

CS & IT ENGINEERING

Discrete Mathematics

Graph Theory

Lecture_ 06



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Topics to be Covered



Topic

Types Of Graphs — II.

Topic

Hypercube.





Topic: Graph Theory

Isomorphic Graph: (same graph but diff representation)

G_1, G_2 are isomorphic to each other, if G_1, G_2 will have 1:1 correspondence

Same no. of vertices.
Same no. of edges.
Same degree sequence.



Topic: Graph Theory

Isomorphic Graph: (same graph but diff representation)



$$G_1 = (V_1, E_1)$$

$$G_2 = (V_2, E_2)$$

$$f: G_1 \rightarrow G_2$$

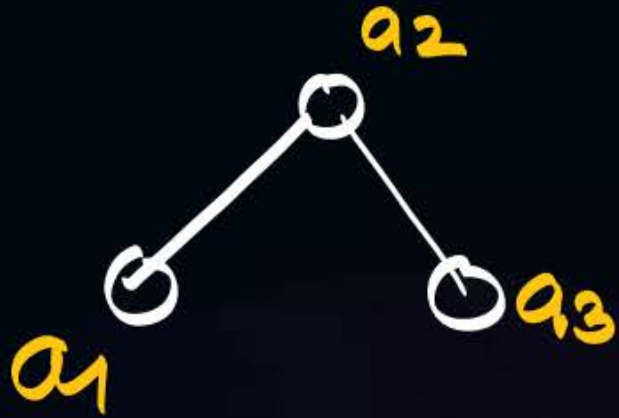
(1:1 correspondence)

$$f: V_1 \rightarrow V_2 \quad f: E_1 \rightarrow E_2$$



Topic: Graph Theory

Isomorphic Graph: (same graph but diff representation)



