



## C interfaces to GALAHAD QPB

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# Chapter 1

## GALAHAD C package qpb

### 1.1 Introduction

#### 1.1.1 Purpose

This package uses a primal-dual interior-point trust-region method to solve the **quadratic programming problem**

$$\text{minimize } q(x) = \frac{1}{2}x^T Hx + g^T x + f$$

subject to the general linear constraints

$$c_i^l \leq a_i^T x \leq c_i^u, \quad i = 1, \dots, m,$$

and the simple bound constraints

$$x_j^l \leq x_j \leq x_j^u, \quad j = 1, \dots, n,$$

where the  $n$  by  $n$  symmetric matrix  $H$ , the vectors  $g$ ,  $a_i$ ,  $c^l$ ,  $c^u$ ,  $x^l$ ,  $x^u$  and the scalar  $f$  are given. Any of the constraint bounds  $c_i^l$ ,  $c_i^u$ ,  $x_j^l$  and  $x_j^u$  may be infinite. Full advantage is taken of any zero coefficients in the matrix  $H$  or the matrix  $A$  of vectors  $a_i$ .

If the matrix  $H$  is positive semi-definite, a global solution is found. However, if  $H$  is indefinite, the procedure may find a (weak second-order) critical point that is not the global solution to the given problem.

#### 1.1.2 Authors

N. I. M. Gould and D. P. Robinson, STFC-Rutherford Appleton Laboratory, England, and Philippe L. Toint, University of Namur, Belgium.

C interface, additionally J. Fowkes, STFC-Rutherford Appleton Laboratory.

#### 1.1.3 Originally released

December 1999, C interface January 2022.

### 1.1.4 Terminology

The required solution  $x$  necessarily satisfies the primal optimality conditions

$$(1a) \quad Ax = c$$

and

$$(1b) \quad c^l \leq c \leq c^u, \quad x^l \leq x \leq x^u,$$

the dual optimality conditions

$$(2a) \quad Hx + g = A^T y + z$$

where

$$(2b) \quad y = y^l + y^u, \quad z = z^l + z^u, \quad y^l \geq 0, \quad y^u \leq 0, \quad z^l \geq 0 \quad \text{and} \quad z^u \leq 0,$$

and the complementary slackness conditions

$$(3) \quad (Ax - c^l)^T y^l = 0, \quad (Ax - c^u)^T y^u = 0, \quad (x - x^l)^T z^l = 0 \quad \text{and} \quad (x - x^u)^T z^u = 0,$$

where the diagonal matrix  $W^2$  has diagonal entries  $w_j^2$ ,  $j = 1, \dots, n$ , where the vectors  $y$  and  $z$  are known as the Lagrange multipliers for the general linear constraints, and the dual variables for the bounds, respectively, and where the vector inequalities hold component-wise.

### 1.1.5 Method

Primal-dual interior point methods iterate towards a point that satisfies these conditions by ultimately aiming to satisfy (1a), (2a) and (3), while ensuring that (1b) and (2b) are satisfied as strict inequalities at each stage. Appropriate norms of the amounts by which (1a), (2a) and (3) fail to be satisfied are known as the primal and dual infeasibility, and the violation of complementary slackness, respectively. The fact that (1b) and (2b) are satisfied as strict inequalities gives such methods their other title, namely interior-point methods.

The problem is solved in two phases. The goal of the first "initial feasible point" phase is to find a strictly interior point which is primal feasible, that is that {1a} is satisfied. The GALAHAD package LSQP is used for this purpose, and offers the options of either accepting the first strictly feasible point found, or preferably of aiming for the so-called "analytic center" of the feasible region. Having found such a suitable initial feasible point, the second "optimality" phase ensures that {4.1a} remains satisfied while iterating to satisfy dual feasibility (2a) and complementary slackness (3). The optimality phase proceeds by approximately minimizing a sequence of barrier functions

$$\frac{1}{2} x^T H x + g^T x + f - \mu \left[ \sum_{i=1}^m \log(c_i - c_i^l) + \sum_{i=1}^m \log(c_i^u - c_i) + \sum_{j=1}^n \log(x_j - x_j^l) + \sum_{j=1}^n \log(x_j^u - x_j) \right],$$

for an appropriate sequence of positive barrier parameters  $\mu$  converging to zero while ensuring that (1a) remain satisfied and that  $x$  and  $c$  are strictly interior points for (1b). Note that terms in the above summations corresponding to infinite bounds are ignored, and that equality constraints are treated specially.

