



ALGORITHMS TASK





Task Number 7
K-th Element of Two Sorted Arrays

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Introduction



Problem Statement

Given two sorted arrays of sizes m and n , find the element that would be at the k -th position in their merged sorted array.

****Examples**:**

Input:

Array 1 = [2, 3, 6, 7, 9]

Array2 = [1, 4, 8, 10]

$k = 5$

Output: 6 (Merged array: [1, 2, 3, 4, 6, 7, 8, 9, 10])





PROBLEMS SOLUTION



Solution 1

**NON-RECURSIVE
APPROACH**

Solution 2

**Merge Sort
(RECURSIVE APPROACH)**

Solution 3

**Bubble Sort
(RECURSIVE APPROACH)**

Solution 4

**Insertion Sort
(RECURSIVE APPROACH)**





NON-RECURSIVE APPROACH

1. Pseudo-code

```
function NonRec_Fun(arr1, size1, arr2, size2, k):  
    i ← 0, j ← 0, count ← 0  
    while i < size1 and j < size2:  
        if arr1[i] < arr2[j]:  
            element ← arr1[i]  
            i ← i + 1  
        else:  
            element ← arr2[j]  
            j ← j + 1  
        count ← count + 1  
        if count == k:  
            return element  
  
    while i < size1:  
        element ← arr1[i]  
        i ← i + 1  
        count ← count + 1  
        if count == k:  
            return element  
  
    while j < size2:  
        element ← arr2[j]  
        j ← j + 1  
        count ← count + 1  
        if count == k:  
            return element  
  
    return -1 // k is out of bounds
```





02 Source code

→ GitHub Link :

03 Sample of output

```
NonRecursive_Task.cpp - Code::Blocks 20.03
File Edit View Search Project Build Debug Fortran wxSmith Tools Tools+ Plugins DoxyBlocks Settings Help
<global>
1 #include <iostream>
2 using namespace std;
3 const int MAX_SIZE = 1000;
4
5 int NonRecursiveTask(int arr1[], int arr2[], int k)
6 {
7     int n1 = arr1[0], n2 = arr2[0];
8     int i = 1, j = 1;
9     while (i <= n1)
10     {
11         arr1[i] = i;
12         i++;
13     }
14     while (j <= n2)
15     {
16         arr2[j] = j;
17         j++;
18     }
19     int kthElement = 0;
20     while (i <= n1 || j <= n2)
21     {
22         if (i <= n1 && (j > n2 || arr1[i] < arr2[j]))
23             kthElement = arr1[i++];
24         else
25             kthElement = arr2[j++];
26     }
27     return kthElement;
28 }
```

C:\Users\del\Downloads\Algo x + v

```
Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
2
3
6
7
9
Enter the size of Array 2: 4
Enter the elements of Array 2 'sorted':
1
4
8
10
Enter k position: 5
The 5-th element in the merged array is: 6

Process returned 0 (0x0)   execution time : 27.543 s
Press any key to continue.
```





04

Time Complexity

$O(n)$

05

Space Complexity

$O(1)$





06

Advantages

- 1- Simple and straightforward implementation
- 2- Easy to understand and verify correctness
- 3- No additional memory allocation needed
- 4- Handles edge cases well (empty arrays, k out of bounds)

07

Disadvantages

- 1- Not optimal for large arrays when k is large
- 2- Linear time complexity could be improved with a binary search approach





1. Pseudo-code

Merge Sort (RECURSIVE APPROACH)

```
function mergeRecursive(arr1, size1, arr2, size2, merged, i = 0, j = 0, k = 0):  
    if i == size1 AND j == size2: // Base case: both arrays fully processed  
        return  
  
    if i == size1: // arr1 exhausted, take from arr2  
        merged[k] = arr2[j]  
        mergeRecursive(arr1, size1, arr2, size2, merged, i, j + 1, k + 1)  
  
    else if j == size2: // arr2 exhausted, take from arr1  
        merged[k] = arr1[i]  
        mergeRecursive(arr1, size1, arr2, size2, merged, i + 1, j, k + 1)  
  
    else if arr1[i] < arr2[j]: // arr1 has smaller element  
        merged[k] = arr1[i]  
        mergeRecursive(arr1, size1, arr2, size2, merged, i + 1, j, k + 1)  
  
    else: // arr2 has smaller or equal element  
        merged[k] = arr2[j]  
        mergeRecursive(arr1, size1, arr2, size2, merged, i, j + 1, k + 1)
```





