

## C interfaces to GALAHAD LPB

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## **Chapter 1**

# GALAHAD C package lpb

## 1.1 Introduction

## 1.1.1 Purpose

This package uses a primal-dual interior-point method to solve the linear programming problem

minimize 
$$q(x) = g^T x + f$$

subject to the general linear constraints

$$c_i^l \le a_i^T x \le c_i^u, \quad i = 1, \dots, m,$$

and the simple bound constraints

$$x_j^l \le x_j \le x_j^u, \quad j = 1, \dots, n,$$

where the vectors  $g, w, x^0, a_i, c^l, c^u, x^l, x^u$  and the scalar f are given. Any of the constraint bounds  $c_i^l, c_i^u, x_j^l$  and  $x_j^u$  may be infinite. Full advantage is taken of any zero coefficients in the matrix A whose rows are the transposes of the vectors  $a_i$ .

#### 1.1.2 Authors

N. I. M. Gould, STFC-Rutherford Appleton Laboratory, England.

C interface, additionally J. Fowkes, STFC-Rutherford Appleton Laboratory.

### 1.1.3 Originally released

August 2018, C interface September 2021.

### 1.1.4 Terminology

The required solution x necessarily satisfies the primal optimality conditions

$$(1a) Ax = c$$

and

$$(1b) c^l \le c \le c^u, \quad x^l \le x \le x^u,$$

the dual optimality conditions

(2a) 
$$q = A^T y + z$$

where

(2b) 
$$y = y^l + y^u, z = z^l + z^u, y^l \ge 0, y^u \le 0, z^l \ge 0 \text{ and } z^u \le 0,$$

and the complementary slackness conditions

(3) 
$$(Ax - c^l)^T y^l = 0$$
,  $(Ax - c^u)^T y^u = 0$ ,  $(x - x^l)^T z^l = 0$  and  $(x - x^u)^T z^u = 0$ ,

where the vectors y and z are known as the Lagrange multipliers for the general linear constraints, and the dual variables for the bounds, respectively, and where the vector inequalities hold component-wise.

#### 1.1.5 Method

Primal-dual interior point methods iterate towards a point that satisfies these conditions by ultimately aiming to satisfy (1a), (2a) and (3), while ensuring that (1b) and (2b) are satisfied as strict inequalities at each stage. Appropriate norms of the amounts by which (1a), (2a) and (3) fail to be satisfied are known as the primal and dual infeasibility, and the violation of complementary slackness, respectively. The fact that (1b) and (2b) are satisfied as strict inequalities gives such methods their other title, namely interior-point methods.

The method aims at each stage to reduce the overall violation of (1a), (2a) and (3), rather than reducing each of the terms individually. Given an estimate  $v=(x,c,y,y^l,y^u,z,z^l,z^u)$  of the primal-dual variables, a correction  $\Delta v=\Delta(x,c,y,y^l,y^uz,z^l,z^u)$  is obtained by solving a suitable linear system of Newton equations for the nonlinear systems (1a), (2a) and a parameterized `residual trajectory" perturbation of (3); residual trajectories proposed by Zhang (1994) and Zhao and Sun (1999) are possibilities. An improved estimate  $v+\alpha\Delta v$  is then used, where the step-size  $\alpha$  is chosen as close to 1.0 as possible while ensuring both that (1b) and (2b) continue to hold and that the individual components which make up the complementary slackness (3) do not deviate too significantly from their average value. The parameter that controls the perturbation of (3) is ultimately driven to zero.

The Newton equations are solved by applying the GALAHAD matrix factorization package SBLS, but there are options to factorize the matrix as a whole (the so-called "augmented system" approach), to perform a block elimination first (the "Schur-complement" approach), or to let the method itself decide which of the two previous options is more appropriate. The "Schur-complement" approach is usually to be preferred when all the weights are nonzero or when every variable is bounded (at least one side), but may be inefficient if any of the columns of A is too dense.

Optionally, the problem may be pre-processed temporarily to eliminate dependent constraints using the GALAHAD package FDC. This may improve the performance of the subsequent iteration.

