

Introduction

14 September 2020 20:03

Amortized Analysis

- Seq of operations

↓ time

average Over all
operations

- Probability is not involved

- "average performance of each operation in worst case"

1. Aggregate analysis

Seq. of 'n'

$T(n)$ - Total time
Complexity
'n' operations
Worst case

$$\text{amortized cost per } \left\{ = \underline{T(n)} \right.$$

operation) n Stack operations

Push (S, 'X')

POP (S)

Actual Cost of

Push - $O(1)$

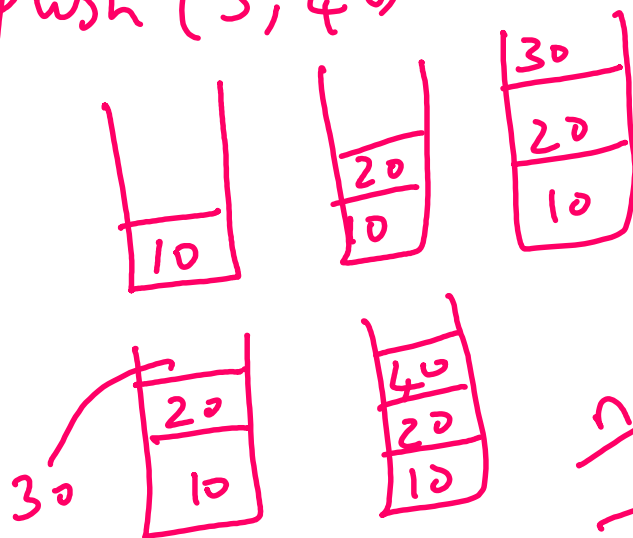
POP - $O(1)$

Seq of operations

Push (S, 10), Push (S, 20)

Push (S, 30), POP (S),

Push (S, 40)



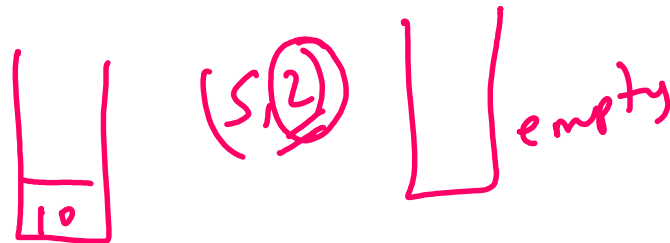
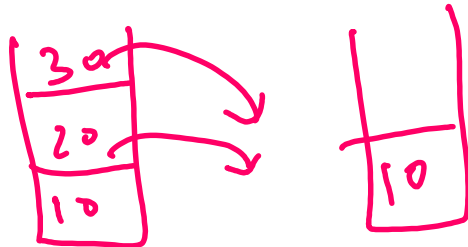
$n = 5$
'n' operations

Total cost of } = $O(n)$
 push & pop

3) multi pop (s, 'k')

k - elements
are popped

multi pop (s, 2)



$$\text{multi pop}(s, k) = \min(s, k)$$

s

no. of elements
in stack

Total cost of
'single' multi pop } = $\min(s, k)$

Worst case total
cost of multi pop } = $O(s)$



(c 2)

<u>3</u>	5
20	
10	

$$s = n$$

$$\text{multiply}(s, -)$$

$$= 2$$

$$\text{multiply}(\underline{s}, \underline{3}) = \underline{\underline{3}}$$

$$= \boxed{O(n)}$$

$$\left(\text{multiply}^{O(n)}(s, k), \text{multiply}^{O(n)}(s, k) \right)$$

$$\dots n \text{ operations}$$

$$\text{Total cost} = O(n^2)$$

$$\text{of 'n' seq of multiply}$$

$$= n \times O(n)$$

$$= O(n^2)?$$

Not
tight

<u>n=3</u>	30
	20
	10

$$\text{multiply}(\underline{s}, \underline{3})$$

$$\text{multiply}(,)$$

X

$$\left. \begin{array}{l} \text{Closest} \\ \text{upper bound} \end{array} \right\} = O(n)$$

$$\left. \begin{array}{l} \text{Seq of operations} \\ \text{push, pop, multiply} \end{array} \right\} = O(n)$$

Total cost

$$T(n) = O(n)$$

$$\left. \begin{array}{l} \text{Amortized} \\ \text{cost of} \\ \text{push or pop} \\ \text{or multiply} \end{array} \right\} = \frac{O(n)}{n} = \underline{\underline{O(1)}}$$

2) Accounting Method

Seq. operations

Operation \rightarrow cost is assigned.

'Charge',
'Credit'

more Credit
less Credit

amortized cost

'... bit

'more credit' \rightarrow 'less' credit

Stack operation

C_i - Actual cost
of i^{th} operation

\hat{C}_i - Amortized cost
of i^{th} operation

$$\sum_{i=1}^n \hat{C}_i \geq \sum_{i=1}^n C_i$$

Amor cost Act. cost

$$\begin{aligned} \text{Total Credits} &= \sum_{i=1}^n \hat{C}_i - \sum_{i=1}^n C_i \\ &\geq 0 \end{aligned}$$

Stack

push

pop

min/max

Actual

$O(1)$

$O(1)$

min/max

$n = n$

$O(n)$

Minimax method

Amortized cost

push 2

pop 0

multiplier 0

constant

following seq

	<u>Amo. cost</u>	<u>Act. cost</u>	<u>Crack</u>
① push(S, 10)	2	1	$2 - 1 = 1$
② push(S, 20)	2	1	$4 - 2 = 2$
③ push(S, 30)	2	1	$6 - 3 = 3$
④ pop(S)	0	1	$3 - 1 = 2$
⑤ pop(S)	0	1	$2 - 1 = 1$
⑥ multiplier(S, 3)	0 $\frac{0}{6=n}$	$\min(1, 3)$ $= 1$ $\frac{1}{6=n}$	$1 - 1 = 0$

n = 6

Total amortized cost } = $O(n)$
= Total act. cost