Assignment No. 01

| Rubric | Score (0 to 4) |
|---------------|---|
| Delivery | MATERIAL PROPERTY OF THE PARTY |
| Understanding | |
| Readability | |
| Discipline | |
| Total | |

Performed On: 03 | 62 | 25

Sign: _____

Q.1 Differentiate between RPC/RMI Remote Method Invocation Remote Procedure Call OA Java-Speciafic O A protocol that allows mechanism to invoke a functions procedure to methods on remote be executed in a remote objects. Sygtem Specific @ Language - independent 1 Java 3 object orientes 3) Procedural programming programming @ Uses Java remote (1) Typically implemented method protocol for Usika lower-level communication. protocols (5) Limited to Java (3) can be used across, Envisonments. different languages plattorms. (6) Automatically Regulars manual serializes Java objects. Sexialization of data Strictures. 1) Depends on underlying 1) provides built-in Security mechanisms nervone process like SSL for encryption (8) More complex to (8) Easier to use in Java applications due Implement in heterogeneous to built-in support Environments a) can be optimized based @Slightly slower due on Protocol and to Viava object Implementation sexialization and deserialization overhead

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0.2] What are services offered by middlewarel. Ans. Middleware is software that connects different application in a distributed System. It helps in communication data management, and security between computers intake were acts bridge between operating system and application, providing essential services to ensure smooth operations.

the sound of the adaption of the

1. Communication Sexuices! application to communicate enables different application to lemmunicate efficiently by handling message exchange and remote calls.

· Remote Procedure Call (RPC); - Allows calling

a function on another machine

· Message oriented Middleware: Uses message queves for asynchrones communication

2. Semony sexuice: - dollars

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· Security is Critical in distributed System
Middleware
. Authentication: - Ensure that only authorized
users can access resources.

Users can access resources.

· Encouption! - Protect data dunny transmission . Alless control: Restricts uses permission bases on voles.

3. Transaction Management!

o In distributed System, multiple operation need to be executed together.
Atomicity: A toansaction is either fully completed or rolled back.
- Consistency: Keeps data accrate across

· Durability: Ensures data remains even after a system failure.

- 4. Resource Management!

 Middleware helps manages system resources efficiently by load balancing distributes processing across multiple sources.
 - · Scalability) Allow System to handle Moreases workload
 - · fault tolarance: Ensures System stability
 even if a part of system fails
- 5. Data Management & Synchoonization;
 Middleware ensure data Consistency guross
 distributed System by.
 - Data replication: Create multiple copses
 - Data Caching! Stores frequently acrosses duta to reduce load terms.
 - 6 Transpasency sexure:
 - -Middleware hides the complexity of distribute System, making them easier to use.
 - resources without knowing their Physical location.
 - · Access toansparrency: Ensures different data formats do not affect system usage

The program of the Control of the Co V Cl.3) Explain Raymond's toce-based algorithm V S Ans. Raymond's Algorithm is a distributed mutual exclusion algorithm that organizes
processes in a logical tree structure to
efficiently manage access to a shared
resources. It is a token - based algorithm
meaning that a process must hold a token
to entex it contrical section. S 1 1 -* Key Features of Raymond's Algorithm 1 1. Tree Stoucture! The processes are arranged in a logical tree, where each node has a parent and potentially multiple children N N 2 Single Token: Only One token exists
In the System. A process must hold
the token to enter critical section. > > 3. Prisorty to closest Requests: The token is always help by the process that last uses its ensuring a form of locality based alcess > 2 he Efficient Request forwarding. Requests
for the token are forwarded up the
free toward token holder. 2 2 2 Example 1- Consider a tree with P2 (2007)
howing Children P2 pand P3, where
P3 has a child P4 2 2 1. Suppose Pu requests the token. 2000000 2. P4 -> P3 -> P1 3. PI sends the token to BB; then to P4

4. P4 executes Ar corrical sections and after completion, passes the token the hext requester. dt Advantagest

- · Less message over head L · Efficient token passing · Scalability

- + Disadvatages:

 Single point of failure
 - Delay due to tree traversal.

