

1.Easy\01.Largest_element_in_array.cpp

```
1  /*
2  QUESTION:-
3  Given an array A[] of size n. The task is to find the largest element in it.
4
5  Example:
6
7  Input:
8  n = 5
9  A[] = {1, 8, 7, 56, 90}
10 Output:
11 90
12 Explanation:
13 The largest element of given array is 90
14 */
15
16 /*
17 APPROACH:-
18 -> Intialize the ans with starting element
19 -> Traverse the entire array and update the ans if the element is greater then ans
20 -> Finally, return the ans
21 */
22
23 // CODE:-
24 int largest(int arr[], int n)
25 {
26     int ans = arr[0];
27     for (int i = 1; i < n; i++)
28     {
29         if (arr[i] > ans)
30             ans = arr[i];
31     }
32     return ans;
33 }
34
35 // TIME COMPLEXITY = O(N)
36 // SPACE COMPLEXITY = O(0)
```

1.Easy\02.Second_largest_element_in_array.cpp

```
1  /*
2  QUESTION:-
3  Given an array Arr of size N, print second largest distinct element from an array.
4
5  Example:
6
7  Input:
8  N = 6
9  Arr[] = {12, 35, 1, 10, 34, 1}
10 Output: 34
11 Explanation: The largest element of the
12 array is 35 and the second largest element
13 is 34.
14 */
15
16 /*
17 APPROACH
18 -> If the current element is larger than 'large' then update second_large and large variables
19 -> Else if the current element is larger than 'second_large' then we update the variable
    second_large.
20 -> Once we traverse the entire array, we would find the second largest element in the
    variable second_large.
21 */
22
23 // CODE:-
24 int print2largest(int arr[], int n)
25 {
26     int prev = -1, curr = arr[0];
27     for (int i = 1; i < n; i++)
28     {
29         if (arr[i] > curr)
30         {
31             prev = curr;
32             curr = arr[i];
33         }
34         else if (arr[i] > prev && arr[i] != curr)
35             prev = arr[i];
36     }
37     return prev;
38 }
39
40 // TIME COMPLEXITY = O(N)
41 // SPACE COMPLEXITY = O(0)
```

1.Easy\03.Check_if_array_is_sorted_and_rotated.cpp

```
1  /*
2  QUESTION:-
3  Given an array nums, return true if the array was originally sorted in non-decreasing order,
4  then rotated some number of positions (including zero). Otherwise, return false.
5  There may be duplicates in the original array.
6  Example 1:
7  Input: nums = [3,4,5,1,2]
8  Output: true
9  Explanation: [1,2,3,4,5] is the original sorted array.
10 You can rotate the array by x = 3 positions to begin on the the element of value 3:
11 [3,4,5,1,2].
12 Example 2:
13 Input: nums = [2,1,3,4]
14 Output: false
15 Explanation: There is no sorted array once rotated that can make nums.
16
17 */
18
19 /*
20 APPROACH:-
21 Compare all neighbour elements (a,b) in A,
22 the case of a > b can happen at most once.
23
24 Note that the first element and the last element are also connected.
25
26 If all a <= b, A is already sorted so answer is true.
27 If all a <= b but only one a > b, and the first element is greater than equal to last element
28 we can rotate and make b the first element so answer is true.
29 Other case, return false.
30 */
31
32 // CODE:-
33 bool check(vector<int> &nums)
34 {
35     int cnt = 0;
36     int n = nums.size();
37     for (int i = 0; i < n - 1; i++)
38     {
39         if (nums[i] > nums[i + 1])
40             cnt++;
41     }
42     if (cnt == 0)
43         return true;
44     else if (cnt == 1 && nums[0] >= nums[n - 1])
45         return true;
46     return false;
47 }
48
49 // TIME COMPLEXITY = O(N)
50 // SPACE COMPLEXITY = O(0)
```

