# ALGORITHMS TASK

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# Task Number 7 K-th Element of Two Sorted Arrays

# **NAMES**

روان ربیع عباس محمد20230069اروي محمد ابراهیم فهیمجانا اشرف ابوسریع مصطف20230419محمد محمد محمد محمدعبدالرحمن محمد احمد شفیق20230318

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# Introduction

#### **Problem Statement**

Given two sorted arrays of sizes m and n, find the element that would be at the k-th position in their merged sorted array.

\*\*Examples\*\*:

Input:

Array 1 = 
$$[2, 3, 6, 7, 9]$$
  
Array2 =  $[1, 4, 8, 10]$   
 $k = 5$ 

Output: 6 (Merged array: [1, 2, 3, 4, 6, 7, 8, 9, 10])







# PROBLEMS SOLUTION

Solution 1

NON-RECURSIVE APPROACH

Solution 2

Binary Search (RECURSIVE APPROACH)

Solution 3

Bubble Sort (RECURSIVE APPROACH)

Solution 4

Insertion Sort (RECURSIVE APPROACH)





# NON-RECURSIVE APPROACH



```
function NonRec_Fun(arr1, size1, arr2, size2, k):
   i \leftarrow 0, j \leftarrow 0, count \leftarrow 0
    while i < size1 and j < size2:
        if arr1[i] < arr2[j]:</pre>
             element ← arr1[i]
             i \leftarrow i + 1
        else:
             element ← arr2[j]
             j \leftarrow j + 1
         count ← count + 1
        if count == k:
             return element
    while i ≺ size1:
        element ← arr1[i]
        i \leftarrow i + 1
        count ← count + 1
        if count == k:
             return element
    while j < size2:
        element ← arr2[j]
        j ← j + 1
        count ← count + 1
        if count == k:
             return element
    return -1 // k is out of bounds
```





## 92 Source code

#### → GitHub Link:

O3 Sample of output

```
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Projects Files FSy
                            #include <iostream>

    ₩orkspace

                           using namespace std;
 - - Algo
                           const int MAX SIZE = 1000;
 int Non! 🖾 C:\Users\dell\Downloads\Alga × + v
                                    Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
                      10
                      11
                      12
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                                    Enter the size of Array 2: 4
Enter the elements of Array 2 'sorted':
                      14
                      15
                      16
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                                    Enter k position: 5
                      20
                                whire The 5-th element in the merged array is: 6
                      21
                     23
                                    Process returned 0 (0x0) execution time : 27.543 \text{ s}
                                    Press any key to continue.
```





# O4 Time Complexity

Best / Worst / Average Case: O(k)

O5 Space Complexity

0(1)





# 

- 1- Simple and straightforward implementation
- 2- Easy to understand and verify correctness
- 3- No additional memory allocation needed
- 4- Handles edge cases well (empty arrays, k out of bounds)

## 07 Disadvantages

- 1- Not optimal for large arrays when k is large
- 2- Linear time complexity could be improved with a binary search approach





# Binary Search (RECURSIVE APPROACH)



```
function mergeRecursive(arr1, size1, arr2, size2, merged, i = 0, j = 0, k = 0):
    if i == size1 AND j == size2: // Base case: both arrays fully processed
        return
   if i == size1: // arr1 exhausted, take from arr2
       merged[k] = arr2[j]
        mergeRecursive(arr1, size1, arr2, size2, merged, i, j + 1, k + 1)
   else if j == size2: // arr2 exhausted, take from arr1
       merged[k] = arr1[i]
        mergeRecursive(arr1, size1, arr2, size2, merged, i + 1, j, k + 1)
   else if arr1[i] < arr2[j]: // arr1 has smaller element
        merged[k] = arr1[i]
        mergeRecursive(arr1, size1, arr2, size2, merged, i + 1, j, k + 1)
   else: // arr2 has smaller or equal element
        merged[k] = arr2[j]
```

mergeRecursive(arr1, size1, arr2, size2, merged, i, j + 1, k + 1)





## 92 Source code

#### → GitHub Link:

# O3 Sample of output

```
e Edit View Search Project Build Debug Fortran wxSmith Tools Tools+ Plugins DoxyBlocks Settings Help
                              Management × ☐ main[1].cpp X
                    1 #include <iostream>

    ₩orkspace

                        using namespace std;
Algo
 ■ Sources
                    4 □void mergeRecursive(int arr1[], int size1, int arr2[], int size2,
                                         int merged[], int i = 0, int j = 0, int k = 0) {
                            // Base case
                           if (i == sizel && j == size2) {
                               return;
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                   11
12
                              Enter the size of Array 1: 5
                              Enter the elements of Array 1 'sorted':
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22
23
                              Enter the size of Array 2: 4
                              Enter the elements of Array 2 'sorted':
                             Enter k position: 5
                              The 5-th element in the merged array is: 6
```





# **Time Complexity**

Best Case: O(m+n) - When all elements need to be processed

Average Case: O(m+n)

Worst Case: O(m+n)

# 5 Space Complexity

O(m+n) for the merged array
O(m+n) for recursion stack in worst case
(though tail recursion optimization may reduce this)





