# MCAC 301: Design and Analysis of Algorithms

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August 26, 2020

#### Linear Search: Correctness

Let us rewrite the linear search algorithm of chapter 2 as follows:

```
input: Array: A[1...n], Key
output: index of first occurrence of key if it is found, 0
        otherwise
i = 1
while i \le n do
   if A[i] = key then
       return i
   end
   else i++
end
return 0
```

### Steps to Prove

- ► Loop Invariance
- ▶ Prove the Loop Invariance using Mathematical Induction
- Use the Loop Invariance to prove the correctness

#### Loop Invarinace

```
\overline{i=1}
while i \leq n do
| \quad \text{if } A[i] = \text{key then}
| \quad \text{return } i 
end
| \quad \text{else } i++
```

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

1. i = r

return 0

2. key  $\notin A[1 \dots r-1]$ 

$$\forall 1 \leq r \leq n+1$$

# Proof Of Loop Invariance

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

- 1. i = r
- 2. key  $\notin A[1 \dots r-1]$

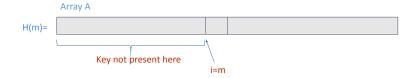
Proof by induction on r.

- 1. Base Case: when r = 1, the claim holds vacuously.
- 2. Induction Hypothesis:  $H(m) \Rightarrow H(m+1)$ .

## Induction Hypothesis

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

- 1. i = r
- 2. key  $\notin A[1 \dots r-1]$



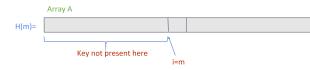
## Induction Hypothesis

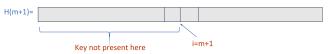
```
\overline{i=1} while i \leq n do | \begin{array}{c} \textbf{if } A[i] = \textit{key then} \\ | \begin{array}{c} \textbf{return } i \\ \textbf{end} \\ | \begin{array}{c} \textit{else } i++ \end{array} \end{aligned} end
```

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

- 1. i = r
- 2. key  $\notin A[1 \dots r-1]$

#### return 0





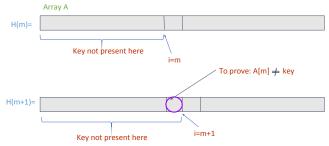
## Induction Hypothesis

```
\overline{i=1}
while i \leq n do
| \begin{array}{c} \textbf{if} \ A[i] = \textit{key then} \\ | \ \textbf{return} \ i \\ | \ \textbf{end} \\ | \ \textit{else} \ i++ \\ \\ \textbf{end} \\ \hline \end{array} |
```

#### <u>return 0</u>

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

- 1. i = r
- 2. key  $\notin A[1 \dots r-1]$



#### **Proof of Correctness**

return 0

Hypothesis H(r): When the control reaches the "While" statement for the  $r^{th}$  time,

- 1. i = r
- 2. key  $\notin A[1 \dots r-1]$

Suppose that the test condition in "While" statement is executed exactly k times. i.e. body of the loop is executed k-1 times.

- 1. Case 1:  $k \le n$ . In this case, i = k and  $A[1 \dots k-1]$  does not contain the key. Also, the while loop terminated because A[i](i.eA[k]) = key. Thus the value returned in statement 4 is the position of the first occurrence of the key in the array.
- 2. Case 2: k = n + 1. By loop invariant hypothesis, i = n + 1 and A[1 ... n] does not contain the key. Since i = n + 1, control goes to statement 8 and the algorithm returns 0.2 0.00