$$f: G_1 \rightarrow G_2$$

$$f: V_1 \rightarrow V_2$$

$$f: E_1 \rightarrow E_2$$

$$f: V_1 \rightarrow V_2$$

$$a_1 \rightarrow b_1$$

$$a_2 \rightarrow b_2$$

$$a_3 \rightarrow b_3$$

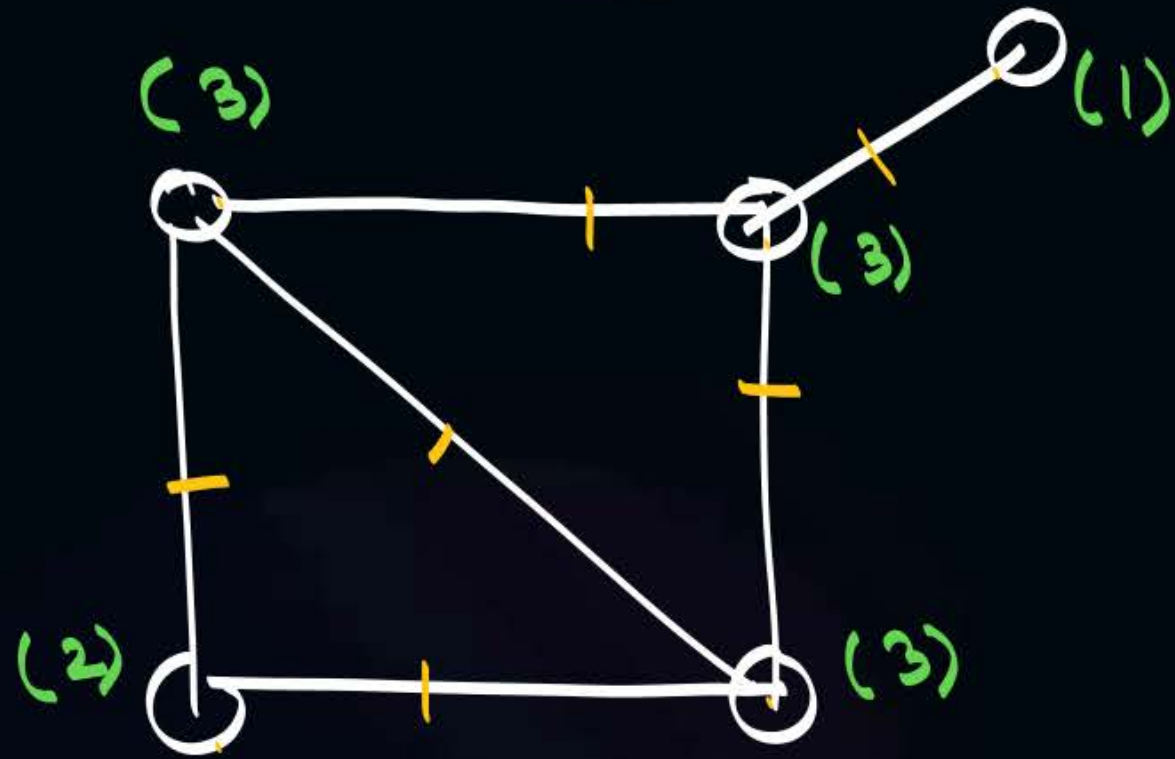
$$f: E_1 \rightarrow E_2$$

$$(a_1, a_2) \rightarrow (b_1, b_2)$$

$$(a_2, a_3) \rightarrow (b_2, b_3)$$



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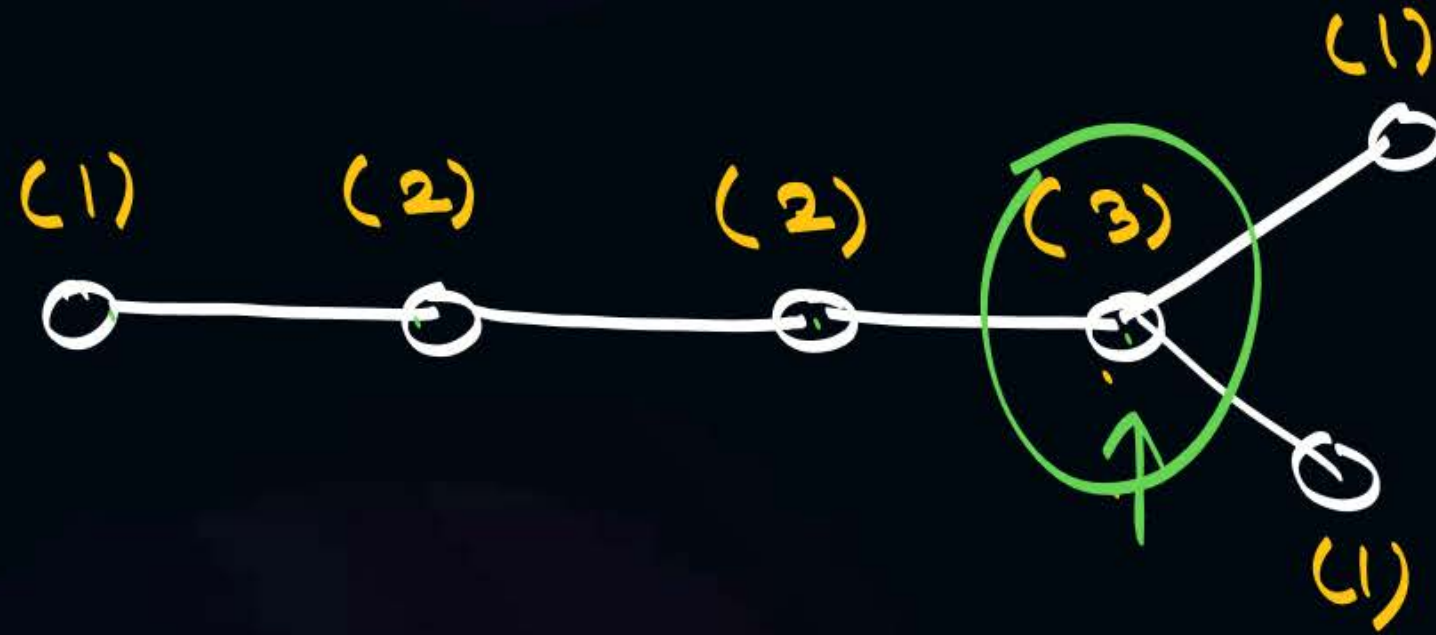
$$\begin{cases} n = 5 \\ e = 6 \\ 3 \ 3 \ 3 \ 2 \ 1 \end{cases}$$



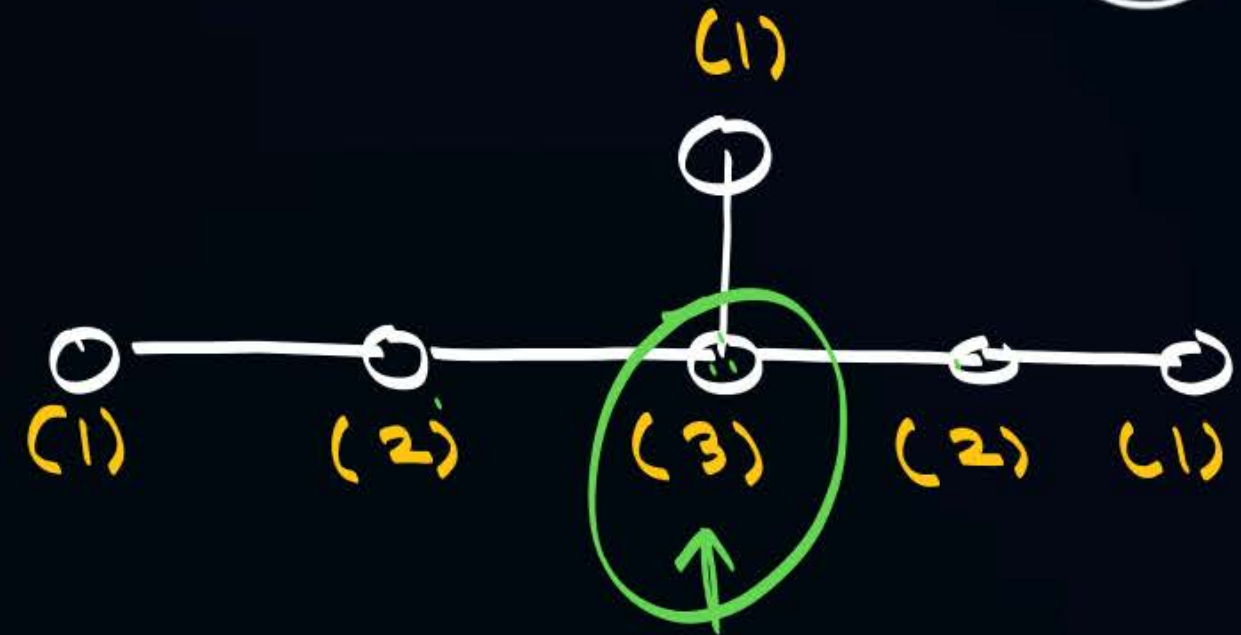
$$\begin{cases} n = 5 \\ e = 6 \\ 3 \ 3 \ 3 \ 2 \ 1 \end{cases}$$



