

Fig. 3.27. NFA N for $(a \mid b)*abb$.

$$C = \epsilon \text{-closure}(\{5\}) = \{1, 2, 4, 5, 6, 7\}$$

Thus, Dtran(A, b) = C.

If we continue this process with the now unmarked sets B and C, we eventually reach the point where all sets that are states of the DFA are marked. This is certain since there are "only" 2^{11} different subsets of a set of eleven states, and a set, once marked, is marked forever. The five different sets of states we actually construct are:

$$A = \{0, 1, 2, 4, 7\}$$
 $D = \{1, 2, 4, 5, 6, 7, 9\}$
 $B = \{1, 2, 3, 4, 6, 7, 8\}$ $E = \{1, 2, 4, 5, 6, 7, 10\}$
 $C = \{1, 2, 4, 5, 6, 7\}$

State A is the start state, and state E is the only accepting state. The complete transition table Dtran is shown in Fig. 3.28.

STATE	INPUT SYMBOL	
	a	ь
A	В	С
В	В	D
c	В	C
D	В	E
E	В	C

Fig. 3.28. Transition table *Dtran* for DFA.

Also, a transition graph for the resulting DFA is shown in Fig. 3.29. It should be noted that the DFA of Fig. 3.23 also accepts $(a \mid b)*abb$ and has one