Each of the barrier subproblems is solved using a trust-region method. Such a method generates a trial correction step  $\Delta(x, c)$  to the current iterate  $(x, c)$  by replacing the nonlinear barrier function locally by a suitable quadratic model, and approximately minimizing this model in the intersection of {4.1a} and a trust region  $\|\Delta(x, c)\| \leq \Delta$  for some appropriate strictly positive trust-region radius  $\Delta$  and norm  $\|\cdot\|$ . The step is accepted/rejected and the radius adjusted on the basis of how accurately the model reproduces the value of barrier function at the trial step. If the step proves to be unacceptable, a linesearch is performed along the step to obtain an acceptable new iterate. In practice, the natural primal "Newton" model of the barrier function is frequently less successful than an alternative primal-dual model, and consequently the primal-dual model is usually to be preferred.

Once a barrier subproblem has been solved, extrapolation based on values and derivatives encountered on the central path is optionally used to determine a good starting point for the next subproblem. Traditional Taylor-series extrapolation has been superceded by more accurate Puiseux-series methods as these are particularly suited to deal with degeneracy.

The trust-region subproblem is approximately solved using the combined conjugate-gradient/Lanczos method implemented in the GALAHAD package GLTR. Such a method requires a suitable preconditioner, and in our case, the only flexibility we have is in approximating the model of the Hessian. Although using a fixed form of preconditioning is sometimes effective, we have provided the option of an automatic choice, that aims to balance the cost of applying the preconditioner against the needs for an accurate solution of the trust-region subproblem. The preconditioner is applied using the GALAHAD matrix factorization package SBLS, but options at this stage are to factorize the preconditioner as a whole (the so-called "augmented system" approach), or to perform a block elimination first (the "Schur-complement" approach). The latter is usually to be preferred when a (non-singular) diagonal preconditioner is used, but may be inefficient if any of the columns of  $A$  is too dense.

In order to make the solution as efficient as possible, the variables and constraints are reordered internally by the GALAHAD package QPP prior to solution. In particular, fixed variables, and free (unbounded on both sides) constraints are temporarily removed.

### 1.1.6 Reference

The method is described in detail in

A. R. Conn, N. I. M. Gould, D. Orban and Ph. L. Toint (1999). A primal-dual trust-region algorithm for minimizing a non-convex function subject to general inequality and linear equality constraints. *Mathematical Programming* **87** 215-249.

### 1.1.7 Call order

To solve a given problem, functions from the qpb package must be called in the following order:

- `qpb_initialize` - provide default control parameters and set up initial data structures
- `qpb_read_specfile` (optional) - override control values by reading replacement values from a file
- `qpb_import` - set up problem data structures and fixed values
- `qpb_reset_control` (optional) - possibly change control parameters if a sequence of problems are being solved
- `qpb_solve_qp` - solve the quadratic program
- `qpb_information` (optional) - recover information about the solution and solution process
- `qpb_terminate` - deallocate data structures

See Section 4.1 for examples of use.

### 1.1.8 Unsymmetric matrix storage formats

The unsymmetric  $m$  by  $n$  constraint matrix  $A$  may be presented and stored in a variety of convenient input formats.

Both C-style (0 based) and fortran-style (1-based) indexing is allowed. Choose `control.f_indexing` as `false` for C style and `true` for fortran style; the discussion below presumes C style, but add 1 to indices for the corresponding fortran version.

Wrappers will automatically convert between 0-based (C) and 1-based (fortran) array indexing, so may be used transparently from C. This conversion involves both time and memory overheads that may be avoided by supplying data that is already stored using 1-based indexing.

### 1.1.8.1 Dense storage format

The matrix  $A$  is stored as a compact dense matrix by rows, that is, the values of the entries of each row in turn are stored in order within an appropriate real one-dimensional array. In this case, component  $n * i + j$  of the storage array  $A\_val$  will hold the value  $A_{ij}$  for  $0 \leq i \leq m - 1, 0 \leq j \leq n - 1$ .

### 1.1.8.2 Sparse co-ordinate storage format

Only the nonzero entries of the matrices are stored. For the  $l$ -th entry,  $0 \leq l \leq ne - 1$ , of  $A$ , its row index  $i$ , column index  $j$  and value  $A_{ij}$ ,  $0 \leq i \leq m - 1, 0 \leq j \leq n - 1$ , are stored as the  $l$ -th components of the integer arrays  $A\_row$  and  $A\_col$  and real array  $A\_val$ , respectively, while the number of nonzeros is recorded as  $A\_ne = ne$ .

### 1.1.8.3 Sparse row-wise storage format

Again only the nonzero entries are stored, but this time they are ordered so that those in row  $i$  appear directly before those in row  $i+1$ . For the  $i$ -th row of  $A$  the  $i$ -th component of the integer array  $A\_ptr$  holds the position of the first entry in this row, while  $A\_ptr(m)$  holds the total number of entries plus one. The column indices  $j$ ,  $0 \leq j \leq n - 1$ , and values  $A_{ij}$  of the nonzero entries in the  $i$ -th row are stored in components  $l = A\_ptr(i), \dots, A\_ptr(i+1)-1$ ,  $0 \leq i \leq m - 1$ , of the integer array  $A\_col$ , and real array  $A\_val$ , respectively. For sparse matrices, this scheme almost always requires less storage than its predecessor.

