Operating Systems Simulation Project Report

CSEN 602 – Milestone 2 German University in Cairo

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Abstract

This report summarizes the design, implementation, and testing of an Operating Systems simulation developed in C with a GTK3 front-end. The simulator implements process control, memory management, synchronization, and multiple scheduling algorithms (FCFS, Round Robin, and MLFQ).

Introduction

The goal of this project is to build a user-friendly OS simulator that demonstrates core operating system concepts: process lifecycle management, scheduling policies, memory operations, and inter-process synchronization. A graphical interface allows step-by-step execution, parameter tuning, and real-time visualization of process state and memory usage.

Objectives

- **Simulate process execution** by parsing and interpreting instruction streams.
- Manage memory in a simulated heap, including variable allocation and access.
- **Implement synchronization** primitives (mutexes) to control access to shared resources.
- **Support multiple scheduling algorithms** (FCFS, Round Robin with configurable quantum, and Multi-Level Feedback Queue).
- Provide a GTK3 GUI for control, logging, and visualization.

System Architecture

The simulator is organized into modular components:W

- memory** module**: Manages a byte-array heap and variable lookup by name.
- pcb** and process modules**: Define and manipulate Process Control Blocks (PCBs) tracking process IDs, state (NEW, READY, RUNNING, BLOCKED, EXIT), and memory bounds.
- **interpreter**** module**: Reads program files, decodes instructions (assignment, I/O, synchronization calls), and executes them against the memory and PCB layers.
- mutex** module**: Provides semWait/semSignal operations to implement mutual exclusion for named semaphores (mutexes).
- **scheduler**** module**: Implements FCFS, Round Robin, and MLFQ policies. The API initScheduler(algo, quantum) sets the active policy; schedule() and onProcessUnblocked() handle queue management and preemption.
- **gui**** module******: A GTK3 application presenting controls (scheduler selection, quantum spinner, step/run buttons), log output, memory view, and process state table.

Implementation Details

1 Memory Management

The memory.c file defines a contiguous simulated memory array. Variables are inserted with storeVariable(name, value), and looked up by loadVariable(name). Bounds checking and initialization ensure safe access.

2 Process Control Block

Each process is represented by a PCB struct holding its ID, program counter index, state enum, and memory offset ranges. PCBs are allocated dynamically on process creation and tracked in the scheduler's ready and blocked queues.

3 Instruction Interpreter

interpreter.c tokenizes each line of the input program. It supports assignments (e.g. a = b + 1), print statements, blocking reads, and mutex operations. The interpreter returns control to the scheduler after each instruction or on I/O waits.

4 Scheduling Policies

• FCFS enqueues processes in arrival order and runs them to completion.

- Round Robin (RR) uses a time-quantum; the scheduler preempts processes after quantum steps and requeues them.
- MLFQ maintains multiple queues of decreasing priority; each higher-priority queue uses shorter quanta. Processes demote on expiry and promote on I/O blocking to prevent starvation.

The GUI exposes a combo-box and spin-button to switch policies at runtime via on_scheduler_changed().

5 Synchronization

The mutex module implements named semaphores with counters and blocking queues. semWait(name) decrements the semaphore or blocks the calling process, while semSignal(name) wakes one blocked process, integrating with the scheduler via onProcessUnblocked().

6 GUI and Integration

A single GTK3 window shows:

- **Control bar**: Scheduler selection, quantum value, step/run/stop buttons.
- Execution log: Live text output of debug messages and instruction traces.
- Process table: Lists PCBs with PID, state, and PC.
- Memory view: Shows current variables and values.

Testing and Validation

Three sample programs (Program_1, Program_2, Program_3) were used to validate:

- 1. **Basic arithmetic and I/O**, ensuring correct interpreter operation and GUI prompts.
- 2. Mutex contention, verifying blocking and unblocking semantics.
- 3. **Scheduling behavior**, confirming RR preemption and MLFQ priority movement.

Conclusion

This project successfully demonstrates a miniature OS environment, complete with memory, processes, synchronization, and scheduling, all managed within a GTK3 graphical interface. The modular design permits easy extension (e.g., additional algorithms) and offers clear visualization for teaching core OS concepts.