

**2.Medium\05.Number\_of\_subarray\_sum\_equal\_k.cpp**

```
1  /*
2  QUESTION:
3  Given an array of integers nums and an integer k, return the total number of subarrays whose
   sum equals to k.
4
5  Example:
6  Input: nums = [1,1,1], k = 2
7  Output: 2
8
9  APPROACH:
10 To find the total number of subarrays with sum equal to k, we can use the technique of prefix
   sum along with a hashmap.
11 1. Initialize a variable `count` to keep track of the count of subarrays with sum equal to k.
12 2. Initialize a variable `prefixSum` to keep track of the prefix sum while iterating through
   the array.
13 3. Initialize a hashmap `sumCount` to store the frequency of prefix sums encountered so far.
14 4. Set the initial prefix sum to 0 and set its count to 1 in the `sumCount` hashmap.
15 5. Iterate through the array and update the prefix sum by adding each element.
16 6. Check if the current prefix sum minus k exists in the `sumCount` hashmap. If it does, add
   the count of that prefix sum to the `count` variable.
17 7. Increment the count of the current prefix sum in the `sumCount` hashmap.
18 8. Finally, return the `count` variable as the total number of subarrays with sum equal to k.
19
20 CODE:
21 */
22
23 int subarraySum(vector<int> &nums, int k)
24 {
25     int pref_sum = 0;
26     unordered_map<int, int> mp;
27     int ans = 0;
28
29     for (int i = 0; i < nums.size(); i++)
30     {
31         pref_sum += nums[i];
32
33         if (pref_sum == k)
34             ans++;
35
36         if (mp.find(pref_sum - k) != mp.end())
37         {
38             ans += mp[pref_sum - k];
39         }
40
41         mp[pref_sum]++;
42     }
43
44     return ans;
45 }
46
47 /*
48 TIME COMPLEXITY: O(n), where n is the size of the input array nums.
```

```
49 | SPACE COMPLEXITY:  $O(n)$ , as we are using a hashmap to store the prefix sums and their  
   | corresponding counts.  
50 | */  
51 |
```