Project of LATEX Page 97 to 100

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January 12, 2021

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Solution:

| Next State | | | |
|-------------------|-----------------|----------------------------|--|
| Present State | 0 | 1 | |
| $[q_0]$ | $[q_0, q_3]$ | $[q_0, q_1]$ | |
| $[q_0,q_3]$ | $[q_0,q_3,q_4]$ | $[q_0,q_1]$ | |
| $[q_0, q_3, q_4]$ | $[q_0,q_3,q_4]$ | $\left[q_0,q_1,q_4\right]$ | |
| $[q_0,q_1,q_4]$ | $[q_0,q_3,q_4]$ | $[q_0,q_1,q_4]$ | |
| $[q_0,\!q_1]$ | $[q_0,q_3]$ | $[q_0,q_1]$ | |

For simplification, let us replace $[q_0]$ by A, $[q_0,q_3,]$ by B, $[q_0,q_3,q_4]$ by C, $[q_0,q_1,q_4]$ by D, and $[q_0,q_1]$ by E. Here, A is the initial state, and C and D are the final states as they contain the state q_0 . The simplified DFA is

| | Next State | | |
|---------------|--------------|--|--------------|
| Present State | 0 | | 1 |
| A | В | | E |
| В | \mathbf{C} | | \mathbf{E} |
| \mathbf{C} | \mathbf{C} | | D |
| D | \mathbf{C} | | D |
| E | В | | \mathbf{E} |

6. Convert the following NFA to an equivalent DFA. Solution :

| Σ | | |
|--------|-------|-----------|
| States | 0 | 1 |
| q_0 | q_0 | q_0,q_1 |
| q_1 | q_2 | q_2 |
| q_0 | _ | q_1 |

($[q_0]$ is the initial state and $[q_1]$ is the final state) **Solution:** Conversion is done in the following ways:

| \sum | | |
|-------------------|----------------------|----------------------------|
| States | 0 | 1 |
| $[q_0]$ | $[q_0]$ | $[q_0,q_1]$ |
| $[q_0,q_1]$ | $[q_0, q_2]$ | $[q_0, q_1, q_2]$ |
| $[q_0, q_1, q_2]$ | $[q_0, q_2]$ | $[q_0, q_1, q_2]$ |
| $[q_0,q_2]$ | $\lfloor q_0 floor$ | $\left[q_0,q_1,q_2\right]$ |

Rename $[q_0]$ as A, $[q_0,q_1]$ as B, $[q_0,q_1,q_2]$ as C, and $[q_0,q_2]$ as D. The beginning state is A, and final states are B and C.

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| Σ | | |
|--------------|---|--------------|
| States | 0 | 1 |
| A | A | В |
| В | D | \mathbf{C} |
| \mathbf{C} | D | \mathbf{C} |
| D | A | С |

7. Convert the following NFA to an equivalent DFA. [UPTU 2005]

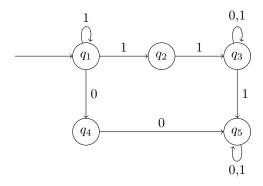
| $\sum_{}$ | | |
|--------------|---------|-----------|
| States | 0 | 1 |
| p | {q,s} | {q} |
| q | $\{e\}$ | $\{q,r\}$ |
| \mathbf{r} | $\{s\}$ | $\{p\}$ |
| \mathbf{s} | Ø | $\{p\}$ |

where p is the initial state and q and s are the final states. Solution:

| Σ | | |
|-------------|-------------|-------------|
| States | 0 | 1 |
| {p} | $\{q,s\}$ | {q} |
| $\{q\}$ | $\{r\}$ | $\{q,r\}$ |
| $\{r\}$ | $\{s\}$ | $\{p\}$ |
| $\{s\}$ | Ø | {p} |
| $\{q,r\}$ | $\{r,s\}$ | $\{p,q,r\}$ |
| $\{r,s\}$ | $\{s\}$ | $\{p\}$ |
| $\{p,q,r\}$ | $\{q,r,s\}$ | $\{p,q,r\}$ |
| $\{q,r,s\}$ | $\{r,s\}$ | $\{p,q,r\}$ |
| $\{q,s\}$ | $\{r\}$ | $\{p,q,r\}$ |
| {Ø} | {Ø} | {Ø} |

Here $\{p\}$ is the beginning state and $\{q\}$, $\{s\}$, $\{q, r\}$, $\{r, s\}$, $\{p, q, r\}$, $\{q, r, s\}$, and $\{q, s\}$ are the final states. \emptyset is the dead state.

