

ICS 2020 Problem Sheet #6:

Problem 6.1:

We know that { AND, NOT } is a complete set of connectives (universal) so (nor) \downarrow and (nand) \uparrow are universal too. We can simply prove that \rightarrow (implication) and \neg (negation) are universal by expressing \downarrow or \uparrow with \rightarrow and \neg .

p	q	$p \rightarrow q$	$\neg p$	$\neg q$	$p \downarrow q$	$p \uparrow q$
F	F	T	T	T	T	T
F	T	T	T	F	F	T
T	F	F	F	T	F	T
T	T	T	F	F	F	F

From the table we deduce that:

$$(p \uparrow q) \equiv (p \rightarrow (\neg q))$$

$$(p \downarrow q) \equiv (\neg((\neg p) \rightarrow q))$$

Thus, implication and negation are universal.

Problem 6.2:

$$a) \phi(A, B) = (\neg A \vee \neg B) \wedge (\neg A \vee B) \wedge (A \vee \neg B)$$

$$\phi(A, B) = (\neg A \vee \neg B) \wedge (\neg A \vee B) \wedge (A \vee \neg B)$$

$$\phi(A, B) = \neg A \vee (\neg B \wedge B) \wedge (A \vee \neg B) \text{ (distributivity)}$$

$$\phi(A, B) = \neg A \vee (0) \wedge (A \vee \neg B) \text{ (identity)}$$

$$\phi(A, B) = \neg A \wedge (A \vee \neg B) \text{ (distributivity)}$$

$$\begin{aligned}\phi(A, B) &= (\neg A \wedge A) \vee (\neg A \wedge \neg B) \\ \phi(A, B) &= 0 \vee (\neg A \wedge \neg B) \text{ (identity)} \\ \phi(A, B) &= (\neg A \wedge \neg B) \text{ (de Morgan's laws)} \\ \phi(A, B) &= \neg(A \wedge B)\end{aligned}$$

$$b) \phi(A, B, C) = (A \wedge \neg B) \vee (A \wedge \neg B \wedge C)$$

$$\begin{aligned}\phi(A, B, C) &= (A \wedge \neg B) \vee (A \wedge \neg B \wedge C) \\ \phi(A, B, C) &= (A \wedge \neg B) \vee ((A \wedge \neg B) \wedge C) \text{ (associativity)} \\ \phi(A, B, C) &= (A \wedge \neg B) \text{ (absorption laws)}\end{aligned}$$

$$c) \phi(A, B, C, D) = (A \vee \neg(B \wedge A)) \wedge (C \vee (D \vee C))$$

$$\begin{aligned}\phi(A, B, C, D) &= (A \vee \neg(B \wedge A)) \wedge (C \vee (D \vee C)) \\ \phi(A, B, C, D) &= (A \vee (\neg B \wedge \neg A)) \wedge (C \vee (D \vee C)) \text{ (de Morgan's laws)} \\ \phi(A, B, C, D) &= ((A \vee \neg B) \wedge (A \vee \neg A)) \wedge (C \vee (D \vee C)) \text{ (distributivity)} \\ \phi(A, B, C, D) &= ((A \vee \neg B) \wedge 1) \wedge (C \vee (D \vee C)) \\ \phi(A, B, C, D) &= (A \vee \neg B) \wedge (C \vee (D \vee C)) \text{ (identity)} \\ \phi(A, B, C, D) &= (A \vee \neg B) \wedge (D \vee (C \vee C)) \text{ (associativity)} \\ \phi(A, B, C, D) &= (A \vee \neg B) \wedge (D \vee C) \text{ (idempotency)}\end{aligned}$$

$$d) \phi(A, B, C) = (\neg(A \wedge B) \vee \neg C) \wedge (\neg A \vee B \vee \neg C)$$

$$\begin{aligned}\phi(A, B, C) &= (\neg(A \wedge B) \vee \neg C) \wedge (\neg A \vee B \vee \neg C) \\ \phi(A, B, C) &= \neg((A \wedge B) \vee C) \wedge (\neg A \vee B \vee \neg C) \text{ (de Morgan's laws)} \\ \phi(A, B, C) &= (\neg((A \wedge B) \vee C) \wedge \neg(A \vee \neg B \vee C)) \text{ (double negation)} \\ \phi(A, B, C) &= \neg(((A \wedge B) \vee C) \wedge (A \vee B \vee C)) \text{ (de Morgan's laws)} \\ \phi(A, B, C) &= \neg(((A \wedge B) \vee C) \wedge (A \vee B)) \vee ((A \wedge B) \vee C) \wedge C) \\ &\quad \text{(distributivity)} \\ \phi(A, B, C) &= \neg((A \wedge B) \vee C) \vee ((A \wedge B) \vee C) \text{ (idempotency)} \\ \phi(A, B, C) &= \neg((A \wedge B) \vee C) \text{ (idempotency)}\end{aligned}$$

$$e) \phi(A, B) = (A \vee B) \wedge (\neg A \vee B) \wedge (A \vee \neg B) \wedge (\neg A \vee \neg B)$$

$$\phi(A, B) = (A \vee B) \wedge (\neg A \vee B) \wedge (A \vee \neg B) \wedge (\neg A \vee \neg B)$$

$$\phi(A, B) = (A \vee B) \wedge (\neg A \vee \neg B) \wedge (\neg A \vee B) \wedge (A \vee \neg B)$$

(commutativity)

$$\phi(A, B) = (A \vee B) \wedge \neg(A \vee B) \wedge (\neg A \vee B) \wedge (A \vee \neg B)$$

(de Morgan's laws)

$$\phi(A, B) = (A \vee B) \wedge \neg(A \vee B) \wedge (\neg A \vee B) \wedge \neg(\neg A \vee B)$$

(double negation)

$$\phi(A, B) = 0 \wedge 0$$

$$\phi(A, B) = 0$$

Problem 6.3:

a) $\phi(P, Q, R, S) = (\neg P \vee Q) \wedge (\neg Q \vee R) \wedge (\neg R \vee S) \wedge (\neg S \vee P)$

P	Q	R	S	($\neg P \vee Q$)	($\neg Q \vee R$)	($\neg R \vee S$)	($\neg S \vee P$)	ϕ
F	F	F	F	T	T	T	T	T
F	F	F	T	T	T	T	F	F
F	F	T	F	T	T	F	T	F
F	F	T	T	T	T	T	F	F
F	T	F	F	T	F	T	F	F
F	T	F	T	T	F	T	F	F
F	T	T	T	T	T	T	F	F
F	T	T	T	T	T	T	F	F
T	F	F	F	F	T	T	T	F
T	F	F	T	F	T	T	T	F
T	F	T	F	F	T	F	T	F
T	F	T	T	F	T	T	T	F
T	T	F	F	T	F	T	T	F

T	T	F	T	T	F	T	T	F
T	T	T	F	T	T	F	T	F
T	T	T	T	T	T	T	T	T

From the truth table we see that 2 interpretation satisfy ϕ

b)

The two interpretations that satisfies ϕ require that P,Q,R,S are all TRUE or ALL false

$$\text{DNF} = (P \wedge Q \wedge R \wedge S) \vee (\neg P \wedge \neg Q \wedge \neg R \wedge \neg S)$$