

Experimental verification : log-barrier interior point methods are not strongly polynomial

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Abstract

Interior-point methods are a very import class of algorithms that solve linear optimization problems. They are widely believed to be polynomial. In this report, we follow on a previous work proving that log-barrier interior point methods are not strongly polynomial presenting a counter-example. We analyzed numerically the behavior of a predictor-corrector algorithm on the counter-example to confirm the result shown previously. In particular, we present an explicit method of constructing a suitable initial point for the algorithm over Puiseux series.

Notations

We denote the scalar product between x and y by $x.y := x^T y$.

In all what follow, we take $n, m \in \mathbb{N}$.

$e := (1 \dots 1)^T \in \mathbb{R}^n$

For a given $x, s, X := \text{Diag}(x_1, \dots, x_n), S := \text{Diag}(s_1, \dots, s_n)$.

1 Interior point method

In all the following we are interested in solving the following linear problem:

$$(P) \begin{cases} \inf_{x \in \mathbb{R}^n} c.x \\ Ax = b \\ x \geq 0 \end{cases}$$

where $c \in \mathbb{R}^n$ called the *cost function*, $b \in \mathbb{R}^m$ (usually $m \leq n$), and $A \in \mathbb{R}^{n \times m}$. $x \geq 0$ means x is in the positive orthan, i.e that $x_i \geq 0$ for all $1 \leq i \leq n$. The admissible set of our problem is noted:

$$\mathcal{F}_P = \{x \in \mathbb{R}^n : Ax = b, x \geq 0\}$$

The strictly admissible set is noted:

$$\mathcal{F}_P^s = \{x \in \mathbb{R}^n : Ax = b, x > 0\}$$

The dual of (P) is written as follow:

$$(D) \begin{cases} \sup_y b.y \\ A^T y + s = c \\ s \geq 0 \end{cases}$$

We also note respectively his admissible set and strictly admissible set:

$$\begin{aligned}\mathcal{F}_D &= \{(y, s) \in \mathbb{R}^m \times \mathbb{R}^n : A^T y + s = c, s \geq 0\} \\ \mathcal{F}_D^s &= \{(y, s) \in \mathbb{R}^m \times \mathbb{R}^n : A^T y + s = c, s > 0\}\end{aligned}$$

It has been shown [2] that $x \in \mathbb{R}^n$ is solution to (P) if and only if there exist a couple $(y, s) \in \mathbb{R}^m \times \mathbb{R}^n$ such that we have:

$$\begin{cases} A^T y + s = c, s \geq 0 \\ Ax = b, x \geq 0 \\ x.s = 0 \end{cases} \quad (1)$$

The same result holds for the dual: (y, s) is an optimal solution of (D) if and only if there exist $x \in \mathbb{R}^n$ such that 1 holds. We denote a primal-dual solution by $z = (x, y, s)$.

1.1 The central path

The positivity constrains over the solution are one of the main difficulties of linear programs. Interior point methods consist of transforming the problem into several problems easier to solve and leading to the optimal solution. Instead of solving directly our linear problem with positivity constrains, we *penalize* the constrains by adding a *log-barrier* to the objective function as follow:

$$(P_\mu) \begin{cases} \inf_x c.x + \mu \text{lb}(x) \\ Ax = b \end{cases} \quad (2)$$

where μ is a positive parameter and $\text{lb} : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$ is defined by:

$$\text{lb}(x) = \begin{cases} -\sum_{i=1}^n \log(x_i), & \text{if } x > 0 \\ +\infty, & \text{else} \end{cases}$$

The *log-barrier* function *pushes* the solution into the positive orthant by *penalizing* the objective function when we get close to the borders, as $\text{lb}(x) \rightarrow +\infty$ when $x_i \rightarrow 0$. The dual of this problem is written as:

$$(D_\mu) \begin{cases} \sup b.y - \mu \text{lb}(s) \\ A^T y + s = c \end{cases} \quad (3)$$

Solving $(P_\mu) - (D_\mu)$ for all $\mu > 0$ gives an application:

$$C_\mu : \mu \in]0, +\infty[\rightarrow z(\mu) := (x(\mu), y(\mu), s(\mu)) \in \mathcal{F}^s$$

where $\mathcal{F}^s := \mathcal{F}_P^s \times \mathcal{F}_D^s$ is the primal-dual strictly feasible set.

We notice that (P_μ) and (D_μ) becomes (P) and (D) when $\mu \rightarrow 0$. We show under reasonable hypothesis that we have as we may expect $z(\mu) \rightarrow z = (x, y, s)$ optimal solution of (P) and (D) . Therefore, the application C_μ construct a path converging to the optimal solution of the problem. Interior point methods consist of following this path to the optimal solution.

The equivalent of equation 1 for a penalized problem can be rewritten in the case of a penalized problem with parameter μ as:

$$\begin{cases} A^T y + s = c, s > 0 \\ Ax = b, x > 0 \\ Xs = \mu e \end{cases} \quad (4)$$

1.2 Predictor-Corrector

Predictor-Corrector algorithm is an example of interior point methods. It is based on an opportunistic way of following the central path towards the optimal solution. On every step of the algorithm, we start by following a Newton direction to the optimal solution by making sure that we don't deviate too much from the central path (prediction) and then correct the path by getting closer to the central path (correction).

Let's start by defining the *duality measure* that quantify how close we are from the optimal solution. For a given $z \in \mathcal{F}^s$, we define the duality measure as:

$$\bar{\mu}(z) = \frac{x \cdot s}{n}$$

Notice that z is in the central path, if and only if $Xs = \bar{\mu}(z)e$.