O.G. Explain: (ii) Any one Physical clock synchronization algorithm! (1) Any one Election Algorithm; Bully Algorithm: The Bully Algorithm is used in distributed Systems the existing the existing one fails. It is called a bolly algorithm because the process with the highest because the process with the highest becomes the reader.

Steps of Bully Algorithm: Steps of Bully Algorithm? 1. Failure Desection: When a process desects
that the coordinator is not responding
it initiates the election. 2 Election process:

The process that detects failure sends an electron message to all higher-ID an electron message is received, it declares itself the leader.

Itself the leader.

That process takes over the electron. 55556666666 3. New Leader seleptions? The process with the highest ID with the election and broadcasts a coordinator message to all other processes. mer recor of the new leases. Color Shared Spread Control of the

(ii) Physical Clock Synthron12ahon Algenthme cristian's Algerithm:

chistian's Algorithm is used to Synchromlann a distributes System.

Steps of Consider's Algorithm: 1. The client sends a time request message to the time server.

- 2 The server seeeives the request and responds with 18 current time.
- 3. The client estimates the round-trip delay (RTT) and adjusts its Clock Using!
 Tadjusted = Tserver & RTT

1. The client sets its clock to the agusted time

Advantages:
Simple & efficient
Nooks well in LAN

Disadvantage:

, Assumes symmetal network delay

· single accurate the server.

Short note on! O.5) Write a Compare Issues related to distorbute 77 (1) Goals and System. NOS DOS 1) Network operating 1) Mos stand () DOS stand to System is full for Middleware Distributed operating system form NOS Operating System QA software @ An Os That 1) An Os That layer that marges across manges multiple networkes compute enables communication connected and resources while mantang computers as Shanna between their independence a Single distributed application Sy stem. > (3) Bridges different 3 Each System 3 Appear as a Os & application operates independently single system to users to enables but network mteropera 6 1ity communication is possible. 2) Transparent, @ Explicit, usess (4) Provides APIS >> for seamless must manually all resonnés access shared Sharing and appear as part of the system resources. in tegounon 3 tay 1+ Tolerance & fault to letance Fault tolerance is High due is low system is VIa redundany does no affect and exor others. 2 process handing. management 2 (6) Example: CORBA , Java (Example: whows server, Amogba, RMI, NET LMUX - based Spride, Remoting. plan 9. network OS. 2 2 2 2

(ii) Goals and Issues related to distributed Chapter 1954 System: -Goals of Distributed Systems. A distributed S designed to achieve the forlowing System is Objectives! 1. Transpadency - Useos should experience the System as a single entity rather than multiple interconnected components. · Access Transparency · Location Transparency · Replication Transparency · Concurrency Transparency · failure Transpavency. 2. Scalability: The System should efficiently handle godinth in terms of wers, data, en 3. fault Tolexance - The system should continue functioning even it some components to seem hadgen fails. 4. Resources Sharing - Rejources like files printers, and databases should be easily shares among users. 5. Concurrency - Multiple users/ processes should be able to operate shultaneously 333 without interference 6. Openness: The system should support Interoperability and be able to integrate different hardware and software component 20000 7. security! Mechanisms must be in place po protect data, control access and energe privary.

- Issues in Distributed Systems: Despite their advantages, distributed Systems face serveral Challenges:
 - 2. Synchoonization Issues: Iceeping clocks Synchronized across different mathines is difficult.
 - 2. Communication failures: Nectworks failure:
 can disrupt interactions between hodes,
 requiring fault-tolerant meachanisms.
 - 3. Comcurrency control: Ensuring data consistency when multiple processes across shared sesources is complex.
 - 4 Deadlocks- Processes wasting indefinitely for resources held by others can lead to system hangs.
 - 5. Security Risks Distributed Systems are more vulcerable to cyberattacks due to remove alless and data withing showing
 - 6. Load Sharing Efficiently distributing workload alooss multiple hodes to prevent bothleneurs is Challenging.