#### 1.1.6 Reference

The basic algorithm is a generalisation of those of

Y. Zhang (1994), On the convergence of a class of infeasible interior-point methods for the horizontal linear complementarity problem, SIAM J. Optimization 4(1) 208-227,

and

G. Zhao and J. Sun (1999). On the rate of local convergence of high-order infeasible path-following algorithms for the  $P_*$  linear complementarity problems, Computational Optimization and Applications 14(1) 293-307,

with many enhancements described by

N. I. M. Gould, D. Orban and D. P. Robinson (2013). Trajectory-following methods for large-scale degenerate convex quadratic programming, Mathematical Programming Computation 5(2) 113-142.

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#### 1.1.7 Call order

To solve a given problem, functions from the lpb package must be called in the following order:

- lpb initialize provide default control parameters and set up initial data structures
- · lpb read specfile (optional) override control values by reading replacement values from a file
- lpb import set up problem data structures and fixed values
- lpb reset control (optional) possibly change control parameters if a sequence of problems are being solved
- lpb solve lp solve the linear program
- lpb information (optional) recover information about the solution and solution process
- lpb\_terminate deallocate data structures

See Section 4.1 for examples of use.

#### 1.1.8 Unsymmetric matrix storage formats

The unsymmetric m by n constraint matrix A may be presented and stored in a variety of convenient input formats.

Both C-style (0 based) and fortran-style (1-based) indexing is allowed. Choose control.f\_indexing as false for C style and true for fortran style; the discussion below presumes C style, but add 1 to indices for the corresponding fortran version.

Wrappers will automatically convert between 0-based (C) and 1-based (fortran) array indexing, so may be used transparently from C. This conversion involves both time and memory overheads that may be avoided by supplying data that is already stored using 1-based indexing.

#### 1.1.8.1 Dense storage format

The matrix A is stored as a compact dense matrix by rows, that is, the values of the entries of each row in turn are stored in order within an appropriate real one-dimensional array. In this case, component n\*i+j of the storage array A\_val will hold the value  $A_{ij}$  for  $0 \le i \le m-1$ ,  $0 \le j \le n-1$ .

#### 1.1.8.2 Sparse co-ordinate storage format

Only the nonzero entries of the matrices are stored. For the l-th entry,  $0 \le l \le ne-1$ , of A, its row index i, column index j and value  $A_{ij}$ ,  $0 \le i \le m-1$ ,  $0 \le j \le n-1$ , are stored as the l-th components of the integer arrays A\_row and A\_col and real array A\_val, respectively, while the number of nonzeros is recorded as A\_ne = ne.

#### 1.1.8.3 Sparse row-wise storage format

Again only the nonzero entries are stored, but this time they are ordered so that those in row i appear directly before those in row i+1. For the i-th row of A the i-th component of the integer array A\_ptr holds the position of the first entry in this row, while A\_ptr(m) holds the total number of entries plus one. The column indices j,  $0 \le j \le n-1$ , and values  $A_{ij}$  of the nonzero entries in the i-th row are stored in components I = A\_ptr(i), ..., A\_ptr(i+1)-1,  $0 \le i \le m-1$ , of the integer array A\_col, and real array A\_val, respectively. For sparse matrices, this scheme almost always requires less storage than its predecessor.

# **Chapter 2**

# File Index

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Here is a list of all files with brief descriptions:	
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## **Chapter 3**

## **File Documentation**

## 3.1 lpb.h File Reference

```
#include <stdbool.h>
#include "galahad_precision.h"
#include "fdc.h"
#include "sbls.h"
#include "fit.h"
#include "roots.h"
#include "cro.h"
#include "rpd.h"
```