To limit the prediction phase and curb it from going too far from the central path, we need to define a neighborhood around the central path that we shouldn't leave. We will use two types of neighborhoods in the following depending on a parameter $\theta > 0$, the norm-2 neighborhood:

$$V_2(\theta) := \{z \in \mathcal{F}^s : \|Xs - \bar{\mu}(z)e\|_2 \leq \theta \bar{\mu}(z)\}$$

and the wide neighborhood:

$$V_\infty^-(\theta) := \{z \in \mathcal{F}^s : \|Xs - \bar{\mu}(z)e\|_\infty \leq \theta \bar{\mu}(z)\}$$

where $\|v\|_\infty = \max(0, \max_i(-v_i))$. The latter can be rewritten as:

$$V_\infty^-(\theta) := \{z \in \mathcal{F}^s : Xs \geq (1 - \theta)\bar{\mu}(z)e\}$$

In both cases, we will note the neighborhood as $V(\theta)$.

Algorithm 1 Predictor-corrector algorithm

We consider two sets $V(\theta)$ and $V(\theta')$ with $0 < \theta' < \theta$.

Suppose that the actual point is $z := (x, y, s) \in V(\theta') \subset \mathcal{F}^s$. The following point $z_+ \in V(\theta')$ is obtained by two steps.

1. Prediction

- (a) Compute Newton direction d from z toward the optimal solution such that we stay in the feasible set \mathcal{F}^s .
- (b) Compute the largest jump $\alpha \in [0, 1]$ such that $z' = z + \alpha d \in V(\theta)$.
- (c) Prediction point: z' .

2. Correction

- (a) Compute Newton direction d' from z' toward the central path such that we stay in the feasible set \mathcal{F}^s .
 - (b) Correction point: $z_+ = z' + d'$.
-

Step 1.a We compute $d = \begin{pmatrix} dx \\ dy \\ ds \end{pmatrix}$ as a solution of [2]:

$$\begin{pmatrix} 0 & A^T & I \\ A & 0 & 0 \\ S & 0 & X \end{pmatrix} \begin{pmatrix} dx \\ dy \\ ds \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ -Xs \end{pmatrix}$$

Step 1.b When $V = V_2$, we compute α by solving:

$$\|X(\alpha)s(\alpha) - \bar{\mu}(z(\alpha))e\|_2^2 = \theta'^2 \bar{\mu}(z(\alpha))^2$$

where $z(\alpha) = z + \alpha d$.

When $V = V_\infty^-$, we need to compute $\alpha^* = \max\{\alpha \in [0, 1] \mid X(\alpha)s(\alpha) \geq (1 - \theta)\bar{\mu}(z(\alpha))e\} = \max\{\alpha \in [0, 1] \mid \min_i x_i(\alpha)s_i(\alpha) = (1 - \theta')\bar{\mu}(z(\alpha))\}$. In order to do that, we compute $\alpha^{**} = \min_i \min\{\alpha_i \in [0, 1] \mid x_i(\alpha_i)s_i(\alpha_i) = (1 - \theta')\bar{\mu}(z(\alpha_i))\}$ which is easier to compute (solving n second degree polynomial equations). Given lemma 2 in [1], we have $z + \alpha d \in V$ for all $\alpha \in [0, \alpha^*]$ and therefore $\alpha^* = \alpha^{**}$.

Step 2.b We compute $d' = \begin{pmatrix} d'x \\ d'y \\ d's \end{pmatrix}$ as a solution of [2]:

$$\begin{pmatrix} 0 & A^T & I \\ A & 0 & 0 \\ S & 0 & X \end{pmatrix} \begin{pmatrix} d'x \\ d'y \\ d's \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ \bar{\mu}(z)e - Xs \end{pmatrix}$$

The correction of the algorithm is shown in the case of $\frac{\theta'^2}{1-\theta'} \leq \sqrt{8}\theta$ [2]. We chose in the following $\theta = \frac{1}{2}$ and $\theta' = \frac{1}{4}$.

2 Log-barrier interior point methods are not strongly polynomial

In their article, Stéphane Gaubert, Xavier Allamigeon and their co-authors prove that log-barrier Interior-Point methods are not strongly polynomial by presenting a counter example instance inspired by a Game-Theory problem.

The counter example is the following:

Given an integer number $r > 1$ and an Archimedean parameter t we consider the following Linear Program:

$$\begin{aligned} & \text{Minimize} && x_1 \\ & \text{Subject to} && x_1 \leq t^2 \\ & && x_2 \leq t \\ & && x_{2j+1} \leq t x_{2j-1}, \quad x_{2j+1} \leq t x_{2j}, \quad 1 \leq j < r \\ & && x_{2j+2} \leq t^{1-1/2^j} (x_{2j-1} + x_{2j}), \quad 1 \leq j < r \\ & && x_{2r-1} \geq 0, \quad x_{2r} \geq 0 \end{aligned} \quad \text{LW}_r(t)$$

$\text{LW}_r(t)$ has $2r$ variables and $3r + 1$ constraints and can be transformed to fit the form of (P) by adding slack variables $(w_j)_{1 \leq j \leq r}$. When transformed, $\text{LW}_r(t)$ has $5r + 1$ variables and $3r + 1$ constraints.