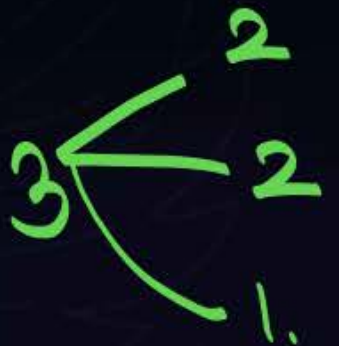
Topic: Graph Theory



$$\begin{cases} 3 & 2 & 2 & 1 & 1 & 1 \\ n = 6 \\ e = 5 \end{cases}$$



$$\begin{cases} 3 & 2 & 2 & 1 & 1 & 1 \\ n = 6 \\ e = 5 \end{cases}$$





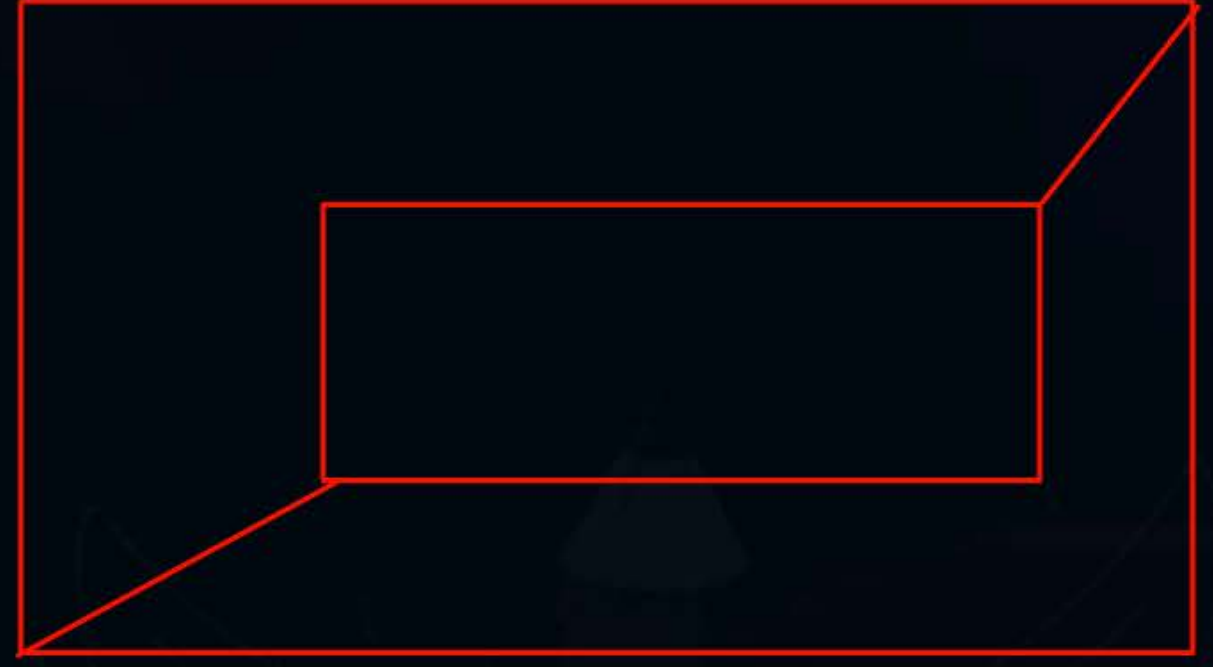
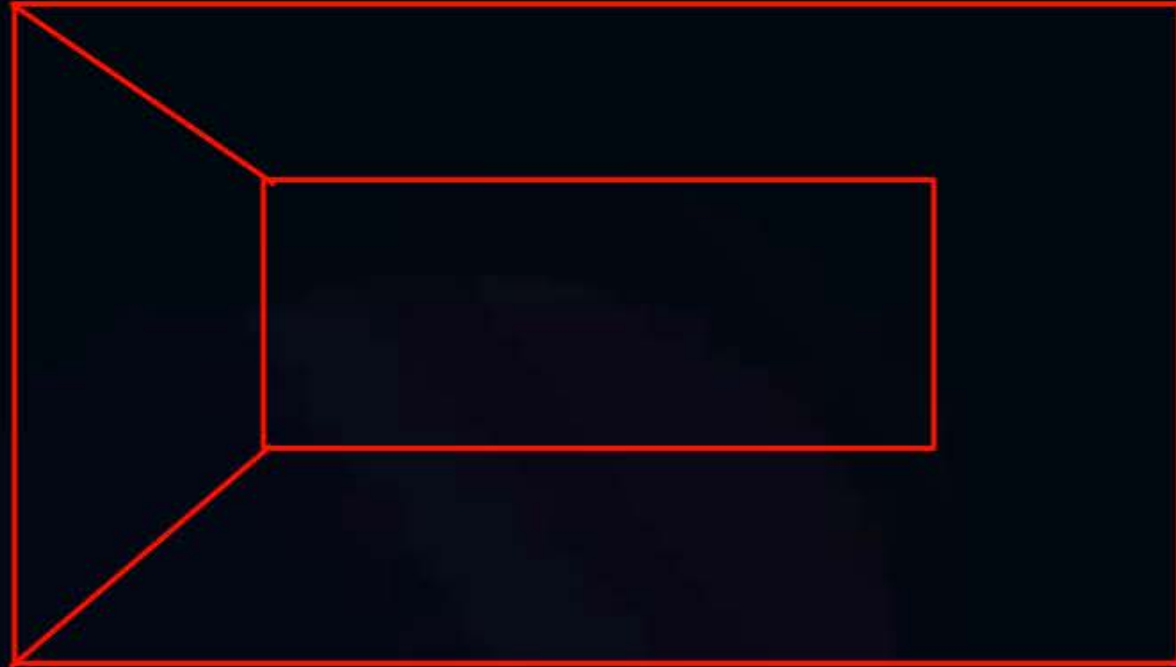
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G_1, G_2 are isomorphic then they will have same degree sequence, same no. of vertices & edges
but vice versa is not true.