## 1.1.9 Symmetric matrix storage formats

Likewise, the symmetric  $n$  by  $n$  objective Hessian matrix  $H$  may be presented and stored in a variety of formats. But crucially symmetry is exploited by only storing values from the lower triangular part (i.e. those entries that lie on or below the leading diagonal).

### 1.1.9.1 Dense storage format

The matrix  $H$  is stored as a compact dense matrix by rows, that is, the values of the entries of each row in turn are stored in order within an appropriate real one-dimensional array. Since  $H$  is symmetric, only the lower triangular part (that is the part  $h_{ij}$  for  $0 \leq j \leq i \leq n - 1$ ) need be held. In this case the lower triangle should be stored by rows, that is component  $i * i/2 + j$  of the storage array  $H\_val$  will hold the value  $h_{ij}$  (and, by symmetry,  $h_{ji}$ ) for  $0 \leq j \leq i \leq n - 1$ .

### 1.1.9.2 Sparse co-ordinate storage format

Only the nonzero entries of the matrices are stored. For the  $l$ -th entry,  $0 \leq l \leq ne - 1$ , of  $H$ , its row index  $i$ , column index  $j$  and value  $h_{ij}$ ,  $0 \leq j \leq i \leq n - 1$ , are stored as the  $l$ -th components of the integer arrays  $H\_row$  and  $H\_col$  and real array  $H\_val$ , respectively, while the number of nonzeros is recorded as  $H\_ne = ne$ . Note that only the entries in the lower triangle should be stored.

### 1.1.9.3 Sparse row-wise storage format

Again only the nonzero entries are stored, but this time they are ordered so that those in row  $i$  appear directly before those in row  $i+1$ . For the  $i$ -th row of  $H$  the  $i$ -th component of the integer array  $H\_ptr$  holds the position of the first entry in this row, while  $H\_ptr(n)$  holds the total number of entries plus one. The column indices  $j$ ,  $0 \leq j \leq i$ , and values  $h_{ij}$  of the entries in the  $i$ -th row are stored in components  $l = H\_ptr(i), \dots, H\_ptr(i+1)-1$  of the integer array  $H\_col$ , and real array  $H\_val$ , respectively. Note that as before only the entries in the lower triangle should be stored. For sparse matrices, this scheme almost always requires less storage than its predecessor.



#### 1.1.9.4 Diagonal storage format

If  $H$  is diagonal (i.e.,  $H_{ij} = 0$  for all  $0 \leq i \neq j \leq n - 1$ ) only the diagonal entries  $H_{ii}$ ,  $0 \leq i \leq n - 1$  need be stored, and the first  $n$  components of the array `H_val` may be used for the purpose.

#### 1.1.9.5 Multiples of the identity storage format

If  $H$  is a multiple of the identity matrix, (i.e.,  $H = \alpha I$  where  $I$  is the  $n$  by  $n$  identity matrix and  $\alpha$  is a scalar), it suffices to store  $\alpha$  as the first component of `H_val`.

#### 1.1.9.6 The identity matrix format

If  $H$  is the identity matrix, no values need be stored.

#### 1.1.9.7 The zero matrix format

The same is true if  $H$  is the zero matrix.



## Chapter 2

# File Index

### 2.1 File List

Here is a list of all files with brief descriptions:

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## Chapter 3

# File Documentation

### 3.1 qpb.h File Reference

```
#include <stdbool.h>
#include "galahad_precision.h"
#include "lsqp.h"
#include "fdc.h"
#include "sbls.h"
#include "gltr.h"
#include "fit.h"
```

#### Data Structures

- struct [qpb\\_control\\_type](#)
- struct [qpb\\_time\\_type](#)
- struct [qpb\\_inform\\_type](#)

#### Functions

- void [qpb\\_initialize](#) (void \*\*data, struct [qpb\\_control\\_type](#) \*control, int \*status)
- void [qpb\\_read\\_specfile](#) (struct [qpb\\_control\\_type](#) \*control, const char specfile[])
- void [qpb\\_import](#) (struct [qpb\\_control\\_type](#) \*control, void \*\*data, int \*status, int n, int m, const char H\_type[], int H\_ne, const int H\_row[], const int H\_col[], const int H\_ptr[], const char A\_type[], int A\_ne, const int A\_row[], const int A\_col[], const int A\_ptr[])
- void [qpb\\_reset\\_control](#) (struct [qpb\\_control\\_type](#) \*control, void \*\*data, int \*status)
- void [qpb\\_solve\\_qp](#) (void \*\*data, int \*status, int n, int m, int h\_ne, const real\_wp\_ H\_val[], const real\_wp\_ g[], const real\_wp\_ f, int a\_ne, const real\_wp\_ A\_val[], const real\_wp\_ c\_l[], const real\_wp\_ c\_u[], const real\_wp\_ x\_l[], const real\_wp\_ x\_u[], real\_wp\_ x[], real\_wp\_ c[], real\_wp\_ y[], real\_wp\_ z[], int x\_stat[], int c\_stat[])
- void [qpb\\_information](#) (void \*\*data, struct [qpb\\_inform\\_type](#) \*inform, int \*status)
- void [qpb\\_terminate](#) (void \*\*data, struct [qpb\\_control\\_type](#) \*control, struct [qpb\\_inform\\_type](#) \*inform)

