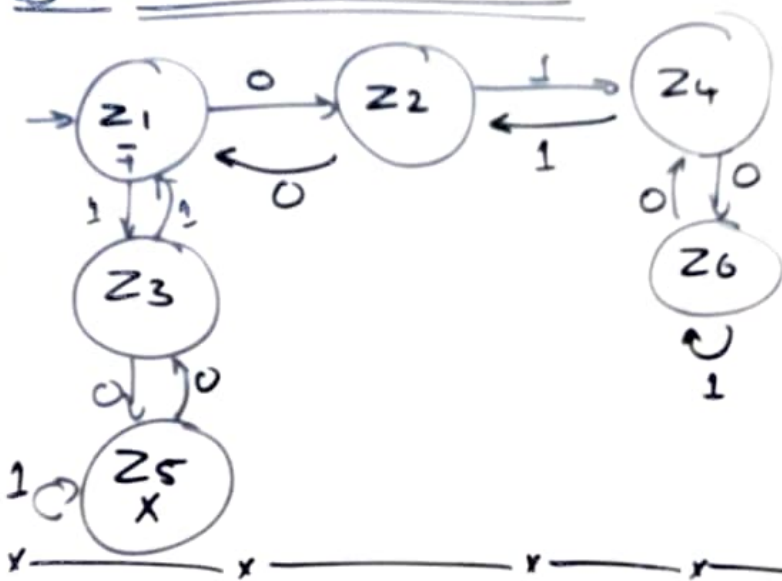


# Q1 DFA Minimization:-



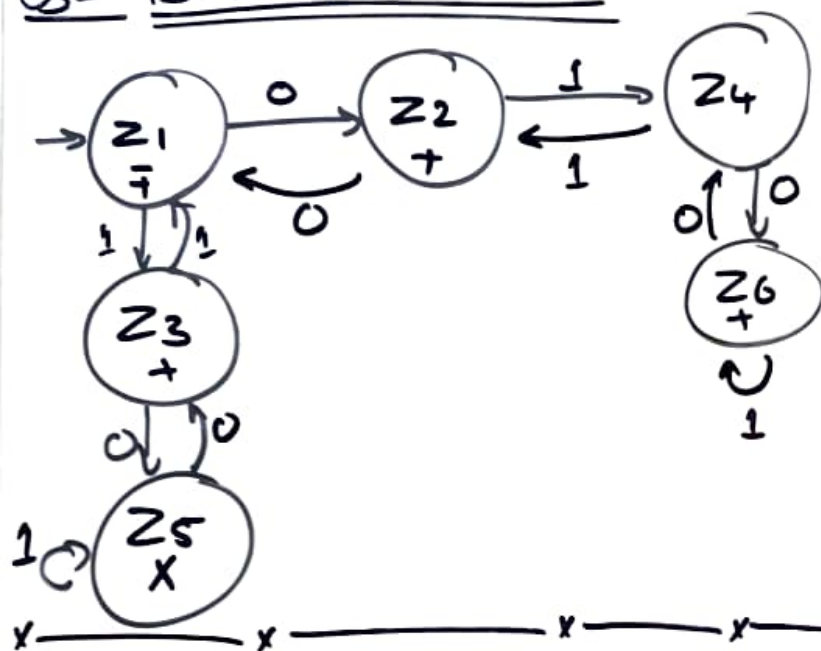
| states | 0  | 1  |
|--------|----|----|
| z1     | z2 | z3 |
| z2     | z1 | z4 |
| z3     | z5 | z1 |
| z4     | z6 | z2 |
| z5     | z3 | z5 |
| z6     | z4 | z6 |

Partition:-

$z_1 | z_2 z_3 z_4 z_5 z_6 |$   
 $z_1 | z_2 | z_3 z_4 z_5 z_6 |$   
 $z_1 | z_2 | z_3 | z_4 z_5 z_6 |$   
 $z_1 | z_2 | z_3 | z_4 | z_5 z_6 |$   
 $z_1 | z_2 | z_3 | z_4 | z_5 | z_6$

DFA is already  
Minimized

# Q1 DFA Minimization:-



| states           | 0              | 1              |
|------------------|----------------|----------------|
| - z <sub>1</sub> | z <sub>2</sub> | z <sub>3</sub> |
| + z <sub>2</sub> | z <sub>1</sub> | z <sub>4</sub> |
| + z <sub>3</sub> | z <sub>5</sub> | z <sub>1</sub> |
| z <sub>4</sub>   | z <sub>6</sub> | z <sub>2</sub> |
| z <sub>5</sub>   | z <sub>3</sub> | z <sub>5</sub> |
| + z <sub>6</sub> | z <sub>4</sub> | z <sub>6</sub> |