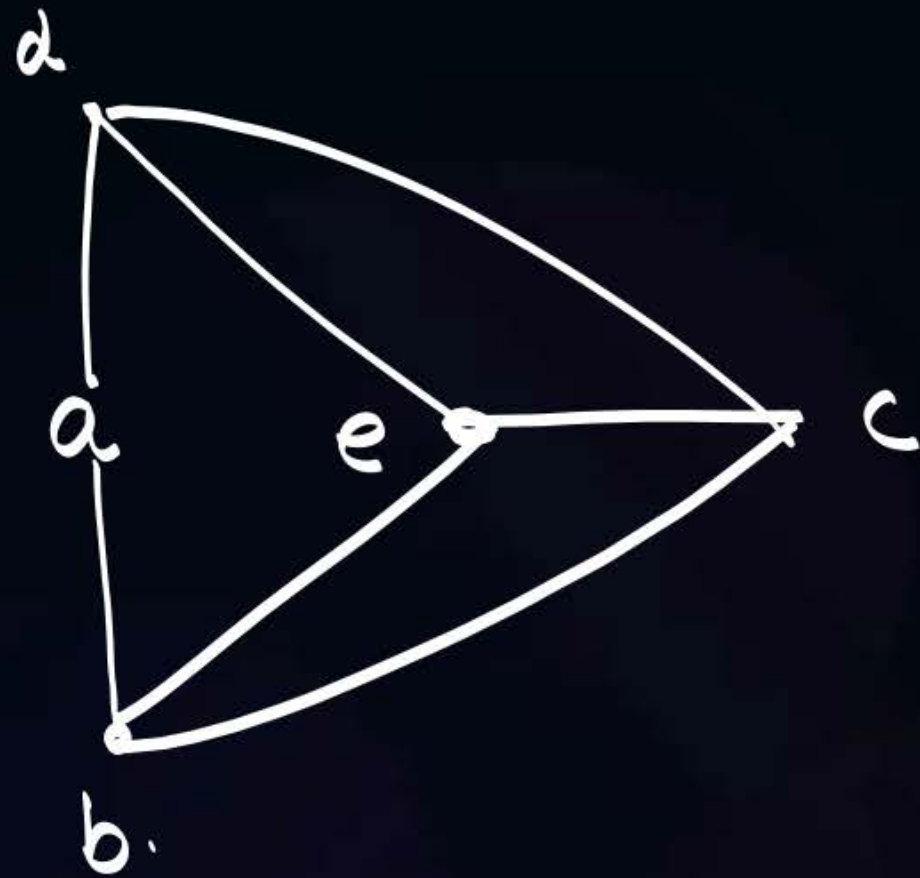


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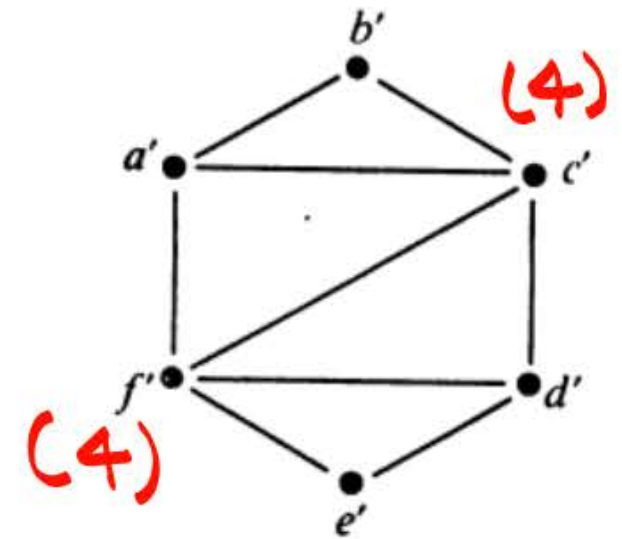
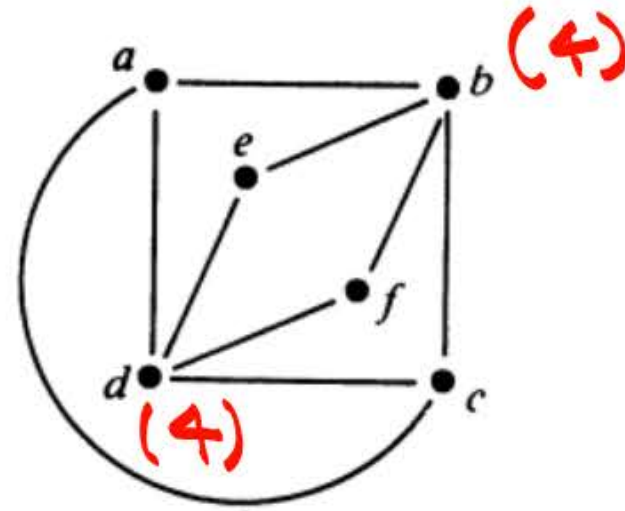




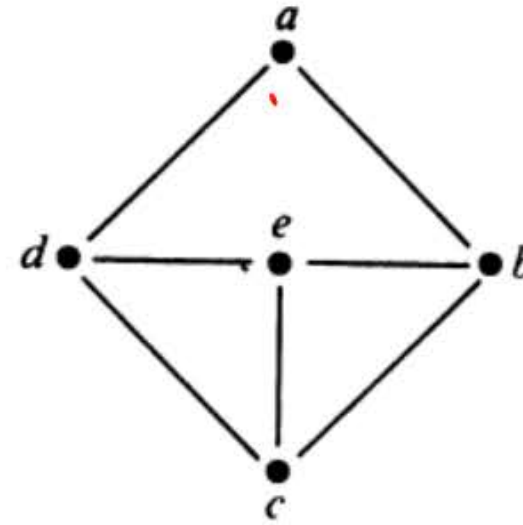
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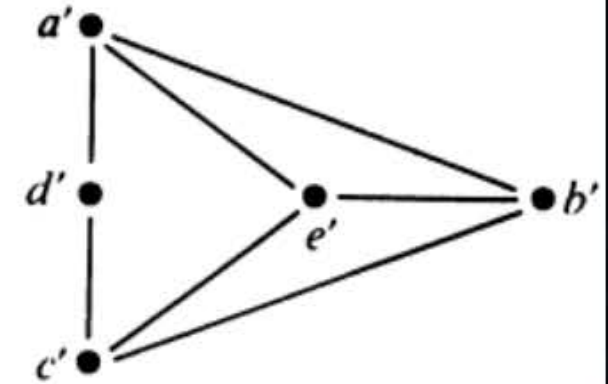
(a)