1.Easy\04.Remove_duplicates_from_sorted_array.cpp

```

1  /*
2  QUESTION:-
3
4  Given an integer array nums sorted in non-decreasing order, remove the duplicates in-place
  such that each unique element appears only once. The relative order of the elements should be
  kept the same. Then return the number of unique elements in nums.
5  Consider the number of unique elements of nums to be k, to get accepted, you need to do the
  following things:
6  Change the array nums such that the first k elements of nums contain the unique elements in
  the order they were present in nums initially. The remaining elements of nums are not
  important as well as the size of nums.
7  Return k.
8
9  Example 1:
10
11 Input: nums = [1,1,2]
12 Output: 2, nums = [1,2,_]
13 Explanation: Your function should return k = 2, with the first two elements of nums being 1
  and 2 respectively.
14 It does not matter what you leave beyond the returned k (hence they are underscores).
15 Example 2:
16
17 Input: nums = [0,0,1,1,1,2,2,3,3,4]
18 Output: 5, nums = [0,1,2,3,4,_,_,_,_,_]
19 Explanation: Your function should return k = 5, with the first five elements of nums being 0,
  1, 2, 3, and 4 respectively.
20 It does not matter what you leave beyond the returned k (hence they are underscores).
21 */
22
23 /*
24 APPROACH:-
25 -> The idea, is to use keep a pointer k which signifies that upto here the array is sorted
26 -> Now traverse the entire array and if arr[k]!=arr[j] that is arr[j] is a unique value hence
  it should be included
27     so increment the k and swap arr[k] with arr[j]
28 -> Return k+1, +1 is because of 0 based indexing
29 */
30
31 // CODE:-
32 int removeDuplicates(vector<int> &nums)
33 {
34     int k = 0; // upto k array contains unique elements
35     for (int j = 1; j < nums.size(); j++)
36     {
37         if (nums[k] != nums[j])
38         {
39             k++;
40             swap(nums[k], nums[j]);
41         }
42     }
43     return k + 1;
44 }

```

1.Easy\05.Rotate_array_left_by_1place.cpp

```
1  /*
2  QUESTION:-
3  Given an array "ARR" containing 'N' elements, rotate this array Left by once means to shift
   all elements by one place to the left and move the first element to the last position in the
   array.
4
5  Example:
6  Input: 'N' 5, 'ARR' = [1, 2, 3, 4, 5]
7  Output: [2, 3, 4, 5, 1]
8
9  Explanation:
10 We moved the 2nd element to the 1st position and 3rd element to the 2nd position and 4th
   element to the 3rd position and 5th element to the 4th position and move oth element to the
   5th position.
11 */
12
13 /*
14 APPROACH:-
15 -> By observing we can the ans is the arr where arr[i] = arr[i+1] and at last place we will
   have arr[0]
16 -> Before traversing store the arr[0] in temp and then traverse the array and make arr[i] =
   arr[i+1]
17 -> Make arr[n-1] = arr[0], where n is the size of the array
18 */
19
20 // CODE:-
21 vector<int> rotateArray(vector<int> &arr, int n)
22 {
23     int temp = arr[0];
24     for (int i = 0; i < n - 1; i++)
25     {
26         arr[i] = arr[i + 1];
27     }
28     arr[n - 1] = temp;
29
30     return arr;
31 }
32
33 // TIME COMPLEXITY = O(N)
34 // SPACE COMPLEXITY = O(0)
```

1.Easy\06.Rotate_array_left&right_by_k_places.cpp

```
1  /*
2  QUESTION:-
3
4  Given an integer array nums, rotate the array to the right by k steps, where k is non-
negative.
5
6  Example 1:
7
8  Input: nums = [1,2,3,4,5,6,7], k = 3
9  Output: [5,6,7,1,2,3,4]
10 Explanation:
11 rotate 1 steps to the right: [7,1,2,3,4,5,6]
12 rotate 2 steps to the right: [6,7,1,2,3,4,5]
13 rotate 3 steps to the right: [5,6,7,1,2,3,4]
14 Example 2:
15
16 Input: nums = [-1,-100,3,99], k = 2
17 Output: [3,99,-1,-100]
18 Explanation:
19 rotate 1 steps to the right: [99,-1,-100,3]
20 rotate 2 steps to the right: [3,99,-1,-100]
21 */
22
23 /*
24 APPROACH:-
25 To rotate the array k places to right follow below steps
26 -> Reverse first n-k elements
27 -> Reverse last k elements
28 -> Reverse the entire array
29
30 To rotate the array k places to left follow below steps
31 -> Reverse first k elements
32 -> Reverse last n-k elements
33 -> Reverse the entire array
34 */
35
36 // CODE:-
37
38 // RIGHT ROTATE:-
39 void rightRotate(int arr[], int n, int k)
40 {
41     k = k % n; // to keep k within the range
42     reverse(arr, arr + (n - k));
43     reverse(arr + (n - k), arr + n);
44     reverse(arr, arr + n);
45 }
46
47 // LEFT ROTATE:-
48 void leftRotate(int arr[], int n, int k)
49 {
50     k = k % n; // to keep k within the range
51     reverse(arr, arr + k);
```

```
52     reverse(arr + k, arr + n);  
53     reverse(arr, arr + n);  
54 }  
55  
56 // TIME COMPLEXITY = O(N)  
57 // SPACE COMPLEXITY = O(0)
```

