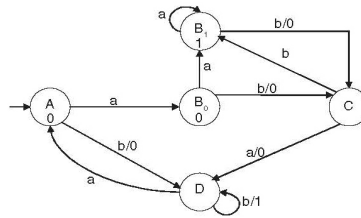
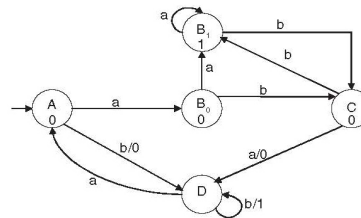


## 1 picture



For the state A, the incoming edges to this state are from  $B_{\{0\}}$  to C with label  $b/0$  and  $B_{\{1\}}$  to C with label  $b/0$ . There is no difference in the outputs of the incoming edges to this state, and so in the constructing Moore machine the output for this state will be 0.

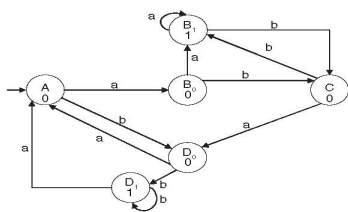
## 2 picture



For the state D, the incoming edges are A to D with label  $b/0$ , from C to D with label  $a/0$ , and from D to D with label  $b/1$ .

We get two different outputs for two incoming edges (D to D output 1, C to D output 0). So, the state D will be divided into two, namely,  $D_{\{0\}}$  and  $D_{\{1\}}$ . The outgoing edges are duplicated for both the states generated from D. The modified machine is

### 3 picture

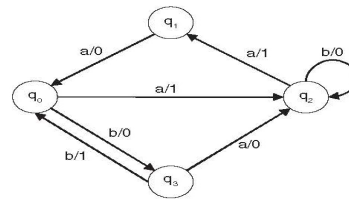


21. Convert the following Mealy machine into an equivalent Moore machine. [UPTU 2004]

Solution: The state  $q_0$  has two incoming edges: from  $q_1$  with label  $a/0$  and from  $q_3$  with label  $b/1$ . As there is a difference in output, the state  $q_0$  is divided into  $q_00$  and  $q_01$  with outputs 0 and 1, respectively. The states  $q_1$  and  $q_3$  have only one incoming edge each, and so there is no need of division. The state  $q_2$  has three incoming edges; among those, two are of output '0'

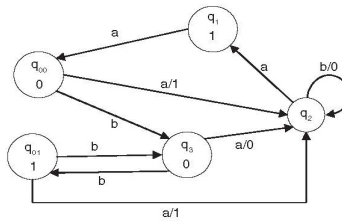
and another is of output '1'. Thus, it is divided into  $q_20$  and  $q_21$  with outputs 0 and 1, respectively.

#### 4 picture



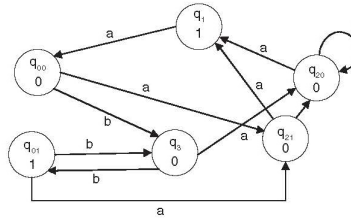
From  $q_1$  input with label 'a' ends on  $q_00$ , and from  $q_3$  input with label 'b' ends on  $q_01$ . The outputs from old  $q_1$  state are duplicated from  $q_00$  and  $q_01$ . The state  $q_1$  and  $q_3$  are not divided.  $q_1$  gets output '1' and  $q_3$  gets output '0'. Dividing the state  $q_0$  and placing  $q_1$  and  $q_3$ , the intermediate machine becomes as follows

#### 5 picture



The state  $q_24$  is divided into  $q_20$  and  $q_21$ . From  $q_00$  and  $q_01$  input with label 'a' ends on  $q_21$ . From  $q_3$  input with label 'a' ends on  $q_20$ . There is a loop on  $q_2$ . That loop will be on  $q_20$  with label 'b'. Another transition with label 'b' is drawn from  $q_21$  to  $q_20$ . The final Moore machine is as follows

## 6 picture



22. Minimize the following finite automata.

Next State			
	Present State	I/P = a	I/P = b
→	A	B	F
	B	A	F
	C	G	A
	D	H	B
	E	A	G
	F	H	C
	G	A	D
	H	A	C

Here F, G, and H are the final states.

Solution: In the finite automata, the states are  $\{A, B, C, D, E, F, G, H\}$ . Name this set as  $S_0$ .

$S_0 : \{A, B, C, D, E, F, G, H\}$

All of the states are 0 equivalents.

In the finite automata, there are two types of states: final state and non-final states.

So, divide the set of states into two parts, Q1 and Q2.

$Q_1 = \{F, G, H\}$   $Q_2 = \{A, B, C, D, E\}$

$S_1 : \{\{F, G, H\}, \{A, B, C, D, E\}\}$

The states belonging to same subset are 1-equivalent because they are in the same set for string length 1. The states belonging to different subsets are 1-distinguishable.

The next states of F are H and C. The next states of G and H are A, D and A, C, respectively.

A, D and A, C belong to the same subset but H and C belong to a different subset.

So, F, G, and H are divided into  $\{F\}, \{G, H\}$ .

For input 0, the next states of A, B, C, D, and E are B, A, G, H, and A, respectively. For input 1, the next states of A, B, C, D, and E are F, F, A, B, and G, respectively. So, the set  $\{A, B, C, D, E\}$  is divided into  $\{A, B, E\}$  and  $\{C, D\}$ .

$S_2 : \{\{F\}\{G, H\}\{A, B, E\}\{C, D\}\}$

By the same process,  $\{A, B, E\}$  is divided into  $\{A, B\}, \{E\}$ .

$S_3 : \{\{F\}\{G, H\}\{A, B\}\{E\}\{C, D\}\} = \{\{A, B\}, \{C, D\}, \{E\}, \{F\}, \{G, H\}\}$

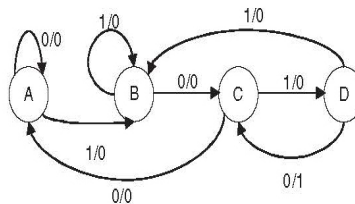
The set is not dividable further. So, these are the states of minimized DFA. Let us rename the subsets as  $q_0, q_1, q_2, q_3$ , and  $q_4$ . The initial state was A, and so here the initial state is  $\{A, B\}$ , i.e.,  $q_0$ . The final state was F, G, and H, and so here the final states are  $\{F\}$ , i.e.,  $q_3$  and  $\{G, H\}$ , i.e.,  $q_4$ . The tabular representation of minimized DFA is

		Next State	
	Present State	I/P = 0	I/P = 1
→	$q_0$	$q_0$	$q_0$
	$q_0$	$q_0$	$q_0$
	$q_0$	$q_0$	$q_0$
	$q_0$	$q_0$	$q_0$
	$q_0$	$q_0$	$q_0$

23. Design a Mealy and Moore machine for detecting a sequence 1010 where overlapping sequences are also accepted. Convert the Moore machine that you have got into a Mealy machine. Are there any differences? How will you prove that the two Mealy machines are equivalent?

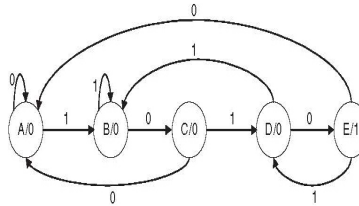
Solution: The Mealy machine is

## 7 picture



The Moore machine is

## 8 picture



The converted Mealy machine from the given Moore machine is (by using the transactional format)

## 9 picture

