

Assignment-4

Question 1-B:

Let us define the equivalence of high-order function g and its CPS version $g^{\$}$ like this, for any $0 \leq i \leq n$ the CPS-equivalent parameters $f_1 \dots f_n$ and $f_1^{\$} \dots f_n^{\$}$ ($g^{\$} f_1^{\$} \dots f_n^{\$} \text{cont}$) is CPS-equivalent to $(\text{cont } (g f_1 \dots f_n))$

Now, we are going to show that pipe is equivalent to pipe\$, using induction on the size of the List.

1- Base: $N=1$

$(\text{cont } (\text{pipe}(f_1^{\$}))) = (\text{cont } f_1^{\$}) \ \& \ (\text{pipe}^{\$} f_1^{\$} \text{cont}) = (\text{cont } (\lambda (x \ c2) (f_1^{\$} x \ c2))) = (\text{cont } f_1^{\$})$

2-Induction step: we assume that $(\text{pipe}^{\$} f_1^{\$} \dots f_n^{\$} \text{cont}) = (\text{cont } (\text{pipe } f_1^{\$} \dots f_n^{\$}))$

3-Prove it for $n+1$:

$$\begin{aligned} & (\text{pipe}^{\$} (f_1^{\$} \dots f_n^{\$} f_{n+1}^{\$} \text{cont})) \\ & (\text{pipe}^{\$} f_2^{\$} \dots f_{n+1}^{\$} (\lambda (f_2^{\$} \dots f_n^{\$}) (\text{cont } (\lambda (x \ c2) (f_1^{\$} x (\lambda (\text{res}) (f_{n+1}^{\$} \text{res } c2)))))) = () \\ & (\lambda (f_2^{\$} \dots f_n^{\$}) (\text{cont } (\lambda (x \ c2) (f_1^{\$} x (\lambda (\text{res}) (f_{n+1}^{\$} \text{res } c2)))))) (\text{pipe } f_2^{\$} \dots f_{n+1}^{\$}) \\ & = (\text{cont } (\lambda (x \ c2) ((\text{pipe } f_2^{\$} \dots f_{n+1}^{\$}) x (\lambda (\text{res}) (f_{n+1}^{\$} \text{res } c2)))) \\ & = (\text{cont } (f_2^{\$} \dots (\text{pipe } f_1^{\$} f_2^{\$} \dots f_{n+1}^{\$}))) = (\text{cont } (\text{pipe } f_1^{\$} \dots f_{n+1}^{\$})) \end{aligned}$$

Question 2-b:

We will use reduce 1-lzl when we want to get a reduce for a finite lazy list.

We will use reduce 1-lzl when we want to get a reduce of one specific prefix of a given infinite lazy list.

We will use reduce 1-lzl when we want to get a reduce of each prefix of an infinite lazy list (as

we use it in Q2e)

Question 2-g:

Advantage: unlike the pi-sum method, which is fixed to a single given 'b' limit, this method can be applied to any approximation level.

Disadvantage: Using this method generates a lot of closures.

Question 3-1:

1. unify $[x(y(y), T, y, z, k(K), y), x(y(T), T, y, z, k(K), L)]$

$$S = \{\} \circ \{T=y, L=y\} = \{T=y, L=y\}$$

$$\boxed{\begin{matrix} y(y) = y(T) \\ L = y \end{matrix}}$$

unify $[x(y(y), y, y, z, k(K), y), x(y(y), y, y, z, k(K), y)]$

$$S = \{T=y, L=y\}$$

2. unify $[f(a, M, f, F, Z, f, x(M)), f(a, x(Z), f, x(M), x(F), f, x(M))]$

כאן נקבע

מכיוון שיש לנו

$$\boxed{\begin{matrix} M = x(Z) & M = x(x(x(M))) \\ F = x(M) & \Rightarrow F = x(x(x(F))) \\ Z = x(F) & Z = x(x(x(Z))) \end{matrix}}$$

3. unify $[t(A, B, C, n(A, B, C), x, y), t(a, b, c, m(A, B, C), X, Y)]$

$$\boxed{\begin{matrix} A=a & X=x \\ B=b & Y=y \\ C=c \end{matrix}}$$

$$S = \{A=a, B=b, C=c, X=x, Y=y\}$$

unify $[t(a, b, c, n(a, b, c), x, y), t(a, b, c, m(a, b, c), x, y)]$

4. unify $[z(a(A, x, Y), D, g), z(a(d, x, g), g, Y)]$

$$S = \{\} \circ \{D=g, Y=g\}$$

$$\boxed{\begin{matrix} a(A, x, Y) = a(d, x, g) \\ D = g \\ Y = g \end{matrix}}$$