1.Easy\07.Move_0's_to_end.cpp

```
1  /*
2  QUESTION:-
3  Given an integer array nums, move all 0's to the end of it while maintaining the relative
   order of the non-zero elements.
4
5  Note that you must do this in-place without making a copy of the array.
6
7  Example 1:
8
9  Input: nums = [0,1,0,3,12]
10 Output: [1,3,12,0,0]
11 Example 2:
12
13 Input: nums = [0]
14 Output: [0]
15 */
16
17 /*
18 APPROACH:-
19 -> The idea is while traversing the array if we found any zero then we have to swap it with
   next non-zero
20 */
21
22 // CODE:-
23 // function to find the next non-zero element
24 int next_nonzero(vector<int> &a, int &j)
25 {
26     while (j < a.size())
27     {
28         if (a[j] != 0)
29             return j;
30         j++;
31     }
32     return -1;
33 }
34 void moveZeroes(vector<int> &nums)
35 {
36     int j = -1; // is to find the next non zero element
37     // i signifies that upto here all elements are non-zero
38     for (int i = 0; i < nums.size(); i++)
39     {
40         if (nums[i] != 0)
41             continue;
42         if (j == -1)
43             j = i + 1;
44         int nxt_non0 = next_nonzero(nums, j);
45         if (nxt_non0 == -1)
46             return;
47         swap(nums[i], nums[nxt_non0]);
48     }
49 }
50
```


1.Easy\09.Union_of_2_sorted_arrays.cpp

```

1  /*
2  QUESTION:-
3  Union of two arrays can be defined as the common and distinct elements in the two arrays.
4  Given two sorted arrays of size n and m respectively, find their union.
5
6
7  Example 1:
8
9  Input:
10 n = 5, arr1[] = {1, 2, 3, 4, 5}
11 m = 3, arr2 [] = {1, 2, 3}
12 Output: 1 2 3 4 5
13 Explanation: Distinct elements including
14 both the arrays are: 1 2 3 4 5.
15
16
17 Example 2:
18
19 Input:
20 n = 5, arr1[] = {2, 2, 3, 4, 5}
21 m = 5, arr2[] = {1, 1, 2, 3, 4}
22 Output: 1 2 3 4 5
23 Explanation: Distinct elements including
24 both the arrays are: 1 2 3 4 5.
25 */
26
27 /*
28 APPROACH:-
29 -> Take two pointer i and j where i is for arr1 and j is for arr2 and traverse
30 -> While traversing 3 cases arises
31     -> arr1[ i ] == arr2[ j ]
32         Here we found a common element, so insert only one element in the union.
33         Let's insert arr[i] in union and whenever we insert element we increment pointer
34         while pointer is not equal to the inserted element
35     -> arr1[i]<arr2[j]
36         Here insert arr[i]
37     -> arr1[i]>arr2[j]
38         Here insert arr2[j]
39 -> Now check if elements of any array is left to traverse then traverse that array
40 */
41 // CODE:-
42 vector<int> findUnion(int arr1[], int arr2[], int n, int m)
43 {
44     int i = 0; // i to keep track in arr1
45     int j = 0; // j to keep track in arr2
46     vector<int> ans;
47
48     while (i < n && j < m)
49     {
50
51         if (arr1[i] < arr2[j])

```

```
52     {
53         ans.push_back(arr1[i++]);
54         while (i < n && arr1[i] == arr1[i - 1])
55             i++;
56     }
57     else if (arr2[j] < arr1[i])
58     {
59         ans.push_back(arr2[j++]);
60         while (j < m && arr2[j] == arr2[j - 1])
61             j++;
62     }
63     // means arr1[i] = arr2[j] in that case we can insert anyone
64     else
65     {
66         ans.push_back(arr1[i++]);
67         j++;
68         while (i < n && arr1[i] == arr1[i - 1])
69             i++;
70         while (j < m && arr2[j] == arr2[j - 1])
71             j++;
72     }
73 }
74
75 while (i < n)
76 {
77     ans.push_back(arr1[i++]);
78     while (i < n && arr1[i] == arr1[i - 1])
79         i++;
80 }
81 while (j < m)
82 {
83     ans.push_back(arr2[j++]);
84     while (j < m && arr2[j] == arr2[j - 1])
85         j++;
86 }
87
88 return ans;
89 }
90
91 // TIME COMPLEXITY = O(N+M)
92 // SPACE COMPLEXITY = O(0)
```

1.Easy\10.Missing_number.cpp

```
1  /*
2  QUESTION:-
3  Given an array nums containing n distinct numbers in the range [0, n], return the only number
  in the range that is missing from the array.
4
5  Example 1:
6
7  Input: nums = [3,0,1]
8  Output: 2
9  Explanation: n = 3 since there are 3 numbers, so all numbers are in the range [0,3]. 2 is the
  missing number in the range since it does not appear in nums.
10 Example 2:
11
12 Input: nums = [0,1]
13 Output: 2
14 Explanation: n = 2 since there are 2 numbers, so all numbers are in the range [0,2]. 2 is the
  missing number in the range since it does not appear in nums.
15 */
16
17 /*
18 APPROACH:-
19 -> Calculate the optimum sum i.e. sum when all elements were present
20 -> Calculate the actual array's sum
21 -> Return the optimum sum - actual sum
22 */
23
24 // CODE:-
25 int missingNumber(vector<int> &nums)
26 {
27     int n = nums.size();
28     long long optimum_sum = (n * (n + 1)) / 2; // the sum if no number is absent
29     long long actual_sum = 0;
30     for (auto it : nums)
31     {
32         actual_sum += it;
33     }
34     return optimum_sum - actual_sum;
35 }
36
37 // TIME COMPLEXITY = O(N)
38 // SPACE COMPLEXITY = O(0)
```