Partition:-

$z_1^x z_2^x z_3^x z_6^x | z_4^y z_5^y$

$z_1 | z_2^x z_3^x z_6^x | z_4 | z_5$

$z_1 | z_2 | z_3^x z_6^x | z_4 | z_5$

$z_1 | z_2 | z_3 | z_4 | z_5 | z_6$

DFA Already Minimized

## Q2 Regular Grammars

a)  $a^n b^m : (n+m) \text{ is even}$

$n \neq m \text{ odd}, n \neq m \text{ even}$

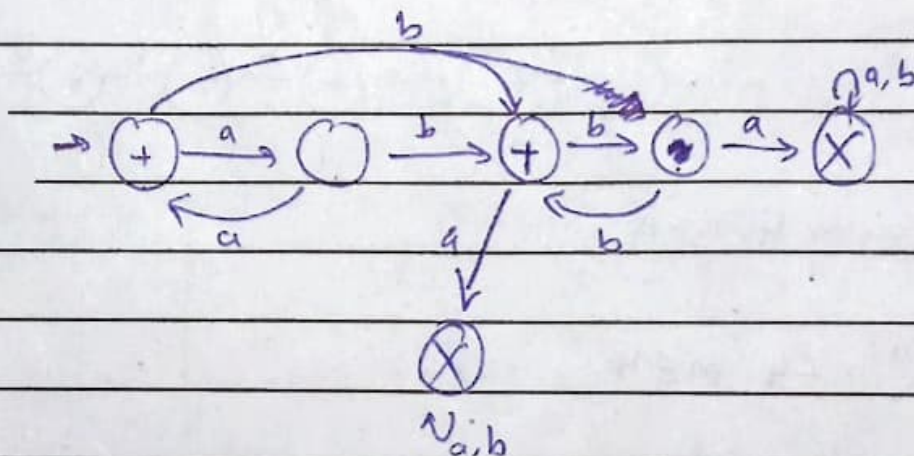
RE:  $a(aa)^* b(bb)^* + (aa)^* (bb)^*$

CFC:  $S \rightarrow aAB \mid AB \mid A \mid B$

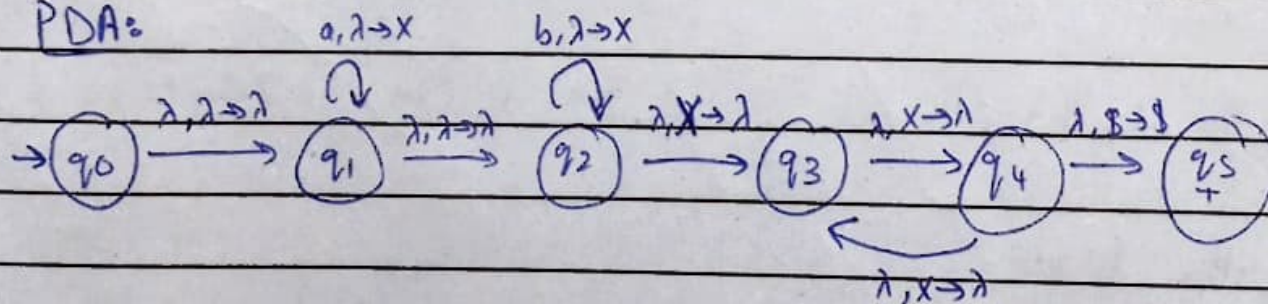
$A \rightarrow aaA \mid \lambda$

$B \rightarrow Bbb \mid \lambda$

DFA:



PDA:



~~RE: App~~

b)  $a^n b^m$ ,  $n \geq 4$ ,  $m \leq 3$

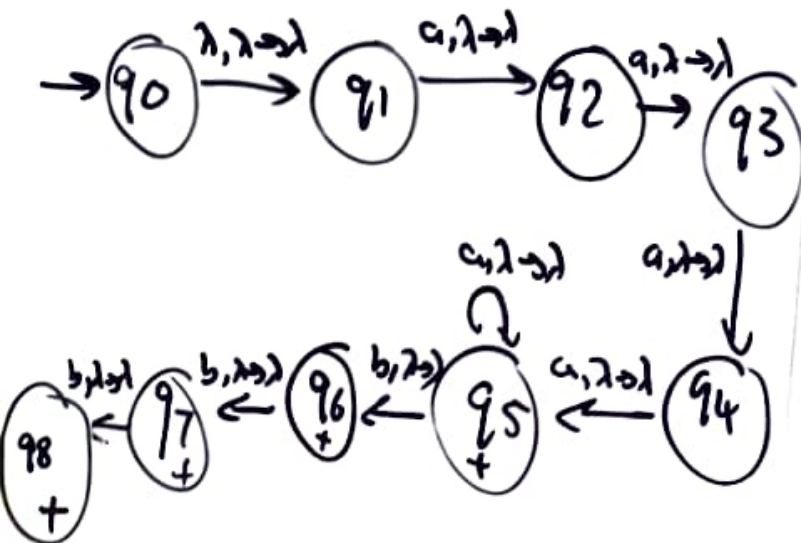
CFG:

$S \rightarrow AB$

$A \rightarrow aaaa|aA$

$B \rightarrow b|bb|bbb|\lambda$

PDA:



