

# CS 5010: Programming Design Paradigms

## Fall 2019

### Lecture 7: Data Structures and Algorithms

Acknowledgement: lecture notes inspired by course material prepared by UW faculty members Z. Fung and H. Perkins.

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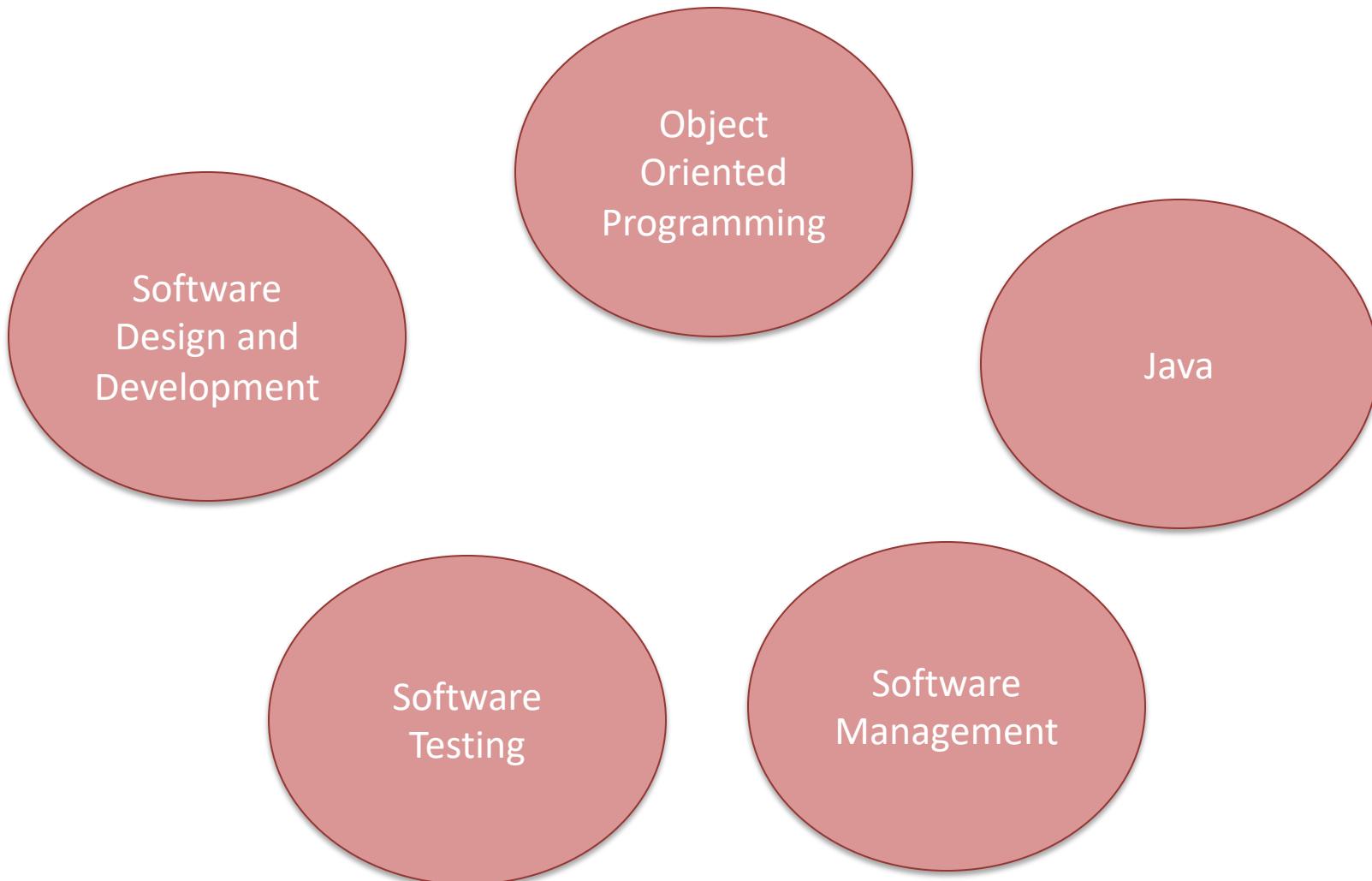
# Administrivia

- Assignment 4 due on Monday, October 21 by 6pm
- Codewalk 4 next week on Tue and Wed, Oct 22 and 23
- Lab 3 – nothing to submit, but please do it before A4 deadline
- (new work process --> we're not using Bottlenose anymore)

# Agenda – Algorithms and Data Structures 1

- **Data collections**
  - Iterating over data collections
    - Interface Iterable
    - Iterator
  - Ordering of objects
    - Interface Comparable
    - Interface Comparator
- **List ADT**
  - Doubly-linked List
  - Inner and nested classes
  - Algorithm: Recursion
  - Recursive data collections
- **Stack ADT**
- **Queue ADT**

# Course so Far...



# Objects and Classes

- **Object** – an entity consisting of states and behavior
  - States stored in variables/fields
  - Behavior represented through methods
- **Class** – template/blueprint describing the states and the behavior that an object of that type supports

# Interface vs. Abstract Class vs. Concrete Class?



[Pictures credit:  
<http://www.shakespearesglobal.com/>  
<https://www.vexels.com/>  
<https://www.amazon.com/Concrete-Leonard-Koren/dp/0714863548>]

# Abstract Class vs. Abstract Data Type?



Vs.



[Pictures credit: <https://www.vexels.com/>, <https://getyouralarm.com>]

- **Abstract class** – contains one or more abstract methods
- **ADT** – a high-level model of a data type, where the data type is defined from a point of view of the user (focus on operations (behavior), not on implementation)

# Review: Abstract Data Type

- **Abstract Data Type (ADT)** - model that describes data by specifying the operations that we can perform on them
- **Clients** care about the ADT → we need to capture the clients expectations in terms of the operations on the ADT
- For each operation, we need to describe:
  - The expected inputs, and any conditions that need to hold for our inputs and/or our ADT
  - The expected outputs and any conditions that need to hold for our output and/or our ADT
  - Invariants about our ADT

# Object Oriented Design Principles

- Encapsulation
- Abstraction
- Information hiding
- Polymorphism
- Inheritance

# Review: Polymorphism

**Polymorphism** – the ability of one instance to be viewed/used as different types (the ability to take many shapes/forms/views)

# Three Types of Polymorphism

Polymorphism – the ability to define different classes and methods as having the same name but taking different data types

Three types of polymorphism:

1. Subtype polymorphism
2. Ad hoc polymorphism
3. Parametric polymorphism

# Review: Subtype Polymorphism

- **Subtypes are substitutable for supertypes**
  - Instance of subtypes won't surprise a client by failing to satisfy the supertype's specification
  - Instance of subtype won't surprise a client by having more expectations than the supertypes' specifications

# Review: Compile Time and Run Time Types

```
Person emily = new Person();
```

```
Singer adele = new Singer();
```

```
Person flora = new Singer();
```

- **Static (compile time) type** – the declared type of a reference variable. Used by a compiler to check syntax
- **Dynamic (run time) type** – the type of an object that the reference variable currently refers to (it can change as the program execution progresses)

## Review: Ad Hoc Polymorphism

- Overloading allows us to create methods that share the same method name but differ in their signature
- Ad hoc polymorphism – another name for function and operator overloading
- Ad hoc polymorphism – a type of polymorphism where a polymorphic functions can be applied to arguments of different types
  - Polymorphic (overloaded) function can denote a number of distinct and potentially heterogeneous implementations, depending on the type of argument(s) to which it is applied

# Parametric Polymorphism

- Parametric polymorphism - ability for a function or type to be written in such a way that it handles values identically without depending on knowledge of their types
  - Such a function or type is called a generic function or data type
- Motivation - parametric polymorphism allows us to write flexible, general code without sacrificing type safety
  - Most commonly used in Java with collections

Algorithms and Data Structures 1

**ASIDE: BASIC JAVA I/O**

# I/O Streams

- Concept of an I/O stream – a **communication channel** (“**pipe**”) between a **source** and a **destination** that allows us to create a flow of data
- I/O Stream has:
  - An input source (a file, another program, device)
  - An output destination (file, another program, device)
  - The kind of data streamed (bytes, ints, objects)

# How to Do I/O

```
import java.io.*;
```

- Open the stream
- Use the stream (read, write, or both)
- Close the stream

open  
use  
close

# Opening a Stream

- **Problem:**
  - There exists some external data that we want to get, or
  - We want to put data somewhere outside your program
- **Solution:** open a stream
  - When we open a stream, we are making a connection to that external place
  - Once the connection is made, we can forget about the external place, and just use the stream

open  
use  
close

## Example of Opening a Stream

- `FileReader` - used to connect to a file that will be used for input:

```
FileReader fileReader = new FileReader(fileName);
```

- `Filename` - specifies where to find the (external) file  
(Note: after instantiating `FileReader` object, we never use `fileName` again; instead, we use `fileReader` object)

open  
use  
close

# Using a Stream

- Using a stream means doing input from it, or output to it
- Some streams can be used only for input, others only for output, still others for both
- But it is not usually that simple - we need to manipulate the data in some way as it comes in, or goes out

open  
use  
close

## Example of Using a Stream

```
int ch;  
ch = fileReader.read();
```

- The `fileReader.read()` method reads one character, and returns it as an integer, or -1 if there are no more characters to read
- The meaning of the integer depends on the file encoding (ASCII, Unicode, other)

open  
use  
close

# Manipulating the Input Data

- Reading characters as integers is not usually what we want to do
- A BufferedReader will convert integers to characters; it can also read whole lines
- The constructor for BufferedReader takes a FileReader parameter:

```
BufferedReader bufferedReader = new  
BufferedReader(fileReader);
```

open  
use  
close

# Reading Lines

```
String s;  
s = bufferedReader.readLine();
```

- A BufferedReader will return null if there is nothing more to read

# BufferedReader

- Reads text from a character-input stream, buffering characters to provide efficient reading of characters, arrays, and lines
- Buffer size may be specified, or the default size may be used (the default is large enough for most purposes)
- Why use BufferedReader?
  - In general, each read request causes a corresponding read request to be made of the underlying character or byte stream
  - These operations may be costly → each invocation of `read()` or `readLine()` in a `FileReader` or `InputStreamReader` causes bytes to be read from the file, converted into characters, and then returned
  - BufferedReader increases efficiency by buffering the input from the specified file

open  
use  
close

# Closing a Stream

- A stream is an **expensive resource!**
- There is a limit on the number of streams that you can have open at one time
- You should not have more than one stream open on the same file
- You should close a stream before you can open it again
- Always close your streams!

Algorithms and Data Structures 1

# DATA COLLECTIONS

# Data Collections?

Collection of chewed gums



Collection of pens



Collection of cassette tapes



Collection of old radios



## What is a data collection?

Shoes collection



Star wars collection



Cars collection



[Pictures credit: <http://www.smosh.com/smash-pit/articles/19-epic-collections-strange-things> ]

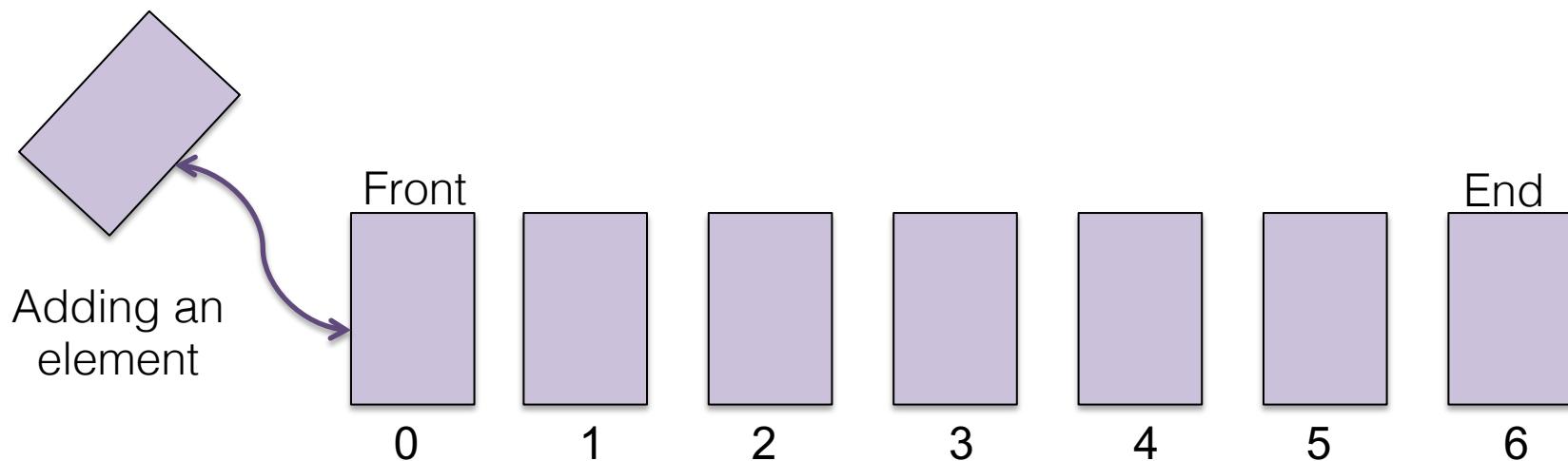
# Data Collections?