#### 3.1.1 Data Structure Documentation

##### 3.1.1.1 struct qpb\_control\_type

control derived type as a C struct

## Data Fields

bool	f_indexing	use C or Fortran sparse matrix indexing
int	error	error and warning diagnostics occur on stream error
int	out	general output occurs on stream out
int	print_level	the level of output required is specified by print_level
int	start_print	any printing will start on this iteration
int	stop_print	any printing will stop on this iteration
int	maxit	at most maxit inner iterations are allowed
int	itref_max	the maximum number of iterative refinements allowed (OBSOL
int	cg_maxit	the maximum number of CG iterations allowed. If cg_maxit < 0, this number will be reset to the dimension of the system + 1
int	indicator_type	specifies the type of indicator function used. Possible values are 1 primal indicator: constraint active $\leq$ distance to nearest bound $\leq$ .indicator_p_tol 2 primal-dual indicator: constraint active $\leq$ distance to nearest bound $\leq$ .indicator_tol_pd * size of corresponding multiplier 3 primal-dual indicator: constraint active $\leq$ distance to nearest bound $\leq$ .indicator_tol_tapia * distance to same bound at previous iteration
int	restore_problem	indicate whether and how much of the input problem should be restored on output. Possible values are 0 nothing restored 1 scalar and vector parameters 2 all parameters
int	extrapolate	should extrapolation be used to track the central path? Possible values 0 never 1 after the final major iteration 2 at each major iteration
int	path_history	the maximum number of previous path points to use when fitting the data
int	factor	the factorization to be used. Possible values are 0 automatic 1 Schur-complement factorization 2 augmented-system factorization (OBSOLETE
int	max_col	the maximum number of nonzeros in a column of A which is permitted with the Schur-complement factorization (OBSOLE
int	indmin	an initial guess as to the integer workspace required by SBLS (OBSOL
int	valmin	an initial guess as to the real workspace required by SBLS (OBSOL
int	infeas_max	the number of iterations for which the overall infeasibility of the problem is not reduced by at least a factor .reduce_infeas before the problem is flagged as infeasible (see reduce_infeas)

## Data Fields

int	precon	the preconditioner to be used for the CG is defined by precon. Possible values are 0 automatic 1 no preconditioner, i.e, the identity within full factorization 2 full factorization 3 band within full factorization 4 diagonal using the barrier terms within full factorization (OBSOLETE)
int	nsemib	the semi-bandwidth of a band preconditioner, if appropriate (OBSOL
int	path_derivatives	the maximum order of path derivative to use
int	fit_order	the order of (Puiseux) series to fit to the path data: $\leq 0$ to fit all da
int	sif_file_device	specifies the unit number to write generated SIF file describing the current problem
real_wp_	infinity	any bound larger than infinity in modulus will be regarded as infinite
real_wp_	stop_p	the required accuracy for the primal infeasibility
real_wp_	stop_d	the required accuracy for the dual infeasibility
real_wp_	stop_c	the required accuracy for the complementarity
real_wp_	theta_d	tolerances used to terminate the inner iteration (for given mu): dual feasibility $\leq \text{MAX}(\text{theta\_d} * \mu ** \text{beta}, 0.99 * \text{stop\_d})$ complementarity $\leq \text{MAX}(\text{theta\_c} * \mu ** \text{beta}, 0.99 * \text{stop\_d})$
real_wp_	theta_c	see theta_d
real_wp_	beta	see theta_d
real_wp_	prfeas	initial primal variables will not be closer than prfeas from their bound
real_wp_	dufeas	initial dual variables will not be closer than dufear from their bounds
real_wp_	muzero	the initial value of the barrier parameter. If muzero is not positive, it will be reset to an appropriate value
real_wp_	reduce_infeas	if the overall infeasibility of the problem is not reduced by at least a factor reduce_infeas over .infeas_max iterations, the problem is flagged as infeasible (see infeas_max)
real_wp_	obj_unbounded	if the objective function value is smaller than obj_unbounded, it will b flagged as unbounded from below.
real_wp_	pivot_tol	the threshold pivot used by the matrix factorization. See the documentation for SBLS for details (OBSOLE
real_wp_	pivot_tol_for_dependencies	the threshold pivot used by the matrix factorization when attempting to detect linearly dependent constraints. See the documentation for FDC for details (OBSOLE
real_wp_	zero_pivot	any pivots smaller than zero_pivot in absolute value will be regarded to zero when attempting to detect linearly dependent constraints (OBSOLE