8. Construct a DFA equivalent to the following NDFA given in the following figure. [UPTU 2004]



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Solution: The tabular representation of the NDFA is

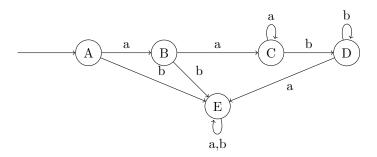
| Next State | | |
|---------------|-------|---------------|
| Present State | 0 | 1 |
| q_0 | q_3 | $\{q_0,q_1\}$ |
| q_1 | Ø | q_2 |
| q_2 | q_2 | $\{q_2,q_4\}$ |
| q_3 | q_4 | Ø |
| q_4 | q_4 | q_4 |

 $(q_0$ is the initial state and q_4 is the final state) The corresponding DFA is

| \sum | | |
|-------------------|-------------------|-------------------|
| States | 0 | 1 |
| q_0 | $\{q_3\}$ | $\{q_0,q_1\}$ |
| $\{q_3\}$ | $\{q_4\}$ | $\{\varnothing\}$ |
| $\{q_4\}$ | $\{q_4\}$ | $\{q_4\}$ |
| $\{q_0,q_1\}$ | $\{q_2\}$ | $\{q_2,q_4\}$ |
| $\{q_2\}$ | $\{q_2\}$ | $\{q_2, q_4\}$ |
| $\{q_2,q_4\}$ | $\{q_2,q_4\}$ | $\{q_2,q_4\}$ |
| $\{\varnothing\}$ | $\{\varnothing\}$ | $\{\varnothing\}$ |

Here $\{q_0\}$ is the beginning state, and $\{q_4\}$, and $\{q_0,q_1\}$ are the final states. (Draw a transitional diagram to complete the answer.)

9. Find the minimal DFAs for the language $L = \{a^n b^m, n \ge 2, m \ge 1\}$ Solution: All 'a' will appear before 'b'. There is at least 2 'a' and 1 'b'. The DFA is the following.



10. Design a DFA for the language L = $\{0^m1^n, m \geq 0, n \geq 1\}$ [JNTU 2007]

Solution: All '0's will appear before '1'. There is at least one '1', but the number of '0's may be zero. The DFA is shown in Fig. 3.58.

11. Construct a DFA which accepts the set of all binary strings that, interpreted as the binary representation of an unsigned decimal integer, is divisible by 5.

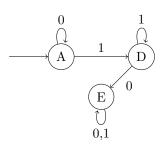


Fig. 3.58

[WBUT 2008]

Solution: For Mod 5, the remainders are 0, 1, 2, 3, and 4. We can assign the states as q_0 , q_1 , q_2 , q_3 , and q_4 . For any binary string, if we add a bit at LSB, then the previous value becomes doubled. (Let the string be 101. The decimal value is 5. If we add another 1 at LSB, the string becomes 1011. The decimal value of the previous 101 becomes 10.) In general, we can write that 'n' becomes 2n + b, where n is the previous number and b is the added bit.

$$(2n + b) \bmod 5 = 2n \bmod 5 + b \bmod 5$$

As b is either 0 or 1, b mod 5 = b. 2n mod 5 is any one of 0, 1, 2, 3, or 4, i.e., 2 X (state number) + a. For this machine, the input alphabets are 0 and 1.

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\begin{array}{l} \delta(q_0,\,0) \to 2\times 0 + 0 = 0 \; \text{means} \; q_0 \\ \delta(q_0,\,1) \to 2\times 0 + 1 = 1 \; \text{means} \; q_1 \\ \delta(q_1,\,0) \to 2\times 1 + 0 = 2 \; \text{means} \; q_2 \\ \delta(q_1,\,1) \to 2\times 1 + 1 = 3 \; \text{means} \; q_3 \\ \delta(q_2,\,0) \to 2\times 2 + 0 = 4 \; \text{means} \; q_4 \\ \delta(q_2,\,1) \to 2\times 2 + 1 = 5\%5 = 0 \; \text{means} \; q_0 \\ \delta(q_3,\,0) \to 2\times 3 + 0 = 6\%5 = 1 \; \text{means} \; q_1 \\ \delta(q_3,\,1) \to 2\times 3 + 1 = 7\%5 = 2 \; \text{means} \; q_2 \\ \delta(q_4,\,0) \to 2\times 4 + 0 = 8\%5 = 3 \; \text{means} \; q_3 \\ \delta(q_4,\,1) \to 2\times 4 + 1 = 9\%5 = 4 \; \text{means} \; q_4 \end{array}
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