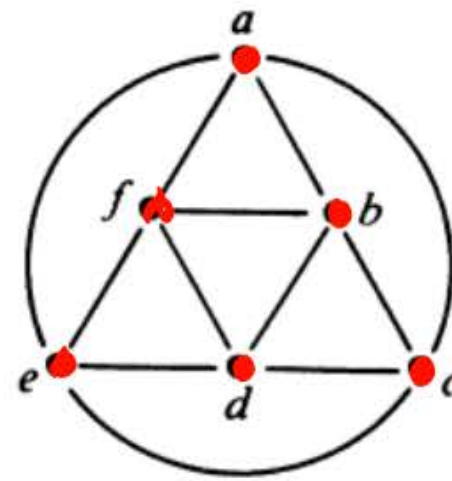
(b) ✓



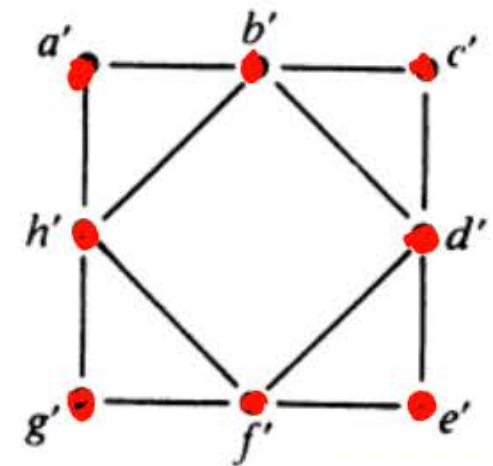
ISO



(c) ✓



not ISO

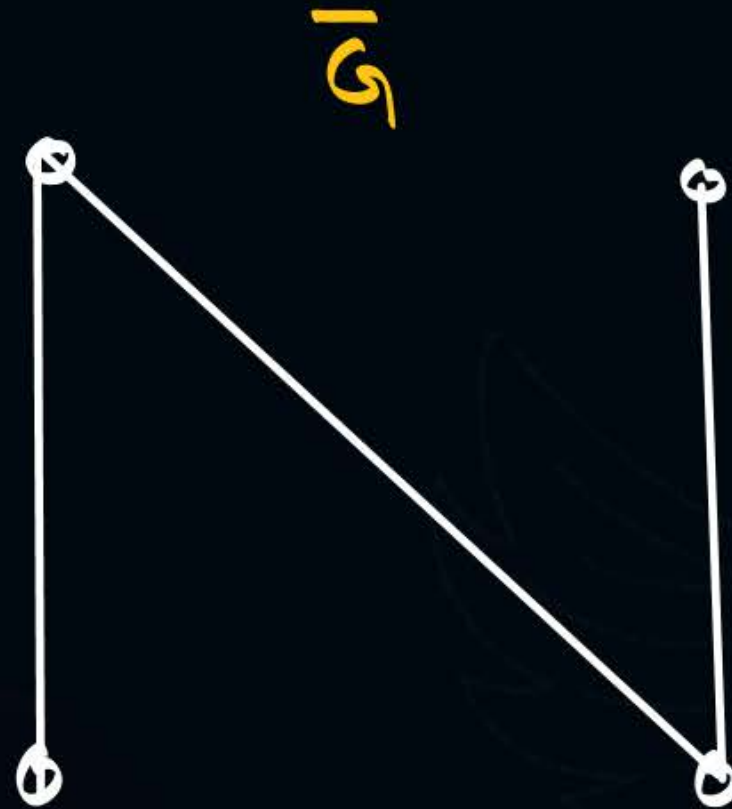
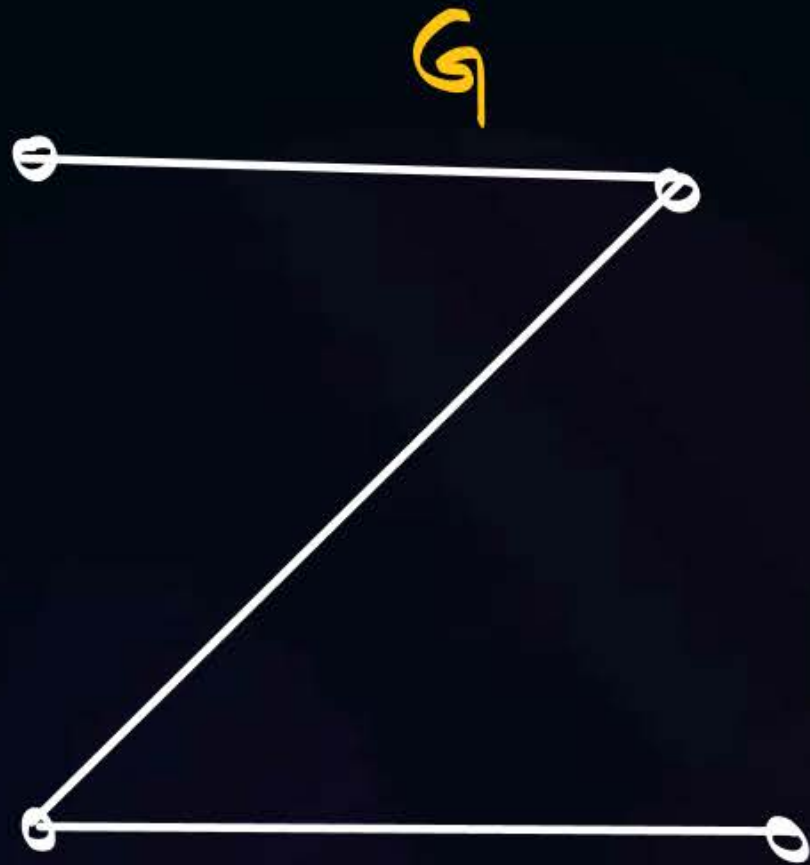




Topic: Graph Theory

Self-complement ($G \cong \bar{G}$)

when a graph is **isomorphic** to its own complement





Topic: Graph Theory

$$V(G) = V(\bar{G})$$

$$E(G) = E(\bar{G})$$

$$\begin{array}{cc} e(G) + e(\bar{G}) = \frac{n(n-1)}{2} \\ \downarrow \quad \quad \downarrow \end{array}$$

$$e + e = \frac{n(n-1)}{2}$$

$$2e = \frac{n(n-1)}{2}$$

$$e = \frac{n(n-1)}{4}$$

$$n=3 \quad e = \frac{n(n-1)}{4} = \frac{3 \cdot 2}{4} \times$$

$$n=4 \quad e = \frac{n(n-1)}{4} = \frac{4 \cdot 3}{4} = 3$$

$$n=5 \quad e = \frac{n(n-1)}{4} = \frac{5 \cdot 4}{4} = 5$$

$$n=6 \quad e = \frac{n(n-1)}{4} = \frac{6 \cdot 5}{4} \times$$



Topic: Graph Theory

$$C_n \quad \begin{matrix} (n=e) \\ (n \geq 3) \end{matrix}$$

↓

C_5

$$n = 5$$

↓

[Self-complement]