c)  $a^n b^m$ ,  $n < 4$ ,  $m \leq 4$

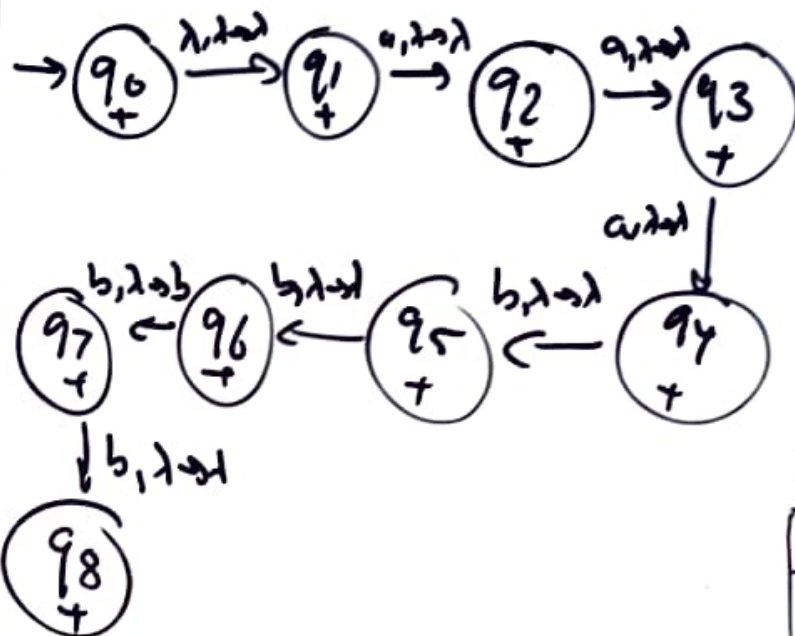
CFG:

$S \rightarrow AB$

$A \rightarrow a|aa|aaa$

$B \rightarrow b|bb|bbb|bbbb$

PDA:





d)

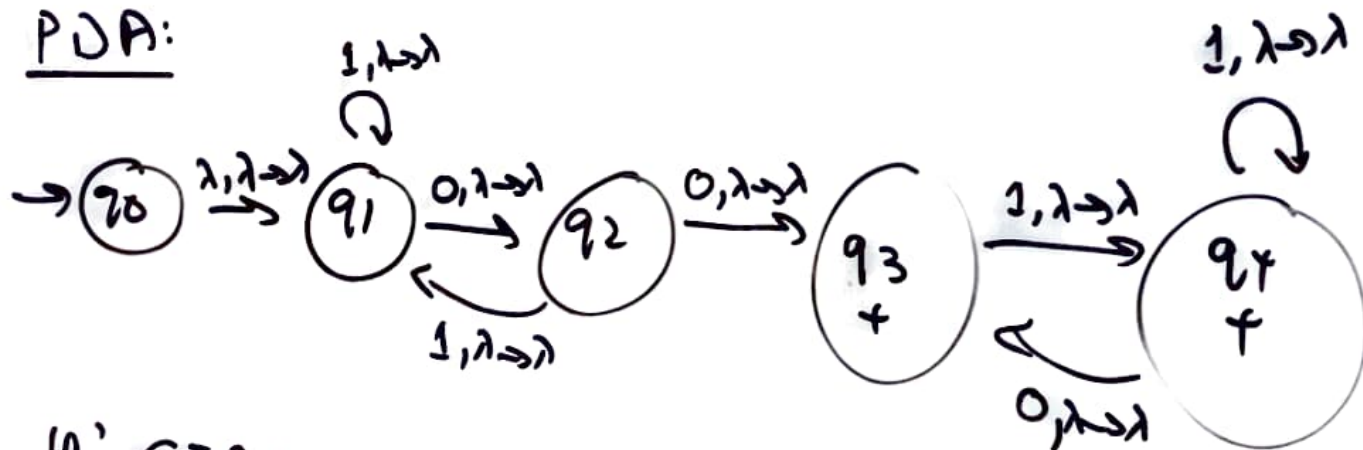
CFC:

$S \rightarrow A00B$

$A \rightarrow 1A \mid 01A \mid \lambda$

$B \rightarrow B1 \mid B10 \mid \lambda$

PDA:



'p' CFC:  $S \rightarrow \lambda$

PDA:



Q) All strings that contain at least 1 occurrence of each symbol in alphabet.

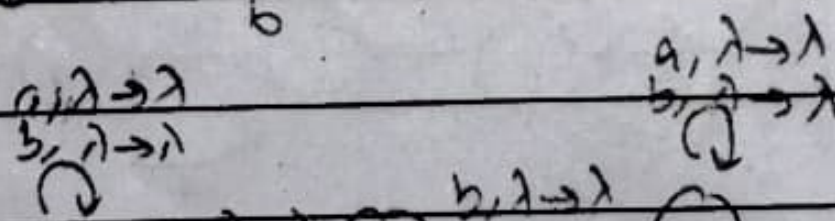
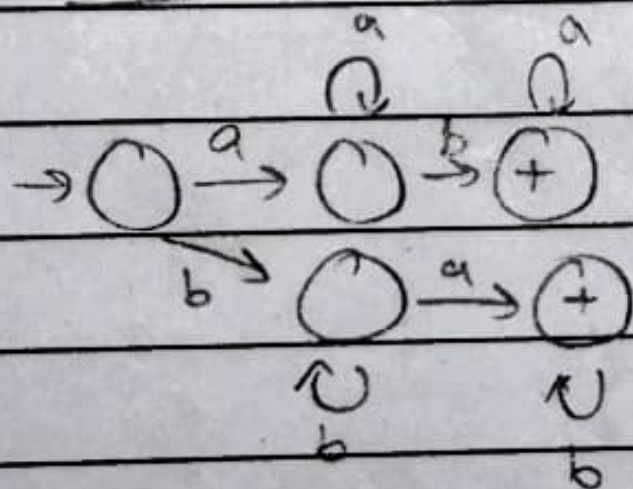
RE:  $(a+b)^* (ab+ba) (a+b)^*$

