

Design & Analysis of Algorithms

①

Tutorial - 7

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Section \rightarrow D

Class R.No. \rightarrow 29

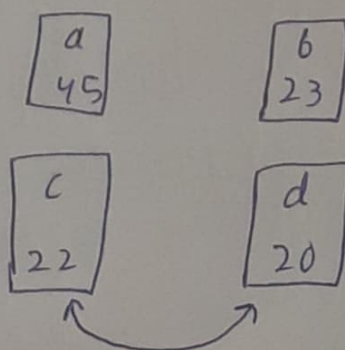
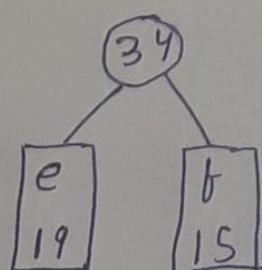
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- ① It is an algorithmic paradigm that builds up a solution by adjoining smaller pieces together, always choosing the next piece that offers the most obvious and immediate benefit.

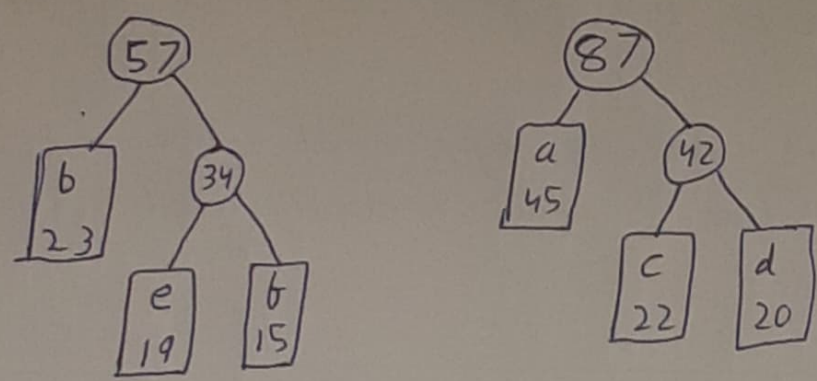
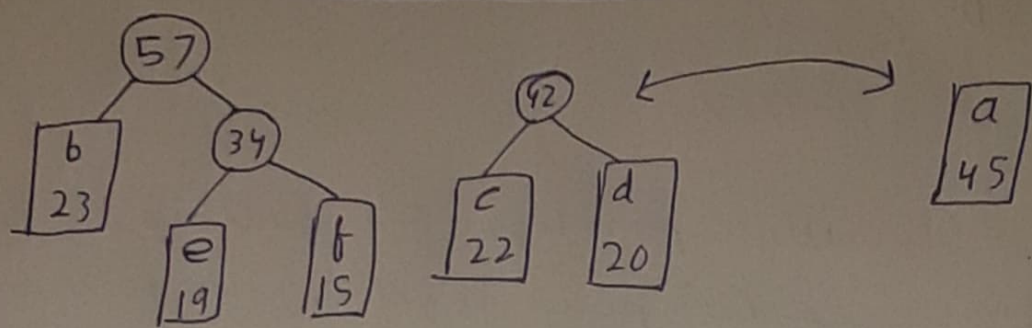
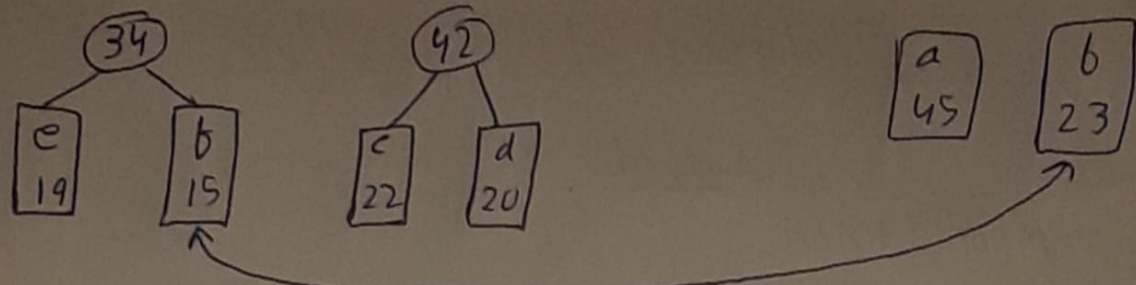
We should use greedy approach whenever a locally optimal solution is also globally optimal.