02 Source code

→ GitHub Link :

03 Sample of output

The screenshot displays the Code::Blocks IDE with a C++ project named 'main[1].cpp'. The source code implements a recursive merge sort algorithm. The output window shows the program's execution, where two arrays are input, sorted, and then merged. The 5th element of the merged array is identified as 6.

```
#include <iostream>
using namespace std;

void mergeRecursive(int arr1[], int size1, int arr2[], int size2,
int merged[], int i = 0, int j = 0, int k = 0) {
    // Base case
    if (i == size1 && j == size2) {
        return;
    }

    if (i < size1 && (j >= size2 || arr1[i] < arr2[j])) {
        merged[k] = arr1[i];
        i++;
    } else {
        merged[k] = arr2[j];
        j++;
    }
    k++;
    mergeRecursive(arr1, size1, arr2, size2, merged, i, j, k);
}

int main() {
    int size1, size2, k;
    int arr1[100], arr2[100], merged[100];

    cout << "Enter the size of Array 1: ";
    cin >> size1;
    cout << "Enter the elements of Array 1 'sorted': ";
    for (int i = 0; i < size1; i++) {
        cin >> arr1[i];
    }

    cout << "Enter the size of Array 2: ";
    cin >> size2;
    cout << "Enter the elements of Array 2 'sorted': ";
    for (int i = 0; i < size2; i++) {
        cin >> arr2[i];
    }

    k = 0;
    mergeRecursive(arr1, size1, arr2, size2, merged, 0, 0, 0);

    cout << "Enter k position: ";
    int k_pos;
    cin >> k_pos;
    cout << "The " << k_pos << "-th element in the merged array is: " << merged[k_pos] << endl;

    return 0;
}
```

Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
2
3
6
7
9
Enter the size of Array 2: 4
Enter the elements of Array 2 'sorted':
1
4
8
10
Enter k position: 5
The 5-th element in the merged array is: 6
Process returned 0 (0x0) execution time : 20.121 s
Press any key to continue.





04 Time Complexity

Best Case: $O(m+n)$ – When all elements need to be processed

Average Case: $O(m+n)$

Worst Case: $O(m+n)$

05 Space Complexity

$O(m+n)$ for the merged array

$O(m+n)$ for recursion stack in worst case

(though tail recursion optimization may reduce this)





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Advantages

- 1- Simple recursive implementation
- 2- Clearly demonstrates the divide-and-conquer approach
- 3- No need for explicit loop management

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Disadvantages

- 1- Recursion overhead for large arrays
- 2- Potential stack overflow for very large arrays
- 3- Slightly more difficult to understand than iterative version



1. Pseudo-code

Bubble Sort (RECURSIVE APPROACH)