CFG:

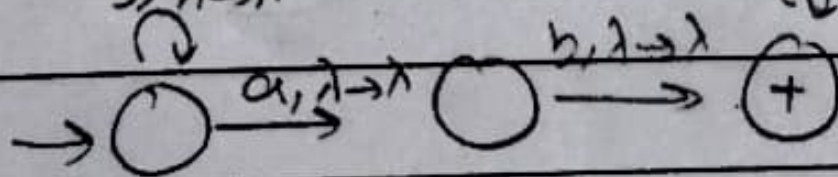
$S \rightarrow AabA \mid AbaA$

$A \rightarrow \cancel{aA} \mid aA \mid bA \mid \lambda$

DFA:



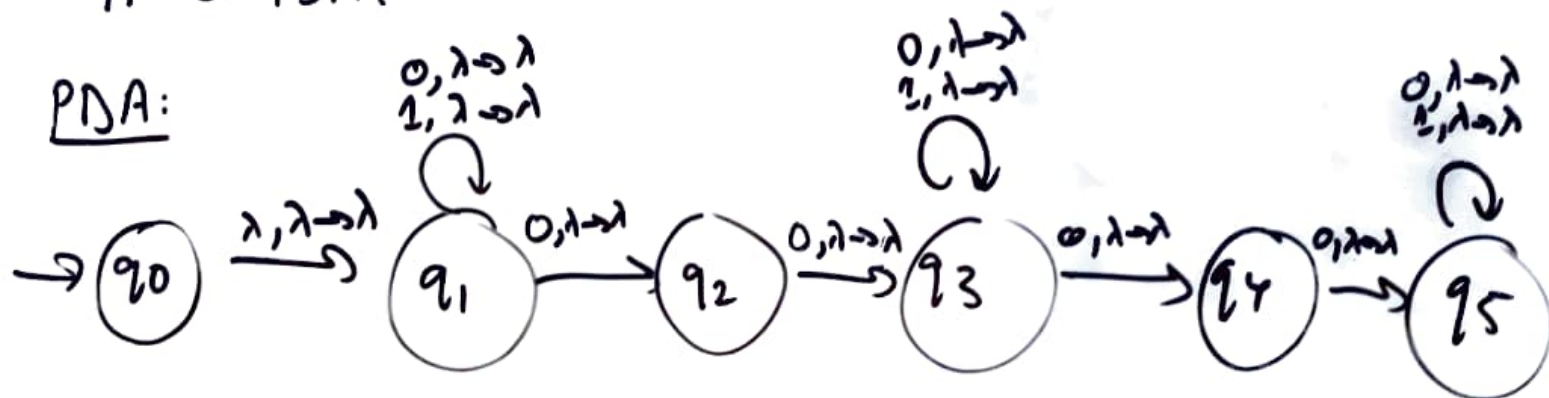
PDA:



g)  $S \rightarrow A00A00A$

$A \rightarrow 0A \mid 1A \mid \lambda$

PDA:

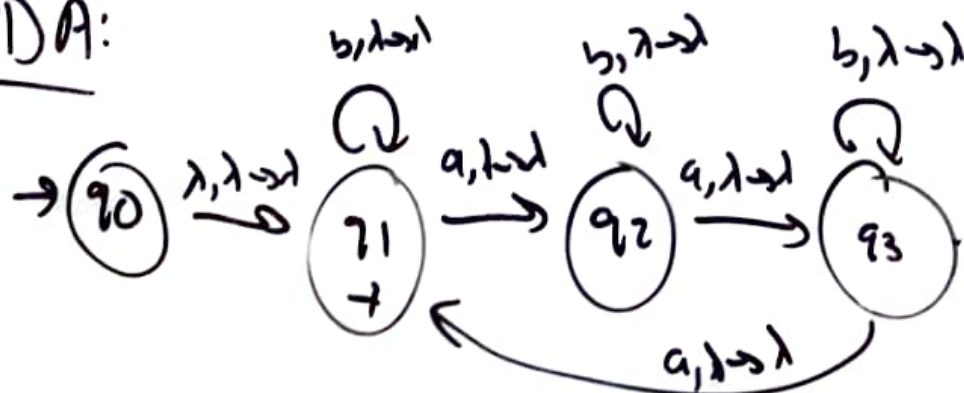


h) CFG:

$S \rightarrow AaAaAaAS \mid \lambda$

$A \rightarrow bA \mid \lambda$

PDA:





(83)

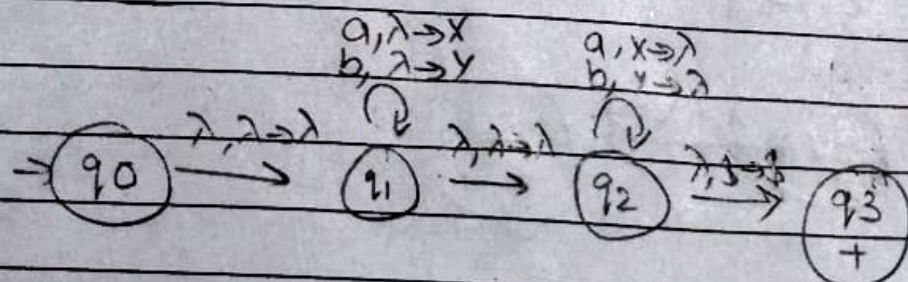
a) CFG:

