

A2

1. Array-based Data Structures

1.1 Implement two Stacks in one Array

- *From 10.1-3 of CLRS*

Explain how to implement two stacks in one array $A[1 : n]$ in such a way that neither stack overflows unless the total number of elements in both stacks together is n .

The `PUSH` and `POP` operations should run in $O(1)$ time.

1.1 Solution [code/TwoStackinOneArray.cpp](#)

2. Linked List

2.1 Reverse a Linked List

- *From 10.2-5 of CLRS*

Give a $\Theta(n)$ -time non-recursive procedure that reverses a singly linked list of elements.

The procedure should use no more than constant storage beyond that needed for the list itself.

2.1 Solution [code/Reverse_Linked_List.cpp](#)

3. Divide-and-Conquer & Recursion

3.1 Multiply Matrices with Squaring Algorithm

- *From 4.2-6 of CLRS*

Suppose that you have a $\Theta(n^\alpha)$ -time algorithm for squaring $n \times n$ matrices, where $\alpha \geq 2$ is a constant. Show how

to use that algorithm to multiply two different $n \times n$ matrices in $\Theta(n^\alpha)$ time.

3.1 Solution

假设 A, B 为两个 $n \times n$ 的矩阵, 我们需要得到 AB

$$\text{假设 } C = \begin{bmatrix} O & A \\ B & O \end{bmatrix} \quad (1)$$

则 C 为一个 $2n \times 2n$ 的矩阵, 且:

$$C^2 = \begin{bmatrix} AB & O \\ O & BA \end{bmatrix} \quad (2)$$

由于 我们对于 $n \times n$ 矩阵有 $\Theta(n^\alpha)$ 时间复杂度的平方算法 ($\alpha \geq 2$ 为一个常数),

那么 得到 C^2 的时间复杂度为 $\Theta((2n)^\alpha) = \Theta(2^\alpha \times n^\alpha) = \Theta(n^\alpha)$

又因为 AB 在 C^2 的左上角, 可以由 C^2 直接得到, 故而 $A \times B$ 的时间复杂度为 $\Theta(n^\alpha)$.

3.2 Sort the Pancakes

- From 1.9 of Algorithms by Jeff Erickson

Suppose you are given a stack of n pancakes of different sizes. You want to sort the pancakes so that smaller pancakes are on top of larger pancakes. The only operation you can perform is a *flip*—insert a spatula under the top k pancakes, for some integer k between 1 and n , and flip them all over.

1. Describe an algorithm to sort an arbitrary stack of n pancakes using $O(n)$ flips. Exactly how many flips does your algorithm perform in the worst case?^[1]
2. Now suppose one side of each pancake is burned. Describe an algorithm to sort an arbitrary stack of n pancakes, so that the burned side of every pancake is facing down, using $O(n)$ flips. Exactly how many flips does your algorithm perform in the worst case?

3.2.1. Solution [code/Sort_Pancakes.cpp](#)

// 核心代码:

```
int SortPancakes() {
    for (int _size = pancakes.size(); _size > 0; _size--) {
        if (pancakes[_size - 1] == _size) continue; // 已经在所在位置了
        flip(search(_size)); // 翻转到头
        flip(_size - 1); // 翻转到尾
    }
    return cnt;
}
```

最坏的情况是每一次都要翻转到头 再翻转到尾巴, 也就是要翻转 $2 \times n = 2n$ 次
而每一次翻转 复杂度为 $_{size}$, 故而总时间复杂度为 $O(n^2)$

3.2.2. Solution [code/Sort_Pancakes2.cpp](#)

// 核心代码:

```
int SortPancakes() {
    for (int _size = pancakes.size(); _size > 1; _size--) {
        if (pancakes[_size - 1].size == _size && pancakes[_size - 1].is_facedown == true) {
            continue; // 已经在所在位置了 且 头朝下
        }
        flip(search(_size)); // 翻转到头

        // 翻转到尾 check一下是否是 burned朝下,如果朝下 翻转至上
        if (pancakes[0].is_facedown) {
            cnt++;
            pancakes[0].flip();
        }
        flip(_size - 1);
    }
    if (!pancakes[0].is_facedown) {
        cnt++;
        pancakes[0].flip();
    }
    return cnt;
}
```

最坏的情况是每一个煎饼都要翻转到头,然后翻面,再翻转到尾巴, 也就是要翻转 $3 \times n = 3n$ 次
而每一次翻转 复杂度为 $_{size}$, 故而总时间复杂度为 $O(n^2)$

3.3 Calculate the Inversions

- From 1.13 of *Algorithms* by Jeff Erickson

An **inversion** in an array $A[1\dots n]$ is a pair of indices (i, j) such that $i < j$ and $A[i] > A[j]$.

The number of inversions in an n -element array is between 0 (if the array is sorted) and $\binom{n}{2}$ (if the array is sorted backward).

Describe and analyze an algorithm to count the number of inversions in an n -element array in $O(n \log n)$ time. [Hint: Modify mergesort.]

Solution [code/count_ReversedPairs.cpp](#)

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1. The exact worst-case optimal number of flips required to sort pancakes (either burned or unburned) is an long-standing open problem; just do the best you can. ↩