## Data Fields

real_wp_	identical_bounds_tol	any pair of constraint bounds (c_l,c_u) or (x_l,x_u) that are closer than identical_bounds_tol will be reset to the average of their values
real_wp_	inner_stop_relative	the search direction is considered as an acceptable approximation to the minimizer of the model if the gradient of the model in the preconditioning(inverse) norm is less than $\max(\text{inner\_stop\_relative} * \text{initial preconditioning(inverse) gradient norm}, \text{inner\_stop\_absolute})$
real_wp_	inner_stop_absolute	see inner_stop_relative
real_wp_	initial_radius	the initial trust-region radius
real_wp_	mu_min	start terminal extrapolation when mu reaches mu_min
real_wp_	inner_fraction_opt	a search direction which gives at least inner_fraction_opt times the optimal model decrease will be found
real_wp_	indicator_tol_p	if .indicator_type = 1, a constraint/bound will be deemed to be active $\Leftrightarrow$ distance to nearest bound $\leq$ .indicator_tol_p
real_wp_	indicator_tol_pd	if .indicator_type = 2, a constraint/bound will be deemed to be active $\Leftrightarrow$ distance to nearest bound $\leq$ .indicator_tol_pd * size of corresponding multiplier
real_wp_	indicator_tol_tapia	if .indicator_type = 3, a constraint/bound will be deemed to be active $\Leftrightarrow$ distance to nearest bound $\leq$ .indicator_tol_tapia * distance to same bound at previous iteration
real_wp_	cpu_time_limit	the maximum CPU time allowed (-ve means infinite)
real_wp_	clock_time_limit	the maximum elapsed clock time allowed (-ve means infinite)
bool	remove_dependencies	the equality constraints will be preprocessed to remove any linear dependencies if true
bool	treat_zero_bounds_as_general	any problem bound with the value zero will be treated as if it were a general value if true
bool	center	if .center is true, the algorithm will use the analytic center of the feasible set as its initial feasible point. Otherwise, a feasible point as close as possible to the initial point will be used. We recommend using the analytic center
bool	primal	if .primal, is true, a primal barrier method will be used in place of the primal-dual method
bool	puiseux	If extrapolation is to be used, decide between Puiseux and Taylor series.
bool	feasol	if .feasol is true, the final solution obtained will be perturbed so that variables close to their bounds are moved onto these bounds
bool	array_syntax_worse_than_do_loop	if .array_syntax_worse_than_do_loop is true, f77-style do loops will be used rather than f90-style array syntax for vector operations (OBSOLETE)



## Data Fields

bool	space_critical	if .space_critical true, every effort will be made to use as little space as possible. This may result in longer computation time
bool	deallocate_error_fatal	if .deallocate_error_fatal is true, any array/pointer deallocation error will terminate execution. Otherwise, computation will continue
bool	generate_sif_file	if .generate_sif_file is .true. if a SIF file describing the current problem is to be generated
char	sif_file_name[31]	name of generated SIF file containing input problem
char	prefix[31]	all output lines will be prefixed by .prefix(2:LEN(TRIM(.prefix))-1) where .prefix contains the required string enclosed in quotes, e.g. "string" or 'string'
struct lsqp_control_type	lsqp_control	control parameters for LSQP
struct fdc_control_type	fdc_control	control parameters for FDC
struct sbbs_control_type	sbbs_control	control parameters for SBLS
struct gltr_control_type	gltr_control	control parameters for GLTR
struct fit_control_type	fit_control	control parameters for FIT

## 3.1.1.2 struct qpb\_time\_type

time derived type as a C struct

## Data Fields

real_wp_	total	the total CPU time spent in the package
real_wp_	preprocess	the CPU time spent preprocessing the problem
real_wp_	find_dependent	the CPU time spent detecting linear dependencies
real_wp_	analyse	the CPU time spent analysing the required matrices prior to factorization
real_wp_	factorize	the CPU time spent factorizing the required matrices
real_wp_	solve	the CPU time spent computing the search direction
real_wp_	phase1_total	the total CPU time spent in the initial-point phase of the package
real_wp_	phase1_analyse	the CPU time spent analysing the required matrices prior to factorization in the initial-point phase
real_wp_	phase1_factorize	the CPU time spent factorizing the required matrices in the initial-point phase
real_wp_	phase1_solve	the CPU time spent computing the search direction in the initial-point phase
real_wp_	clock_total	the total clock time spent in the package
real_wp_	clock_preprocess	the clock time spent preprocessing the problem
real_wp_	clock_find_dependent	the clock time spent detecting linear dependencies
real_wp_	clock_analyse	the clock time spent analysing the required matrices prior to factorization
real_wp_	clock_factorize	the clock time spent factorizing the required matrices
real_wp_	clock_solve	the clock time spent computing the search direction
real_wp_	clock_phase1_total	the total clock time spent in the initial-point phase of the package