 $S \rightarrow aSa \mid bSb \mid \lambda$ 

1a4b a

|   |  |
|---|--|
|   |  |
| y |  |
| x |  |

PDA:



Stack Operation

Test String: abba

|    |  |    |  |    |  |    |  |    |  |
|----|--|----|--|----|--|----|--|----|--|
|    |  |    |  |    |  |    |  |    |  |
|    |  |    |  | y  |  |    |  |    |  |
|    |  | x  |  | x  |  | x  |  | x  |  |
| \$ |  | \$ |  | \$ |  | \$ |  | \$ |  |

 $q_1, a, \lambda \rightarrow x$  $q_1, b, \lambda \rightarrow y$  $q_2, y \rightarrow \lambda$  $q_2, x \rightarrow \lambda$  $\lambda, \$ \rightarrow \$$ Pumping Lemma: $a^n b^n b^n a^n$ let  $n=2$  $|xy| \leq 2$ 

$$\begin{array}{ccccccc} a & a & b & b & b & b & a & a \\ \hline x & y & & & z & & & \end{array}$$
 $xy^2z = abbbbbaa$  reject, therefore not regular.

SOLO

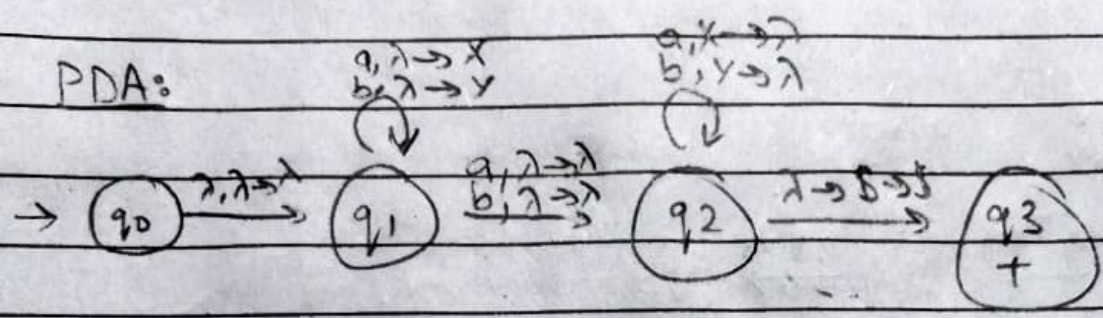


b)

CFG:-

$S \rightarrow aSa \mid bSb \mid a \mid b$

PDA:



Pumping Lemma:

$a^n b^n a^n$  let  $n=3$

$\underbrace{aaa}_x \underbrace{bbb}_y \underbrace{aaa}_z$

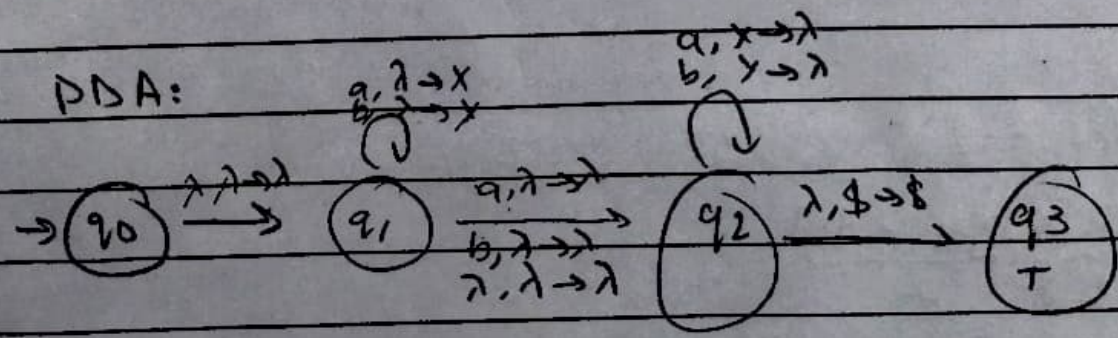
$|xy| \leq 3$

$abbbaaa$  & odd length palindrome therefore not regular.

c' CFG:

$S \rightarrow aSa \mid bSb \mid a \mid b \mid \lambda$

PDA:



Q  $a^n b a^n$

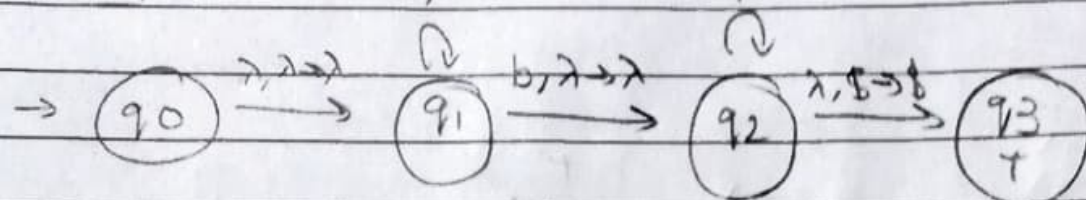
CFG:

$S \rightarrow aSa \mid b$

PDA:

