

**Algorithm** Prim( $E, \text{cost}, n, t$ )

*//  $E$  is the set of edges in  $G$ .*

*//  $\text{cost}[l : n, 1 : n]$  is the cost adjacency matrix of an  $n$  vertex graph such that*

*//  $\text{cost}[i, j]$  is either a positive real number or  $\infty$  if no edge  $(i, j)$  exists.*

*// A minimum spanning tree is computed and stored as a set of edges in the*

*// array  $t[1:n-1, 1:2]$ .*

*//  $(t[i, 1], t[i, 2])$  is an  $i^{\text{th}}$  edge in the minimum-cost spanning tree.*

*// The final cost is returned.*

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Let  $(k, l)$  be an edge of minimum cost in  $E$ ;

$\text{mincost} := \text{cost}[k, l]$ ;

$t[1, 1] := k; t[1, 2] := l$ ;

**for**  $i := 1$  **to**  $n$  **do**   *// Initialize near[ ] array.*

**if**  $(\text{cost}[i, l] < \text{cost}[i, k])$  **then**   *// If  $l$  is nearer to  $i$  than  $k$*

$\text{near}[i] := l$ ;

**else**

$\text{near}[i] := k$ ;

$\text{near}[k] := \text{near}[l] := 0$ ;

**for**  $i := 2$  **to**  $n - 1$  **do**

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*// Find  $n-2$  additional edges for  $t$ .*

        Let  $j$  be an index such that  $\text{near}[j] \neq 0$  and  $\text{cost}[j, \text{near}[j]]$  is  
        minimum;

$t[i, 1] := j$ ;

$t[i, 2] := \text{near}[j]$ ;

$\text{mincost} := \text{mincost} + \text{cost}[j, \text{near}[j]]$ ;

$\text{near}[j] := 0$ ;

**for**  $k := 1$  **to**  $n$  **do** *//update near[ ] array*

**if**  $(\text{near}[k] \neq 0)$  **and**  $((\text{cost}[k, \text{near}[k]] > (\text{cost}[k, j])))$  **then**

$\text{near}[k] := j$ ;

    }

**return**  $\text{mincost}$ ;

}