## Data Fields

real_wp_	clock_phase1_analyse	the clock time spent analysing the required matrices prior to factorizat in the inital-point phase
real_wp_	clock_phase1_factorize	the clock time spent factorizing the required matrices in the inital-poi phase
real_wp_	clock_phase1_solve	the clock time spent computing the search direction in the inital-point

## 3.1.1.3 struct qpb\_inform\_type

inform derived type as a C struct

## Data Fields

int	status	return status. See QPB_solve for details
int	alloc_status	the status of the last attempted allocation/deallocation
char	bad_alloc[81]	the name of the array for which an allocation/deallocation error occurred
int	iter	the total number of iterations required
int	cg_iter	the total number of conjugate gradient iterations required
int	factorization_status	the return status from the factorization
int	factorization_integer	the total integer workspace required for the factorization
int	factorization_real	the total real workspace required for the factorization
int	nfacts	the total number of factorizations performed
int	nbacts	the total number of "wasted" function evaluations during the linesearch
int	nmods	the total number of factorizations which were modified to ensure that th matrix was an appropriate preconditioner
real_wp_	obj	the value of the objective function at the best estimate of the solution determined by QPB_solve
real_wp_	non_negligible_pivot	the smallest pivot which was not judged to be zero when detecting linear dependent constraints
bool	feasible	is the returned "solution" feasible?
struct <a href="#">qpb_time_type</a>	time	timings (see above)
struct <a href="#">lsqp_inform_type</a>	lsqp_inform	inform parameters for LSQP
struct <a href="#">fdc_inform_type</a>	fdc_inform	inform parameters for FDC
struct <a href="#">sbls_inform_type</a>	sbls_inform	inform parameters for SBLS
struct <a href="#">gltr_inform_type</a>	gltr_inform	return information from GLTR
struct <a href="#">fit_inform_type</a>	fit_inform	return information from FIT

## 3.1.2 Function Documentation

## 3.1.2.1 qpb\_initialize()

```
void qpb_initialize (
    void ** data,
```

```

    struct qpb_control_type * control,
    int * status )

```

Set default control values and initialize private data

#### Parameters

in, out	<i>data</i>	holds private internal data
out	<i>control</i>	is a struct containing control information (see <a href="#">qpb_control_type</a> )
out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently): <ul style="list-style-type: none"> <li>• 0. The import was succesful.</li> </ul>

#### 3.1.2.2 qpb\_read\_specfile()

```

void qpb_read_specfile (
    struct qpb_control_type * control,
    const char specfile[] )

```

Read the content of a specification file, and assign values associated with given keywords to the corresponding control parameters

#### Parameters

in, out	<i>control</i>	is a struct containing control information (see <a href="#">qpb_control_type</a> )
in	<i>specfile</i>	is a character string containing the name of the specification file

#### 3.1.2.3 qpb\_import()

```

void qpb_import (
    struct qpb_control_type * control,
    void ** data,
    int * status,
    int n,
    int m,
    const char H_type[],
    int H_ne,
    const int H_row[],
    const int H_col[],
    const int H_ptr[],
    const char A_type[],
    int A_ne,
    const int A_row[],
    const int A_col[],
    const int A_ptr[] )

```

Import problem data into internal storage prior to solution.