#### **Data Structures**

- struct lpb\_control\_type
- struct lpb\_time\_type
- struct lpb\_inform\_type

#### **Functions**

- void lpb\_initialize (void \*\*data, struct lpb\_control\_type \*control, int \*status)
- void lpb\_read\_specfile (struct lpb\_control\_type \*control, const char specfile[])
- void lpb\_import (struct lpb\_control\_type \*control, void \*\*data, int \*status, int n, int m, const char A\_type[], int A\_ne, const int A\_row[], const int A\_col[], const int A\_ptr[])
- void lpb\_reset\_control (struct lpb\_control\_type \*control, void \*\*data, int \*status)
- void lpb\_solve\_lp (void \*\*data, int \*status, int n, int m, const real\_wp\_ g[], const real\_wp\_ f, int a\_ne, const real\_wp\_ A\_val[], const real\_wp\_ c\_l[], const real\_wp\_ c\_u[], const real\_wp\_ x\_l[], const real\_wp\_ x\_l[], real\_wp\_ x[], real\_wp\_ c[], real\_wp\_ z[], int x\_stat[], int c\_stat[])
- void lpb\_information (void \*\*data, struct lpb\_inform\_type \*inform, int \*status)
- void lpb\_terminate (void \*\*data, struct lpb\_control\_type \*control, struct lpb\_inform\_type \*inform)

#### 3.1.1 Data Structure Documentation

#### 3.1.1.1 struct lpb\_control\_type

control derived type as a C struct

**Examples** 

```
lpbt.c, and lpbtf.c.
```

## Data Fields

bool	f_indexing	use C or Fortran sparse matrix indexing
int	error	error and warning diagnostics occur on stream
		error
int	out	general output occurs on stream out
int	print_level	the level of output required is specified by print_level
int	start_print	any printing will start on this iteration
int	stop_print	any printing will stop on this iteration
int	maxit	at most maxit inner iterations are allowed
int	infeas_max	the number of iterations for which the overall infeasibility of the problem is not reduced by at least a factor .reduce_infeas before the problem is flagged as infeasible (see reduce_infeas)
int	muzero_fixed	the initial value of the barrier parameter will not be changed for the first muzero_fixed iterations
int	restore_problem	indicate whether and how much of the input problem should be restored on output. Possible values are
		0 nothing restored     1 scalar and vector parameters
		2 all parameters
		_ all parameters
int	indicator_type	specifies the type of indicator function used. Pssible values are
		<ul> <li>1 primal indicator: constraint active if and only if distance to nearest bound &lt;= .indicator_p_tol</li> </ul>
		<ul> <li>2 primal-dual indicator: constraint active if and only if distance the nearest bound &lt;= .indicator_tol_pd * size of corresponding multiplier</li> </ul>
		3 primal-dual indicator: constraint active if and only if distance to the nearest bound <= .indicator_tol_tapia * distance to same bound at previous iteration
int	arc	which residual trajectory should be used to aim from the current iteration to the solution
		1 the Zhang linear residual trajectory
		2 the Zhao-Sun quadratic residual trajectory
		3 the Zhang arc ultimately switching to the Zhao-Sun residual trajectory
		4 the mixed linear-quadratic residual trajectory
int	series_order	the order of (Taylor/Puiseux) series to fit to the path data

## Data Fields

int	sif_file_device	specifies the unit number to write generated SIF
		file describing the current problem
int	qplib_file_device	specifies the unit number to write generated QPLIB file describing the current problem
real_wp_	infinity	any bound larger than infinity in modulus will be regarded as infinite
real_wp_	stop_abs_p	the required absolute and relative accuracies for the primal infeasibili
real_wp_	stop_rel_p	see stop_abs_p
real_wp_	stop_abs_d	the required absolute and relative accuracies for the dual infeasibility
real_wp_	stop_rel_d	see stop_abs_d
real_wp_	stop_abs_c	the required absolute and relative accuracies for the complementarity
real_wp_	stop_rel_c	see stop_abs_c
real_wp_	prfeas	initial primal variables will not be closer than prfeas from their bound
real_wp_	dufeas	initial dual variables will not be closer than dufeas from their bounds
real_wp_	muzero	the initial value of the barrier parameter. If muzero is not positive, it will be reset to an appropriate value
real_wp_	tau	the weight attached to primal-dual infeasibility compared to complementa when assessing step acceptance
real_wp_	gamma_c	individual complementarities will not be allowed to be smaller than gamma_c times the average value
real_wp_	gamma_f	the average complementarity will not be allowed to be smaller than gamma_f times the primal/dual infeasibility
real_wp_	reduce_infeas	if the overall infeasibility of the problem is not reduced by at least a factor reduce_infeas over .infeas_max iterations, the problem is flagged as infeasible (see infeas_max)
real_wp_	obj_unbounded	if the objective function value is smaller than obj_unbounded, it will be flagged as unbounded from below.
real_wp_	potential_unbounded	if W=0 and the potential function value is smaller than potential_unbounded * number of one-sided bounds, the analytic center will be flagged as unbounded
real_wp_	identical_bounds_tol	any pair of constraint bounds (c_l,c_u) or (x_l,x_u) that are closer than identical_bounds_tol will be reset to the average of their values
real_wp_	mu_lunge	start terminal extrapolation when mu reaches mu_lunge
real_wp_	indicator_tol_p	if .indicator_type = 1, a constraint/bound will be deemed to be active if and only if distance to nearest bound <= .indicator_p_tol