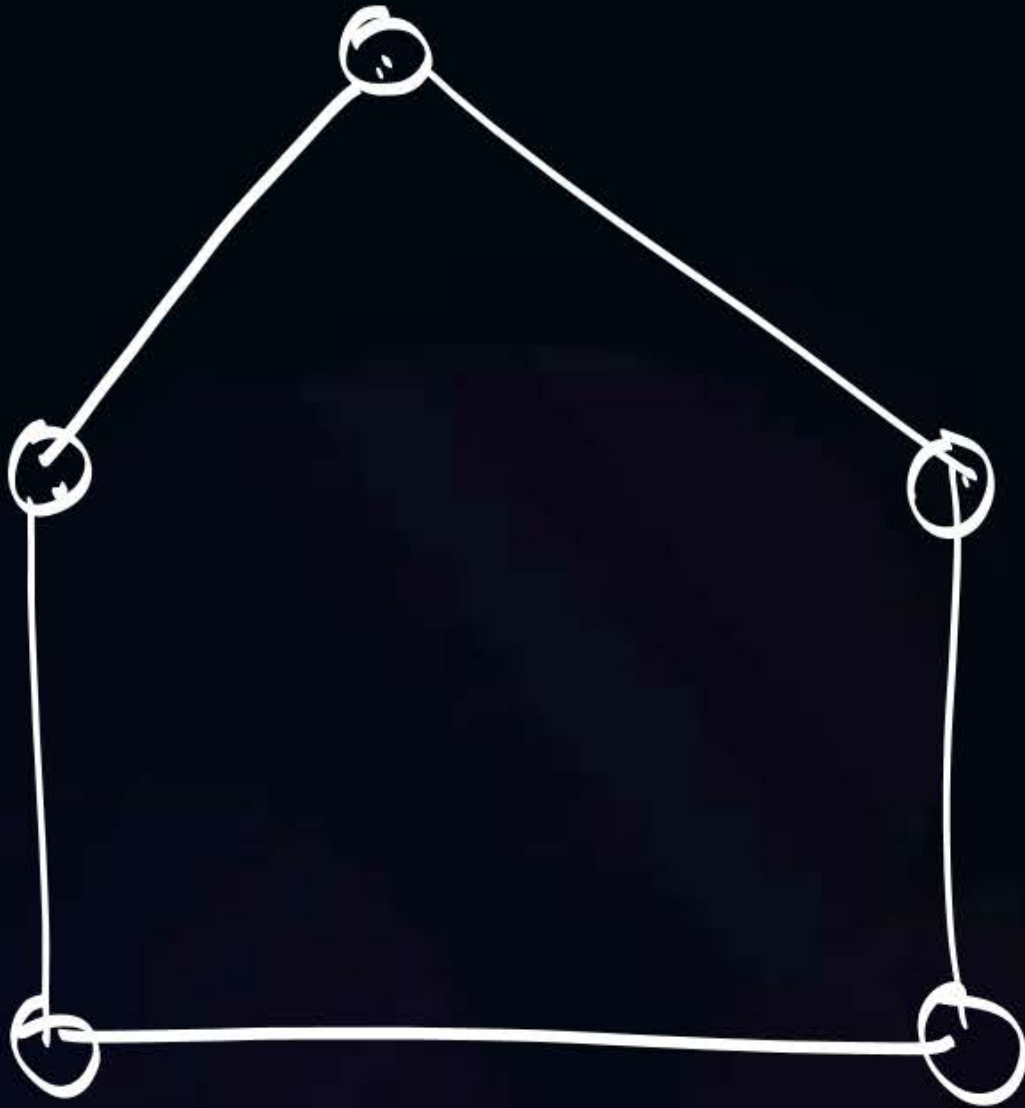
$$n=4 \quad e = \frac{n(n-1)}{2} = \frac{4 \cdot 3}{2} = 3$$

$$n=5 \quad e = \frac{5 \cdot 4}{2} = 5 \rightarrow (C_5)$$



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C_5



$$\overline{C_5} = C_5$$



C_5 is only cycle graph which is self complement.



Topic: Graph Theory



Consider a Graph vertices are represented as n -bit signal.
two vertices are adjacent, if there bit position changes by
1-bit.

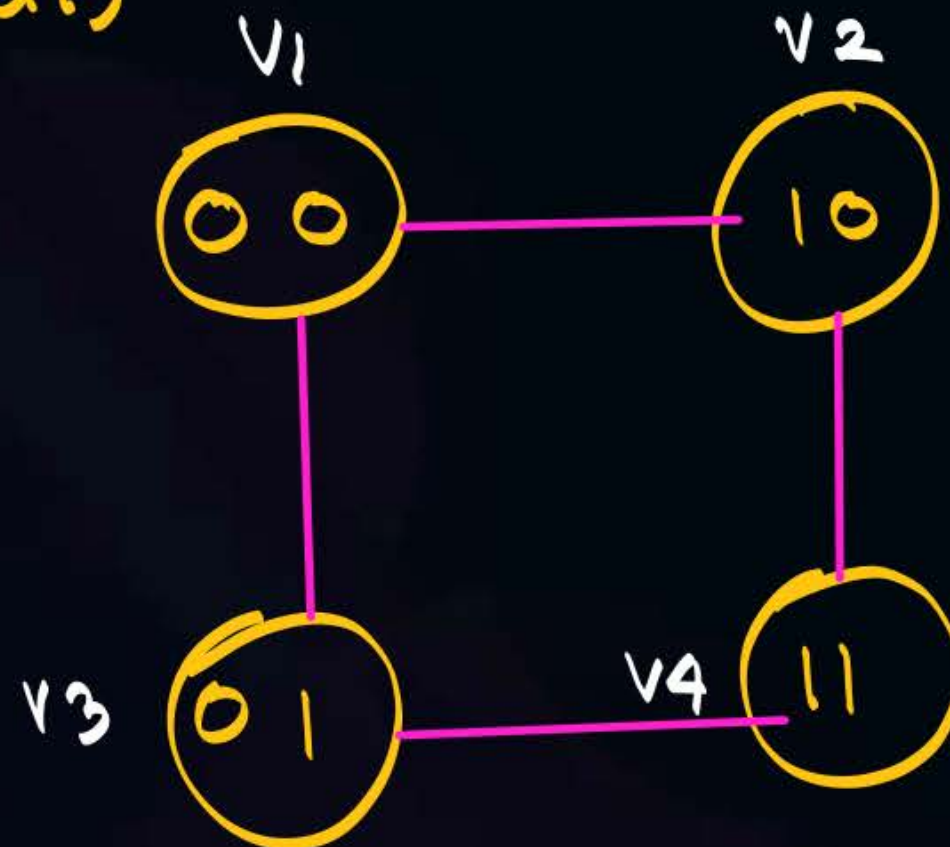
what will be total no. of edges in G ? ($n \geq 1$)



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Graph vertices are represented as n -bit signal.

$n = 2$
(2-bit) Total vertices = $2^2 = 4$. Every vertex of graph is represented as 2-bit signal.



* Degree of each vertex will be 2.