1.Easy\11.Max_consecutive_1's.cpp

```
1  /*
2  QUESTION:-
3  Given an array nums containing n distinct numbers in the range [0, n], return the only number
  in the range that is missing from the array.
4
5  Example 1:
6
7  Input: nums = [3,0,1]
8  Output: 2
9  Explanation: n = 3 since there are 3 numbers, so all numbers are in the range [0,3]. 2 is the
  missing number in the range since it does not appear in nums.
10 Example 2:
11
12 Input: nums = [0,1]
13 Output: 2
14 Explanation: n = 2 since there are 2 numbers, so all numbers are in the range [0,2]. 2 is the
  missing number in the range since it does not appear in nums.
15 */
16
17 /*
18 APPROACH:-
19 -> Traverse the entire array and within it run a loop while element's are equal to 1 and
  store the count
20 -> Update the ans as max(ans,cnt)
21 */
22
23 // CODE:-
24 int findMaxConsecutiveOnes(vector<int> &nums)
25 {
26     int ans = 0;
27     for (int i = 0; i < nums.size(); i++)
28     {
29         int cnt = 0; // to store the count of consecutive 1's
30         while (i < nums.size() && nums[i] == 1)
31         {
32             cnt++;
33             i++;
34         }
35         ans = max(ans, cnt);
36     }
37     return ans;
38 }
39
40 // TIME COMPLEXITY = O(N)
41 // SPACE COMPLEXITY = O(0)
42
```

1.Easy\12.Longest_subarray_with_given_sum.cpp

```
1  /*
2  QUESTION:-
3  You are given an array 'A' of size 'N' and an integer 'K'. You need to print the length of
  the longest subarray of array 'A' whose sum = 'K'.
4  Example:
5  Input: 'N' = 7 'K' = 3
6  'A' = [1, 2, 3, 1, 1, 1, 1]
7  Output: 3
8  Explanation: Subarrays whose sum = '3' are:
9  [1, 2], [3], [1, 1, 1], [1, 1, 1]
10 Here, the length of the longest subarray is 3, which is our final answer.
11 */
12
13 /*
14 APPROACH:-
15 -> Use sliding window approach using two pointers start and end
16 -> Run a loop to traverse the entire array add from end and subtract from start when sum>k
17 -> If sum==k then, update the ans now, window size = end-start+1
18 */
19
20 // CODE:-
21 int longestSubarrayWithSumK(vector<int> a, long long k)
22 {
23     int start = 0;
24     int ans = 0;
25     long long sum = 0;
26     int n = a.size();
27
28     for (int end = 0; end < n; end++)
29     {
30         sum += a[end];
31         while (sum > k)
32         {
33             sum -= a[start];
34             start++;
35         }
36         if (sum == k)
37         {
38             ans = max(ans, end - start + 1);
39         }
40     }
41     return ans;
42 }
43
44 // TIME COMPLEXITY = O(N)
45 // SPACE COMPLEXITY = O(0)
```

1.Easy\13.Find_element_present_only_once.cpp

```
1  /*
2  QUESTION:-
3  Given a non-empty array of integers nums, every element appears twice except for one. Find
   that single one.
4  You must implement a solution with a linear runtime complexity and use only constant extra
   space.
5
6
7  Example 1:
8  Input: nums = [2,2,1]
9  Output: 1
10
11 Example 2:
12 Input: nums = [4,1,2,1,2]
13 Output: 4
14 */
15
16 /*
17 APPROACH:-
18 -> We can use XOR operation as we know xor cancels out the same elements
19 -> Intial xr=0 then traverse the entire array and xor each element with xr
20 -> Since only one element is present once and all other are present twice so the remaining
   element would be the
21     one which is present only once cause all other gets cancels out
22 */
23
24 // CODE:-
25 int singleNumber(vector<int> &nums)
26 {
27     int xr = 0;
28     for (int i = 0; i < nums.size(); i++)
29     {
30         xr = nums[i] ^ xr;
31     }
32     return xr;
33 }
34
35 // TIME COMPLEXITY = O(N)
36 // SPACE COMPLEXITY = O(0)
```