## **Data Fields**

real_wp_	indicator_tol_pd	if .indicator_type = 2, a constraint/bound will be deemed to be active if and only if distance to nearest bound <= .indicator_tol_pd * size of corresponding multiplier
real_wp_	indicator_tol_tapia	if .indicator_type = 3, a constraint/bound will be deemed to be active if and only if distance to nearest bound <= .indicator_tol_tapia * distance to same bound at previous iteration
real_wp_	cpu_time_limit	the maximum CPU time allowed (-ve means infinite)
real_wp_	clock_time_limit	the maximum elapsed clock time allowed (-ve means infinite)
bool	remove_dependencies	the equality constraints will be preprocessed to remove any linear dependencies if true
bool	treat_zero_bounds_as_general	any problem bound with the value zero will be treated as if it were a general value if true
bool	just_feasible	if .just_feasible is true, the algorithm will stop as soon as a feasible point is found. Otherwise, the optimal solution to the problem will be found
bool	getdua	if .getdua, is true, advanced initial values are obtained for the dual variables
bool	puiseux	decide between Puiseux and Taylor series approximations to the arc
bool	every_order	try every order of series up to series_order?
bool	feasol	if .feasol is true, the final solution obtained will be perturbed so tha variables close to their bounds are moved onto these bounds
bool	balance_initial_complentarity	if .balance_initial_complentarity is true, the initial complemetarity is required to be balanced
bool	crossover	if .crossover is true, cross over the solution to one defined by linearly-independent constraints if possible
bool	space_critical	if .space_critical true, every effort will be made to use as little space as possible. This may result in longer computation time
bool	deallocate_error_fatal	if .deallocate_error_fatal is true, any array/pointer deallocation error will terminate execution. Otherwise, computation will continue
bool	generate_sif_file	if .generate_sif_file is .true. if a SIF file describing the current problem is to be generated
bool	generate_qplib_file	if .generate_qplib_file is .true. if a QPLIB file describing the current problem is to be generated
char	sif_file_name[31]	name of generated SIF file containing input problem
char	qplib_file_name[31]	name of generated QPLIB file containing input problem
char	prefix[31]	all output lines will be prefixed by .prefix(2:LEN(TRIM(.prefix))-1) where .prefix contains the required string enclosed in quotes, e.g. "string" or 'string'

### **Data Fields**

struct fdc_control_type	fdc_control	control parameters for FDC
struct sbls_control_type	sbls_control	control parameters for SBLS
struct fit_control_type	fit_control	control parameters for FIT
struct roots_control_type	roots_control	control parameters for ROOTS
struct cro_control_type	cro_control	control parameters for CRO

## 3.1.1.2 struct lpb\_time\_type

time derived type as a C struct

### **Data Fields**

real_wp_	total	the total CPU time spent in the package
real_wp_	preprocess	the CPU time spent preprocessing the problem
real_wp_	find_dependent	the CPU time spent detecting linear dependencies
real_wp_	analyse	the CPU time spent analysing the required matrices prior to factorization
real_wp_	factorize	the CPU time spent factorizing the required matrices
real_wp_	solve	the CPU time spent computing the search direction
real_wp_	clock_total	the total clock time spent in the package
real_wp_	clock_preprocess	the clock time spent preprocessing the problem
real_wp_	clock_find_dependent	the clock time spent detecting linear dependencies
real_wp_	clock_analyse	the clock time spent analysing the required matrices prior to factorization
real_wp_	clock_factorize	the clock time spent factorizing the required matrices
real_wp_	clock_solve	the clock time spent computing the search direction

## 3.1.1.3 struct lpb\_inform\_type

inform derived type as a C struct

## Examples

lpbt.c, and lpbtf.c.