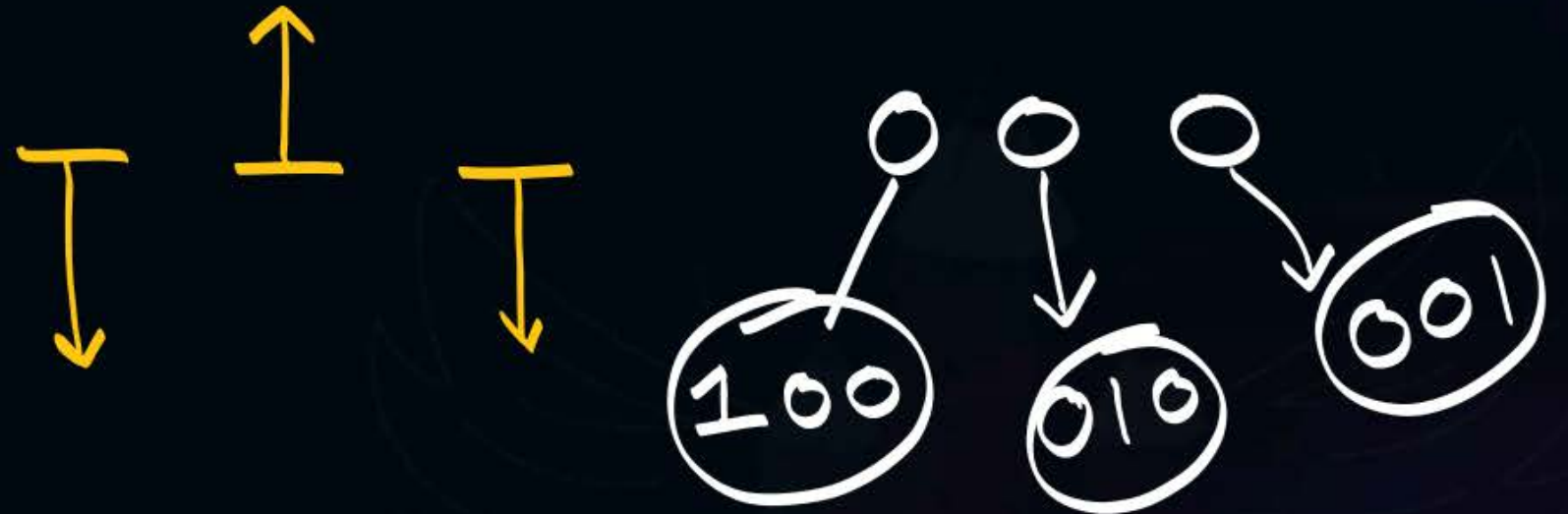
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$n = 3$ (3-bit)

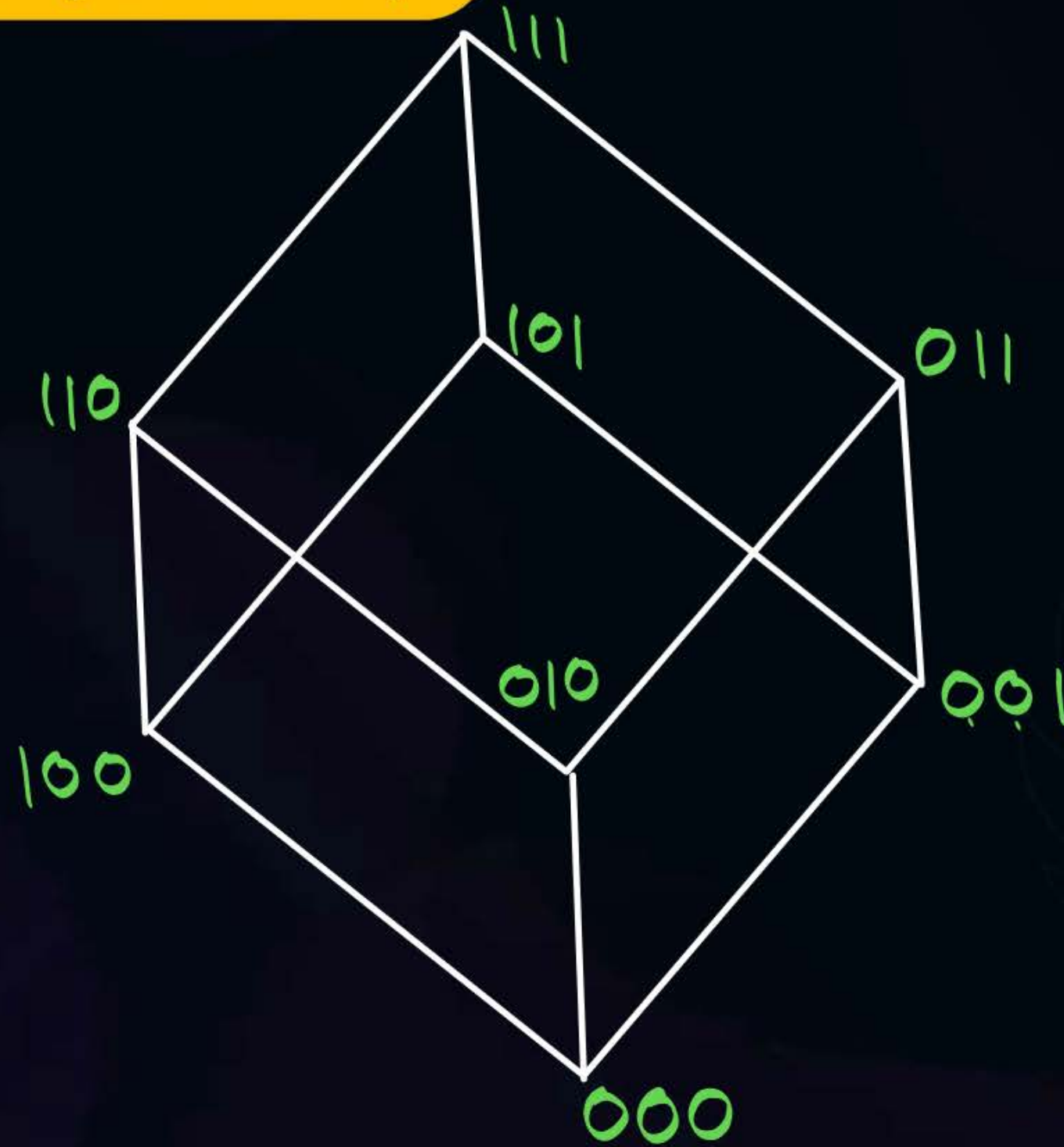
Total vertices = $2^3 = 8$

Degree of each vertex will be 3.





Topic: Graph Theory





Topic: Graph Theory

$n \rightarrow$ bit signal.