# 

- 1- Simple recursive implementation
- 2-Clearly demonstrates the divide-and-conquer approach
- 3- No need for explicit loop management

## O7 Disadvantages

- 1- Recursion overhead for large arrays
- 2- Potential stack overflow for very large arrays
- 3- Slightly more difficult to understand than iterative version





#### 1. Pseudo-code

# Bubble Sort (RECURSIVE APPROACH)

```
function bubbleSortRec(arr, n):

function bubbleSortRec(arr, n):

function bubbleSortRec(arr, n):

function bubbleSortRec(arr, n):

return

// Perform one full pass of bubble sort

for i from 0 to n - 2:

function 0 to n - 2:

with arr[i] > arr[i + 1]:

www.arr[i] > arr[i + 1])

// Recurse on the remaining array

bubbleSortRec(arr, n - 1)

end function

description:
```





O2 Source code

→ GitHub Link:

O3 Sample of output

```
Projects Files FSy
                                                                                                        1 #include <iostream>
using namespace std;
                                                                                                                             const int MAX_SIZE = 1000;
                                                                                                                       pvoid bubbleSortRec(int arr[], int n) {
                                                                                                                                              if (n == 1) return;
                                                                                                                                                 for (int i = 0; i < n - 1; i++) {
                                                                                                                                                                 © C:\Users\dell\Downloads\Alga × + v
                                                                                                                                                        Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
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                                                                                                    21
                                                                                                                                                          Enter the size of Array 2: 4
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23
                                                                                                                                                            Enter the elements of Array 2 'sorted':
                                                                               Code:Blocks X Q Enter k position: 5

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\Tools\Binn;C:\Program
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                                                                               (x86)\Microsoft SQL Se:\Launcher;C:\U3ers\dell\\Process returned 0 (0x0) execution time : 15.850 s\U3ers\dell\AppBatalor\\Press any key to continue.
```





# O4 Time Complexity

Worst-case: O(n²) - When the array is sorted in reverse order

Best-case: O(n) - When the array is already sorted (with optimized version)

Average-case: O(n²)

O5 Space Complexity

O(n)





## 

- 1- Simple to understand and implement
- 2- Stable sorting algorithm
- 3- In-place sorting

## O7 Disadvantages

- 1- Consistently poor performance (even on sorted input)
- 2- No practical advantages over other sorts
- 3- Recursion provides no benefit over iterative version
- 4- Least efficient of all four approaches





# insertion sort (RECURSIVE APPROACH)



```
function insertionSortRec(arr, n):
         if n <= 1:
             return
         insertionSortRec(arr, n-1) // Sort first n-1 elements
         last = arr[n-1]
                                     // Last element to be inserted
         j = n - 2
         while j \ge 0 and arr[j] > last:
10 🗸
             arr[j + 1] = arr[j]
11
             j = j - 1
12
13
         arr[j + 1] = last
14
     end function
15
16
```





92 Source code

→ GitHub Link:

O3 Sample of output

```
1 #include <iostream>

    Workspace

                        using namespace std;
- Algo
 const int MAX SIZE = 1000;
                       void insertionSortRec(int arr[], int n) {
                            inse. C:\Users\dell\Downloads\Alga × + ~
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20
21
22
23
                               Enter the size of Array 1: 5
Enter the elements of Array 1 'sorted':
                               Enter the size of Array 2: 4
Enter the elements of Array 2 'sorted':
                                 Enter k position: 5
                                The 5-th element in the merged array is: 6
               \Users\dell\AppData\Local\
Executing: "D:\CodeBlocks/
```





# 04 Time Complexity

Worst-case: O(n²) - When the array is sorted in reverse order Best-case: O(n) - When the array is already sorted

Average-case: O(n²)

O5 Space Complexity

O(n)





## 

- 1- Efficient for small or nearly-sorted arrays
- 2- Stable algorithm (preserves order of equal elements)
- 3- In-place sorting (no additional space needed beyond recursion stack)

#### O7 Disadvantages

- 1- Quadratic time makes it impractical for large arrays
- 2- Poor performance on reverse-sorted input
- 3- Recursion adds unnecessary overhead vs iterative version



#### PERFORMANCE COMPARISON



Algorithm

Time Complexity Space Complexity

**Best For** 

Non-recursive Merge

O(n)

0(1)

Large sorted arrays

Recursive Merge

O(n)

O(n)

Medium sorted arrays

Recursive Insertion Sort

 $O(n^2)$ 

O(n)

Small arrays

Recursive Bubble Sort

 $O(n^2)$ 

O(n)

Educational purposes





#### Conclusion

The comparison clearly shows that when merging two already sorted arrays, the merge approaches (especially the non-recursive one) are vastly superior to the sorting approaches. The sorting approaches become useful only when the input arrays aren't guaranteed to be sorted, though even then more efficient sorting algorithms than insertion or bubble sort would typically be preferred. This exercise demonstrates how algorithm choice dramatically affects performance, with the merge approaches being O(n) while the sorting approaches are O(n²) for this problem.







# THANK YOU