#### **Data Fields**

int	status	return status. See LPB_solve for details
int	alloc_status	the status of the last attempted
		allocation/deallocation
char	bad_alloc[81]	the name of the array for which an
		allocation/deallocation error ocurred
int	iter	the total number of iterations required
int	factorization_status	the return status from the factorization
int	factorization_integer	the total integer workspace required for the
		factorization
int	factorization_real	the total real workspace required for the
		factorization

### **Data Fields**

int	nfacts	the total number of factorizations performed
int	nbacts	the total number of "wasted" function evaluations during the linesearch
int	threads	the number of threads used
real_wp_	obj	the value of the objective function at the best estimate of the solution determined by LPB_solve
real_wp_	primal_infeasibility	the value of the primal infeasibility
real_wp_	dual_infeasibility	the value of the dual infeasibility
real_wp_	complementary_slackness	the value of the complementary slackness
real_wp_	init_primal_infeasibility	these values at the initial point (needed bg GALAHAD_CLPB)
real_wp_	init_dual_infeasibility	see init_primal_infeasibility
real_wp_	init_complementary_slackness	see init_primal_infeasibility
real_wp_	potential	the value of the logarithmic potential function sum -log(distance to constraint boundary)
real_wp_	non_negligible_pivot	the smallest pivot which was not judged to be zero when detecting linear dependent constraints
bool	feasible	is the returned "solution" feasible?
int	checkpointsIter[16]	checkpoints(i) records the iteration at which the criticality measures first fall below $10^{-i}$ , i = 1,, 16 (-1 means not achieved)
real_wp_	checkpointsTime[16]	see checkpointsIter
struct lpb_time_type	time	timings (see above)
struct fdc_inform_type	fdc_inform	inform parameters for FDC
struct sbls_inform_type	sbls_inform	inform parameters for SBLS
struct fit_inform_type	fit_inform	return information from FIT
struct roots_inform_type	roots_inform	return information from ROOTS
struct cro_inform_type	cro_inform	inform parameters for CRO
struct rpd_inform_type	rpd_inform	inform parameters for RPD

## 3.1.2 Function Documentation

## 3.1.2.1 lpb\_initialize()

Set default control values and initialize private data

#### **Parameters**

Tilly ode data Holde pirrate internal data	in,out	data	holds private internal data
--	--------	------	-----------------------------

#### **Parameters**

out	control	is a struct containing control information (see lpb_control_type)
out	status	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently):
		0. The import was succesful.

#### Examples

lpbt.c, and lpbtf.c.

### 3.1.2.2 lpb\_read\_specfile()

Read the content of a specification file, and assign values associated with given keywords to the corresponding control parameters

#### **Parameters**

in,out	control	is a struct containing control information (see lpb_control_type)
in	specfile	is a character string containing the name of the specification file

### 3.1.2.3 lpb\_import()

Import problem data into internal storage prior to solution.

### **Parameters**

in	control	is a struct whose members provide control paramters for the remaining prcedures (see <a href="https://linear.com/lpb_control_type">lpb_control_type</a> )
in,out	data	holds private internal data
in, out	status	<ul> <li>is a scalar variable of type int, that gives the exit status from the package. Possible values are:</li> <li>0. The import was succesful</li> <li>-1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>-2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>-3. The restrictions n &gt; 0 or m &gt; 0 or requirement that A_type contains its relevant string 'dense', 'coordinate' or 'sparse_by_rows' has been violated.</li> </ul>
in	n	is a scalar variable of type int, that holds the number of variables.
in	m	is a scalar variable of type int, that holds the number of general linear constraints.
in	A_type	is a one-dimensional array of type char that specifies the unsymmetric storage scheme used for the constraint Jacobian, $A$ . It should be one of 'coordinate', 'sparse_by_rows' or 'dense; lower or upper case variants are allowed.
in	A_ne	is a scalar variable of type int, that holds the number of entries in $A$ in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes.
in	A_row	is a one-dimensional array of size A_ne and type int, that holds the row indices of $A$ in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes, and in this case can be NULL.
in	A_col	is a one-dimensional array of size A_ne and type int, that holds the column indices of $A$ in either the sparse co-ordinate, or the sparse row-wise storage scheme. It need not be set when the dense or diagonal storage schemes are used, and in this case can be NULL.
in	A_ptr	is a one-dimensional array of size $n+1$ and type int, that holds the starting position of each row of $A$ , as well as the total number of entries plus one, in the sparse row-wise storage scheme. It need not be set when the other schemes are used, and in this case can be NULL.