Total vertices = 2^n

Degree of each vertex will be n .

$$\sum d(v_i) = 2e$$

$$2^n \times n = 2e$$

$$e = n \cdot 2^{n-1}$$

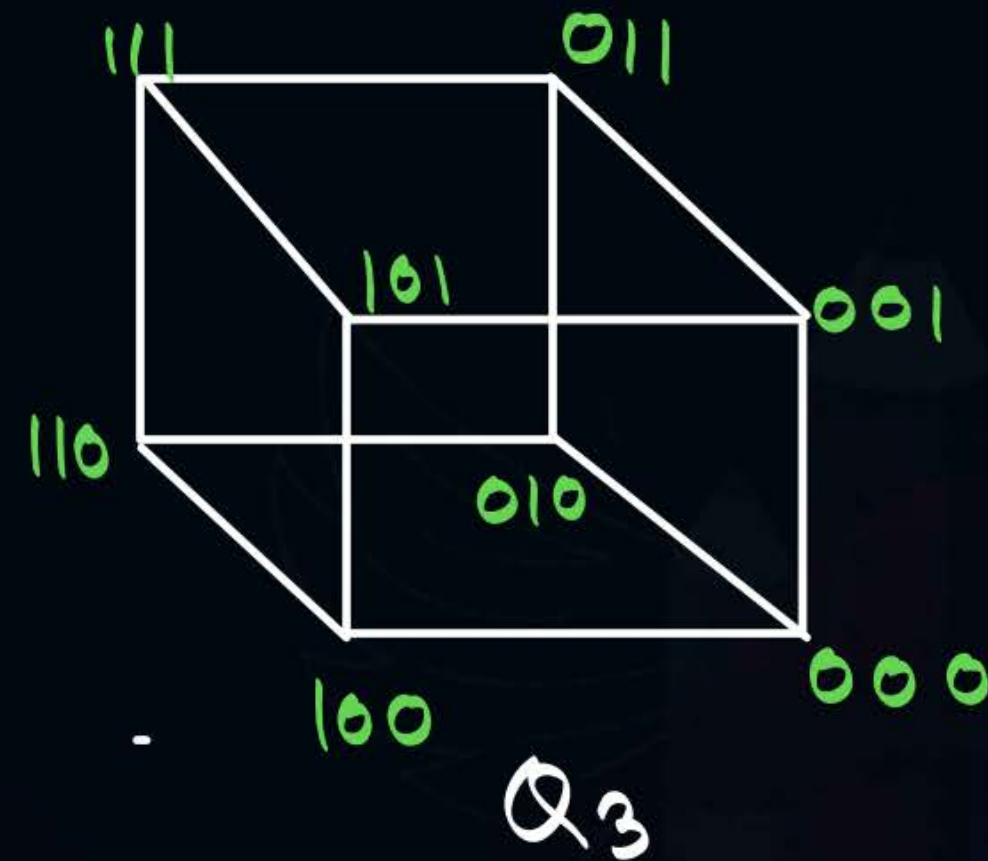
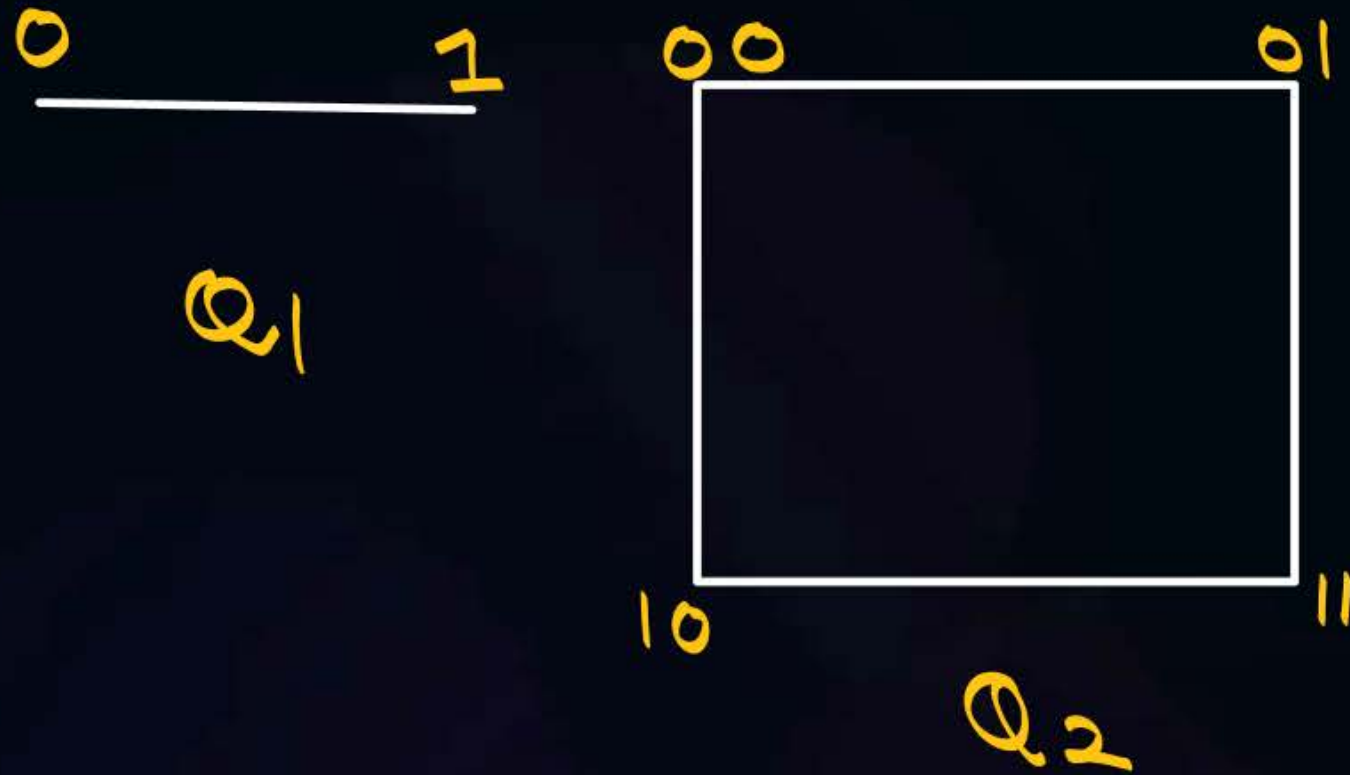


Topic: Graph Theory

Hypercube (Q_n)

↳ n-bit signal

$$V(Q_n) = 2^n$$
$$E(Q_n) = n \cdot 2^{n-1}$$





THANK - YOU