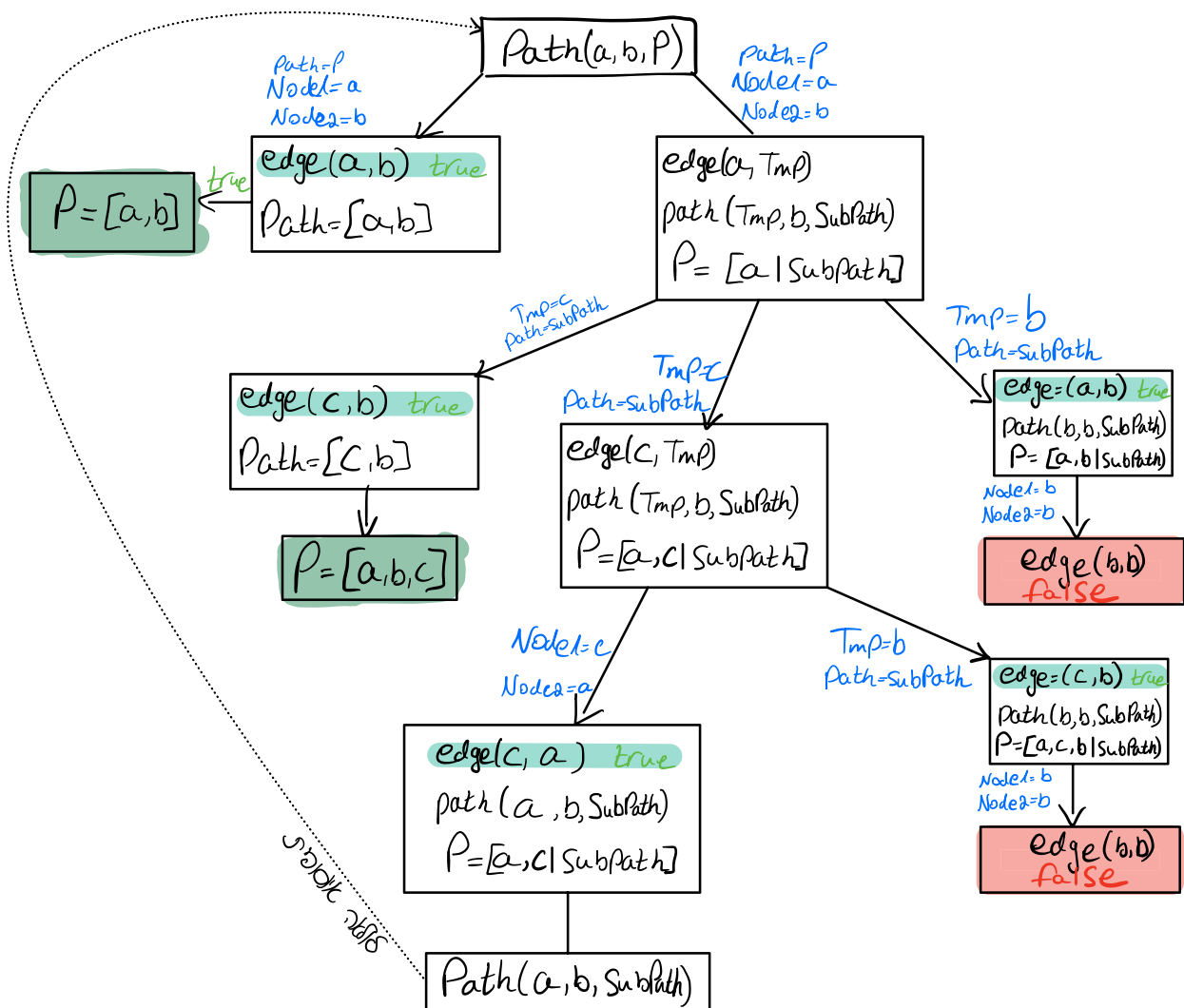
$$\text{unify}[a(A, x, Y), a(d, x, g)]$$

$$S' = \{A=d, Y=g\}$$

unify $[z(a(d, x, g), g, g), z(a(d, x, g), g, g)]$

$$S = S \circ S' = \{D=g, Y=g, A=d\}$$

Question 3-3a



→ according to the given proof tree it's infinite tree.

→ according to the given proof tree we have more than one success path therefore it's success tree.