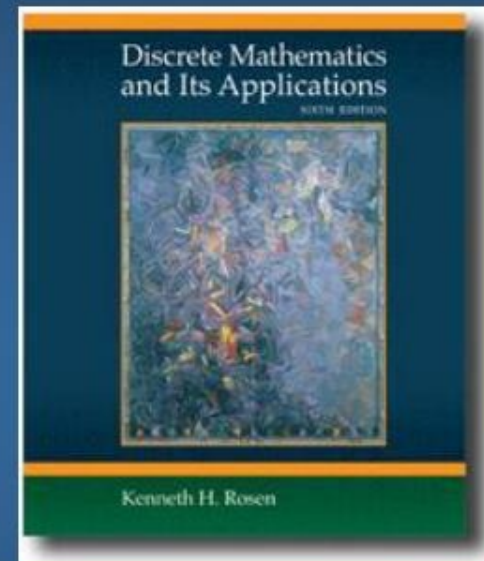


Discrete Mathematics and Its Applications

Sixth Edition
By Kenneth Rosen

Chapter 7 Advanced Counting Techniques



- ◆ 7.1 Recurrence Relations
- ◆ 7.2 Solving Linear Recurrence Relations
- ◆ 7.3 Divide-and-Conquer Algorithms and Recurrence Relations
- ◆ 7.4 Generating Functions
- ◆ 7.5 Inclusion—Exclusion
- ◆ 7.6 Applications of Inclusion—Exclusion

7.1 Recurrence Relations

- Definition 1: A *recurrence relation* for the sequence $\{a_n\}$ is an equation that expresses a_n in terms of one or more of the previous terms of the sequence, namely, a_0, a_1, \dots, a_{n-1} , for all integers n with $n \geq n_0$, where n_0 is a nonnegative integer.
 - A sequence is called a *solution* of a recurrence relation if its terms satisfy the recurrence relation
 - Ex.1-2

Modeling with Recurrence Relations

- Ex.3: Compound interest
- Ex.4: Rabbits and the Fibonacci numbers
- Ex.5: The Tower of Hanoi
- Ex.6
- Ex.7: Codeword enumeration
- Ex.8












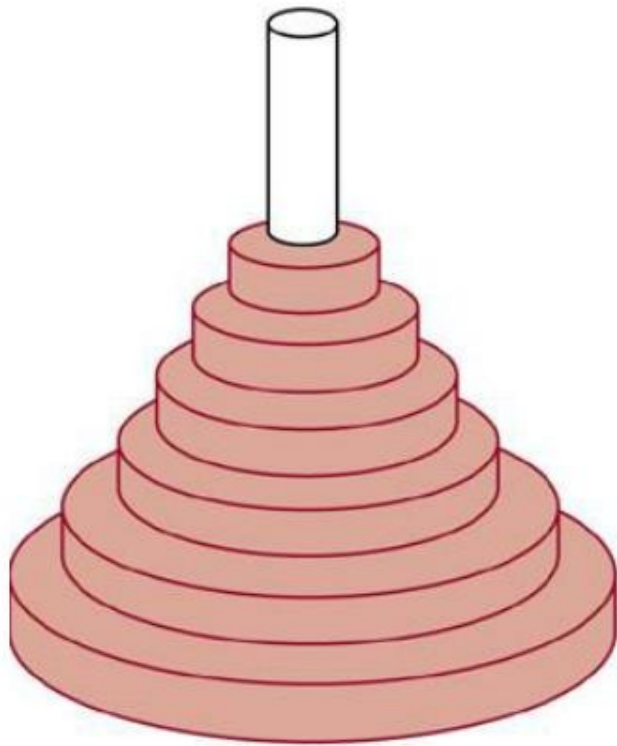
Reproducing pairs (at least two months old)	Young pairs (less than two months old)	Month	Reproducing pairs	Young pairs	Total pairs
		1	0	1	1
		2	0	1	1
		3	1	1	2
		4	1	2	3
		5	2	3	5
	 	6	3	5	8

FIGURE 1 Rabbits on an Island.

