Assignment No. 02

Rubric	Score (0 to 4)
Delivery	
Understanding	
Readability	and the Control of the Control
Discipline	In Buckley
Total	

Performed On: 11 03 25

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Sign:

and process (1) Explain different load estimation and protocology transfer policies used by load bakering algorithms, this. Load balancing ensures that computational tasks are evaly distributed accoss multiple hodes in distributed system to prevent bottleneites and improve efficiency. Load balancing algorithms are different load estimation policies and process transfer policies to decides when and how to distribute tasks. 1. Load Estimation Policies: - Load estimation policies determine how the system measures and evaluates the workload of each node. (i) CPV Queue Length-Based Estimation:
Measyning humber of processes waiting
in the CPV queue.
A higher queue length indicates a higher Load, U · simple but may hot reflect the actual processing speed (ii) CPU Utilization-Based Estimation:
Measures the percentage of CPU usage over a given period.
High CPU Utilization suggests an overloader - can be maccurate it sher resources. (iii) Response Time - Based Estimation: -Evaluates how long it takes for a Projess to execute . Jugh response time suggests system is under heavy load. · Useful for interactive applications.

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2. Process Transfer policies: - Process transfer policies determine when and which process should be toansferred from an overloaded note to and understanding note (1) Threshold - Bases Policy: Defines upper and lower Threshold values for System load.

If a node's load exceeds the upper threshold, it transfers processes to a lower-load hode.

If a load falls below the lower threshold, it can accept processes. (ii) Load Difference - Based Policy: Transfers processes when the sufference between the most loaded and least loades note exceeds à predefines . Ensures and even distribytion of tasks. threshold. (iii) Probabilistic Transfer Policy:-· Instead of fixed thresholds, it tandomly selects processes for transfer bases on prophability. Lecision-making overhead but may be methicient in some cases. (iv) sender-initiated us receives mitiates Policy. · sender-initiates: - An overloades node actively looks for an underloades node to sene processes. · Pecerver - Initiates: An underloaded node requests tasks from overloaded nodes.

0.2] Explain any 5 data central consistency model with examples? Ans. A- data-centor consistency model defines how updates to a shared data store appear to different processes in a distributed system These models determine the order and visibility of read and write operations to ensure consistency while balancing system performance Here are five key consistency models with examples:

1. Strict Consistency: Finsures that any read operation returns The most recent write value Every process observes the same order of updates 'm real - time.

Example:

· Process P1 writes X = 10 at time t1. · Process P2 reads X at time t2

P2 must see X = 10 immediately, enoung

2. Sequential Consistency: - Ensures all processes observe operations in same order, but not necessarily in real-time. The sequence of operations must be consistent actoss all nodes.

· P1 1 x=10, then Y=20

. P2: Reads Y=20, then before Y=20 . P3: Must see X=10 before Y=20

3. Causal Consistency: Ensures that causally related operations appear in the correct order. Unitelated operations can appear in different orders from different processes.

Example:

· P1: Wastes x =10

· P2: Reads X=10, then writes Y=20

· P3: Must see x=10 before Y=20, but another write (Z=30) unrelated to 2 and Y can appear in any order.

4. Eventual Consistency: Ensures that if no updates occur, all nodes eventually converge to the same value.

Example: A user updates their profile picture on a social media platform.

5. Lineariability: A stricter from of sequential consistency where operations appear historianeously across all nodes.

Example! writes x = 10 at time t1.

• p1: Writes x = 10 at time t2 > t1

and must see x = 10 immediately

• P3: Cannot see an old value (x = 5)after p2 has seen x = 10.

0.3) what are the desirable features of good DFS?. Ans. A Distributed file System (DFS) allow users to access and manage files across multiple networked machines as if they were stored locally. A good DFS should have the following desirable features:

1. Transpasency:

. Access toanspasency: users should access
remote files as if they were local

. Location Transparency: file location

should not affect the users operations.

- · Replication Transparency.
- 2. Scalability: It should support dynamic addition of storage nodes.
- 3. fauit tolerance and Reliability 1.

 Data replication ensures availability even it a server fails.

 Redundant storage minimizes data loss vists.
- 4. Performance Efficiency:

 fast file retrieval and minimal network.

 latency.

 Load balancing mechanism should

 prevent bottleneurs.
- 5. Security and Access Control:
 - · Authentication and authorization to control access to files. Secure data toansmission.
- 6. Ease of Use and Manageability:

 Oseo-friendly interface for file operations

 Easy integration with existing application

 and operating systems.
- 1. Support for Heserogeneity:
 Support different file formats and access
 protocols (Eg. NFS, SMB, Hadoop HDF).

O. al Explain various file caching schemes

O. i) Explain various file caching Schemes.

Ans. file caching is a technique used in distributed file System to improve performance by reducing access time and inetwork traiffic. Various Caching schemes exist to optimize file access based on system requirements.

1. Client - Side Caching The client machine store recently accessed files or file tragments locally to reduce the need

tor remove access.

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advantage: - Reduces server load and network latency.

Toproves response the for forguently accessed files.

disadvantages:Risk of stale data if files are modifical at the server but non updated in cacher

2. Server-Size Caching: The Server maintains a Caching of frequently requested files to Serve multiple dients efficiently.

advantages:
Petites disk I/O operations.
Improves performance for shares read operations.

disadvantages:
. Increased memory ugage on the server
. May not benefit cliens with highly
personalizer Lata access patterns.

3. Disk-Based Caching 1- Instead of using memory cached files are stored on disk to handle large datasets.

advantage: · more storage capacity than RAM-bases disadvantage! · Slower than RAM - based caching due to disk read with overhead. h. Write-Thoough Caching! - Every write operation is immediately propagated to the serves before being acknowledged · hudantees data integrity.
· No fish of data loss due to cache failure. disabantages; attency due to frequent heawork communication slower performance compared to other dacking schemes. 5. Write -Back Caching: - write operations are temporgrily stores in the cache and written to the server only at specific intervals. a dvantages: · Resules network touffir and Improves work performance. Useful for applications with frequent updates to the same file. disadvantages: · May cause consistency isoves if multiple clients modify same file.

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0.5] Explain Code migration.

Ans. code migration refers to the process of moving program code, execution state, or both between different machines in a distributed system. This helps in improving performance, load balancing & resources utilization.

Types of Code Migration:

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1. Process Migration
- Moves a running process from one machine to another wall balancing and fault to lexange.

Useful for load balancing and fault to lexange.

Example: Migrating a running application from an overloaded server to less busy one.

2 computation Migration:

Moves only computation while data
remains on the original machine
reduces network toaffic by moving
computation closes to the data.

Example: Running datasets givery on
the server instead of fetching all data
to the client

3. Data Migration:

Move desta from one node to another for performance optimization.

Observed in cloude storage and distributes destabases.

Example: Moving frequently accessed files to a faster storage system.