

3.Hard\03.3_sum.cpp

```
1  /*
2  QUESTION:
3  Given an integer array nums, return all the triplets [nums[i], nums[j], nums[k]] such that i
   != j, i != k, and j != k, and nums[i] + nums[j] + nums[k] == 0.
4
5  Example:
6  Input: nums = [-1,0,1,2,-1,-4]
7  Output: [[-1,-1,2],[-1,0,1]]
8  Explanation:
9  nums[0] + nums[1] + nums[2] = (-1) + 0 + 1 = 0.
10 nums[1] + nums[2] + nums[4] = 0 + 1 + (-1) = 0.
11 nums[0] + nums[3] + nums[4] = (-1) + 2 + (-1) = 0.
12 The distinct triplets are [-1,0,1] and [-1,-1,2].
13 Notice that the order of the output and the order of the triplets does not matter.
14 */
15
16 /*
17 APPROACH:
18 To find all triplets that sum up to zero, we can follow these steps:
19 1. Sort the input array in non-decreasing order.
20 2. Iterate through the array and fix the first element as nums[k] (where k = 0 to n-1).
21 3. Use two pointers (i and j) to find the other two elements such that nums[i] + nums[j] = -
   nums[k].
22 4. Move the pointers accordingly to find all possible triplets.
23 5. Skip duplicate elements to avoid duplicate triplets.
24 6. Return the resulting triplets.
25 */
26
27 vector<vector<int>> threeSum(vector<int> &nums)
28 {
29     vector<vector<int>> ans;
30     sort(nums.begin(), nums.end());
31
32     for (int k = 0; k < nums.size(); k++)
33     {
34         int i = k + 1;
35         int j = nums.size() - 1;
36         int target = -nums[k];
37
38         while (i < j)
39         {
40             int sum = nums[i] + nums[j];
41
42             if (sum == target)
43             {
44                 ans.push_back({nums[k], nums[i], nums[j]});
45                 i++;
46                 j--;
47
48                 // Skip duplicate elements
49                 while (i < j && nums[i] == nums[i - 1])
50                     i++;
```

```
51         while (i < j && nums[j] == nums[j + 1])
52             j--;
53     }
54     else if (sum < target)
55     {
56         i++;
57     }
58     else
59     {
60         j--;
61     }
62 }
63
64 // Skip duplicate elements
65 while (k + 1 < nums.size() && nums[k + 1] == nums[k])
66     k++;
67 }
68
69 return ans;
70 }
71
72 /*
73 TIME COMPLEXITY:  $O(n^2)$ , where  $n$  is the size of the input array.
74 The sorting step takes  $O(n \log n)$ , and the two-pointer traversal takes  $O(n^2)$  in the worst
75 case.
76 Hence, the overall time complexity is  $O(n^2)$ .
77
78 SPACE COMPLEXITY:  $O(1)$ , as we are using a constant amount of extra space for storing the
79 output and variables.
80 */
```