$\text{LW}_r(t)$ can be rewritten in a matricial form as follows:

$$\begin{aligned} & \text{Minimize} && c.x \\ & \text{Subject to} && A(t)x \leq b(t) \\ & && x \geq 0 \end{aligned}$$

2.1 Field of Puiseux series and Tropical Geometry tools

We denote by \mathbb{K} the field of real and absolutely convergent Puiseux series, ie: the field of elements $f = \sum_{\alpha \in \mathbb{R}} a_\alpha t^\alpha$ such that:

$$\begin{cases} (i) a_\alpha \in \mathbb{R} \text{ for all } \alpha \in \mathbb{R}, \\ (ii) \text{supp}(\mathbf{f}) = \{\alpha \in \mathbb{R} : a_\alpha \neq 0\} \text{ is either finite or has } -\infty \text{ as the only accumulation point} \end{cases} \quad (5)$$

The definition of the field \mathbb{K} guarantees that every series $\mathbf{f} = \sum_{\alpha \in \mathbb{R}} a_\alpha t^\alpha$ has a leading term $a_{\alpha_0} t^{\alpha_0}$ where α_0 is defined as $\alpha_0 := \max(\text{supp}(\mathbf{f}))$. This allows us to define a total ordering on \mathbb{K} where an element is positive if its leading term has a positive coefficient α .

The *valuation* map $val : \mathbb{K} \rightarrow \mathbb{R} \cup \{-\infty\}$ that allows us to establish a link between the problem $\mathbf{LW}_r(t)$ and Tropical Geometry is defined by: $val(\mathbf{f})$ for $\mathbf{f} \in \mathbb{K}$ is the greatest element α_0 of $\text{supp}(\mathbf{f})$ or equivalently:

$$val(\mathbf{f}) = \lim_{t \rightarrow +\infty} \log_t |\mathbf{f}(t)|$$

This *valuation* map has the following particular properties :

$\forall \mathbf{f}, \mathbf{g} \in \mathbb{K}$:

$$val(\mathbf{f} + \mathbf{g}) \leq \max(val(\mathbf{f}), val(\mathbf{g})) \quad \text{and} \quad val(\mathbf{f}\mathbf{g}) = val(\mathbf{f}) + val(\mathbf{g}). \quad (6)$$

where the first inequality becomes an equality if the leading terms of \mathbf{f} and \mathbf{g} don't cancel which is the case when \mathbf{f}, \mathbf{g} are positive Puiseux series (i.e in \mathbb{K}_+). $val : (\mathbb{K}_+, +, \times) \rightarrow (\mathbb{T}, \oplus, \odot)$ where $\mathbb{T} := \mathbb{R} \cup \{-\infty\}$ is then some kind of morphisme.

We can then define Puiseux polyhedron as we do in a real numbers setting ([1] for a precise justification)

$$\mathcal{P} = \{\mathbf{x} \in \mathbb{K}^d : \mathbf{A}\mathbf{x} \leq \mathbf{b}\}, \quad (7)$$

where $\mathbf{A} \in \mathbb{K}^{p \times d}$, $\mathbf{b} \in \mathbb{K}^p$ (with $p \geq 0$), and \leq stands for the partial order over \mathbb{K}^p .

As explained in [1], this allows us to tropicalize the problem $\mathbf{LW}_r(t)$ and to compute it's exact tropical central path $\mathcal{C}^{trop}(\lambda)$ which in some sense describes the behaviour of the real central path when t is sufficiently large, and again according to [1]. The combinatorial analysis of the tropical central path gives us an yields a lower bound on the curvature of the central path of the parametric family of linear problems over the reals $(\mathbf{LW}_r(t))_{t>0}$

The tropical central path is given by the equations:

for x :

$$\begin{aligned} x_1^\lambda &= \min(\lambda, 2) \\ x_2^\lambda &= 1 \\ x_{2j+1}^\lambda &= 1 + \min(x_{2j-1}^\lambda, x_{2j}^\lambda), \quad 1 \leq j < r \\ x_{2j+2}^\lambda &= (1 - 1/2^j) + \max(x_{2j-1}^\lambda, x_{2j}^\lambda), \quad 1 \leq j < r \end{aligned}$$

for the slack variables w :

$$\begin{aligned} w_1^\lambda &= 2, \quad w_2^\lambda = 1, \quad w_3^\lambda = 1 + \min(\lambda, 2) \\ \forall 1 < j < r \quad w_{3j}^\lambda &= 1 + \min(w_{3j-3}^\lambda, w_{3j-2}^\lambda) \\ \forall 1 \leq j < r \quad w_{3j+1}^\lambda &= 1 + w_{3j-1}^\lambda \\ w_{3j+2}^\lambda &= (1 - \frac{1}{2^j}) + \max(w_{3j}^\lambda - 1, w_{3j-1}^\lambda) \\ w_{3r}^\lambda &= \min(w_{3r-3}^\lambda, w_{3r-2}^\lambda) \\ w_{3r+1}^\lambda &= \max(w_{3r-3}^\lambda, w_{3r-2}^\lambda) - \frac{1}{2^{r-1}} \end{aligned}$$

For $z = (x, w, s, y)$ in the central path, we have $\begin{pmatrix} xs \\ wy \end{pmatrix} = \mu e$, and since the valuation of μ is λ , we deduce y^λ by tropicalizing this equality which yields:

$$y_1^\lambda = \lambda - 2, \quad y_2^\lambda = \lambda - 1, \quad y_3^\lambda = \lambda - 1 - \min(\lambda, 2) \quad (8a)$$

$$\forall 1 < j < r \quad y_{3j}^\lambda = \max(y_{3j-3}^\lambda, y_{3j-2}^\lambda) - 1 \quad (8b)$$

$$\forall 1 \leq j < r \quad y_{3j+1}^\lambda = y_{3j-1}^\lambda - 1 \quad (8c)$$

$$y_{3j+2}^\lambda = \min(y_{3j}^\lambda + 1, y_{3j-1}^\lambda) - (1 - \frac{1}{2^j}) \quad (8d)$$

$$y_{3r}^\lambda = \max(y_{3r-3}^\lambda, y_{3r-2}^\lambda) \quad (8e)$$

$$y_{3r+1}^\lambda = \min(y_{3r-3}^\lambda, y_{3r-2}^\lambda) + \frac{1}{2^{r-1}} \quad (8f)$$

When plotting in a plane the points $(x_{2r-1}^\lambda, x_{2r}^\lambda)$ for different values of the tropical parameter λ we expect to get a staircase form with an exponential number of steps.