## Examples

lpbt.c, and lpbtf.c.

### 3.1.2.4 | lpb\_reset\_control()

Reset control parameters after import if required.

#### **Parameters**

in	control	is a struct whose members provide control paramters for the remaining prcedures (see
		lpb_control_type)
in,out	data	holds private internal data
in,out	status	is a scalar variable of type int, that gives the exit status from the package. Possible values are:
		0. The import was succesful.

### 3.1.2.5 lpb\_solve\_lp()

```
void lpb_solve_lp (
             void ** data,
             int * status,
             int n_{i}
             int m_{,}
             const real_wp_ g[],
             const real_wp_ f,
             int a_ne,
             const real_wp_ A_val[],
             const real_wp_ c_1[],
             const real_wp_ c_u[],
             const real_wp_ x_1[],
             const real_wp_ x_u[],
             real_wp_ x[],
             real_wp_ c[],
             real_wp_ y[],
             real_wp_ z[],
             int x_stat[],
             int c_stat[] )
```

Solve the linear program.

#### **Parameters**

in, out   <i>data</i>   holds private internal data
---

## **Parameters**

in,out	status	is a scalar variable of type int, that gives the entry and exit status from the package.  Possible exit are:
		0. The run was succesful.
		<ul> <li>-1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> </ul>
		<ul> <li>-2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> </ul>
		<ul> <li>-3. The restrictions n &gt; 0 and m &gt; 0 or requirement that A_type contains its relevant string 'dense', 'coordinate' or 'sparse_by_rows' has been violated.</li> </ul>
		-5. The simple-bound constraints are inconsistent.
		<ul> <li>-7. The constraints appear to have no feasible point.</li> </ul>
		<ul> <li>-9. The analysis phase of the factorization failed; the return status from the factorization package is given in the component inform.factor_status</li> </ul>
		<ul> <li>-10. The factorization failed; the return status from the factorization package is given in the component inform.factor_status.</li> </ul>
		<ul> <li>-11. The solution of a set of linear equations using factors from the factorization package failed; the return status from the factorization package is given in the component inform.factor_status.</li> </ul>
		-16. The problem is so ill-conditioned that further progress is impossible.
		-17. The step is too small to make further impact.
		-18. Too many iterations have been performed. This may happen if control.maxit is too small, but may also be symptomatic of a badly scaled problem.
		<ul> <li>-19. The CPU time limit has been reached. This may happen if control.cpu_time_limit is too small, but may also be symptomatic of a badly scaled problem.</li> </ul>
in	n	is a scalar variable of type int, that holds the number of variables
in	m	is a scalar variable of type int, that holds the number of general linear constraints.
in	g	is a one-dimensional array of size n and type double, that holds the linear term $g$ of the objective function. The j-th component of g, j = 0,, n-1, contains $g_j$ .
in	f	is a scalar of type double, that holds the constant term $f$ of the objective function.
in	a_ne	is a scalar variable of type int, that holds the number of entries in the constraint Jacobian matrix ${\cal A}.$
in	A_val	is a one-dimensional array of size a_ne and type double, that holds the values of the entries of the constraint Jacobian matrix ${\cal A}$ in any of the available storage schemes.
in	c_I	is a one-dimensional array of size m and type double, that holds the lower bounds $c^l$ on the constraints $Ax$ . The i-th component of c_l, i = 0,, m-1, contains $c^l_i$ .
in	c_u	is a one-dimensional array of size m and type double, that holds the upper bounds $c^l$ on the constraints $Ax$ . The i-th component of c_u, i = 0,, m-1, contains $c^u_i$ .
in	x_I	is a one-dimensional array of size n and type double, that holds the lower bounds $x^l$ on the variables $x$ . The j-th component of x_l, j = 0,, n-1, contains $x^l_j$ .
	1	1