- **Data collection** - an object used to store data (think *data structures*)
  - Stored objects called **elements**
  - Some typically operations:
    - `add()`
    - `remove()`
    - `clear()`
    - `size()`
    - `contains()`
- Some examples: `ArrayList`, `LinkedList`, `Stack`, `Queue`, `Maps`, `Sets`, `Trees`

## Why do we need different data collections?

# Example: ArrayList vs. LinkedList

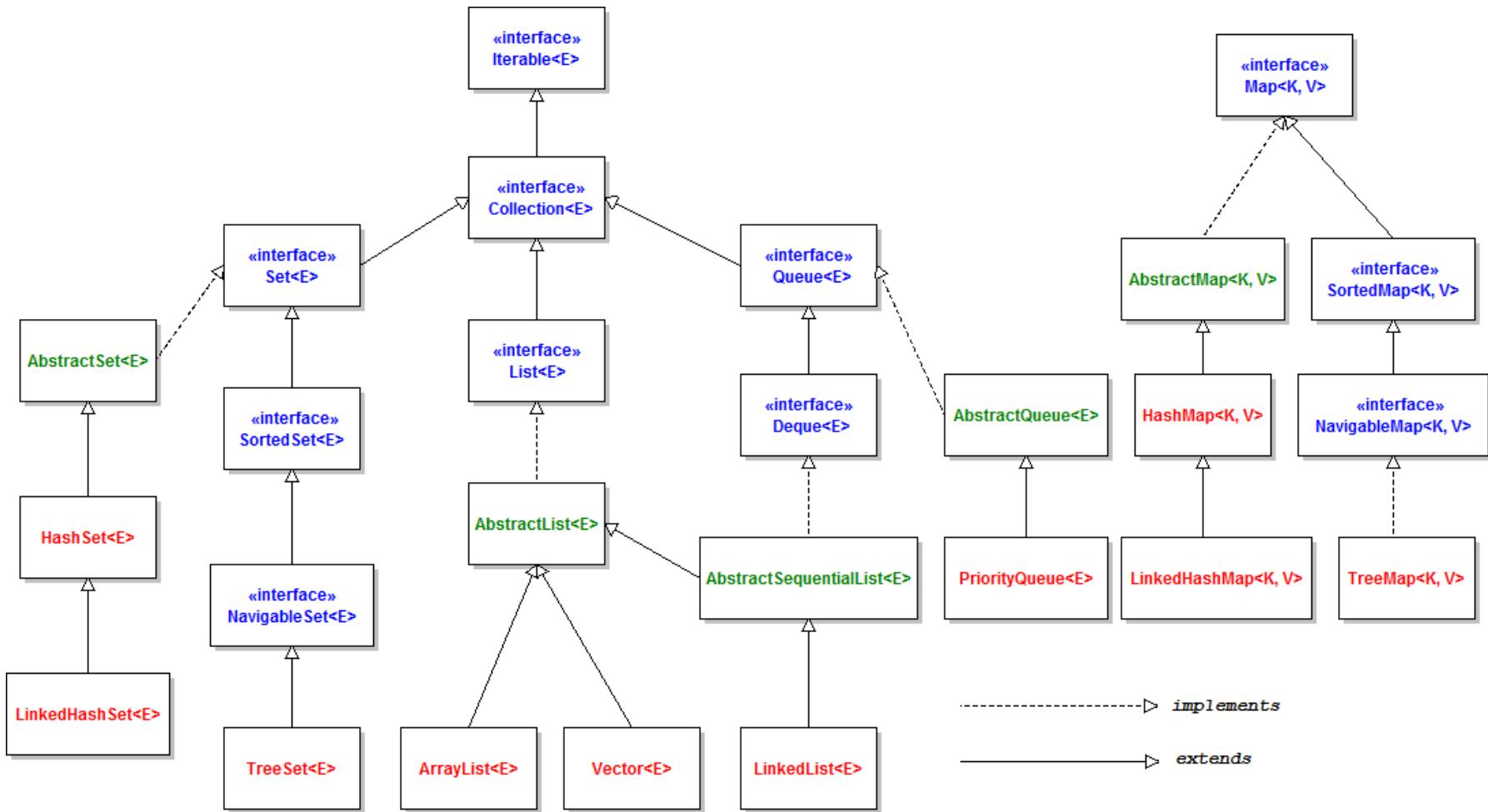
- List - a collection of elements with 0-based **indexes**
  - Elements can be added to the front, back, or in the middle
  - Can be implemented as an [ArrayList](#) or as a [LinkedList](#)
  - What is the complexity of adding an element to the front of an:
    - [ArrayList](#)?
    - [LinkedList](#)?



Algorithms and Data Structures 1

# JAVA COLLECTIONS FRAMEWORK

# Java Collections API



[Pictures credit: <http://www.codejava.net>]

# Java Collections Framework

- Under the Java Collections Framework, we have:
  - Interfaces that define the behavior of various data collections
  - Abstract classes that implement the interface(s) of the collection framework, that can then be extended to create a specialized data collection
  - Concrete classes that provide a general-purpose implementation of the interface(s)

# Java Collections Framework

- Goals of the Java Collections Framework:
  1. Reduce programming effort by providing the most common data structures
  2. Provide a set of types that are easy to use and extend
  3. Provide flexibility through defining a standard set of interfaces for collections to implement
  4. Improve program quality through the use and reuse of tested implementations of common data structures

# Java Collections Framework

- Part of the `java.util` package
- Interface `Collection<E>`:
  - Root interface in the collection hierarchy
  - Extended by four interfaces:
    - `List<E>`
    - `Set<E>`
    - `Queue<E>`
    - `Map<K, V>`
  - Extends interface `Iterable<T>`

# Interfaces Iterable<T>

- Super-interface for interface Collection<T>
- Implementing interface Iterable<T> allows an object to be traversed using the **for each loop**
- Every object that implements Iterable<T> must provide a method **Iterator iterator()**

Modifier and Type	Method and Description
default void	<b>forEach(Consumer&lt;? super T&gt; action)</b> Performs the given action for each element of the Iterable until all elements have been processed or the action throws an exception.
<b>Iterator&lt;T&gt;</b>	<b>iterator()</b> Returns an iterator over elements of type T.
default <b>Spliterator&lt;T&gt;</b>	<b>spliterator()</b> Creates a <b>Spliterator</b> over the elements described by this Iterable.

# Interfaces Iterable<T> and forEach Loop

- Super-interface for interface Collection<T>
- Implementing interface Iterable<T> allows an object to be traversed using the **for each loop**
- Every object that implements Iterable<T> must provide a method **Iterator iterator()**
- **ForEach loop:**  
default void forEach(Consumer<? super T> action)
- Performs the action for each element of the interface Iterable, until:
  - All elements have been processed or
  - The action throws an exception
- Actions are performed in the order of iteration (unless otherwise specified by the implementing class)
- Exceptions thrown by the action are relayed to the caller

## Interfaces Iterable<T> and Iterator<E>

- Super-interface for interface Collection<T>
- Implementing interface Iterable<T> allows an object to be traversed using the **for each loop**
- Every object that implements Iterable<T> must provide a method **Iterator iterator()**
- Interface Iterator – an iterator over a collection

```
public interface Iterator<E> {  
    boolean hasNext();  
    E next();  
    void remove();  
}
```

- Iterator remove() method – removes the last item returned by method next()

# Direct Use of an Iterator<E>

- Interface Iterator – an iterator over a collection

```
public interface Iterator<E> {  
    boolean hasNext();  
    E next();  
    void remove();  
}
```
- Careful when an iterator is used directly (not via a for each loop) → if you make any structural changes to a collection being iterated (add, remove, clear), the iterator is no longer valid (ConcurrentModificationException thrown)

Algorithms and Data Structures 1

# NATURAL ORDERING OF OBJECTS

# Comparing Objects

- **Problem:** How do we compare `String` in some list of `Strings`?
  - Operators like `<` and `>` do not work with `String` objects
  - But we do think of `Strings` as having an alphabetical ordering
- **Natural ordering:** Rules governing the relative placement of all values of a given type
- **Comparison function:** Code that, when given two values A and B of a given type, decides their relative ordering:
  - `A < B`, `A == B`, `A > B`

# The compareTo method

- A standard way for a Java class to define a comparison function for its objects is to define a `compareTo` method
  - Example: in the `String` class, there is a method:  
`public int compareTo(String other)`
- A call of `A.compareTo(B)` will return:
  - a value  $< 0$  if A comes "before" B in the ordering,
  - a value  $> 0$  if A comes "after" B in the ordering,
  - 0 if A and B are considered "equal" in the ordering.

# Using compareTo

- compareTo can be used as a test in an if statement

```
String a = "alice";
String b = "bob";
if (a.compareTo(b) < 0) { // true
    ...
}
```

Primitives	Objects
if (a < b) { ... }	if (a.compareTo(b) < 0) { ... }
if (a <= b) { ... }	if (a.compareTo(b) <= 0) { ... }
if (a == b) { ... }	if (a.compareTo(b) == 0) { ... }
if (a != b) { ... }	if (a.compareTo(b) != 0) { ... }
if (a >= b) { ... }	if (a.compareTo(b) >= 0) { ... }
if (a > b) { ... }	if (a.compareTo(b) > 0) { ... }

# compareTo() and Java Collections

- We can use an array or list of Strings with Java's included binary search method because it calls `compareTo` internally

```
String[] a = {"al", "bob", "cari", "dan", "mike"};  
int index = Arrays.binarySearch(a, "dan"); // 3
```

- Java's TreeSet/Map use `compareTo` internally for ordering

```
Set<String> set = new TreeSet<String>();  
for (String s : a) {  
    set.add(s);  
}  
System.out.println(s);  
// [al, bob, cari, dan, mike]
```

# Ordering Our Own Types

- **Problem:** we cannot binary search or make a TreeSet/Map of **arbitrary types**, because Java doesn't know how to order the elements
- **Example:** the program compiles but crashes when we run it

```
Set<HtmlTag> tags = new TreeSet<HtmlTag>();  
tags.add(new HtmlTag("body", true));  
tags.add(new HtmlTag("b", false));  
...  
Exception in thread "main"
```

```
java.lang.ClassCastException  
    at  
java.util.TreeSet.add(TreeSet.java:238)
```

# Comparable Template

```
public class name implements Comparable<name> {  
    ...  
  
    public int compareTo(name other) {  
        ...  
    }  
}
```

# Comparable Example

```
public class Point implements Comparable<Point> {
    private int x;
    private int y;
    ...

    // sort by x and break ties by y
    public int compareTo(Point other) {
        if (x < other.x) {
            return -1;
        } else if (x > other.x) {
            return 1;
        } else if (y < other.y) {
            return -1;      // same x, smaller y
        } else if (y > other.y) {
            return 1;      // same x, larger y
        } else {
            return 0;      // same x and same y
        }
    }
}
```

