

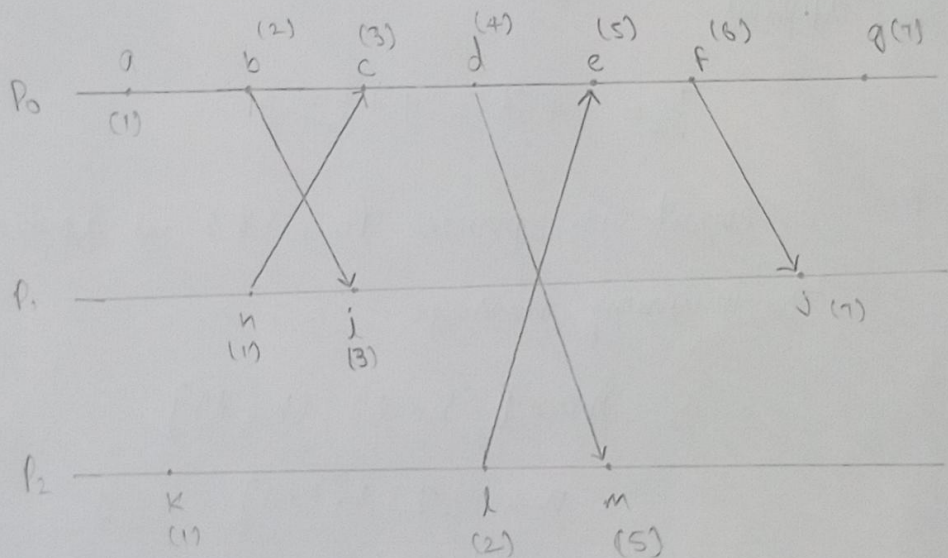
DIGITAL ASSIGNMENT - II

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Q.1. Consider the Figure, assume that initial logical clock values are all initialized to 0.

a) List the Lamport timestamps for each event and provide timestamps for each labeled event.



Assuming $d=1$

Implementations Rules [IR]:

$$IR1: C_i(\text{event } 2) = C_i(\text{event } 1) + d$$

Dealing with only events of one process.

$$IR2: C_j(\text{event } 2) = \max[C_j(\text{event } 1) + d, t_m + d]$$

Dealing with process events having communication with each other.

i) Since $d=1$, for all first events of all processes will be 1.

$$\begin{array}{c} P_0 \\ \underline{a=1} \end{array}$$

$$\begin{array}{c} P_1 \\ \underline{h=1} \end{array}$$

$$\begin{array}{c} P_2 \\ \underline{k=1} \end{array}$$

ii) For internal process b to P_0 , $IR1$ will be applied.

$$\underline{b} = 1 + d = 1 + 1 = 2$$

iii) c event in process P_0 , $IR2$ is applied due to incoming message.

$$\begin{aligned} \underline{c} &= \max[(b+d), (h+d)] \\ &= \max[2+1, 1+1] = 3 \end{aligned}$$

iv) i event in process P_1 , $IR2$ is applied due to incoming message

$$\begin{aligned} \underline{i} &= \max[(h+d), (b+d)] \\ &= \max[1+1, 2+1] = 3 \end{aligned}$$

v) For internal event d in P_0 , $IR1$ will be applied.

$$\underline{d} = 3 + 1 = 4$$

vi) For internal event l in process P_2 , IR_1 will be applied.

$$\underline{l} = l + 1 = 2$$

vii) e event, IR_2 will be applied due to incoming message

$$\begin{aligned}\underline{e} &= \max[(l+d), (1+d)] \\ &= \max[4+1, 2+1] \\ &= 5\end{aligned}$$

viii) m event in P_2 has incoming message, so IR_2 is applied.