② Name	TC	SC
Activity Selection	$O(n \log n) \leftrightarrow O(n)$	$O(n)$
Job Sequencing	$O(n^2) \leftrightarrow O(n \log n)$	$O(n)$
Fractional knapsack	$O(n \log n) \leftrightarrow O(n)$	$O(n)$
Huffman Encoding	$O(n \log n) \leftrightarrow O(\log n)$	$O(n)$

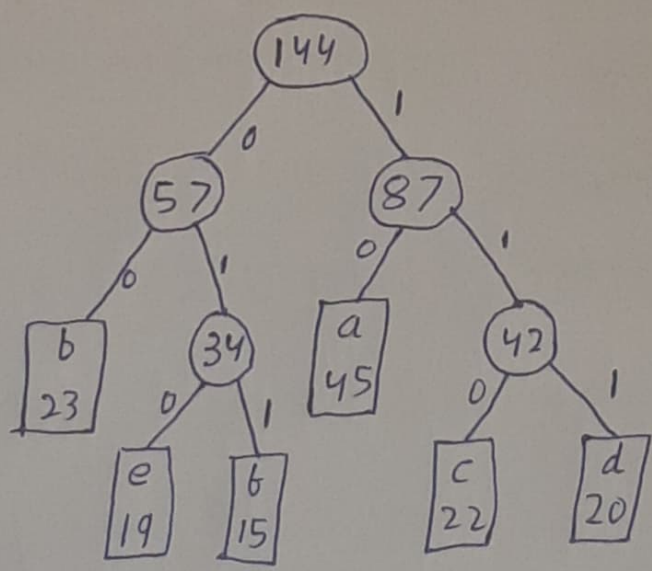
- ③ $a = 45, d = 20, b = 23, e = 19, c = 22, f = 15$



2



Final Huffman Tree ⇒



$a = 10$
 $b = 00$
 $c = 110$
 $d = 111$
 $e = 010$
 $f = 111$

Total bits Used ⇒ $(45 \times 2) + (23 \times 2) + (22 \times 3) + (20 \times 3)$
 $+ (19 \times 3) + (15 \times 3)$
 $= 364 \text{ bits}$

- ④ A 2-tree is used to implement Huffman encoding algorithm. It is a binary tree where every node has either 2-child or no child.

Applications of Huffman Encoding \Rightarrow

- Data compression in long files without any loss.
- To implement traffic routes with traffic magnitude

Ans-5

V	10	5	15	7	6	18	3
w	2	3	5	7	1	4	1
V/w	5	1.67	3	1	6	4.5	3

$$k = 15 - 1 - 2 - 4 - 5 - 1 - 2 = 0$$

$$\text{Profit} = 30 + 10 + 18 + 15 + 3 + 3 - 34 = 79 - 34$$

V	6	10	18	15	3	5
w	1	2	4	5	1	3
V/w	6	5	4.5	3	3	1.67

Ans-6 Fractional Knapsack: It is using a greedy approach as we have divided our profits to the smallest unit possible & then builds upon it.

Huffman Encoding: It is using the greedy approach as we have divided our profits to the smallest unit possible & then builds upon it.

(4)

Huffman Encoding: It is using the greedy approach as it always places the node with the lower frequency further from the parent node.

Ans-7

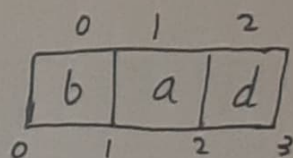
<u>Start</u>	1	2	0	6	9	10
<u>End</u>	3	5	7	8	11	12
<u>Index</u>	0	1	2	3	4	5

Jobs to do = [0], [3], [4] or [5]

i.e. \Rightarrow Max = 4

Ans-8

	Profit	Deadline
a	20	2
b	15	2
c	10	1
d	5	3
e	1	3

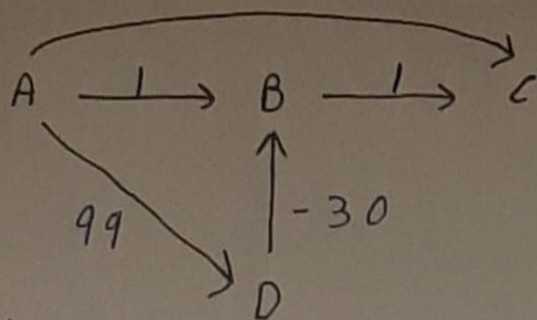


$$\text{Profit} = 20 + 15 + 5 = 40$$

Ans-9 Times when not to use greedy approach

- When the approach involves a lot of consumptions, a such as "pick always the -----"
- We should avoid greedy approach on complex implementation.
- When we are making performance-critical applications.

Eg \Rightarrow Dijkstra's algorithm is very unoptimized for graphs with negative edges.



We can't find the distance of the pair $[n, c] \rightarrow$ it gives 0, though it is -200

Ans - 10 Normally, the time complexity of Job sequencing is $O(n^2)$ but we can improve it using a Priority Queue, made of Max Heap.

Algorithm :-

1. Sort the job based on deadlines
2. Iterate the end & calculate the available slots b/w 2 consecutive deadlines include all data in Max-Heap
3. If there are slots available & there are jobs in the max heap, include the job ID with max profit & deadlines in the result.

Sort the array based on deadlines

Time Complexity $\Rightarrow O(n \log(n))$

Space Complexity $\Rightarrow O(n)$