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Transition graph

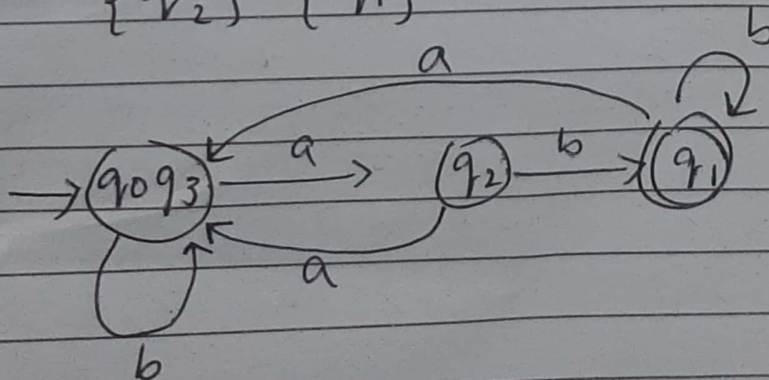
	a	b
→ q ₀	q ₂	q ₃
q ₁	q ₀	q ₁
q ₂	q ₀	q ₁
q ₃	q ₂	q ₃

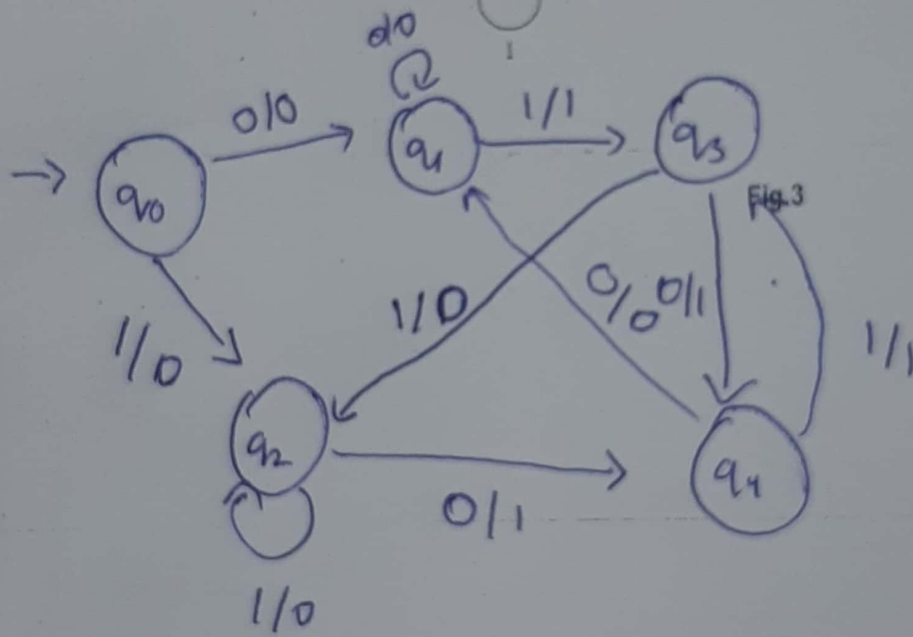
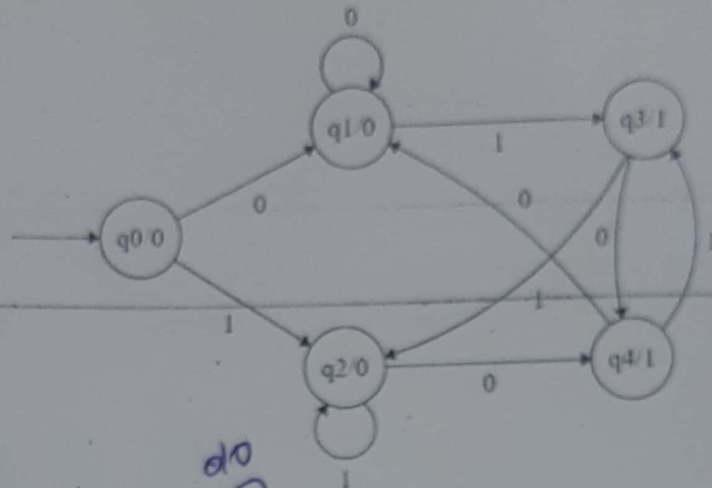
#1

{q₀, q₂, q₃} {q₁}

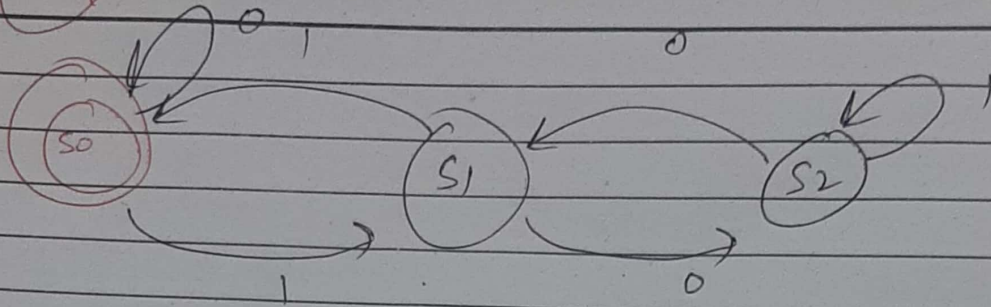
{q₀, q₃}

{q₂} {q₁}





#2



old states	0	1
$S_0 \equiv Z_1$	$S_0 \equiv Z_1$	$S_1 \equiv Z_2$
$S_2 \equiv Z$		
$S_0 \equiv Z_1$		

old states	0	1
$S_0 \equiv S_0$	$S_0 \equiv S_0$	$S_1 \equiv S_1$
$S_1 \equiv S_1$	$S_2 \equiv S_2$	$S_0 \equiv S_0$
$S_2 \equiv S_2$	$S_1 \equiv S_1$	$S_2 \equiv S_2$
Same	DFA	would be generated.

Q#3

www, $w \in \{a, b\}^*$

lets take $a^n b^n a^n b^n a^n b^n$.

Assume $n \geq 4$ as pumping length.

the sample word would be -
 $aaaa bbbb aaaa bbbb aaaa bbbb$.

let divide it into x, y, z -

$\underbrace{aaaa}_{x} \underbrace{bbbb}_{y} \underbrace{aaaa bbbb aaaa bbbb}_{z}$

$xy^0z = aabbbb aaaa bbbb aaaa bbbb$

this not part of original language.

hence this is dis proof.

Part 2
 \Rightarrow ~~$(aaaa)^*$~~ aa language of this pattern
 which are regular.

$(aaa)^*$
 $(bbb)^*$