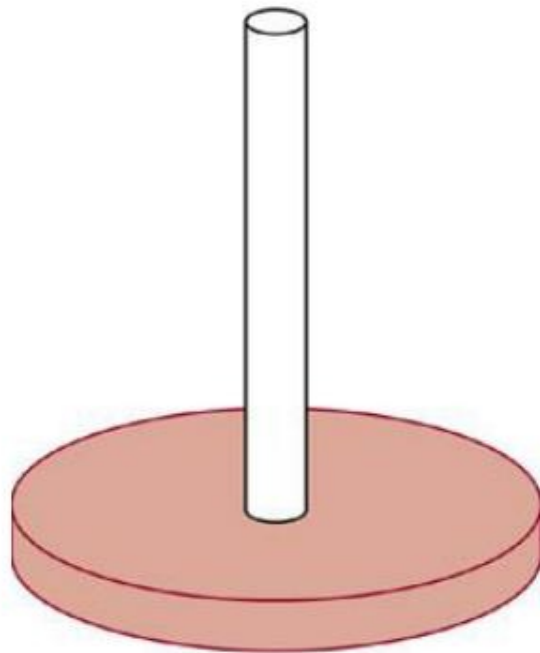


Peg 1

Peg 2

Peg 3

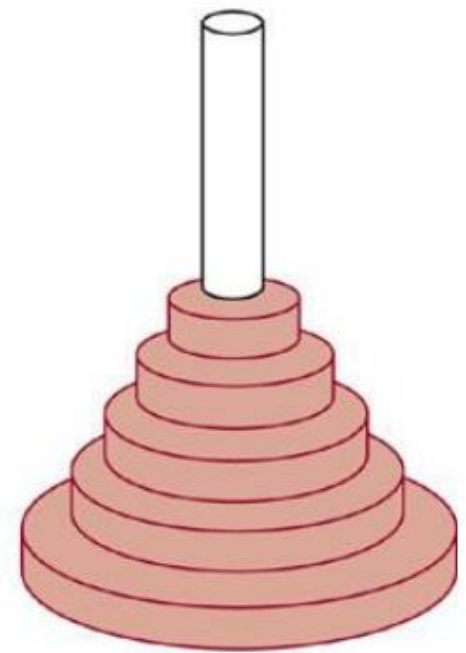
FIGURE 2 The Initial Position in the Tower of Hanoi.



Peg 1



Peg 2



Peg 3

FIGURE 3 An Intermediate Position in the Tower of Hanoi.

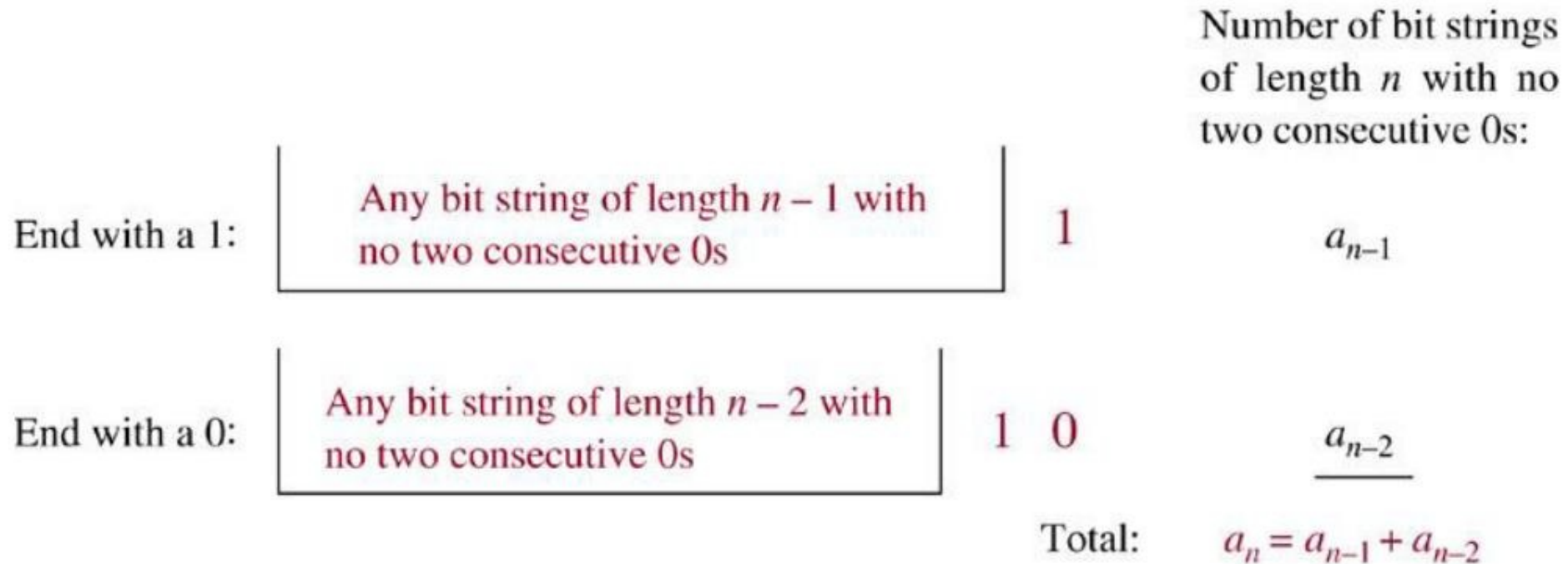


FIGURE 4 Counting Bit Strings of Length n with No Two Consecutive 0s.

7.2 Solving Linear Recurrence Relations

- Definition 1: A *linear homogeneous recurrence relation of degree k with constant coefficients*:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k},$$

where c_1, c_2, \dots, c_k are real numbers, and $c_k \neq 0$.

