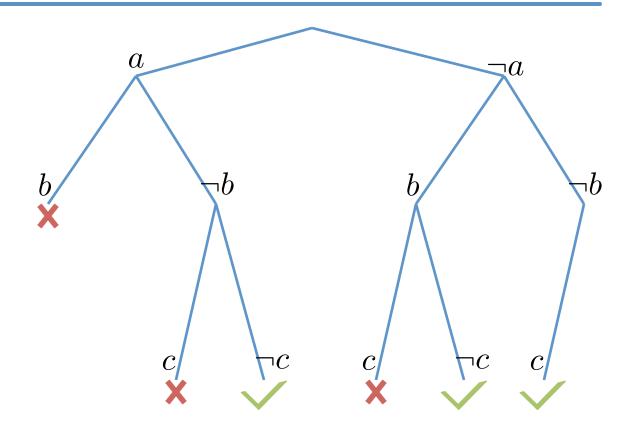


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$$(a \lor b \lor c)$$

$$\wedge (\underline{\neg a} \vee \underline{\neg b})$$

$$\wedge (\underline{\neg b} \vee \underline{\neg c})$$

$$\wedge (\underline{\neg c} \vee \underline{\neg a})$$

Satisfiable 3 solutions total