## Parameters

in	<i>control</i>	is a struct whose members provide control paramters for the remaining prcedures (see <a href="#">qpb_control_type</a> )
in, out	<i>data</i>	holds private internal data
in, out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are: <ul style="list-style-type: none"> <li>• 0. The import was succesful</li> <li>• -1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -3. The restrictions <math>n &gt; 0</math> or <math>m &gt; 0</math> or requirement that a type contains its relevant string 'dense', 'coordinate', 'sparse_by_rows', 'diagonal', 'scaled_identity', 'identity', 'zero' or 'none' has been violated.</li> <li>• -23. An entry from the strict upper triangle of <math>H</math> has been specified.</li> </ul>
in	<i>n</i>	is a scalar variable of type int, that holds the number of variables.
in	<i>m</i>	is a scalar variable of type int, that holds the number of general linear constraints.
in	<i>H_type</i>	is a one-dimensional array of type char that specifies the <a href="#">symmetric storage scheme</a> used for the Hessian, $H$ . It should be one of 'coordinate', 'sparse_by_rows', 'dense', 'diagonal', 'scaled_identity', 'identity', 'zero' or 'none', the latter pair if $H = 0$ ; lower or upper case variants are allowed.
in	<i>H_ne</i>	is a scalar variable of type int, that holds the number of entries in the lower triangular part of $H$ in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes.
in	<i>H_row</i>	is a one-dimensional array of size $H\_ne$ and type int, that holds the row indices of the lower triangular part of $H$ in the sparse co-ordinate storage scheme. It need not be set for any of the other three schemes, and in this case can be NULL.
in	<i>H_col</i>	is a one-dimensional array of size $H\_ne$ and type int, that holds the column indices of the lower triangular part of $H$ in either the sparse co-ordinate, or the sparse row-wise storage scheme. It need not be set when the dense, diagonal or (scaled) identity storage schemes are used, and in this case can be NULL.
in	<i>H_ptr</i>	is a one-dimensional array of size $n+1$ and type int, that holds the starting position of each row of the lower triangular part of $H$ , as well as the total number of entries plus one, in the sparse row-wise storage scheme. It need not be set when the other schemes are used, and in this case can be NULL.
in	<i>A_type</i>	is a one-dimensional array of type char that specifies the <a href="#">unsymmetric storage scheme</a> used for the constraint Jacobian, $A$ . It should be one of 'coordinate', 'sparse_by_rows' or 'dense'; lower or upper case variants are allowed.
in	<i>A_ne</i>	is a scalar variable of type int, that holds the number of entries in $A$ in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes.
in	<i>A_row</i>	is a one-dimensional array of size $A\_ne$ and type int, that holds the row indices of $A$ in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes, and in this case can be NULL.
in	<i>A_col</i>	is a one-dimensional array of size $A\_ne$ and type int, that holds the column indices of $A$ in either the sparse co-ordinate, or the sparse row-wise storage scheme. It need not be set when the dense or diagonal storage schemes are used, and in this case can be NULL.

## Parameters

in	<i>A_ptr</i>	is a one-dimensional array of size $n+1$ and type <code>int</code> , that holds the starting position of each row of $A$ , as well as the total number of entries plus one, in the sparse row-wise storage scheme. It need not be set when the other schemes are used, and in this case can be <code>NULL</code> .
----	--------------	--

## 3.1.2.4 qpb\_reset\_control()

```
void qpb_reset_control (
    struct qpb_control_type * control,
    void ** data,
    int * status )
```

Reset control parameters after import if required.

## Parameters

in	<i>control</i>	is a struct whose members provide control parameters for the remaining procedures (see <a href="#">qpb_control_type</a> )
in, out	<i>data</i>	holds private internal data
in, out	<i>status</i>	is a scalar variable of type <code>int</code> , that gives the exit status from the package. Possible values are: <ul style="list-style-type: none"> <li>• 0. The import was successful.</li> </ul>

## 3.1.2.5 qpb\_solve\_qp()

```
void qpb_solve_qp (
    void ** data,
    int * status,
    int n,
    int m,
    int h_ne,
    const real_wp_ H_val[],
    const real_wp_ g[],
    const real_wp_ f,
    int a_ne,
    const real_wp_ A_val[],
    const real_wp_ c_l[],
    const real_wp_ c_u[],
    const real_wp_ x_l[],
    const real_wp_ x_u[],
    real_wp_ x[],
    real_wp_ c[],
    real_wp_ y[],
    real_wp_ z[],
```

```
int x_stat[],
int c_stat[] )
```

Solve the quadratic program when the Hessian  $H$  is available.

#### Parameters

in, out	<i>data</i>	holds private internal data
in, out	<i>status</i>	<p>is a scalar variable of type int, that gives the entry and exit status from the package. Possible exit are:</p> <ul style="list-style-type: none"> <li>• 0. The run was succesful.</li> <li>• -1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -3. The restrictions <math>n &gt; 0</math> and <math>m &gt; 0</math> or requirement that a type contains its relevant string 'dense', 'coordinate', 'sparse_by_rows', 'diagonal', 'scaled_identity', 'identity', 'zero' or 'none' has been violated.</li> <li>• -5. The simple-bound constraints are inconsistent.</li> <li>• -7. The constraints appear to have no feasible point.</li> <li>• -9. The analysis phase of the factorization failed; the return status from the factorization package is given in the component inform.factor_status</li> <li>• -10. The factorization failed; the return status from the factorization package is given in the component inform.factor_status.</li> <li>• -11. The solution of a set of linear equations using factors from the factorization package failed; the return status from the factorization package is given in the component inform.factor_status.</li> <li>• -16. The problem is so ill-conditioned that further progress is impossible.</li> <li>• -17. The step is too small to make further impact.</li> <li>• -18. Too many iterations have been performed. This may happen if control.maxit is too small, but may also be symptomatic of a badly scaled problem.</li> <li>• -19. The CPU time limit has been reached. This may happen if control.cpu_time_limit is too small, but may also be symptomatic of a badly scaled problem.</li> <li>• -23. An entry from the strict upper triangle of <math>H</math> has been specified.</li> </ul>
in	<i>n</i>	is a scalar variable of type int, that holds the number of variables
in	<i>m</i>	is a scalar variable of type int, that holds the number of general linear constraints.
in	<i>h_ne</i>	is a scalar variable of type int, that holds the number of entries in the lower triangular part of the Hessian matrix $H$ .
in	<i>H_val</i>	is a one-dimensional array of size h_ne and type double, that holds the values of the entries of the lower triangular part of the Hessian matrix $H$ in any of the available storage schemes.
in	<i>g</i>	is a one-dimensional array of size n and type double, that holds the linear term $g$ of the objective function. The $j$ -th component of $g$ , $j = 0, \dots, n-1$ , contains $g_j$ .