$a, \lambda \rightarrow x$

$a, x \rightarrow \lambda$



Pumping lemma:

Let  $n=2$

$\underline{aa} \underline{b} \underline{aa}$   
 $\underline{xy} \quad \underline{z}$

$|xy| \leq n, |xy| \leq 2$

Pump  $y$  0 times:  $abaa$  doesn't exist in  $a^n b a^n$ ,  
therefore not Regular

Q' WW:  $W \in \{a, b\}^*$

Not Possible Not CFG



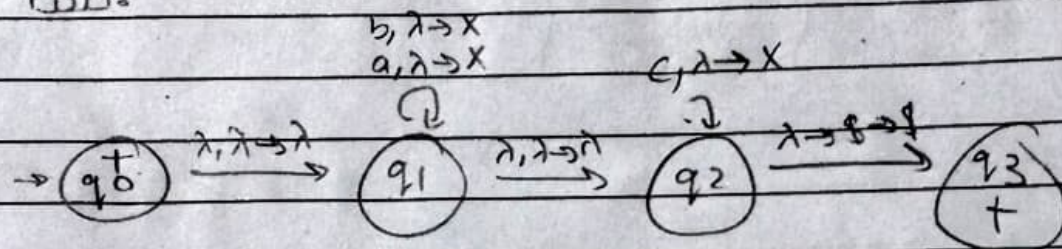
$\{a^n b^m c^{n+m} \mid n \geq 0, m \geq 0\}$

CFG:

$S \rightarrow aSc \mid \lambda \mid B$

$B \rightarrow bBc \mid \lambda$

PDA:



Stack Operations:

Test String abcc

|    |  |    |  |    |  |    |  |    |  |
|----|--|----|--|----|--|----|--|----|--|
|    |  |    |  |    |  |    |  |    |  |
|    |  |    |  | X  |  |    |  |    |  |
|    |  | X  |  | X  |  | X  |  | X  |  |
| \$ |  | \$ |  | \$ |  | \$ |  | \$ |  |

$q_1, a, \lambda \rightarrow X$

$q_1, b, \lambda \rightarrow X$

$q_2, c, X \rightarrow \lambda$

$q_2, c, X \rightarrow \lambda$

$\lambda, \$ \rightarrow \lambda$



Date \_\_\_\_\_

Pumping Lemma:

Let  $n=2$   $m=3$

$aabbcccc$   
 $\underbrace{aa}_{x} \underbrace{bb}_{y} \underbrace{cccc}_{z}$

$|xy| \leq 2$

pump  $y$  0 times  $aabbcccc \notin a^n b^m c^{n+m}$   
 therefore not regular.

$\exists a^{n+m} b^n c^m$

CFG:

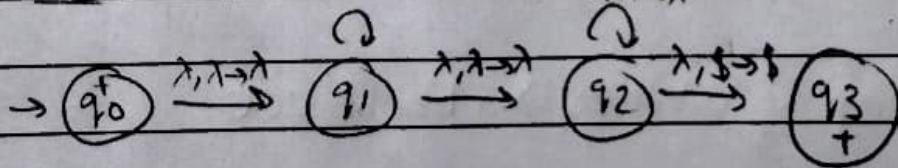
$S \rightarrow aSc \mid B \mid \lambda$

$B \rightarrow aBb \mid \lambda$

PDA:

$a, \lambda \rightarrow x$

$b, x \rightarrow \lambda$   
 $c, x \rightarrow \lambda$



Pumping Lemma:

Let  $n=2$ ,  $m=3$

$aaaaabbccc$

$aaaaabbccc$   
 $\underbrace{aaaa}_{xy} \underbrace{abb}_{z} ccc$

$|xy| \leq 2$

pump  $y$  0 times :  $aaaaabbccc \notin a^{n+m} b^n c^m$   
 Not Regular

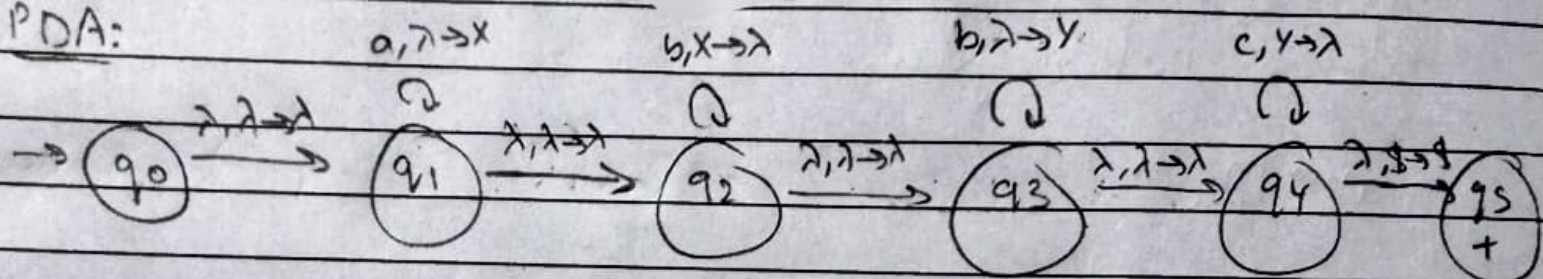


$$h \quad a^n b^{n+m} c^m$$

CFG:

$$S \rightarrow AB$$
$$A \rightarrow aAb \mid \lambda$$
$$B \rightarrow bBc \mid \lambda$$