# Interface Comparator

- **Problem:** we may want to be able to order instances of some classes by more than one property (for example, by first name and by last name). What can we do?
- **Solution:** use interface `Comparator<T>` - a comparison function, which imposes a total ordering on some collection of objects
  - Can be passed to a sort method, to allow precise control over the sort order
  - Can also be used to control the order of certain data structures (tree sets and tree maps)
  - Can also be used to provide an ordering for collections of objects that don't have a natural ordering

# Method `compare(T o1, T o2)`

- `int compare(T o1, T o2)` - compares its two arguments for order and returns:
  - A negative integer if the first argument is less than the second
  - Zero, if the first argument is equal to the second
  - A positive integer if the first argument is greater than the second
- The implementor must also ensure that the relation is **transitive**:  
 $((\text{compare}(x, y) > 0) \&\& (\text{compare}(y, z) > 0)) \text{ implies } \text{compare}(x, z) > 0$
- Not strictly required, but good rule to follow: **compare() method should be consistent with equals() method**:

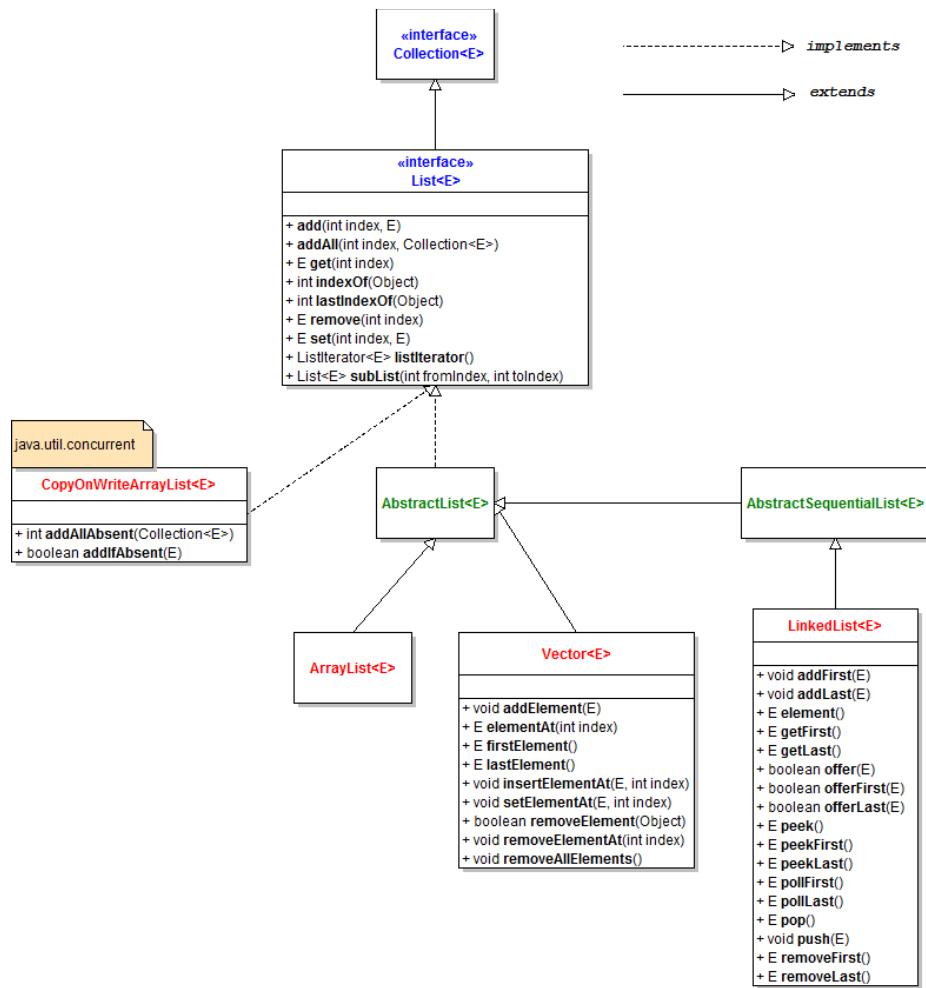
`(compare(x, y) == 0) == (x.equals(y))`

(Any comparator that violates this condition should clearly indicate this fact. The recommended language is "Note: this comparator imposes orderings that are inconsistent with equals.") **(don't do this!)**

Algorithms and Data Structures 1

**LIST ADT**

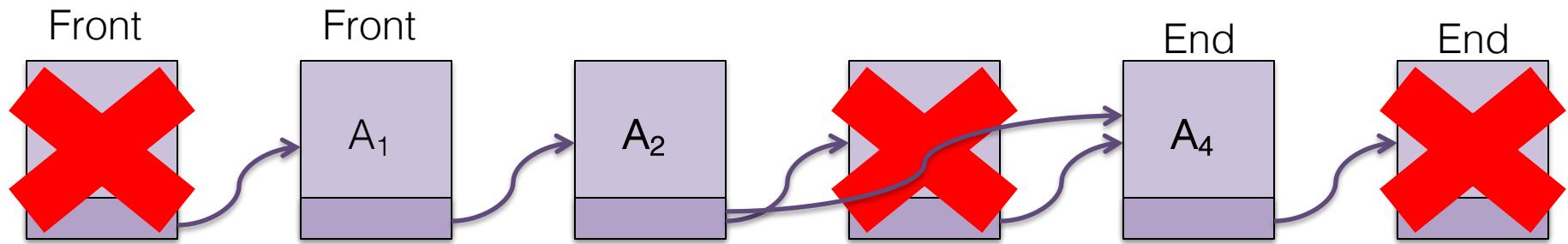
# Java List API



[Pictures credit: <http://www.codejava.net>]

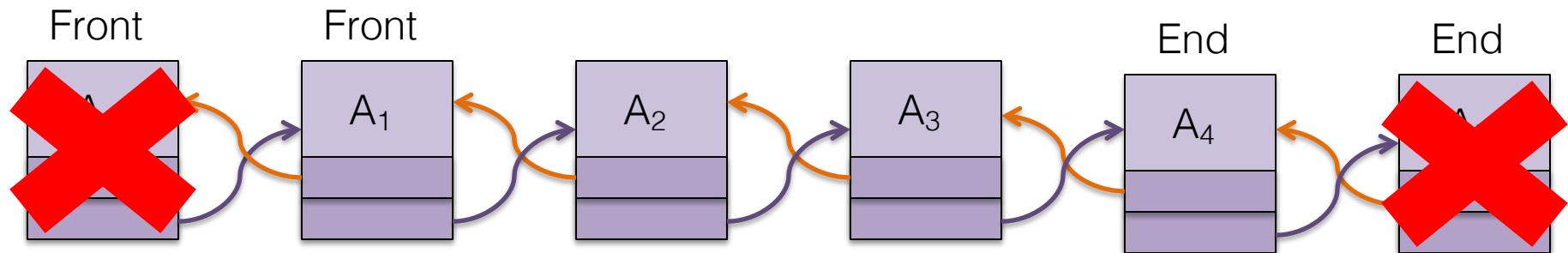
- `List<E>` - the base interface
- **Abstract subclasses:**
  - `AbstractList<E>`
  - `AbstractSequentialList<E>`
- **Concrete classes:**
  - `ArrayList<E>`
  - `LinkedList<E>`
  - `Vector<E>` (legacy collection)
  - `CopyOnWriteArrayList<E>` (class under `java.util.concurrent` package)
- **Main methods:**
  - `E get(int index);`
  - `E set(int index, E newValue);`
  - `Void add(int index, E x);`
  - `Void remove(int index);`
  - `ListIterator<E> listIterator();`

# Example: Removing Elements from a LinkedList



- Remove the first element of the list
- Remove element A<sub>3</sub> from the list
- Remove the last element in the list
- What's the tricky part about removing elements A<sub>3</sub> and A<sub>5</sub>?
- Having to find their predecessors (elements A<sub>2</sub> and A<sub>4</sub>) and updating their link to last node
- Idea: every node maintains the link to its previous and its next node → doubly linked list

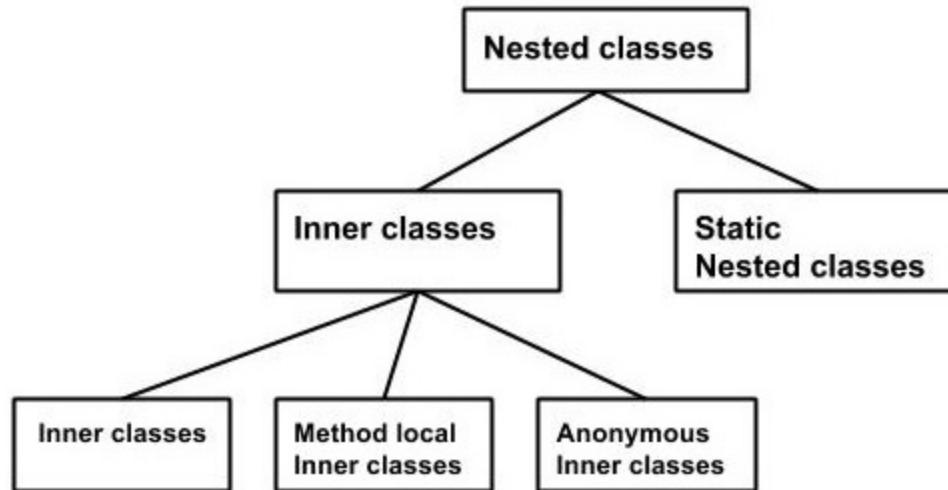
# Doubly Linked List



- Removing the first element of the list
- Removing the last element of the list

# Implementation of a Doubly LinkedList

- Review – inner and nested classes:



[Pictures credit: [https://www.tutorialspoint.com/java/images/inner\\_classes.jpg](https://www.tutorialspoint.com/java/images/inner_classes.jpg)]

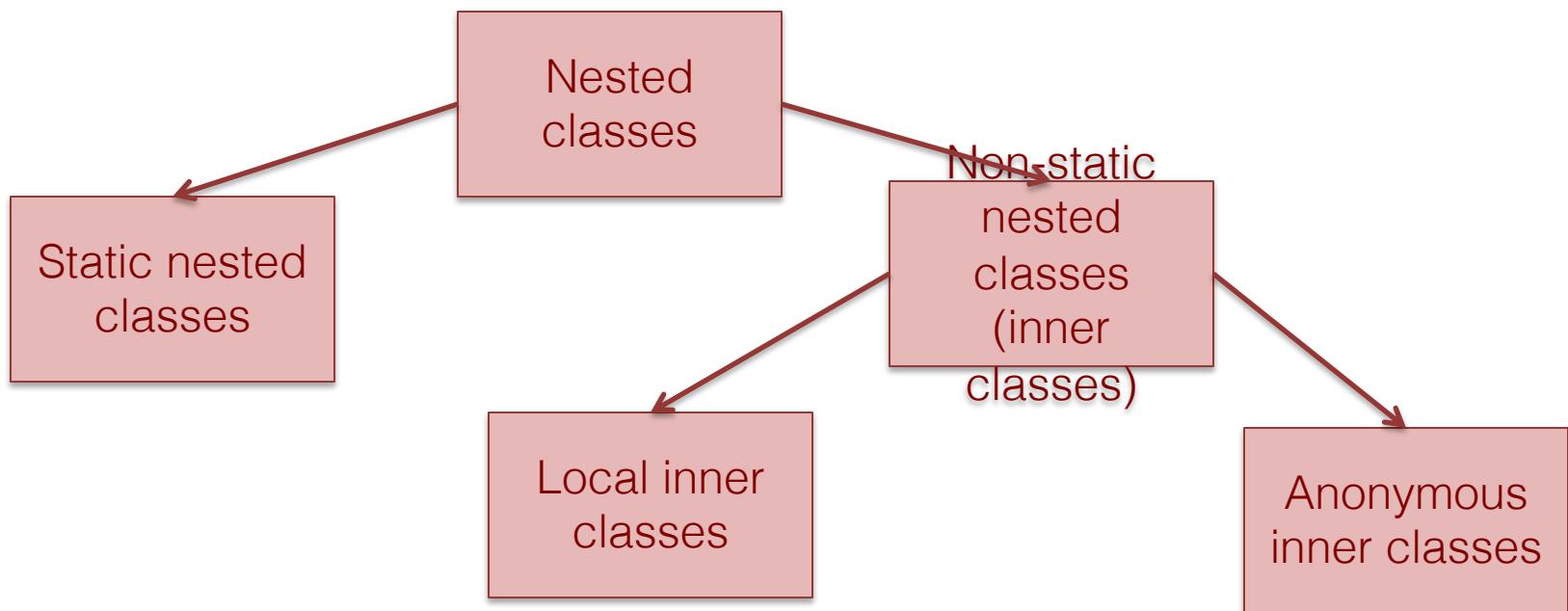
- Java allows us to define a class *within* another class – a **nested class**
  - Static nested classes
  - Non-static nested classes or **inner classes** – have access to other members of the enclosing class, even if they are declared private