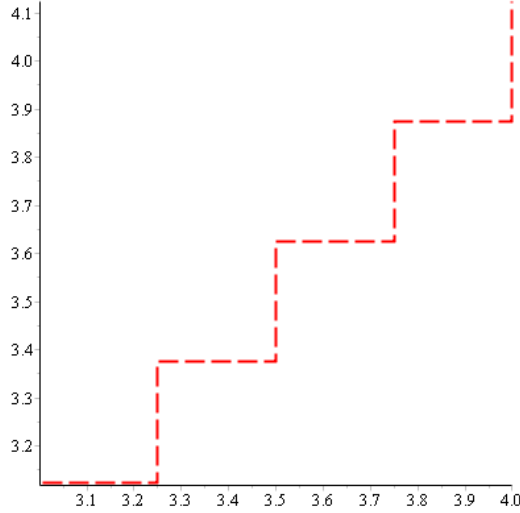


Figure 1: The form of the two last tropical x -coordinates

Our objective is to reproduce experimentally the theoretical results of [1] by running a Predictor-Corrector algorithm on instances of $\mathbf{LW}_r(t)$ and plotting an approximation of the tropical path given by the *valuation* ($\log_t(\cdot)$ when t is large) of the prediction/correction points following the primal-dual central path.

3 Constructing an initial point

To run the predictor corrector algorithm on our counter-example, we need to find an initial point z_0 to start the algorithm that verifies the following conditions:

1. $z_0 \in V(\theta')$. In particular, we need to have $z_0 \in \mathcal{F}^s$.
2. $\bar{\mu}(z_0)$ large enough. We need to be sufficiently far in the central path to be able to see all the steps.

Condition 1 is the most challenging to satisfy. A natural way to proceed is to start from a point from the tropical path – which we know exactly its form for all λ – and try to lift it into a real point z such that $\bar{\mu}(z) \approx t^\lambda$ and z close to the central path. In order to do that, we need to go through Puiseux series:

$$\begin{aligned}
& \text{Tropical} && : z^\lambda = \mathcal{C}^{trop}(\lambda) \\
\longrightarrow & \text{Puisseux series} && : \mathbf{z} \mid \mathbf{z} \in \mathcal{F}^s, \text{ val}(\mathbf{z}) = z^\lambda \\
\longrightarrow & \text{Real point} && : \mathbf{z} = \mathbf{z}(t) \mid \mathbf{z} \in V(\theta), \bar{\mu}(\mathbf{z}) \approx t^\lambda
\end{aligned}$$

for t sufficiently large.

In all what follows, we consider $\lambda \in]0, 2]$ and $z^\lambda = \mathcal{C}^{trop}(\lambda)$. We will start therefore by analyzing the set $\{\mathbf{z} \mid \mathbf{z} \in \mathcal{F}^s, \text{ val}(\mathbf{z}) = z^\lambda\}$.

In order to satisfy $\mathbf{z} \in \mathcal{F}^s$, we will exploit the form of the conditions. We recall that $\mathbf{z} = (\mathbf{x}, \mathbf{w}, \mathbf{s}, \mathbf{y})$ and

$$\mathbf{z} \in \mathcal{F}^s \iff \begin{cases} \mathbf{A}\mathbf{x} + \mathbf{w} = \mathbf{b} \\ \mathbf{s} - \mathbf{A}^T\mathbf{y} = \mathbf{c} \\ \mathbf{x}, \mathbf{w}, \mathbf{s}, \mathbf{y} > 0 \end{cases}$$

Given the form of the conditions, we construct \mathbf{x} and \mathbf{y} from the tropical point and define $\mathbf{s} := \mathbf{A}^T\mathbf{y} + \mathbf{c}$ and $\mathbf{w} := \mathbf{b} - \mathbf{A}\mathbf{x}$ to guarantee the two first feasibility conditions.

In particular we will seek to construct such points by *lifting* the tropical points.

Definition 1. Let z^λ be a tropical point in the tropical path. We say that a Puisseux serie point \mathbf{z} is lifted from z^λ if there exist α and β such that

$$\begin{aligned}
\forall i, \quad \mathbf{x}_i &= \alpha_i t^{x_i^\lambda} \\
\forall i, \quad \mathbf{y}_i &= \beta_i t^{y_i^\lambda}
\end{aligned}$$

and $\mathbf{s} := \mathbf{A}^T\mathbf{y} + \mathbf{c}$, $\mathbf{w} := \mathbf{b} - \mathbf{A}\mathbf{x}$.

Note that a lifted point \mathbf{z} verify in particular $\text{val}(\mathbf{z}) = z^\lambda$, but not necessarily that $\mathbf{z} \in \mathcal{F}^s$. The following result ensures that focusing only on lifted points does not reduce our chances of getting into the feasible set.

For a given vector of Puisseux series \mathbf{x} , we denote by $\bar{\mathbf{x}}$ the Puisseux series vector obtained by taking the leading monomial on each coordinate.

Proposition 2. Let $\mathbf{z} = (\mathbf{x}, \mathbf{w}, \mathbf{y}, \mathbf{s})$ be a feasible Puisseux series.