### **Parameters**

in	x_u	is a one-dimensional array of size n and type double, that holds the upper bounds $x^l$ on the variables $x$ . The j-th component of x_u, j = 0,, n-1, contains $x^l_j$ .
in,out	Х	is a one-dimensional array of size n and type double, that holds the values $x$ of the optimization variables. The j-th component of x, j = 0,, n-1, contains $x_j$ .
out	С	is a one-dimensional array of size m and type double, that holds the residual $c(x)$ . The i-th component of c, i = 0,, m-1, contains $c_i(x)$ .
in,out	у	is a one-dimensional array of size n and type double, that holds the values $y$ of the Lagrange multipliers for the general linear constraints. The j-th component of y, j = 0,, m-1, contains $y_i$ .
in,out	Z	is a one-dimensional array of size n and type double, that holds the values $z$ of the dual variables. The j-th component of z, j = 0,, n-1, contains $z_j$ .
out	x_stat	is a one-dimensional array of size n and type int, that gives the optimal status of the problem variables. If $x\_stat(j)$ is negative, the variable $x_j$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.
out	c_stat	is a one-dimensional array of size m and type int, that gives the optimal status of the general linear constraints. If c_stat(i) is negative, the constraint value $a_i^T x$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.

## Examples

lpbt.c, and lpbtf.c.

## 3.1.2.6 lpb\_information()

## Provides output information

#### **Parameters**

in,out	data	holds private internal data
out	inform	is a struct containing output information (see lpb_inform_type)
out	status	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently):
		0. The values were recorded succesfully

## Examples

lpbt.c, and lpbtf.c.

## 3.1.2.7 lpb\_terminate()

## Deallocate all internal private storage

#### **Parameters**

in,out	data	holds private internal data
out	control	is a struct containing control information (see lpb_control_type)
out	inform	is a struct containing output information (see lpb_inform_type)

## Examples

lpbt.c, and lpbtf.c.

## **Chapter 4**

## **Example Documentation**

## 4.1 lpbt.c

This is an example of how to use the package to solve a linear program. A variety of supported constraint matrix storage formats are shown.

Notice that C-style indexing is used, and that this is flaggeed by setting control.f\_indexing to false.

```
/* lpbt.c */
/* Full test for the LPB C interface using C sparse matrix indexing \star/
#include <stdio.h>
#include <math.h>
#include "lpb.h"
int main(void) {
     // Derived types
     void *data;
     struct lpb_control_type control;
     struct lpb_inform_type inform;
     // Set problem data
     int n = 3; // dimension int m = 2; // number of general constraints double g[] = \{0.0, 2.0, 0.0\}; // linear term in the objective double f = 1.0; // constant term in the objective int A_ne = 4; // Jacobian elements
     int A_ne = 4; // Jacobian elements
int A_row[] = {0, 0, 1, 1}; // row indices
int A_col[] = {0, 1, 1, 2}; // column indices
int A_ptr[] = {0, 2, 4}; // row pointers
double A_val[] = {2.0, 1.0, 1.0, 1.0 }; // values
double c_l[] = {1.0, 2.0}; // constraint lower bound
double c_u[] = {2.0, 2.0}; // constraint upper bound
     double x_1[] = {-1.0, - INFINITY, - INFINITY}; // variable lower bound
     double x_u[] = \{1.0, INFINITY, 2.0\}; // variable upper bound
     // Set output storage
double c[m]; // constraint values
     int x_stat[n]; // variable status
     int c_stat[m]; // constraint status
     int status;
     for ( int d=1; d <= 3; d++) {
           // Initialize LPB
           lpb_initialize( &data, &control, &status );
           // Set user-defined control options
           control.f_indexing = false; // C sparse matrix indexing
           // Start from 0
           double x[] = \{0.0, 0.0, 0.0\};
double y[] = \{0.0, 0.0\};
           double z[] = \{0.0, 0.0, 0.0\};
           switch(d){
                case 1: // sparse co-ordinate storage
   st = 'C';
                      lpb_import( &control, &data, &status, n, m,
                                       "coordinate", A_ne, A_row, A_col, NULL );
                      lpb_solve_lp( &data, &status, n, m, g, f,
```

```
A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
                    x_stat, c_stat );
   printf(" case %1i break\n",d);
   case 2: // sparse by rows
st = 'R';
       lpb_solve_lp( &data, &status, n, m, g, f,
                   A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
                   x_stat, c_stat );
      break:
   case 3: // dense
st = 'D';
       int A_dense_ne = 6; // number of elements of A
       A_dense_ne, A_dense, c_l, c_u, x_l, x_u,
                    x, c, y, z, x_stat, c_stat);
lpb_information( &data, &inform, &status );
if (inform.status == 0) {
  printf("%c:%6i iterations. Optimal objective value = %5.2f status = %1i\n",
          st, inform.iter, inform.obj, inform.status);
   printf("%c: LPB_solve exit status = %li\n", st, inform.status);
//printf("x: ");
//printf("\n");
//printf("\n");
//printf("gradient: ");
//for( int i = 0; i < n; i++) printf("%f ", g[i]); //printf("\n");
// Delete internal workspace
lpb_terminate( &data, &control, &inform );
```