# Nested Classes

- Nested class – a class defined within some other class
- Some nested class features:
  - The scope of a nested class is bounded by the scope of its outer class (i.e., the inner class does not exist independently of its outer class)
    - Class example: node and tree
  - A nested is a member of its outer class
    - Can be declared private, public, protected
    - Can have access to all members (including private) of its outer class
    - Outer class does not have access to the members of the nested class

# Nested Classes

- Nested class – a class defined within some other class
- Nested classes can be:
  - Static classes
  - Non-static (inner classes)



# Static Nested Classes

- Behave the same way as top-level classes

```
class A {  
    //code for A  
    static class B {  
        //code for B  
    }  
}
```

- To access a static nested class, we need to use the name of the outer class:

```
A.B b = new A.B()
```

# Inner Classes

- Object of inner classes exist within an instance of the outer class

```
class X {  
    //code for X  
  
    Class Y {  
        //code for Y  
    }  
}
```

- To access an inner nested class, we need to do so through an instance of the outer class:

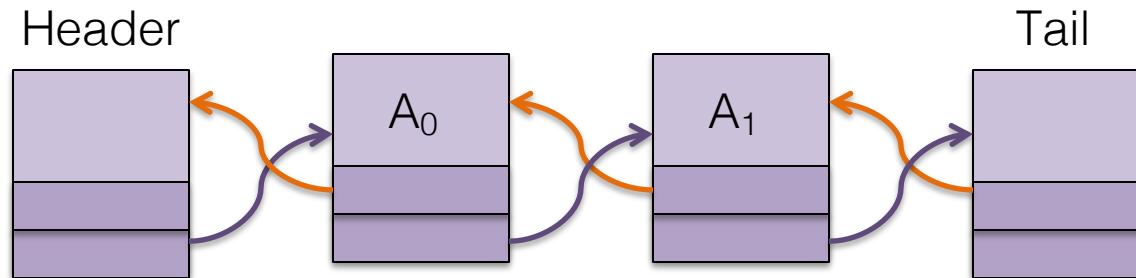
```
X x = new X();  
X.Y y = x.new Y();
```

# Difference Between Static and Inner Classes

- Static nested classes do not have a direct access to the non-static members of the outer class (non-static variables and methods)
  - Static class must access the non-static members of its enclosing class through an object
- Inner classes have access to all members (static and non-static, including private) of its outer class, and may refer to them directly
  - Used more frequently

# Implementation of a Doubly LinkedList

- **Doubly linked list** – need to provide and maintain links to both ends of the list
- Implemented classes:
  - Class MyLinkedList
  - Class Node – private nested class, contains the data and links to previous and next nodes
  - Class LinkedListIterator – private inner class, implementing the interface Iterator
- **Sentinel nodes:**
  - Header and tail nodes, used to logically represent the beginning and the end markers



Algorithms and Data Structures 1

# LIST ADT AND RECURSIVE DATA COLLECTIONS

# Recursion



Did you mean: recursion

Recursion - Wikipedia, the free encyclopedia

A visual form of **recursion** known as the Droste effect. The woman in this image contains a smaller image of herself holding the same ...  
[en.wikipedia.org/w/index.php?title=Recursion&oldid=9000000](https://en.wikipedia.org/w/index.php?title=Recursion&oldid=9000000) - Cached - Similar -

Recursion (computer science) - Wikipedia, the free encyclopedia

**Recursion** in computer science is a way of thinking about and solving

[Pictures credit: <http://www.telegraph.co.uk/technology/google/6201814/Google-easter-eggs-15-best-hidden-jokes.html>]

# Recursion

- **Recursion** – an operation defined in terms of itself
  - Solving a problem recursively means solving smaller occurrences of the same problem
- **Recursive programming** – an object consists of methods that call themselves to solve some problem
- Can you think of some examples of recursions and recursive programs?

# Recursive Algorithm

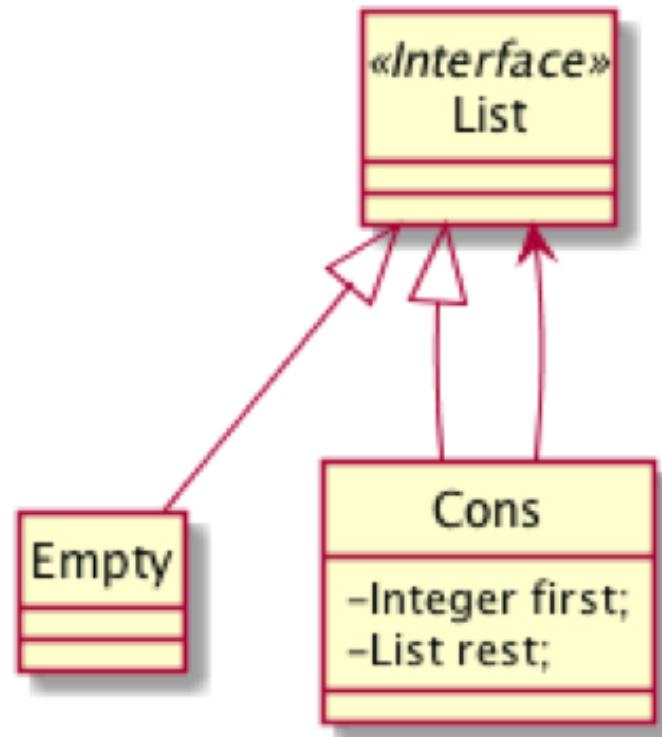
- Every recursive algorithm consists of:
  - **Base case** – at least one simple occurrence of the problem that can be answered directly
  - **Recursive case** - more complex occurrence that cannot be directly answered, but can be described in terms of smaller occurrences of the same problem
- A crucial part of recursive programming is identifying these cases
- What were base cases for our Fibonacci problem in Assignment 1?
- Why are we now talking about recursion???

# Recursive Data Structures

- **Recursive data structure** - a data structure partially composed of smaller or simpler instances of the same data structure
- Is linked list a recursive data structure?
- Let's see - a linked list is either
  - Null (base case)
  - A node whose next field references a list

# Lists as a Recursive Data Collection

- List – an ordered collection (also known as a sequence)
- A linked list is either:
  - Null (base case)
  - A node whose next field references a list

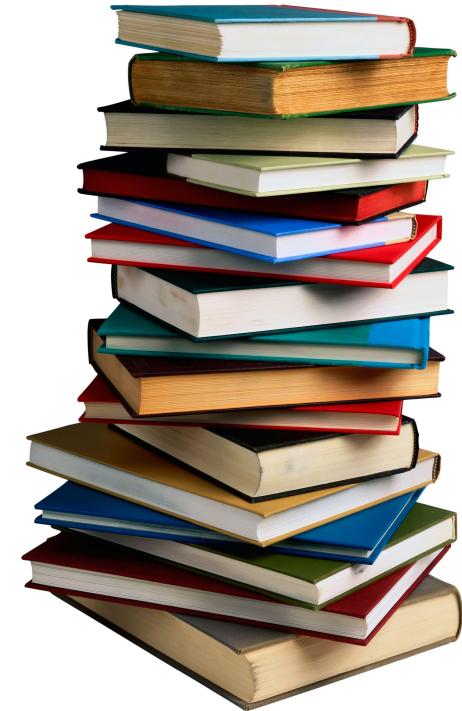


Algorithms and Data Structures 1

**STACK ADT**

# Stacks

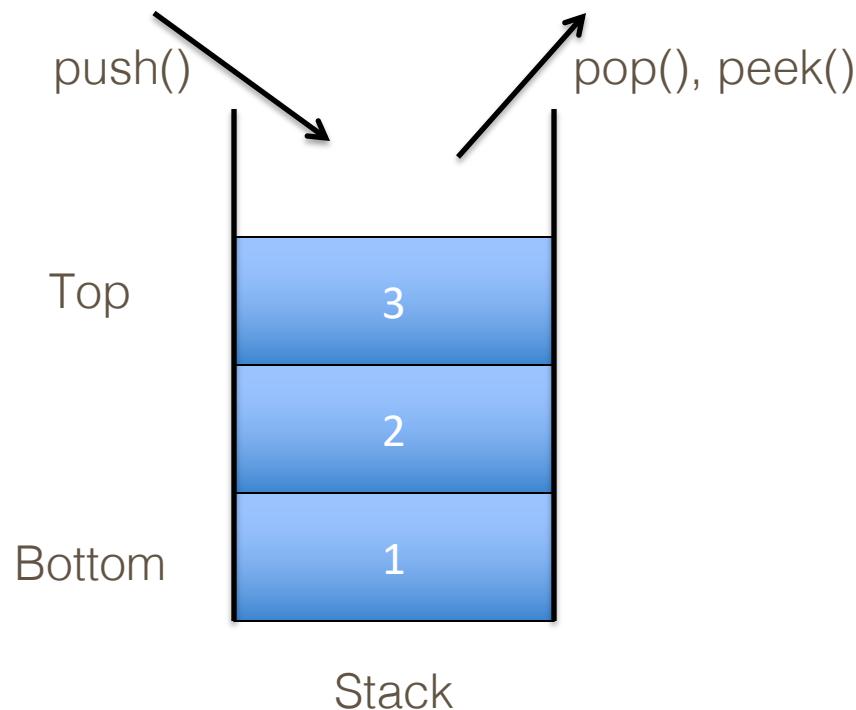
- Some of my favorite stacks:



[Pictures credit: <https://rukminim1.flixcart.com/image/1408/1408/stacking-toy>,  
<http://battellemmedia.com/wp-content/uploads/2014/08/National-Pancake-Day-at-IHOP.jpg>,  
[http://all4desktop.com/data\\_images/original/4245681-book.jpg](http://all4desktop.com/data_images/original/4245681-book.jpg)]

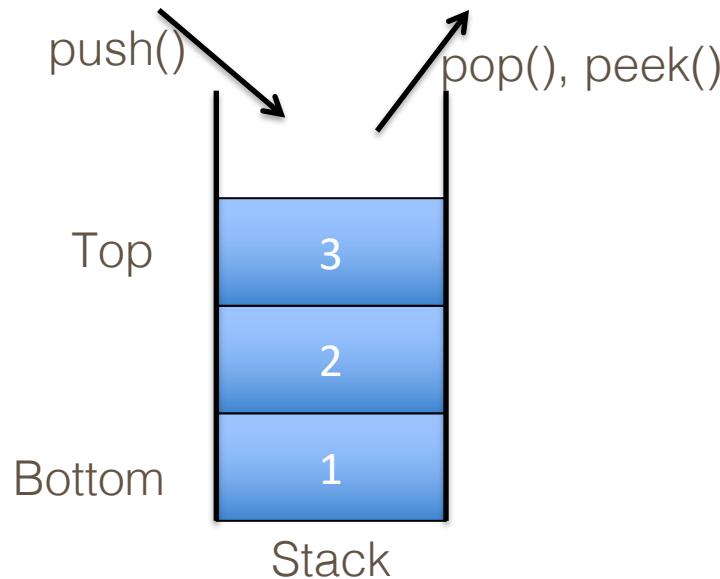
# What is a Stack?

