CS 6410: Compilers

Fall 2019

HW 2 - Context Free Grammars and LR Parsing

Assigned: Tuesday, October 1, 2018, Due: Saturday, October 19, 2019
Instructor: Tamara Bonaci
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Submission Guidelines

- Please push your homework as a single .pdf file through Canvas.
- You do not have to type in your submission hand-written and then scanned, or photographed documents
 are fine, as long as the total size of your document is not too big, and your document is readable.
- This assignment is meant to be worked on individually, and you should submit it by 11:59pm on Saturday, October 19, 2019.

Problem 1 (Problem modified from the Aho, Sethi, Ullman book)

Please consider the following grammar:

```
S ::= cSdS | dScS | \epsilon
```

- 1. Show that this grammar is ambiguous by constructing two leftmost derivations for the sentence cdcd.
- 2. Show that this grammar is ambiguous by constructing **two rightmost derivations** for the sentence cdcd.
- 3. Construct the corresponding parse trees for your derivations.

Problem 2 (Problem modified from the Aho, Sethi, Ullman book)

Consider the following grammar:

```
S ::= ( L ) | x
L ::= L , S | S
```

- 1. Give a left-most derivation of (x, (x, x)).
- 2. Give a right-most derivation of (x, (x, x)).
- 3. Show the steps that a shift-reduce parser goes through when it parses (x, x, x). That is, show the contents of the stack, and the remaining input at each step.
- 4. Suppose that we replace the left-recursive production L::=L,S with a right-recursive one L::=S,L. What general effect does this have on the depth of the stack during a shift-reduce parse? (Hint: you might want to work through the parse of (x, x, x) again to see what changes.)

Problem 3 (Cooper and Torczon, Problem 3.10)

Consider the following grammar:

```
Start ::= S
S ::= A a
A ::= B C | B C f
B ::= b
C ::= c
```

Is the given grammar an LR(1) grammar? Please show your work.

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Problem 4 (Cooper and Torczon, Problem 3.11)

Consider a robotic arm that accepts two commands: ∇ puts an apple in the bag, and \triangle takes an apple out of the bag. Assume that the robotic arm starts with an empty bag.

A valid sequence for the robotic arm should have no prefix that contains more \triangle commands than it contains ∇ commands. For example, $\nabla\nabla\triangle\triangle$ and $\nabla\triangle\nabla$ are valid commands, but $\triangle\triangle\nabla\nabla$ and $\triangle\nabla\triangle$ are not.

- 1. Write a grammar that represents all the valid command sequences for the robotic arm.
- 2. Prove that the grammar it LR(1).