PDA:



### Pumping Lemma:

Let  $n=2$ ,  $m=3$

$$\underline{|xy| \leq 2}$$

aa bbbbbbccc  
x y z

$\underbrace{aa}_{x} \underbrace{bbbb}_{y} \underbrace{ccc}_{z}$  pump  $y$  0 times  $\Rightarrow a bbbbbb ccc$

$a^m b^{n+m} c^m$  therefore not regular



abb

asbb

aB  
↓

Date \_\_\_\_\_

i)  $a^n b^{2m}$

$a(bb)^*$

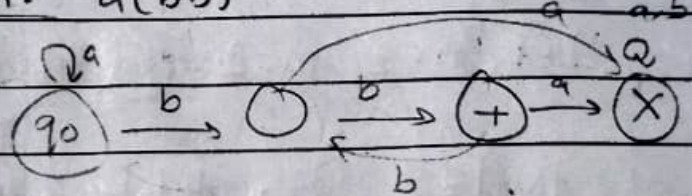
CFG:

$S \rightarrow AB$

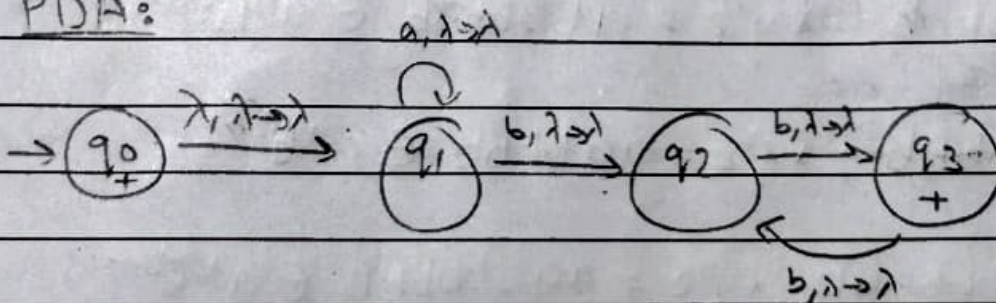
$A \rightarrow aA \mid \lambda$

$B \rightarrow bbB \mid \lambda$

DFA:  $a^*(bb)^*$



PDA:



Stack Operations

Test string: abb

|  |    |  |    |  |    |    |
|--|----|--|----|--|----|----|
|  |    |  |    |  |    |    |
|  |    |  |    |  |    |    |
|  |    |  |    |  |    |    |
|  | \$ |  | \$ |  | \$ | \$ |

$q_1, a, \lambda \rightarrow \lambda$

$q_2, b, \lambda \rightarrow \lambda$

$q_3, b, \lambda \rightarrow \lambda$







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## Stack Operations:-

test string aabbb

|    |  |    |   |    |   |    |  |    |  |
|----|--|----|---|----|---|----|--|----|--|
|    |  |    |   |    |   |    |  |    |  |
|    |  |    | X |    | X |    |  |    |  |
|    |  |    | X |    | X |    |  |    |  |
|    |  |    | X |    | X |    |  |    |  |
| \$ |  | \$ |   | \$ |   | \$ |  | \$ |  |

$q_1, a, \lambda \rightarrow xxx$

$q_2, b, x \rightarrow \lambda$

$q_2, b, X \rightarrow 1$

## Pumping Lemma:-

Let  $n=2$

$\overset{xy}{\text{aaaa}} \overset{z}{\text{bbbbbb}}$

$|xy| \leq 2$

$x y^2 z = \text{aaa bbbbbb} \neq a^n b^{3n}$  therefore not Regular.

Date \_\_\_\_\_

Q4 1)

remove  $B \rightarrow ab$  &  $S \rightarrow abB$  ~~unreachable~~ ~~unreachable~~

$S \rightarrow abS \mid abA \mid abB$

$S \rightarrow abS \mid abA$

$A \rightarrow cd$

$A \rightarrow cd$

$B \rightarrow ab$

$C \rightarrow dc$

$C \rightarrow dc$

remove  $C \rightarrow dc$  unreachable

Let  ~~$ab$~~   $X \rightarrow ab$ ,  $C \rightarrow c$ ,  $D \rightarrow d$ ,  $E \rightarrow a$ ,  $F \rightarrow b$

$S \rightarrow abS \mid abA$

$S \rightarrow XS \mid XA$

$A \rightarrow cd$

$A \rightarrow C \mid D$

$X \rightarrow EF$

$C \rightarrow c$

$D \rightarrow d$

$E \rightarrow a$

$F \rightarrow b$

2)  $S \rightarrow ABC \mid a$

remove  $E \rightarrow e$ ,  $F \rightarrow f$ ,  $G \rightarrow g$  unreachable

Let  $X \rightarrow AB$

$A \rightarrow b$

$B \rightarrow c$

$C \rightarrow d$

$E \rightarrow e$

$F \rightarrow f$

$G \rightarrow g$

$S \rightarrow ABC \mid a$

$A \rightarrow b$

$B \rightarrow c$

$C \rightarrow d$

$S \rightarrow XC \mid a$

$A \rightarrow b$

$B \rightarrow c$

$C \rightarrow d$

$X \rightarrow AB$



Date \_\_\_\_\_

Remove  $x \rightarrow SBD$  unreachable,  $A \rightarrow Bad$ ,  $A \rightarrow bSX$   
 $A \rightarrow a$