- **Stack** – a data collection that retrieves elements in the LIFO order (last-in-first-out)



- Is there another way to think about a stack?

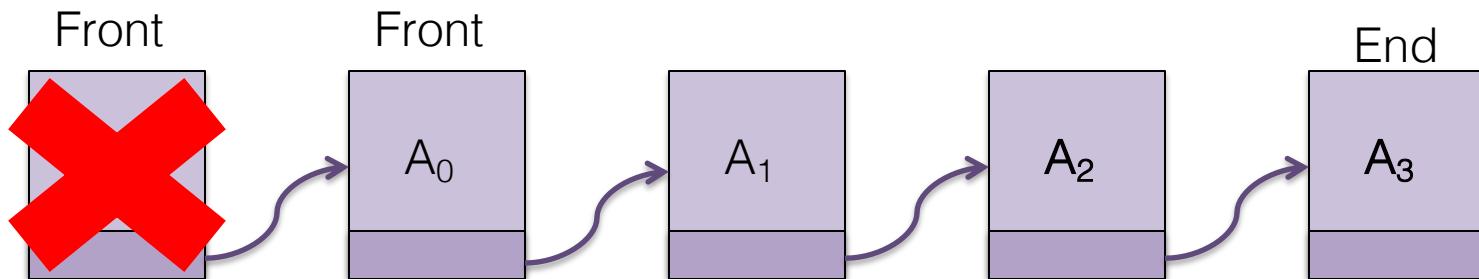
# What is a Stack?



- Is there another way to think about a stack?
- **Stack** – a constrained data collection where clients are limited to use only limited optimized methods (`pop`, `push`, `peek`)
- **Stack** – a list with restriction that insertions and deletions can be performed in only one position, the end of the list, called **the top**

# Implementations of a Stack

- Stack – a list with restriction that insertions and deletions can be performed in only one position, the end of the list, called the top
- Since a stack is a list, any list implementation will do:
  - Example: linked list implementation
    - push ()
    - peek () / top ()
    - pop ()



# Java Class Stack

Stack <E> ( )	Object constructor – constructs a new stack with elements of type E
push (value)	Places given value on top of the stack
pop ()	Removes top value from the stack, and returns it. Throws EmptyStackException if the stack is empty.
peek ()	Returns top value from the stack without removing it. Throws EmptyStackException if the stack is empty.
size ()	Returns the number of elements on the stack.
isEmpty ()	Returns true if the stack is empty.

Example:

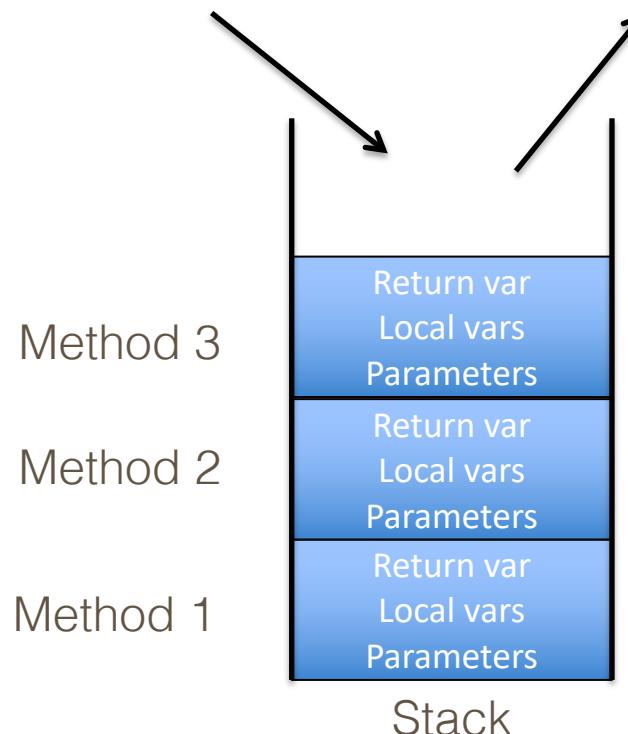
```
Stack<String> s = new Stack<String>();  
s.push("Hello");  
s.push("PDP");  
s.push("Fall 2019"); //bottom ["Hello", "PDP", "Fall 2019"] top  
System.out.println(s.pop()); //Fall 2019
```

# Applications of a Stack

- Programming languages and compilers:
  - Method calls (*call=push, return=pop*)
  - Compilers (parsers)
- Matching up related pairs of things:
  - Find out whether a string is a palindrome
  - Examine a file to see if its braces { } match
  - Convert "infix" expressions to pre/postfix
- Sophisticated algorithms:
  - Searching through a maze with "backtracking"

# Example: Methods Call

- In some system, whenever there is a method call, some information about the current state of the system needs to be stored before the control is transferred to a new method:
  - Parameters
  - Local variables
  - Return address



# Example: Postfix Expressions

- Suppose you are using a calculator to compute the total cost of your groceries
  - Add the costs of individual items
  - Multiply by 1.1 to account for local sales tax
- The natural way to do this with a calculator:  
$$5.5 + 4.5 + 7 + 8 * 1.1$$
- What is the expected result?
  - 27.5 (expected value)
  - 25.8
- That depends on how “smart” is your calculator!

## Example: Postfix Expressions

- The natural way to do this with a calculator:  
$$5.5 + 4.5 + 7 + 8 * 1.1$$
- What if we represent the given expression in the **postfix** or **Reverse Polish notation**:  
$$5.5 \ 1.1 * \ 4.5 \ 1.1 * + \ 7 \ 1.1 * + \ 8 \ 1.1 * +$$
- The easiest way to implement this is with a stack:  
$$(5.5 * 1.1) + (4.5 * 1.1) + (7 * 1.1) + (8 * 1.1)$$

Algorithms and Data Structures 1

# QUEUE ADT

# My Least Favorite Queues



[Pictures credit: <http://airport.blog.ajc.com>, <https://s1cdn.autoevolution.com/images/news/the-longest-traffic-jam-in-history-12-days-62-mile-long-47237-7.jpg>]

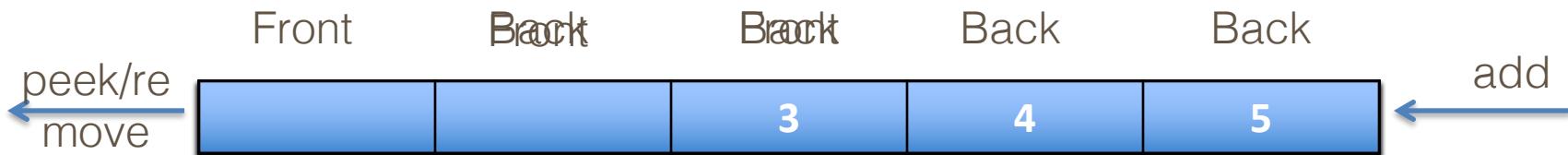
# What is a Queue?

- **Queue** – a data collection that retrieves elements in the FIFO order (first in, first out)
  - Elements are stored in order of insertion, but don't have indexes
    - Client can only:
      - Add to the end of the queue,
      - Examine/remove the front of the queue
- Basic queue operations:
  - **Add** (enqueue) - add an element to the back of the queue
  - **Peek** - examine the front element
  - **Remove** (dequeue) - remove the front element



# Implementations of Queues

- Like stack, queue can be seen as a list with restriction, and can be implemented as a list:
- Example: `ArrayList` implementation
  - Initially, queue only has elements 1 and 2
  - Add another element, 3, to the queue
  - Add another element, 4, to the queue
  - Add another element, 5, to the queue
  - Remove an element from the queue
  - Remove an element from the queue



What happens when we remove two more elements from the queue?

# Implementations of Queues

- What happens when we remove two more elements from the queue?



- Approach – circular array implementation - whenever front or back get to the end of the array, allow them to wrap around to the beginning
- Example:

- Add another element, 6, to the queue

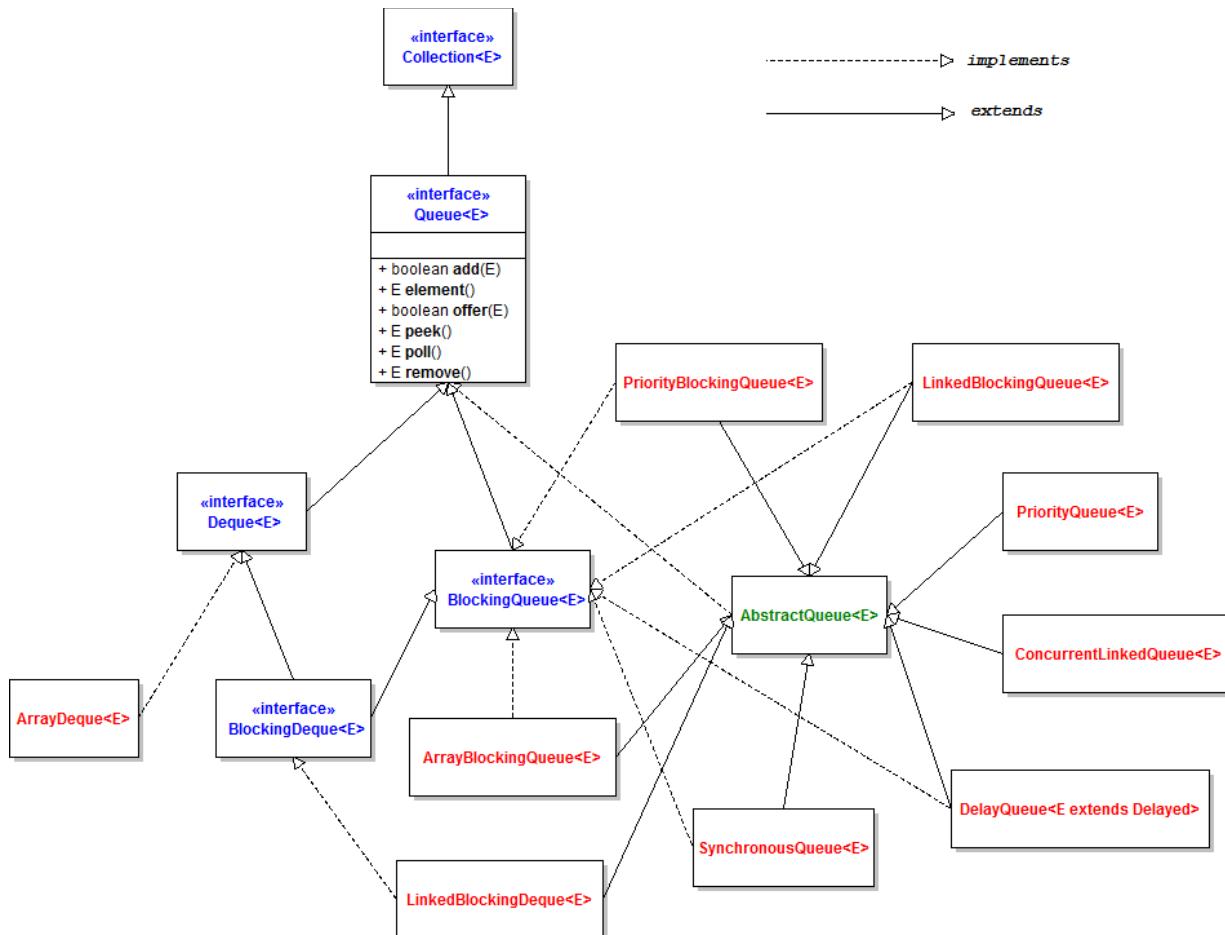


- What is the front and what is the back of the queue now?

# Applications of Queues

- Operating systems:
  - Queue of print jobs to send to the printer
  - Queue of programs / processes to be run
  - Queue of network data packets to send
- Programming:
  - Modeling a line of customers or clients
  - Storing a queue of computations to be performed in order
- Real world examples:
  - People waiting in some line
  - ???