$\tilde{\mathbf{z}} := (\bar{\mathbf{x}}, \tilde{\mathbf{w}} := \mathbf{b} - \mathbf{A}\bar{\mathbf{x}}, \bar{\mathbf{y}}, \tilde{\mathbf{s}} := \mathbf{c} + \mathbf{A}\bar{\mathbf{y}})$ is also a feasible point in \mathbb{K} (i.e $\tilde{\mathbf{z}} \in \mathcal{F}^s$)

Proof. We have by definition of \mathbf{w} and $\bar{\mathbf{w}}$:

$$\begin{aligned}
\mathbf{w}_1 &= t^2 - \mathbf{x}_1 \\
\mathbf{w}_2 &= t - \mathbf{x}_2 \\
\mathbf{w}_{3j} &= t \mathbf{x}_{2j-1} - \mathbf{x}_{2j+1} \\
\mathbf{w}_{3j+1} &= t \mathbf{x}_{2j} - \mathbf{x}_{2j+1} \\
\mathbf{w}_{3j+2} &= t^{1-1/2^j} (\mathbf{x}_{2j-1} + \mathbf{x}_{2j}) - \mathbf{x}_{2j+2}
\end{aligned} \tag{9}$$

and :

$$\begin{aligned}
\tilde{\mathbf{w}}_1 &= t^2 - \bar{\mathbf{x}}_1 \\
\tilde{\mathbf{w}}_2 &= t - \bar{\mathbf{x}}_2 \\
\tilde{\mathbf{w}}_{3j} &= t \bar{\mathbf{x}}_{2j-1} - \bar{\mathbf{x}}_{2j+1} \\
\tilde{\mathbf{w}}_{3j+1} &= t \bar{\mathbf{x}}_{2j} - \bar{\mathbf{x}}_{2j+1} \\
\tilde{\mathbf{w}}_{3j+2} &= t^{1-1/2^j} (\bar{\mathbf{x}}_{2j-1} + \bar{\mathbf{x}}_{2j}) - \bar{\mathbf{x}}_{2j+2}
\end{aligned} \tag{10}$$

since $\bar{\mathbf{x}}$ is the Puisseux series of the leading terms of \mathbf{x} , the positivity of \mathbf{w} yields that of $\tilde{\mathbf{w}}$. The same argument proves, using $\tilde{\mathbf{s}} := \mathbf{c} + \mathbf{A}\bar{\mathbf{y}}$, that $\tilde{\mathbf{s}}$ is a positive Puisseux series.

□

In order for a lifted Puiseux series \mathbf{z} to be feasible, we only need to ensure that $\mathbf{x} > 0$, $\mathbf{w} > 0$, $\mathbf{y} > 0$, $\mathbf{s} > 0$. Notice that the two pair of conditions $\mathbf{x} > 0$, $\mathbf{w} > 0$ and $\mathbf{y} > 0$, $\mathbf{s} > 0$ are independent, therefore, we can state a feasibility condition on each pair independently.

It was showed in [1] that fixing a sequence of positive numbers $\alpha_0 = \frac{1}{2} > \alpha_1 > \dots > \alpha_{r-1} > 0$ and defining \mathbf{x} as

$$\mathbf{x}_{2j+1} := \alpha_j t^{x_{2j+1}^\lambda}, \quad \mathbf{x}_{2j+2} := \alpha_j t^{x_{2j+2}^\lambda}, \quad (0 \leq j < r) \quad (11)$$

ensures the feasibility of the lifted pair \mathbf{x} , \mathbf{w} .

We will now state a necessary and sufficient condition for the feasibility of the pair $\mathbf{s} := \mathbf{c} + \mathbf{A}\mathbf{y}$, \mathbf{y} .

Proposition 3. *Let \mathbf{z} be a lifted Puiseux series lifted from the tropical point. Let β be such that $\mathbf{y}_i := \beta_i t^{y_i^\lambda}$. The lifted pair $(\mathbf{s} := \mathbf{c} + \mathbf{A}\mathbf{y}, \mathbf{y})$ is feasible for every $\lambda \leq 2$ if and only if:*

$$\beta > 0$$

$$\begin{cases} 1 > \beta_3 + \beta_5 \\ \beta_2 > \beta_4 + \beta_5 \end{cases}$$

$$\forall j \begin{cases} \min(\beta_{3j}, \beta_{3j+1}) > \beta_{3j+3} + \beta_{3j+5} \\ \beta_{3j+2} > \beta_{3j+4} + \beta_{3j+5} \end{cases}$$

$$\begin{cases} \min(\beta_{3r-3}, \beta_{3r-2}) > \beta_{3r} \\ \beta_{3r-1} > \beta_{3r+1} \end{cases}$$

Proof. Suppose β verifies the condition. We only have to prove that $\mathbf{s}, \mathbf{y} > 0$ to ensure the feasibility of the lifted pair. $\mathbf{y} > 0$ results directly from the positivity of β . Let us now prove the positivity of \mathbf{s} . By definition of the lifted pair, we have

$$\begin{aligned} \mathbf{s}_1 &= \mathbf{y}_1 - t\mathbf{y}_3 - t^{\frac{1}{2}}\mathbf{y}_5 + 1 \\ \mathbf{s}_2 &= \mathbf{y}_2 - t\mathbf{y}_4 - t^{\frac{1}{2}}\mathbf{y}_5 \end{aligned} \quad (12)$$