## 4.2 lpbtf.c

This is the same example, but now fortran-style indexing is used.

```
/* lpbtf.c */
/\star Full test for the LPB C interface using Fortran sparse matrix indexing \star/
#include <stdio.h>
#include <math.h>
#include "lpb.h
int main(void) {
      // Derived types
      void *data:
      struct lpb_control_type control;
      struct lpb_inform_type inform;
      // Set problem data
      int n = 3; // dimension
int m = 2; // number of general constraints
      double g[] = {0.0, 2.0, 0.0}; // linear term in the objective double f = 1.0; // constant term in the objective int A_ne = 4; // Jacobian elements
      int A_row[] = {1, 1, 2, 2}; // row indices
     int A_row[] = {1, 1, 2, 2}; // row indices
int A_col[] = {1, 2, 2, 3}; // column indices
int A_ptr[] = {1, 3, 5}; // row pointers
double A_val[] = {2.0, 1.0, 1.0, 1.0 }; // values
double c_l[] = {1.0, 2.0}; // constraint lower bound
double c_u[] = {2.0, 2.0}; // constraint upper bound
      double x_{\perp}[] = \{-1.0, -\text{INFINITY}, -\text{INFINITY}; // variable lower bound double <math>x_{\perp}[] = \{1.0, -\text{INFINITY}, 2.0\}; // variable upper bound
      // Set output storage
      double c[m]; // constraint values
int x_stat[n]; // variable status
int c_stat[m]; // constraint status
      char st;
      int status;
     for( int d=1; d <= 3; d++) {</pre>
            // Initialize LPB
            lpb_initialize( &data, &control, &status );
            // Set user-defined control options
```

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```
control.f_indexing = true; // Fortran sparse matrix indexing
// Start from 0
double x[] = {0.0,0.0,0.0};
double y[] = {0.0,0.0};
double z[] = {0.0,0.0,0.0};
switch(d){
    case 1: // sparse co-ordinate storage
    st = 'C';
        lpb_import( &control, &data, &status, n, m,
        printf(" case %1i break\n",d);
    case 2: // sparse by rows
st = 'R';
        A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
                      x_stat, c_stat );
    break;
case 3: // dense
st = 'D';
        int A_dense_ne = 6; // number of elements of A double A_dense[] = {2.0, 1.0, 0.0, 0.0, 1.0, 1.0};
        lpb_information( &data, &inform, &status );
if (inform.status == 0) {
    printf("%c:%6i iterations. Optimal objective value = %5.2f status = %1i\n",
           st, inform.iter, inform.obj, inform.status);
   printf("%c: LPB_solve exit status = %li\n", st, inform.status);
//printf("x: ");
//printf( X. ),
//for( int i = 0; i < n; i++) printf("%f ", x[i]);
//printf("\n");</pre>
//printf("gradient: ");
//for( int i = 0; i < n; i++) printf("%f ", g[i]);
//printf("\n");
// Delete internal workspace
lpb_terminate( &data, &control, &inform );
```

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