# Class Diagram of the Queue API



[Pictures credit: <http://www.codejava.net/java-core/collections/class-diagram-of-queue-api>]

# Java Interface Queue

<code>add (value)</code>	places given value at back of queue
<code>remove ()</code>	removes value from front of queue and returns it; throws a <code>NoSuchElementException</code> if queue is empty
<code>peek ()</code>	returns front value from queue without removing it; r eturns <code>null</code> if queue is empty
<code>size ()</code>	returns number of elements in queue
<code>isEmpty ()</code>	returns <code>true</code> if queue has no elements

- Example:

```
Queue<Integer> myQueue = new LinkedList<Integer>();  
myQueue.add(10);  
myQueue.add(16);  
myQueue.add(2019); // front [10, 16, 2019] back  
System.out.println(myQueue.remove()); // 10
```

# Mixing Queues and Stacks

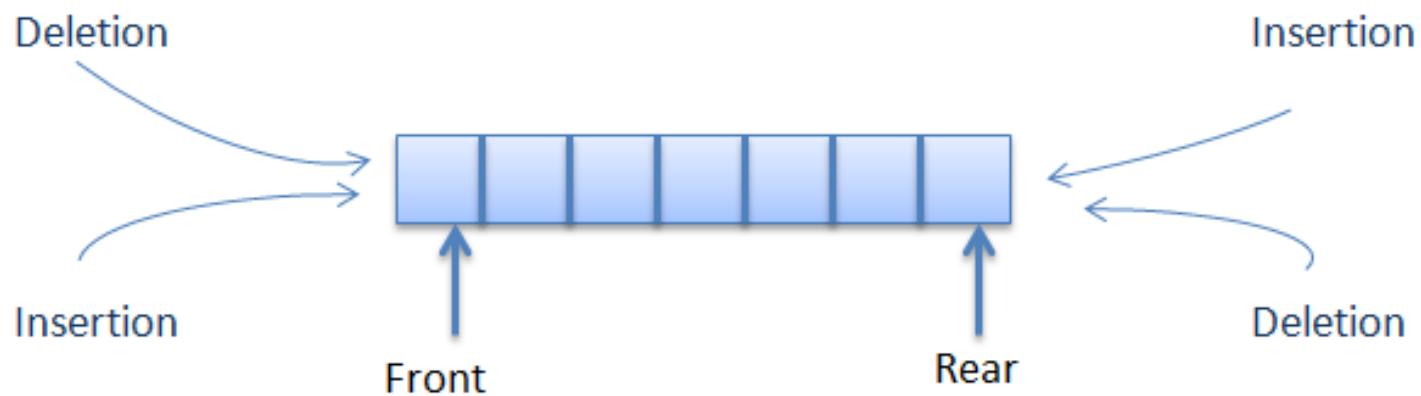
- We often mix stacks and queues to achieve certain effects
- Example: Reverse the order of the elements of a queue

```
Queue<Integer> q = new LinkedList<Integer>();  
q.add(1);  
q.add(2);  
q.add(3); // [1, 2, 3]  
Stack<Integer> s = new Stack<Integer>();  
while (!q.isEmpty()) {  
    s.push(q.remove()); } // Q -> S  
while (!s.isEmpty()) {  
    q.add(s.pop()); } // S -> Q  
System.out.println(q); // [3, 2, 1]
```

Algorithms and Data Structures 1

# DEQUE ADT

# Deque



[Pictures credit: <http://www.java2novice.com/data-structures-in-java/queue/double-ended-queue/>]

# Deque

	<b>First Element (Head)</b>	
	<i>Throws exception</i>	<i>Special value</i>
<b>Insert</b>	<code>addFirst(e)</code>	<code>offerFirst(e)</code>
<b>Remove</b>	<code>removeFirst()</code>	<code>pollFirst()</code>
<b>Examine</b>	<code>getFirst()</code>	<code>peekFirst()</code>

	<b>Last Element (Tail)</b>	
	<i>Throws exception</i>	<i>Special value</i>
<b>Insert</b>	<code>addLast(e)</code>	<code>offerLast(e)</code>
<b>Remove</b>	<code>removeLast()</code>	<code>pollLast()</code>
<b>Examine</b>	<code>getLast()</code>	<code>peekLast()</code>

[Pictures credit:<https://docs.oracle.com/javase/7/docs/api/java/util/Deque.html>

# Deque

Queue Method	Equivalent Deque Method
<code>add(e)</code>	<code>addLast(e)</code>
<code>offer(e)</code>	<code>offerLast(e)</code>
<code>remove()</code>	<code>removeFirst()</code>
<code>poll()</code>	<code>pollFirst()</code>
<code>element()</code>	<code>getFirst()</code>
<code>peek()</code>	<code>peekFirst()</code>

[Pictures credit:<https://docs.oracle.com/javase/7/docs/api/java/util/Deque.html>

# Deque

Stack Method	Equivalent Deque Method
<code>push(e)</code>	<code>addFirst(e)</code>
<code>pop()</code>	<code>removeFirst()</code>
<code>peek()</code>	<code>peekFirst()</code>

Algorithms and Data Structures 1

# SET ADT

# Sets

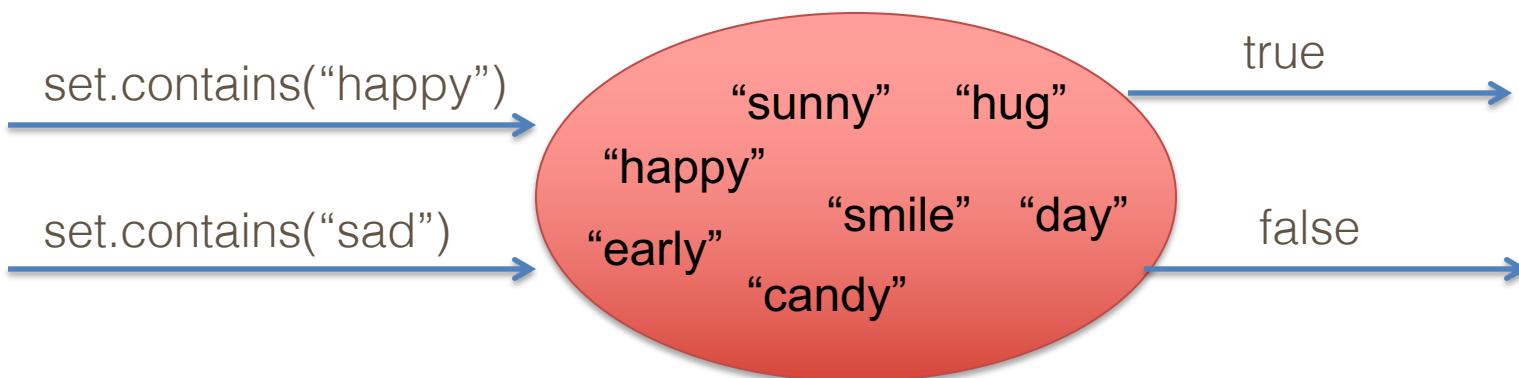


[Pictures credit:

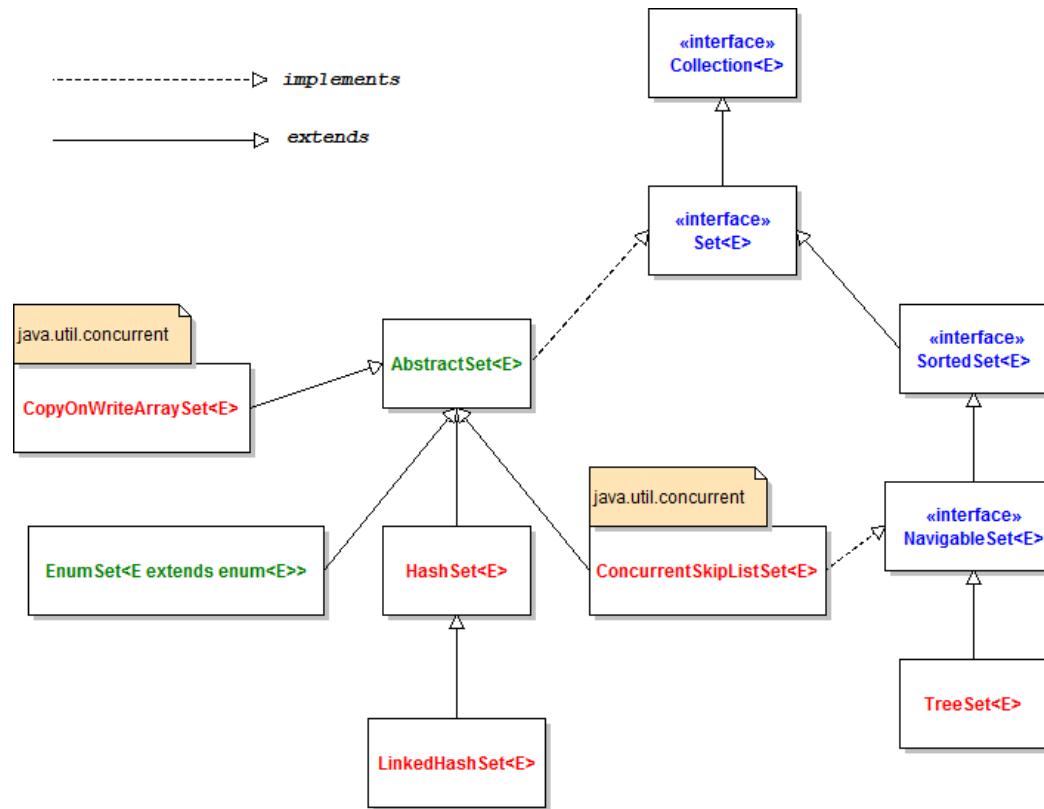
[https://www.beyondtheblackboard.com/components/com\\_virtuemart/shop\\_image/product/full/SET---Box---Transparent-Background---8-22-11\\_0.png](https://www.beyondtheblackboard.com/components/com_virtuemart/shop_image/product/full/SET---Box---Transparent-Background---8-22-11_0.png)]

# Sets

- **Set** - a collection of unique values (no duplicates allowed) that can perform the following operations efficiently:
  - add,
  - remove,
  - search (contains)
- We don't think of a set as having indexes; we just add things to the set in general and don't worry about order



# Set API Class Diagram



[Pictures credit:

<http://www.codejava.net/images/articles/javacore/collections/Set%20API%20class%20diagram.png>]

# Set Implementations

- In Java, sets are represented by Set type in `java.util`
- Set is implemented by `HashSet` and `TreeSet` classes
  - `HashSet`: implemented using a "hash table" array
    - Very fast:  $O(1)$  for all operations
    - Elements are stored in unpredictable order
  - `TreeSet`: implemented using a binary search tree
    - Pretty fast:  $O(\log N)$  for all operations
    - Elements are stored in sorted order
  - `LinkedHashSet`:
    - $O(1)$  but stores in order of insertion, but slightly slower than `HashSet` because of extra info stored

## Set Methods

- We can construct an empty set, or one based on a given collection
- Examples:

```
Set<Integer> set = new TreeSet<Integer>(); // empty
```

```
List<String> list = new ArrayList<String>();
```

```
...
```

```
Set<String> set2 = new HashSet<String>(list);
```

# Set Methods

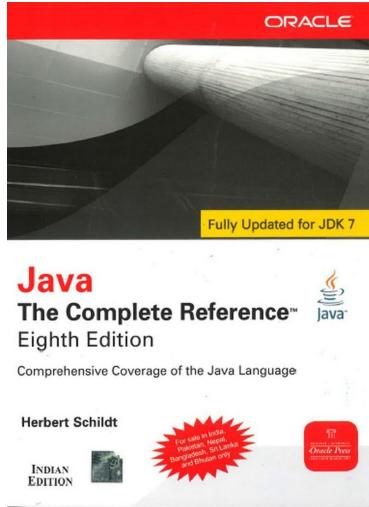
<b>add (value)</b>	adds the given value to the set
<b>contains (value)</b>	returns true if the given value is found in this set
<b>remove (value)</b>	removes the given value from the set
<b>clear ()</b>	removes all elements of the set
<b>size ()</b>	returns the number of elements in list
<b>isEmpty ()</b>	returns true if the set's size is 0
<b>toString ()</b>	returns a string such as "[3, 42, -7, 15]"