$$\forall 1 \leq j < r-1 \quad \mathbf{s}_{2j+1} = \mathbf{y}_{3j} + \mathbf{y}_{3j+1} - t\mathbf{y}_{3j+3} - t^{1-\frac{1}{2j}}\mathbf{y}_{3j+5} \quad (13)$$

$$\mathbf{s}_{2j+2} = \mathbf{y}_{3j+2} - t\mathbf{y}_{3j+4} - t^{1-\frac{1}{2j}}\mathbf{y}_{3j+5} \quad (14)$$

$$\mathbf{s}_{2r-1} = \mathbf{y}_{3r-3} + \mathbf{y}_{3r-2} - \mathbf{y}_{3r} \quad (15)$$

$$\mathbf{s}_{2r} = \mathbf{y}_{3r-1} - \mathbf{y}_{3r+1} \quad (16)$$

We will start by proving the following inequalities on y^λ that derive directly from equations 8:

$$0 \geq \max(y_3^\lambda + 1, y_5^\lambda + \frac{1}{2}) \quad (17a)$$

$$y_2^\lambda \geq \max(1 + y_4^\lambda, \frac{1}{2} + y_5^\lambda) \quad (17b)$$

$$\forall 1 \leq j < r-1 \quad \max(y_{3j}^\lambda, y_{3j+1}^\lambda) \geq \max(1 + y_{3j+3}^\lambda, 1 - \frac{1}{2^j} + y_{3j+5}^\lambda) \quad (17c)$$

$$y_{3j+2}^\lambda \geq \max(1 + y_{3j+4}^\lambda, 1 - \frac{1}{2^j} + y_{3j+5}^\lambda) \quad (17d)$$

$$\max(y_{3r-3}^\lambda, y_{3r-2}^\lambda) \geq y_{3r} \quad (17e)$$

$$y_{3r-1}^\lambda \geq y_{3r+1}^\lambda \quad (17f)$$

Because $\lambda \leq 2$, from 8a and 8d we have $y_3^\lambda + 1 = 0$ and $y_5^\lambda + \frac{1}{2} = \min(0, \lambda - 1) \leq 0$, hence $0 \geq \max(y_3^\lambda + 1, y_5^\lambda + \frac{1}{2})$ wich proves 17a.

8c and 8d yields to $y_4^\lambda + 1 = y_2^\lambda$ and $y_5^\lambda + \frac{1}{2} = \min(0, \lambda - 1) \leq \lambda - 1 = y_2^\lambda$ which proves 17b.

For $1 \leq j < r - 1$, we have thanks to 8c and 8e, $1 + y_{3j+3}^\lambda = \max(y_{3j}^\lambda, y_{3j}^\lambda)$ and $y_{3j+5}^\lambda + 1 - \frac{1}{2^j} = \min(y_{3j+3}^\lambda + 1, y_{3j+2}^\lambda) \leq 1 + y_{3j+3}^\lambda = \max(y_{3j}^\lambda, y_{3j+1}^\lambda)$ proving 17c.

8c and 8e yield to $y_{3j+4}^\lambda + 1 = y_{3j+2}^\lambda$ and $y_{3j+5}^\lambda + 1 - \frac{1}{2^j} = \min(y_{3j+3}^\lambda + 1, y_{3j+2}^\lambda) \leq y_{3j+2}^\lambda$ which leads to 17d.

The two inequalities for 17e have already been proved.

From 8d we have $y_{3r-1}^\lambda = \min(y_{3r-3}^\lambda + 1, y_{3r-4}^\lambda) - (1 - \frac{1}{2^{r-1}})$. To prove 17f, it suffice to show that $y_{3r+1}^\lambda \leq y_{3r-3}^\lambda + \frac{1}{2^{r-1}}$ and $y_{3r+1}^\lambda \leq y_{3r-4}^\lambda - (1 - \frac{1}{2^{r-1}})$. The former is a consequence of 8f and the latter is a consequence of 8c applied to $r - 1$.

17f is a direct consequence of 8f.

The inequalities 17 ensure that there exist a possible β such that $\mathbf{s} > 0$. In fact, it is a bare reformulation of the fact that the leading monomial of every coordinate of \mathbf{s} must have a coefficient of the form $\sum_{j \in J+} \beta_j - \sum_{j \in J-} \beta_j$ where $J+, J- \subset [3r+1]$ and $J+ \neq \emptyset$ (the elements in inequations 17 being the powers of the terms in the expression of \mathbf{s}). Since $\beta > 0$, this is obviously a necessary condition for the positivity of \mathbf{s} . The rest of the conditions on β simply state that for all \mathbf{s} coordinates leading monomial's coefficient of the form $\sum_{j \in J+} \beta_j - \sum_{j \in J-} \beta_j$, we have $\sum_{j \in J+} \beta_j \geq \sum_{j \in J-} \beta_j$ for the worst case scenario of $J+, J-$ (i.e $|J+| = 1$ and $|J-|$ maximal).

Reciprocally, let (\mathbf{s}, \mathbf{y}) be a lifted feasible point. $\mathbf{y} > 0$ implies that $\beta > 0$. We show by choosing particular cases of $\lambda < 2$, that each of the inequalities in 17 is saturated in a given case. Therefore, we reach the worse case scenarios on $J+, J-$ for all the coordinates of \mathbf{s} proving therefore the necessity of the conditions over β .

□

Remark 4. We denote ${}^s A := (\text{sign}(A_{i,j}))_{i,j}$. For a matrix $M \in \mathbb{R}^{m \times n}$ and a vector $v \in \mathbb{R}^n$, we define the product \bullet as $(M \bullet v)_j = \min_k \frac{(M_{jk}+1)}{2} \beta_k + \sum_k \frac{(M_{jk}-1)}{2} \beta_k$ for all j . The previous conditions on β can be rewritten as:

$$\beta > 0, {}^s A \bullet \beta + c > 0$$

Remark 5. The previous formula suggest a generalization of the results to every standard linear problem. While the condition being sufficient is easy to generalize, the reciprocal is not obviously true. We suggest therefore the following proposition.