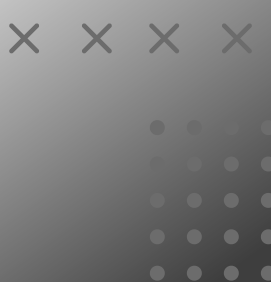
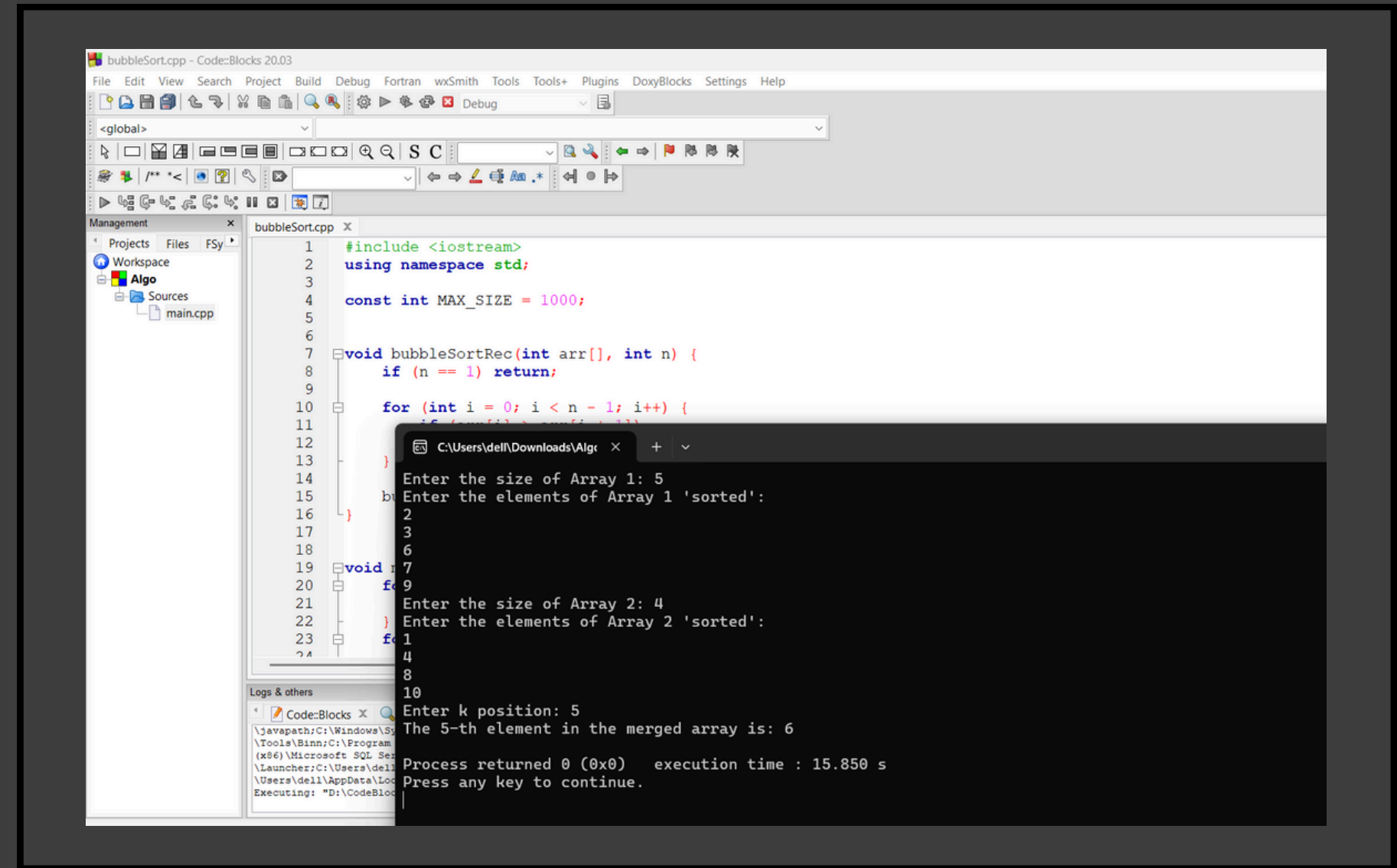
```
1  function bubbleSortRec(arr, n):  
2  ✓    if n == 1:  
3        return  
4  
5        // Perform one full pass of bubble sort  
6        for i from 0 to n - 2:  
7  ✓        if arr[i] > arr[i + 1]:  
8            swap(arr[i], arr[i + 1])  
9  
10       // Recurse on the remaining array  
11       bubbleSortRec(arr, n - 1)  
12  end function  
13
```



02 Source code

→ GitHub Link :

03 Sample of output





04 Time Complexity

Worst-case: $O(n^2)$ – When the array is sorted in reverse order

Best-case: $O(n)$ – When the array is already sorted (with optimized version)

Average-case: $O(n^2)$

05 Space Complexity

$O(n)$





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Advantages

- 1- Simple to understand and implement
- 2- Stable sorting algorithm
- 3- In-place sorting

