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1 Number of different binary trees

(From CLRS 12-4)

Let b_n denote the number of different binary trees with n nodes. In this problem, you will find a formula for b_n , as well as an asymptotic estimate.

1. Show that $b_0 = 1$ and that, for $n \ge 1$,

$$b_n = \sum_{k=0}^{n-1} b_k b_{n-1-k}$$

显然 $b_0 = 1$

那么 b_n 的情况有 左子树有k个节点, 右子树有n-1-k个节点, k可以从0取到n-1, 且两边独立, 使用乘法原则, 故而

$$b_{n} = \sum_{k=0}^{n-1} b_{k} b_{n-1-k}$$

2. Let B(x) be the generating function

$$B(x) = \sum_{n=0}^{\infty} b_n x^n$$

Show that $B(x) = xB(x)^2 + 1$, and hence one way to express B(x) in closed form is

$$B(x) = \frac{1}{2x}(1 - \sqrt{1 - 4x})$$

$$B(x)^{2} = \sum_{0 \le i,j}^{\infty} b_{i} b_{j} x^{i+j} = 1 + \sum_{n=1}^{\infty} \left(x^{n} \sum_{k=0}^{\infty} b_{k} b_{n-k} \right) = 1 + \sum_{n=1}^{\infty} \left(b_{n+1} x^{n} \right)$$
$$xB(x)^{2} + 1 = x \left(1 + \sum_{n=1}^{\infty} \left(b_{n+1} x^{n} \right) \right) + 1 = \sum_{n=1}^{\infty} \left(b_{n} x^{n} \right) + 1$$
$$= \sum_{n=0}^{\infty} \left(b_{n} x^{n} \right) = B(x)$$

当我们得到上述等式后,根据求根公式得

$$B(x) = \frac{1 \pm \sqrt{1 - 4x}}{2x}$$

$$\sum_{x\to 0} B(x) = 0,$$

$$\therefore B(x) = \frac{1}{2x} \left(1 - \sqrt{1 - 4x} \right)$$

3. Show that

$$b_n = \frac{1}{n+1} \binom{2n}{n}$$

(the nth Catalan number) by using the Taylor expansion of $\sqrt{1-4x}$ around x=0.

p.s. If you wish, instead of using the Taylor expansion, you may use the generalization of the binomial theorem (where n can be any real number) to noninteger exponents.

方法一:

$$b_0 = 1 = \frac{1}{1} \binom{0}{0}$$
; $b_1 = 1 = \frac{1}{2} \binom{2}{1}$ 我们假设 $n \le k$: $b_n = \frac{1}{n+1} \binom{2n}{n}$; 现需要证明: $n = k+1$, $b_n = \frac{1}{n+1} \binom{2n}{n}$ 我们有: $b_n = \sum_{i=0}^{n-1} b_i b_{n-1-i}$ 后面暂时没想出来 ...

方法二:

$$\sqrt{1-4x} = 1 + \frac{-4x}{2} - \sum_{n=2}^{\infty} \frac{(2n-3)!!}{2^n * n!} (4x)^n = 1 - 2x - \sum_{n=2}^{\infty} 2 * \frac{(4n-6)!!}{n!} x^n$$

$$B(x) = \frac{1}{2x} \left(1 - \sqrt{1-4x} \right) = \frac{1}{2x} \left(2x + 2 \sum_{n=2}^{\infty} \frac{(4n-6)!!}{n!} x^n \right) = 1 + \sum_{n=2}^{\infty} \frac{(4n-6)!!}{n!} x^{n-1}$$

$$= 1 + \sum_{n=1}^{\infty} \frac{(4n-2)!!}{(n+1)!} x^n = 1 + \sum_{n=1}^{\infty} b_n x^n$$

$$\therefore b_n = \frac{(4n-2)!!}{(n+1)!} (n > 0) = \frac{1}{n+1} \frac{2^n (2n-1)!!}{n!} = \frac{1}{n+1} \frac{(2n)!! (2n-1)!!}{n! n!}$$

$$= \frac{1}{n+1} \frac{(2n)!! (2n-1)!!}{n! n!} = \frac{1}{n+1} \left(\frac{(2n)!}{n! n!} \right) = \frac{1}{n+1} \binom{2n}{n}$$

4. Show that

$$b_n = \frac{4^n}{\sqrt{\pi}n^{3/2}} (1 + O(1/n))$$

$$b_n = \frac{1}{n+1} \binom{2n}{n} = \frac{1}{n+1} \left(\frac{(2n)!}{n! \, n!} \right) = \frac{1}{n+1} \left(\frac{\sqrt{2\pi 2n} \left(\frac{2n}{e} \right)^{2n}}{\sqrt{2\pi n} \left(\frac{n}{e} \right)^n \sqrt{2\pi n} \left(\frac{n}{e} \right)^n} \right)$$

$$= \frac{1}{n+1} \left(\frac{(2)^{2n}}{\sqrt{\pi n}} \right) = \frac{1}{n} \left(\frac{n}{1+n} \right) \left(\frac{(2)^{2n}}{\sqrt{\pi n}} \right) = \frac{1}{n} \left(\frac{4^n}{\sqrt{\pi}n^{\frac{1}{2}}} \right) \left(1 - \left(\frac{1}{n+1} \right) \right) = \frac{4^n}{\sqrt{\pi}n^{\frac{3}{2}}} (1 + O\left(\frac{1}{n} \right))$$

2 AVL trees

(From CLRS 13-3)

An **AVL tree** is a binary search tree that is **height balanced**: for each node x, the heights of the left and right subtrees of differ by at most 1. To implement an AVL tree, maintain an extra attribute h in each node such that x.h is the height of node x. As for any other binary search tree T, assume that node T.root points to the root node. Prove that an AVL tree with n nodes has height $O(\lg n)$.

(Hint: Prove that an AVL tree of height h has at least F_h nodes, where F_h is the hth Fibonacci number.)

对h进行归纳

- 1. h = 1 时, nodeNumber = $1 \ge F_1 = 1$
- 2. 假设 $h \le k$ 时, nodeNumber_h $\ge F_h$
- 3. 下证: h = k + 1 时 $nodeNumber_h \ge F_h$ k+1高的树 至少有一个k高的子树,且不存在高大于k的子树,那么根据ALV的性质,我们知道另一个子 树的高度为k或k-1,

所以 $nodeNumber_{k+1} \ge nodeNumber_k + nodeNumber_{k-1} \ge F_k + F_{k-1} = F_{k+1}$

而
$$F_h = O\left(\frac{\varphi^{\rm h} - (1-\varphi)^{\rm h}}{\sqrt{5}}\right)$$
 其中 $\varphi = \frac{1+\sqrt{5}}{2}$ $F_h = O(a^{\rm h})$ 其中 a 为某个常数 故而当有 n 个节点时,其高度为 $\log(n)$