Proposition 6. *Consider a standard linear problem*

$$(P) \begin{cases} \inf_x c \cdot x \\ Ax = b \\ x \geq 0 \end{cases}$$

where $c \in \mathbb{R}^n$, $b \in \mathbb{R}^m$, and $A \in \mathbb{R}^{n \times m}$ and let $\lambda > 0$. Suppose the corresponding Puiseux series strictly feasible set $\mathcal{F}^s \cap \{z \in \mathbb{K} | \text{val}(z) = \lambda\}$ is non empty. Let $z = (\mathbf{x}, \mathbf{w}, \mathbf{y}, \mathbf{s})$ be a lifted Puiseux series lifted from a tropical point in the central path. Let β be such that $\mathbf{y}_i := \beta_i t^{y_i^\lambda}$ for all i . If $\beta > 0$, ${}^s A \bullet \beta + c > 0$ then (\mathbf{y}, \mathbf{s}) is strictly feasible.

Choosing α and β under the constraints fixed above yields a lifted feasible point $z = (\mathbf{x}, \mathbf{w}, \mathbf{y}, \mathbf{s})$ in \mathcal{F} . For a large enough t , we know therefore that $z(t)$ is a feasible point in $\mathcal{P}(t)$. The problem of feasibility being solved, we still need to get a point in the neighbourhood V . The conditions over β being necessary, we know that we didn't lose many possibilities for the choice of the initial point (The eventual losses being in the choice of α).

α and β must be chosen to ensure that the lifted point is in V . It is very difficult to find a choice of α and β to satisfy this constraint because the equations defining V are not simple.

One way to overcome this problem is to start from a feasible point in $\mathcal{P}(t)$ and get closer to V following Newton directions (correction steps). Once we find a feasible point in V we can launch the Predictor-Corrector on an instance of $\mathbf{LW}_r(t)$.

4 Experimentation

The following plots are outputs of running the predictor corrector algorithm presented in section 1 starting with the initial point constructed in section 3. We run the algorithm till $\bar{\mu}(z) < 1$ where z is the current point.

When t become large enough the image of a the central path of the problem $\mathbf{LW}_r(t)$ under the function \log_t takes the form of the tropical central path which is a staircase (and the convergence is proven to be uniform in [1]). The image below illustrates this convergence and we see that the image of the central path under \log_t starts to stick to the staircase form of the tropical central path (in green) showing (as expected in [1]) that, for a sufficiently large t , the total curvature of the central path is bounded from below exponentially in r .

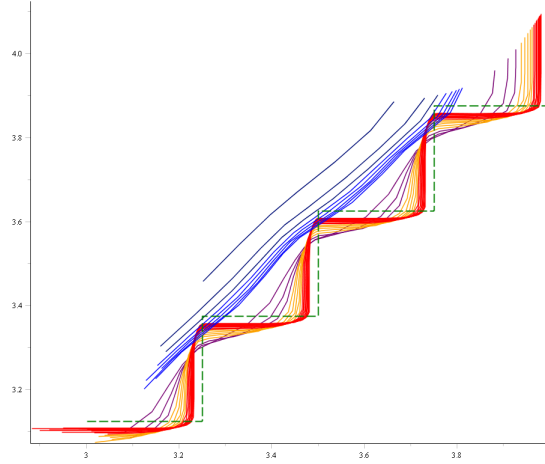


Figure 2: Deformation of the central path when t becomes large ($\log_t - \log_t$ scale of real central path) for $r=4$

The following plots show that the number of steps in the tropical central path is exponential in r and that the threshold $t_0(r)$ above which the staircase form appears is also exponential in r . On each r , we select t such that we can clearly see the staircase form of the path.

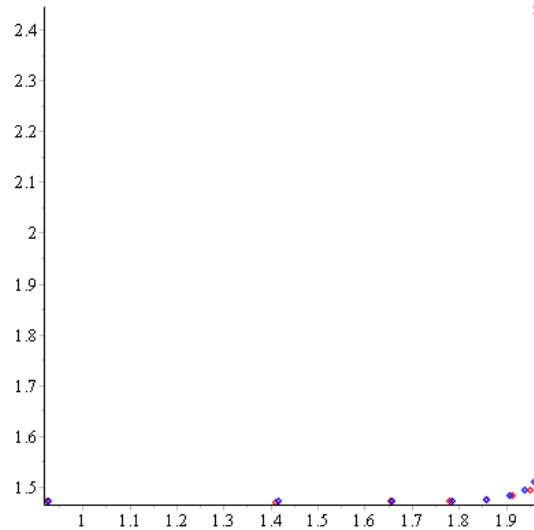


Figure 3: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 2$, $t = 10^{10}$. $\log_t - \log_t$ scale. Red: prediction steps, blue: correction steps.