Algorithms and Data Structures 1

**MAP ADT**

# Maps

- Write a program to count the number of occurrences of every unique word in a large text file (e.g. *Java Reference Manual*)

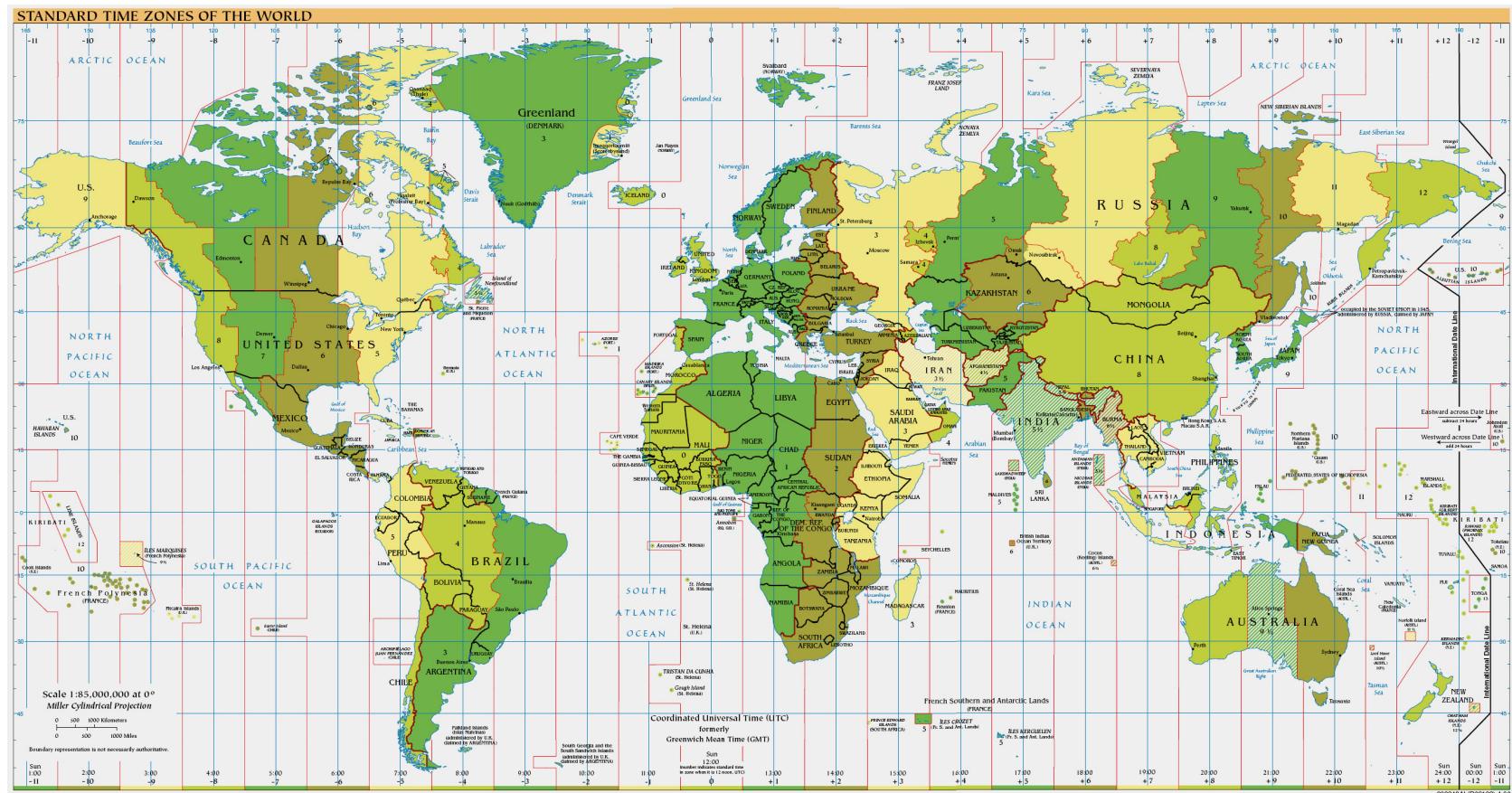


[Pictures credit: <https://images-na.ssl-images-amazon.com/images/I/61rKnDmww9L.jpg>]

# Maps

- Write a program that stores, modifies and retrieves:
  - Assignment grades for every student in this College
  - Financial information for every client of some bank
  - Browsing history for every user of some search engine
  - Searches and transactions for every user of some online retailer
  - Activity and likes of every user of some online platform
- Question: What do these records have in common?
- The way we think about them → every data sample has a unique user → unique ID (key)
- What is the appropriate data collection for this data?
- Maps

# Maps



# Maps

- **Map** – a data collection that holds a set of unique *keys* and a collection of *values*, where each key is associated with one value
- Also known as:
  - Dictionary
  - Associative array
  - Hash
- Basic map operations:
  - **put**(*key, value*) - adds a mapping from a key to a value
  - **get**(*key*) - retrieves the value mapped to the key
  - **remove**(*key*) - removes the given key and its mapped value

# Map Implementations

- In Java, maps are represented by `Map` type in `java.util`
- Map is implemented by the `HashMap` and `TreeMap` classes
  - `HashMap` - implemented using a "hash table"
    - Extremely fast:  $O(1)$
    - Keys are stored in unpredictable order
  - `TreeMap` - implemented as a linked "binary tree" structure
    - Very fast:  $O(\log N)$
    - Keys are stored in sorted order
  - `LinkedHashMap` -  $O(1)$ 
    - Keys are stored in order of insertion

# Map Implementations

- Map requires 2 types of parameters:
  - One for keys
  - One for values

- Example:

```
// maps from String keys to Integer values
Map<String, Integer> votes = new HashMap<String, Integer>();
```

# Map Methods

<code>put(key, value)</code>	adds a mapping from the given key to the given value; if the key already exists, replaces its value with the given one
<code>get(key)</code>	returns the value mapped to the given key (null if not found)
<code>containsKey(key)</code>	returns true if the map contains a mapping for the given key
<code>remove(key)</code>	removes any existing mapping for the given key
<code>clear()</code>	removes all key/value pairs from the map
<code>size()</code>	returns the number of key/value pairs in the map
<code>isEmpty()</code>	returns true if the map's size is 0
<code>toString()</code>	returns a string such as "{a=90, d=60, c=70}"

# keySet and Values

- keySet method returns a **Set** of all keys in the map
  - It can loop over the keys in a `foreach` loop
  - It can get each key's associated value by calling `get` on the map

Example:

```
Map<String, Integer> ages = new TreeMap<String, Integer>();  
ages.put("Marty", 19);  
ages.put("Geneva", 2); // ages.keySet() returns Set<String>  
ages.put("Vicki", 57);  
for (String name : ages.keySet()) {           // Geneva -> 2  
    int age = ages.get(name);                  // Marty -> 19  
    System.out.println(name + " -> " + age); // Vicki -> 57  
}
```

# Methods keySet and values

- values method returns a collection of all values in the map
  - It can loop over the values in a `foreach` loop
  - No easy way to get from a value to its associated key(s)

<code>keySet ()</code>	returns a set of all keys in the map
<code>values ()</code>	returns a collection of all values in the map
<code>putAll (<b>map</b>)</code>	adds all key/value pairs from the given map to this map
<code>equals (<b>map</b>)</code>	returns <code>true</code> if given map has the same mappings as this one

## Example: Maps

- Many words are similar to other words
- For example word wine can become:
  - dine, fine, line, mine, nine, pine, or vine
  - wide, wife, wipe, or wire
  - wind, wing, wink, or wins
- Write a program to find all words that can be changed into at least 15 other words by a single one-character substitution
- Assume that:
  - We have a dictionary consisting of approximately 89,000 different words of varying lengths
  - Most words are between 6 and 11 characters

Algorithms and Data Structures 1

# HASHING AND HASH FUNCTIONS

# Introduction to Hashing

- **Problem:** how much does a new phone cost?
  - e.g., Pixel 3??
  - e.g., iPhone XR??
- **Ways to answer this question:**
  - Look it up in an unsorted list of all phone prices:  $O(n)$
  - Look it up in a sorted list of all phone prices:  $O(\log n)$
  - Ask your buddy if she remembers:  $O(1)$

# Introduction to Hashing

- Does this difference in time complexity matter?

Num Items	Simple Search $O(n)^2$	Binary Search $O(\log n)$	Buddy $O(1)$
100	10 sec	1 sec	Instant
1000	1.6 min	1 sec	Instant
10000	16.6 min	2 sec	Instant

It sure does!

How do we replicate our buddy with data structures???

# Introduction to Hashing

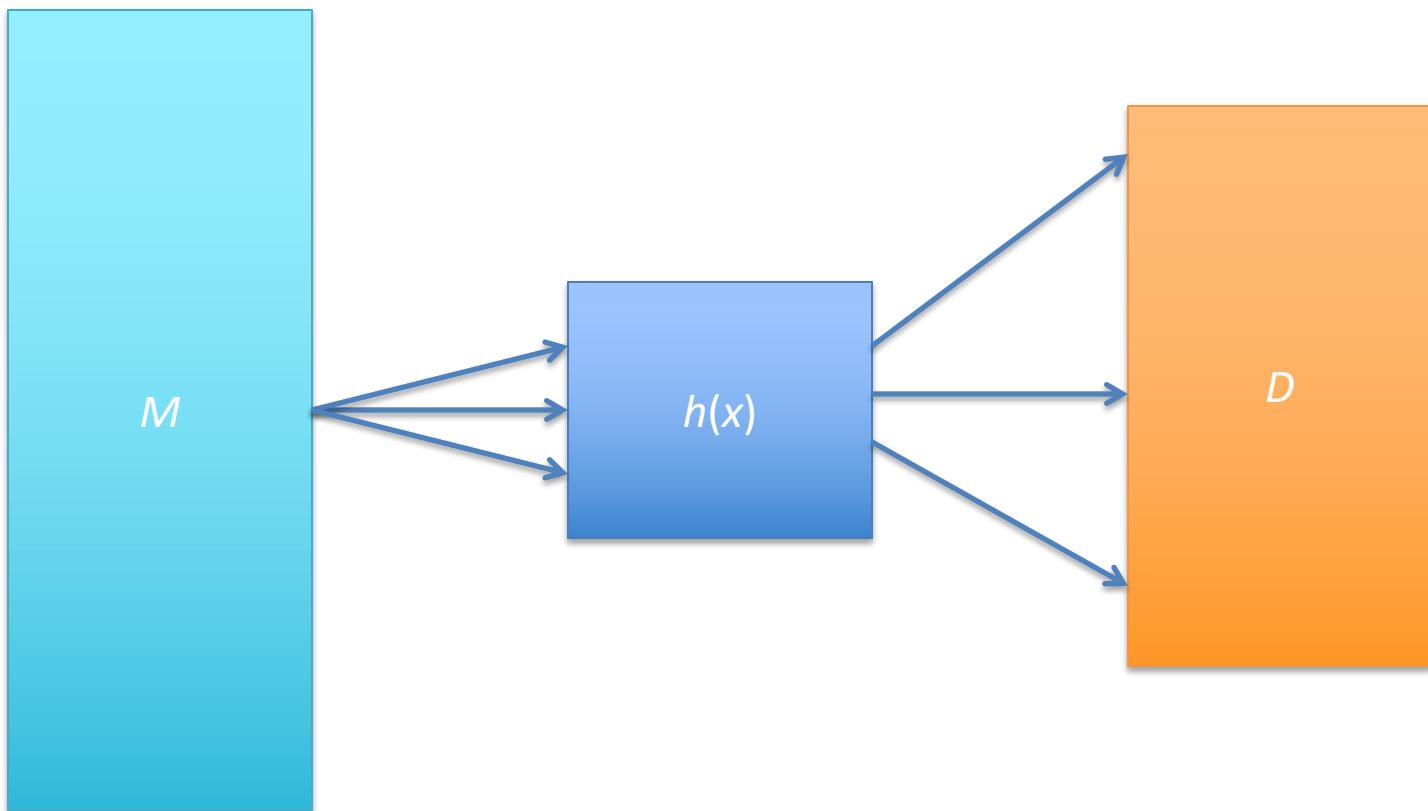
- How do we replicate our buddy with data structures???
- Observation - each item is essentially 2 items:
  - Product name, and
  - Price
- If we sort by name, we can find a product in  $O(\log n)$  time
  - How can we get down to  $O(1)$ ?