## Parameters

in	<i>f</i>	is a scalar of type double, that holds the constant term $f$ of the objective function.
in	<i>a_ne</i>	is a scalar variable of type int, that holds the number of entries in the constraint Jacobian matrix $A$ .
in	<i>A_val</i>	is a one-dimensional array of size <i>a_ne</i> and type double, that holds the values of the entries of the constraint Jacobian matrix $A$ in any of the available storage schemes.
in	<i>c_l</i>	is a one-dimensional array of size <i>m</i> and type double, that holds the lower bounds $c^l$ on the constraints $Ax$ . The <i>i</i> -th component of <i>c_l</i> , $i = 0, \dots, m-1$ , contains $c_i^l$ .
in	<i>c_u</i>	is a one-dimensional array of size <i>m</i> and type double, that holds the upper bounds $c^u$ on the constraints $Ax$ . The <i>i</i> -th component of <i>c_u</i> , $i = 0, \dots, m-1$ , contains $c_i^u$ .
in	<i>x_l</i>	is a one-dimensional array of size <i>n</i> and type double, that holds the lower bounds $x^l$ on the variables $x$ . The <i>j</i> -th component of <i>x_l</i> , $j = 0, \dots, n-1$ , contains $x_j^l$ .
in	<i>x_u</i>	is a one-dimensional array of size <i>n</i> and type double, that holds the upper bounds $x^u$ on the variables $x$ . The <i>j</i> -th component of <i>x_u</i> , $j = 0, \dots, n-1$ , contains $x_j^u$ .
in, out	<i>x</i>	is a one-dimensional array of size <i>n</i> and type double, that holds the values $x$ of the optimization variables. The <i>j</i> -th component of <i>x</i> , $j = 0, \dots, n-1$ , contains $x_j$ .
out	<i>c</i>	is a one-dimensional array of size <i>m</i> and type double, that holds the residual $c(x)$ . The <i>i</i> -th component of <i>c</i> , $i = 0, \dots, m-1$ , contains $c_i(x)$ .
in, out	<i>y</i>	is a one-dimensional array of size <i>n</i> and type double, that holds the values $y$ of the Lagrange multipliers for the general linear constraints. The <i>j</i> -th component of <i>y</i> , $j = 0, \dots, n-1$ , contains $y_j$ .
in, out	<i>z</i>	is a one-dimensional array of size <i>n</i> and type double, that holds the values $z$ of the dual variables. The <i>j</i> -th component of <i>z</i> , $j = 0, \dots, n-1$ , contains $z_j$ .
out	<i>x_stat</i>	is a one-dimensional array of size <i>n</i> and type int, that gives the optimal status of the problem variables. If <i>x_stat</i> ( <i>j</i> ) is negative, the variable $x_j$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.
out	<i>c_stat</i>	is a one-dimensional array of size <i>m</i> and type int, that gives the optimal status of the general linear constraints. If <i>c_stat</i> ( <i>i</i> ) is negative, the constraint value $a_i^T x$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.

## 3.1.2.6 qpb\_information()

```
void qpb_information (
    void ** data,
    struct qpb_inform_type * inform,
    int * status )
```

Provides output information

## Parameters

in, out	<i>data</i>	holds private internal data
out	<i>inform</i>	is a struct containing output information (see <a href="#">qpb_inform_type</a> )
out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently): <ul style="list-style-type: none"> <li>• 0. The values were recorded successfully</li> </ul>

### 3.1.2.7 qpb\_terminate()

```
void qpb_terminate (
    void ** data,
    struct qpb_control_type * control,
    struct qpb_inform_type * inform )
```

Deallocate all internal private storage

#### Parameters

in, out	<i>data</i>	holds private internal data
out	<i>control</i>	is a struct containing control information (see <a href="#">qpb_control_type</a> )
out	<i>inform</i>	is a struct containing output information (see <a href="#">qpb_inform_type</a> )



## Chapter 4

# Example Documentation

### 4.1 qpbt.c

This is an example of how to use the package to solve a quadratic program. A variety of supported Hessian and constraint matrix storage formats are shown.

Notice that C-style indexing is used, and that this is flagged by setting `control.f_indexing` to `false`.

### 4.2 qpbtf.c

This is the same example, but now fortran-style indexing is used.



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