3)  $S \rightarrow aB | bX$

$S \rightarrow aB | bX$

$A \rightarrow Bad | bSX | a$

$B \rightarrow aSB | bBX$

$B \rightarrow aSB | bBX$

$X \rightarrow SBD | aBX | ad$

$X \rightarrow aBX | ad$

Remove  $B \rightarrow aSB | bBX$ ,  $x \rightarrow aBX$ ,  $A \rightarrow Bad$  can't generate anything

$S \rightarrow bX$

$X \rightarrow ad$

Let  $C \rightarrow a$ ,  $D \rightarrow b$ ,  $E \rightarrow d$

$S \rightarrow DX$

$X \rightarrow CE$

$C \rightarrow a$

$D \rightarrow b$

$E \rightarrow d$

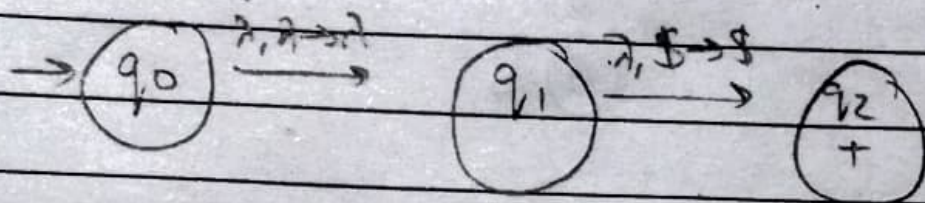


Q5 a)  $S \rightarrow XS \mid \lambda$

$X \notin A$  is unreachable

$A \rightarrow aXb \mid Ab \mid ab$

$S \rightarrow \lambda$



b) PDA Not possible

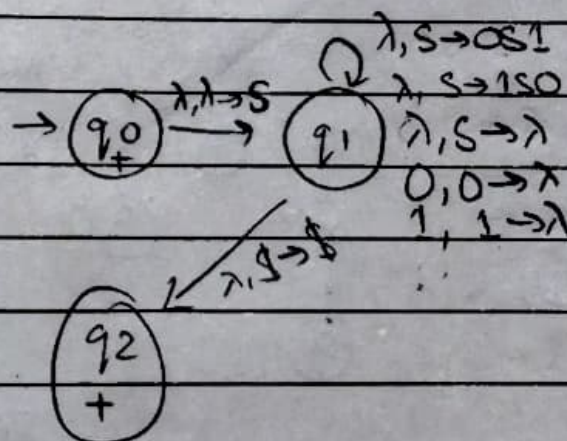
c)  $S \rightarrow OS1 \mid 1SO \mid \lambda$

$S \rightarrow OSX \mid 1SY \mid \lambda$

$X \rightarrow 1$

$Y \rightarrow 0$

} same

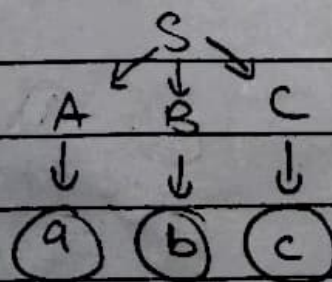


Q6 a)  $S \rightarrow ABC$

$A \rightarrow a$

$B \rightarrow b$

$C \rightarrow c$



No ambiguity because there is only 1 Derivation tree

$\frac{aS}{aaS}$

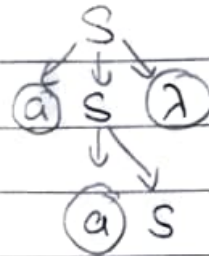
Date \_\_\_\_\_

Q6b

Test string  $aas$

$$S \rightarrow aS \mid \lambda$$

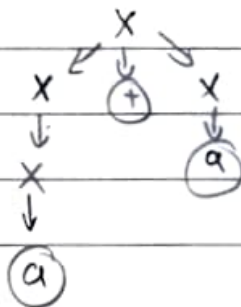
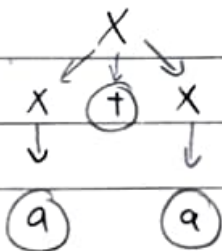
No ambiguity only 1 Derivation Tree



Q6c

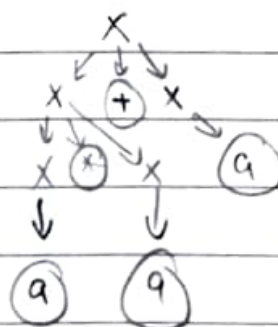
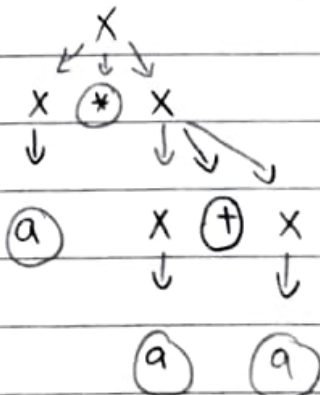
$$X \rightarrow X+X \mid X*X \mid X \mid a$$

Test string  $a+a$



Multiple Derivation Tree's  
Ambiguous

Test string  $a*a+a$



Multiple Derivation Trees  
Ambiguous