## 2017-09-14

- Function: for every  $x \in X$ , there is exactly one  $y \in Y$  such that  $(x, y) \in R$ .
  - Injective (one-to-one;  $|X| \le |Y|$ ): for every  $y \in Y$ , there is at most one  $x \in X$  such that f(x) = y.
  - Surjective (onto;  $|X| \ge |Y|$ ): for every  $y \in Y$ , there is at least one  $x \in X$  such that f(x) = y.
  - **Bijective** (one-to-one and onto; |X| = |Y|): both injective and surjective.
- Countable: there is an injective function from X to  $\mathbb{N}$ .
  - Examples of countably infinite sets:  $\mathbb{Q}$ ,  $\mathbb{Z}$ ,  $\mathbb{N}$ .
  - Examples of uncountably infinite sets:  $\mathbb{C}$ ,  $\mathbb{R}$ ,  $2^{\mathbb{N}}$ .
  - Proof of uncountability: Cantor's diagonalization.
- Alphabet ( $\Sigma$ ): a finite set of symbols.
- Word (w): a finite sequence of symbols over  $\Sigma$ .
  - $\epsilon$  (empty word): the word of length 0.
  - $\Sigma^n$ : the set of all words of length *n* over  $\Sigma$ .
  - $\Sigma^* = \bigcup_{n \ge 0} \Sigma^n$ : the set of all finite words over  $\Sigma$ .
  - $\circ \quad \Sigma^+ = \cup_{n \ge 1} \Sigma^n = \Sigma^* \setminus \{\epsilon\}$
- Language (L): a subset of  $\Sigma^*$ , i.e.  $L \subseteq \Sigma^*$ .
  - Given two languages  $L_1$  and  $L_2$  over  $\Sigma$ ,  $L_1L_2 = \{w_1w_2 | w_1 \in L_1, w_2 \in L_2\}$  (similar to Cartesian product).
- Other examples of countably infinite sets:
  - The set of all finite sequences over  $\Sigma$ , i.e.  $\Sigma^*$ .
  - The set of all computer programs. ∵ computer programs have (1) finite symbols, and (2) finite length.
- Other examples of uncountably infinite sets:
  - The set of all languages over  $\Sigma$ , i.e.  $2^{\Sigma^*}$ .
- Question: For every language, we can write a computer program to detect it. (False)