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Disadvantages

- 1- Consistently poor performance (even on sorted input)
- 2- No practical advantages over other sorts
- 3- Recursion provides no benefit over iterative version
- 4- Least efficient of all four approaches



iINSERTION Sort (RECURSIVE APPROACH)

1. Pseudo-code

```
1  function insertionSortRec(arr, n):  
2  ✓    if n <= 1:  
3        return  
4  
5        insertionSortRec(arr, n-1) // Sort first n-1 elements  
6  
7        last = arr[n-1]           // Last element to be inserted  
8        j = n - 2  
9  
10 ✓    while j >= 0 and arr[j] > last:  
11        arr[j + 1] = arr[j]  
12        j = j - 1  
13  
14        arr[j + 1] = last  
15    end function  
16
```



02 Source code

→ [GitHub Link](#) :

03 Sample of output

```
1 #include <iostream>
2 using namespace std;
3
4 const int MAX_SIZE = 1000;
5
6
7 void insertionSortRec(int arr[], int n) {
8     if (n <= 1) return;
9
10    inse
11
12    int
13    int
14
15    while
16
17
18    }
19
20    arr[
21
22
23
24 void main() {
25     int n1, n2, k;
26     int arr1[MAX_SIZE], arr2[MAX_SIZE], mergedArray[MAX_SIZE];
27
28     cout << "Enter the size of Array 1: ";
29     cin >> n1;
30     cout << "Enter the elements of Array 1 'sorted': ";
31     for (int i = 0; i < n1; i++) {
32         cin >> arr1[i];
33     }
34
35     cout << "Enter the size of Array 2: ";
36     cin >> n2;
37     cout << "Enter the elements of Array 2 'sorted': ";
38     for (int i = 0; i < n2; i++) {
39         cin >> arr2[i];
40     }
41
42     cout << "Enter k position: ";
43     cin >> k;
44
45     insertionSortRec(arr1, n1);
46     insertionSortRec(arr2, n2);
47
48     merge(arr1, 0, n1-1, arr2, 0, n2-1, mergedArray, k);
49
50     cout << "The " << k << "-th element in the merged array is: " << mergedArray[k-1] << endl;
51
52     return 0;
53 }
```

Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
2
3
6
7
9
Enter the size of Array 2: 4
Enter the elements of Array 2 'sorted':
1
4
8
10
Enter k position: 5
The 5-th element in the merged array is: 6
Process returned 0 (0x0) execution time : 25.513 s
Press any key to continue.





04 Time Complexity

Worst-case: $O(n^2)$ – When the array is sorted in reverse order

Best-case: $O(n)$ – When the array is already sorted

Average-case: $O(n^2)$

05 Space Complexity

$O(n)$





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Advantages

- 1- Efficient for small or nearly-sorted arrays
- 2- Stable algorithm (preserves order of equal elements)
- 3- In-place sorting (no additional space needed beyond recursion stack)

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Disadvantages

- 1- Quadratic time makes it impractical for large arrays
- 2- Poor performance on reverse-sorted input
- 3- Recursion adds unnecessary overhead vs iterative version



PERFORMANCE COMPARISON



Algorithm

Time
Complexity

Space
Complexity

Best For

Non-recursive Merge

$O(n)$

$O(1)$

Large sorted arrays

Recursive Merge

$O(n)$

$O(n)$

Medium sorted arrays

Recursive Insertion Sort

$O(n^2)$

$O(n)$

Small arrays

Recursive Bubble Sort

$O(n^2)$

$O(n)$

Educational purposes



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Conclusion

The comparison clearly shows that when merging two already sorted arrays, the merge approaches (especially the non-recursive one) are vastly superior to the sorting approaches. The sorting approaches become useful only when the input arrays aren't guaranteed to be sorted, though even then more efficient sorting algorithms than insertion or bubble sort would typically be preferred.

This exercise demonstrates how algorithm choice dramatically affects performance, with the merge approaches being $O(n)$ while the sorting approaches are $O(n^2)$ for this problem.



THANK YOU