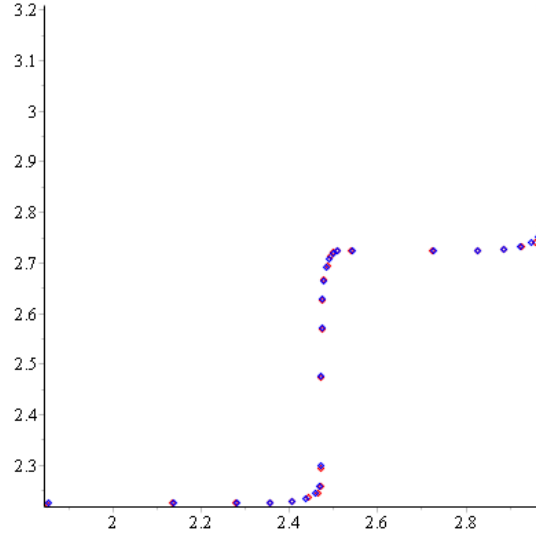


Figure 4: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 3$, $t = 10^{12}$. $\log_t - \log_t$ scale. *Red*: prediction steps, *blue*: correction steps.

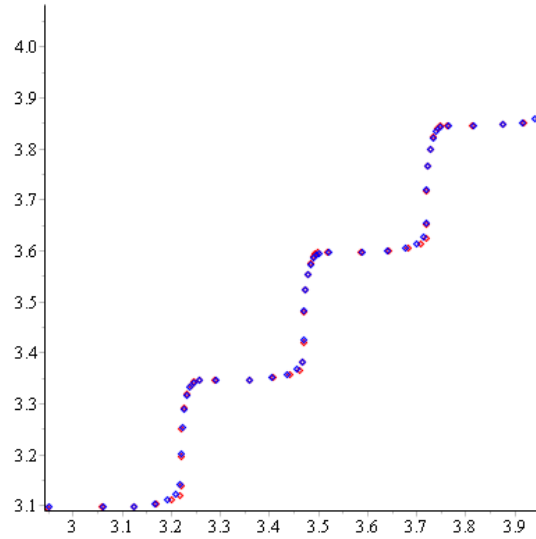


Figure 5: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 4$, $t = 10^{12}$. $\log_t - \log_t$ scale. *Red*: prediction steps, *blue*: correction steps.

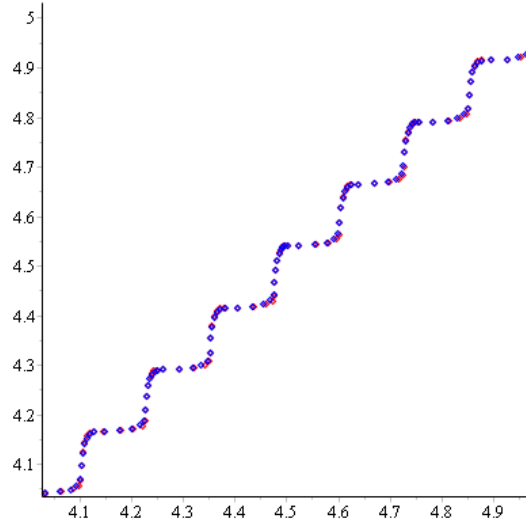


Figure 6: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 5, t = 10^{17}$. $\log_t - \log_t$ scale. *Red*: prediction steps, *blue*: correction steps.

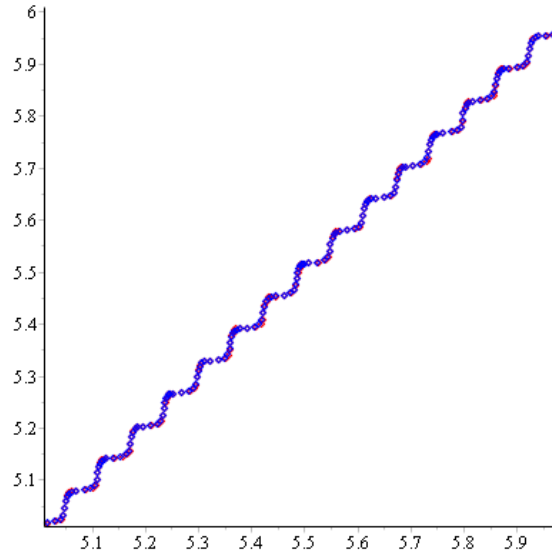


Figure 7: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 6, t = 10^{25}$. $\log_t - \log_t$ scale. *Red*: prediction steps, *blue*: correction steps.

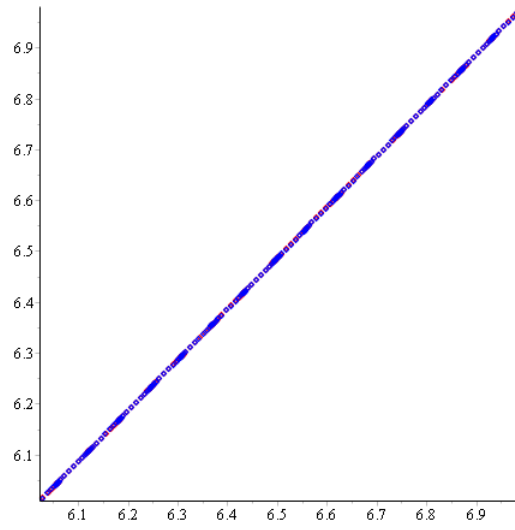


Figure 8: Projection on (x_{2r-1}, x_{2r}) of the Path followed by the predictor corrector on $\mathbf{LW}_r(t)$ for $r = 7, t = 10^{33}$. $\log_t - \log_t$ scale. *Red*: prediction steps, *blue*: correction steps.

Notice that in the last figure we cannot see the staircase form even with $t = 10^{33}$. This is due the fact that the threshold $t_0(r)$ increases exponentially with respect to r .

References

- [1] Log-Barrier Interior Point Methods Are Not Strongly Polynomial X.Allamigeon, P.Benchimol, S.Gaubert, and M.JOSWIG
- [2] <https://who.rocq.inria.fr/Jean-Charles.Gilbert/ensta/optim.html>