# Introduction to Hash Functions

- Use a **hash function** to transform a {string, number, struct, object, . . . } into a number
- **Be consistent** - every time you give it the same input, it returns the same value
  - Every time you give it “iPhone XR”, it returns ‘xxx’
- Make sure that the **output is distinguishing**
  - In the best case, it returns a different value for every distinct input given to it
  - Why: a function that always returns 1 is not helpful

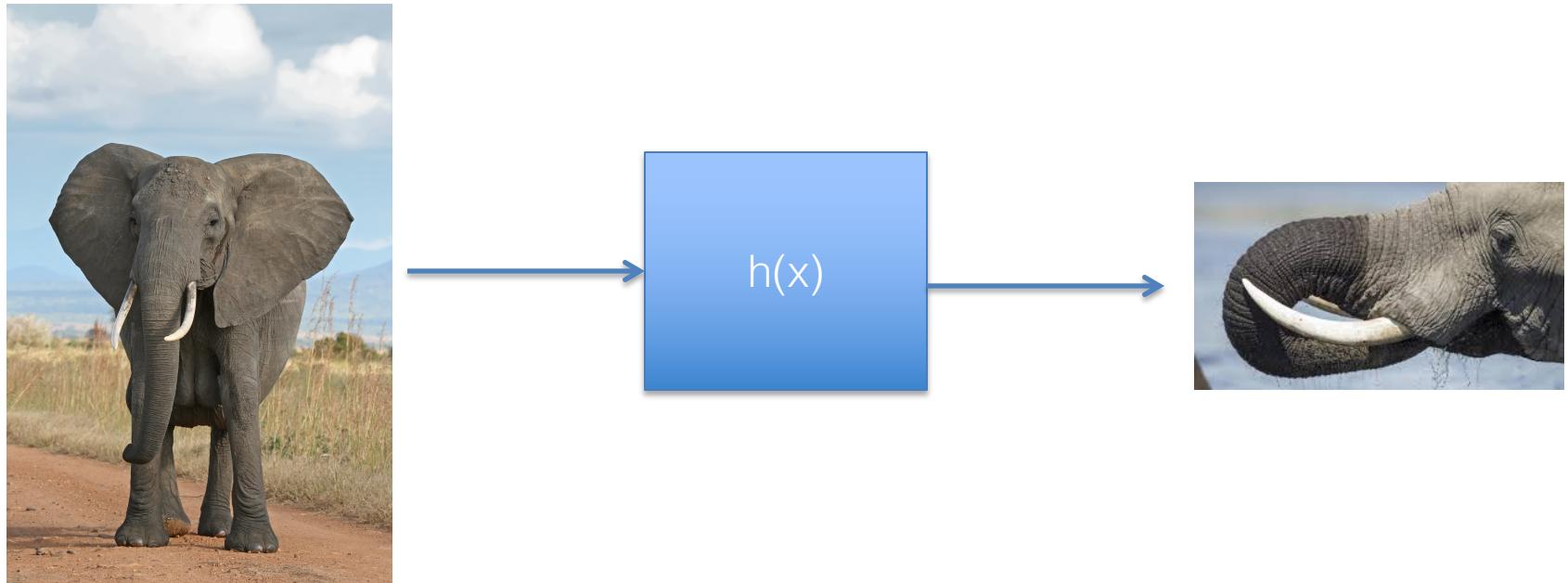
# Introduction to Hash Functions

- Another possible approach – pseudo-random mapping using a hash function  $h(x)$



# Introduction to Hash Functions

- Hash function - maps space M into target space D



[Pictures credit: [https://upload.wikimedia.org/wikipedia/commons/3/37/African\\_Bush\\_Elephant.jpg](https://upload.wikimedia.org/wikipedia/commons/3/37/African_Bush_Elephant.jpg)  
[https://aos.iacpublishinglabs.com/question/aq/1400px-788px/what-are-elephant-tusks-used-for\\_3c174bec-bd85-4fab-a617-7a1a33f14c62.jpg?domain=cx-aos.ask.com](https://aos.iacpublishinglabs.com/question/aq/1400px-788px/what-are-elephant-tusks-used-for_3c174bec-bd85-4fab-a617-7a1a33f14c62.jpg?domain=cx-aos.ask.com)]

# Introduction to Hash Functions

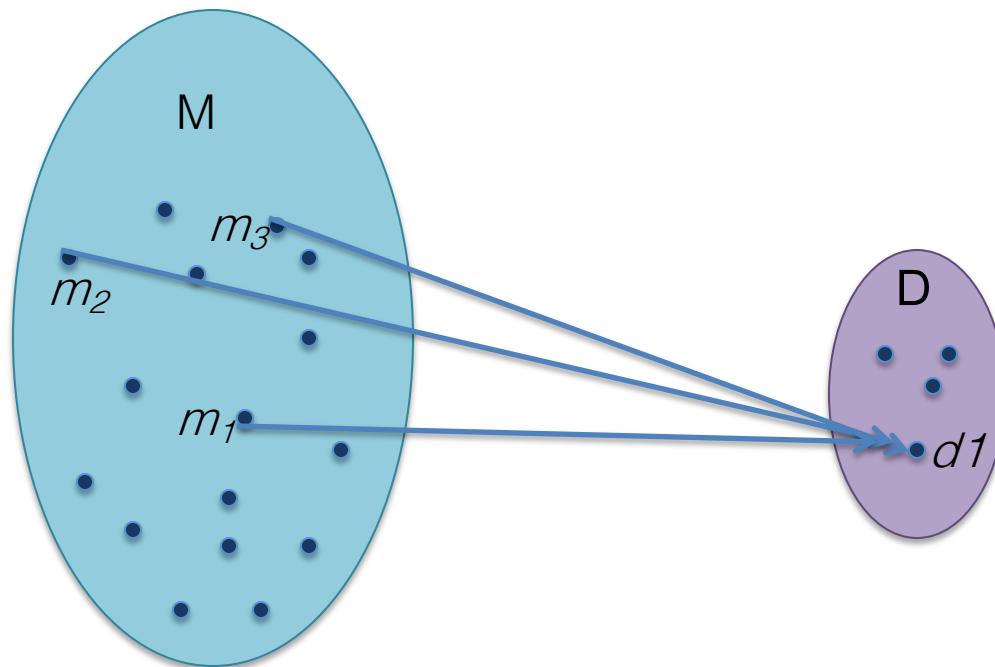
- Hash function - maps the large space  $M$  into target space  $D$
- Desired properties of hash functions:
  - Repeatability:
    - For every  $x$  in  $D$ , it should always be  $h(x) = h(x)$
  - Equally distributed:
    - For some  $y, z$  in  $D$ ,  $P(h(y)) = P(h(z))$
  - Constant-time execution:  $h(x) = O(1)$

# Simple Hash Function

```
/**  
 * A hash method for String objects.  
 * @param key the String to hash.  
 * @param tableSize the size of the hash table.  
 * @return the hash value.  
 */  
public static int hash(String key, int tableSize) {  
    int hashVal = 0;  
  
    for(int i=0; i<key.length(); i++)  
        hashVal = 37 * hashVal + key.charAt(i);  
  
    hashVal %= tableSize;  
    if(hashVal < 0)  
        hashVal += tableSize;  
    return hashVal;  
}
```

# Problems with Hash Functions

- **Hash function** – can be thought of as a “lossy compression function”, since  $|M| \ll |D|$

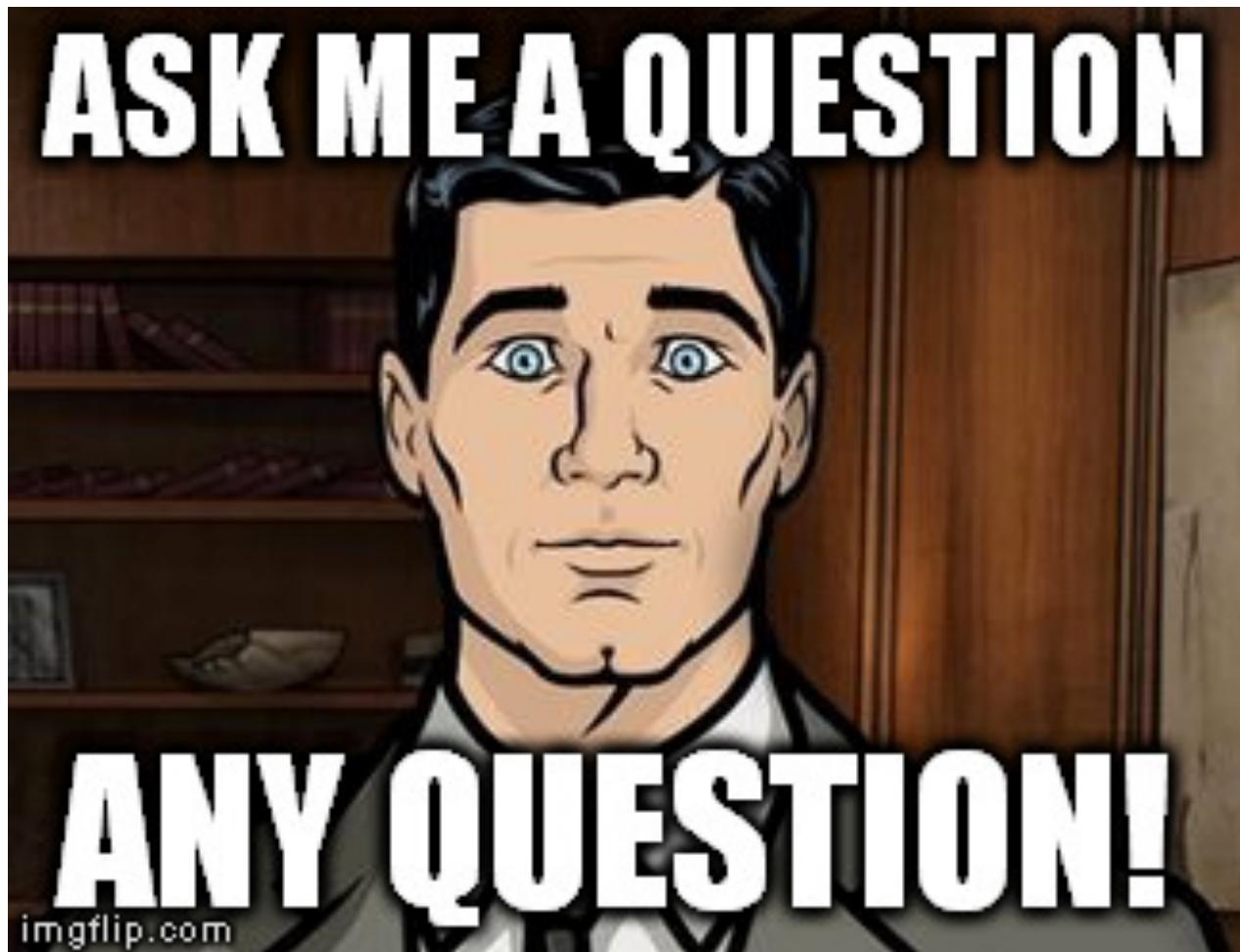


- Do you see any problems here?
- Yes, collision!

# Resolving Collisions

- Hash function – can be thought of as a “lossy compression function”, since  $|M| \ll |D|$
- Problem – collision
- Possible approaches to resolve collisions:
  - Store data in the next available space
  - Store both in the same space
  - Try a different hash
  - Resize the array

# Your Questions



[Meme credit: imgflip.com]

# References and Reading Material

- Mark Allen Weiss, Data Structures and Algorithm Analysis in Java, chapters 1 through 4
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