$$\begin{aligned}m &= \max[(l+d), (e+d)] \\ &= \max[2+1, 4+1] \\ &= 5\end{aligned}$$

ix) For internal events f and g , IR_1 is applied

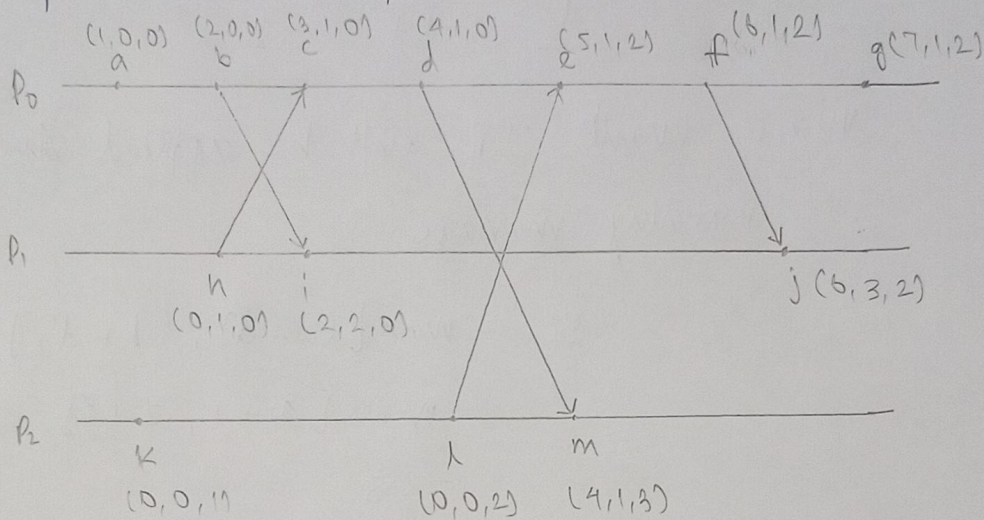
$$\underline{f} = e + d = 5 + 1 = 6$$

$$\underline{g} = f + d = 6 + 1 = 7$$

x) For event j in process P_1 , IR_2 will be applied.

$$\begin{aligned}\underline{j} &= \max[(i+d), (f+d)] \\ &= \max[6+1, 3+1] = 7\end{aligned}$$

b) List the Vector Clock timestamps for each event and provide timestamps for each labeled event.



Assuming $d=1$

Implementation Rules [IR]:

[IR1]: Clock C_i is incremented b/w any two successive events in process P_i :

$$C_i[i] := C_i[i] + d \quad (d > 0)$$

[IR2]: Event a sending message m by process P_i , then message m is assigned a vector timestamp $tm = C_i(a)$; on receiving the same message m by process P_j , C_j is updated as

$$\forall k, C_j[k] := \max(C_j[k], tm[k])$$

i) Events a , h and k are internal events to processes P_0 , P_1 and P_2 respectively. So IR1 will be applied.

$$\underline{a} \rightarrow (1, 0, 0) \quad \underline{h} \rightarrow (0, 1, 0) \quad \underline{k} \rightarrow (0, 0, 1)$$

ii) For internal events b and l in P_0 and P_2 respectively, IR1 will be applied.

$$\underline{b} \rightarrow (2, 0, 0) \quad \underline{l} \rightarrow (0, 0, 2)$$

iii) Events c and i processes there is a incoming messages so IR2 will be applied.

$$c \rightarrow \max[(3, 0, 0), (3, 1, 0)] \neq k$$

$$\underline{c} \rightarrow (3, 1, 0)$$

$$i \rightarrow \max[(2, 0, 0), (0, 2, 0)] \neq k$$

$$\underline{i} \rightarrow (2, 2, 0)$$

iv) Event d is internal event to process P_0 , so IR2 is applied.

$$d \rightarrow (4, 1, 0)$$

v7 Processes P_0 and P_2 have e and m events respectively, having incoming messages, IR2 will be applied.

$$e \rightarrow \max[(5, 1, 0), (0, 0, 2)] \neq k$$

$$\underline{e} \rightarrow (5, 1, 2)$$

$$m \rightarrow \max[(0, 0, 3), (4, 1, 0)] \neq k$$

$$\underline{m} \rightarrow (4, 1, 3)$$

vii7 For internal events f and g to process P_0 , IR1 will be applied.

$$\underline{f} \rightarrow (6, 1, 2)$$

$$\underline{g} \rightarrow (7, 1, 2)$$

viii7 Event j in P_1 has incoming message, IR2 will be applied.

$$j \rightarrow \max[(2, 3, 0), (6, 1, 2)]$$

$$\underline{j} \rightarrow (6, 3, 2)$$