– Ex.1

- $P_n = (1.11)P_{n-1}$
- $f_n = f_{n-1} + f_{n-2}$
- $a_n = a_{n-5}$

Solving Linear Homogeneous Recurrence Relations with Constant Coefficients

- $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$
 - $a_n = r^n$ is a solution iff $r^n = c_1 r^{n-1} + c_2 r^{n-2} + \dots + c_k r^{n-k}$
 - $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_{k-1} r - c_k = 0$
 - Characteristic equation
 - Characteristic roots
- Theorem 1: Let c_1 and c_2 are real numbers. Suppose $r^2 - c_1 r - c_2 = 0$ has two **distinct** roots r_1 and r_2 . Then the sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1 a_{n-1} + c_2 a_{n-2}$ **iff** $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$ for $n = 0, 1, 2, \dots$, where α_1 and α_2 are constants.
 - Proof.

- Ex.3: $a_n = a_{n-1} + 2a_{n-2}$, $a_0 = 2$, $a_1 = 7$.
- Ex.4: Fibonacci numbers.
- Theorem 2: Let c_1 and c_2 are real numbers with $c_2 \neq 0$. Suppose $r^2 - c_1r - c_2 = 0$ has **only one root** r_0 . Then the sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1a_{n-1} + c_2a_{n-2}$ **iff** $a_n = \alpha_1 r_0^n + \alpha_2 n r_0^n$ for $n = 0, 1, 2, \dots$, where α_1 and α_2 are constants.
 - Ex. 5: $a_n = 6a_{n-1} - 9a_{n-2}$, $a_0 = 1$, $a_1 = 6$.

- Theorem 3: Let c_1, c_2, \dots, c_k be real numbers. Suppose $r^k - c_1 r^{k-1} - \dots - c_k = 0$ has k distinct roots r_1, r_2, \dots, r_k . Then the sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ iff $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n + \dots + \alpha_k r_k^n$ for $n = 0, 1, 2, \dots$, where $\alpha_1, \alpha_2, \dots, \alpha_k$ are constants.
 - Ex.6: $a_n = 6a_{n-1} - 11a_{n-2} + 6a_{n-3}$, $a_0 = 2$, $a_1 = 5$, $a_2 = 15$.

- Theorem 4: Let c_1, c_2, \dots, c_k be real numbers. Suppose $r^k - c_1 r^{k-1} - \dots - c_k = 0$ has t distinct roots r_1, r_2, \dots, r_t with multiplicities m_1, m_2, \dots, m_t such that $m_i \geq 1$ and $m_1 + m_2 + \dots + m_t = k$. Then the sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ **iff**

$$a_n = (\alpha_{1,0} + \alpha_{1,1}n + \dots + \alpha_{1,m_1-1}n^{m_1-1})r_1^n + (\alpha_{2,0} + \alpha_{2,1}n + \dots + \alpha_{2,m_2-1}n^{m_2-1})r_2^n + \dots + (\alpha_{t,0} + \alpha_{t,1}n + \dots + \alpha_{t,m_t-1}n^{m_t-1})r_t^n$$
 for $n = 0, 1, 2, \dots$, where $\alpha_{i,j}$ are constants.
 - Ex.7
 - Ex.8: $a_n = -3a_{n-1} - 3a_{n-2} - a_{n-3}$, $a_0 = 1$, $a_1 = -2$, $a_2 = -1$.

Linear Nonhomogeneous Recurrence Relations with Constant Coefficients

- Nonhomogeneous:
 - $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n)$
 - Associated homogeneous recurrence relation:
$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$
 - Ex.9: $a_n = a_{n-1} + 2^n$, $a_n = a_{n-1} + a_{n-2} + n^2 + n + 1$, $a_n = 3a_{n-1} + n3^n$, $a_n = a_{n-1} + a_{n-2} + a_{n-3} + n!$

- Theorem 5: If $\{a_n^{(p)}\}$ is a particular solution of the nonhomogeneous recurrence relation with constant coefficients $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n)$, then every solution is of the form $\{a_n^{(p)} + a_n^{(h)}\}$, where $a_n^{(h)}$ is a solution of the associated homogeneous recurrence relation $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$.
 - Proof.
 - Ex.10. $a_n = 3a_{n-1} + 2n$, $a_1 = 3$.
 - Ex.11: $a_n = 5a_{n-1} - 6a_{n-2} + 7^n$.

- Theorem 6: Suppose $\{a_n\}$ satisfies the nonhomogeneous recurrence relation with constant coefficients $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n)$, and $F(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n$. When s is **not** a root of the associate recurrence relation, there is a particular solution of the form $(p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n$.
When s is a root of the characteristic equation and its multiplicity is m , there is a particular solution of the form $n^m (p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n$.
 - Ex.12: $a_n = 6a_{n-1} - 9a_{n-2} + F(n)$, $F(n) = 3^n, n3^n, n^2 2^n, (n^2 + 1)3^n$.
 - Ex.13: $a_n = a_{n-1} + n$.

7.3 Divide-and-Conquer Algorithms and Recurrence Relations

- Divide-and-conquer algorithms
 - Divide: a problem of size n into a subproblems, each of size n/b
 - Conquer: combine the solutions of subproblems
- Divide-and-conquer recurrence relation
 - $f(n) = af(n/b) + g(n)$
 - Ex.1: binary search
 - Ex.2: finding the maximum and minimum of a sequence
 - Ex.3: merge sort
 - Ex.4: fast multiplication of integers
 - Ex.5: fast matrix multiplication

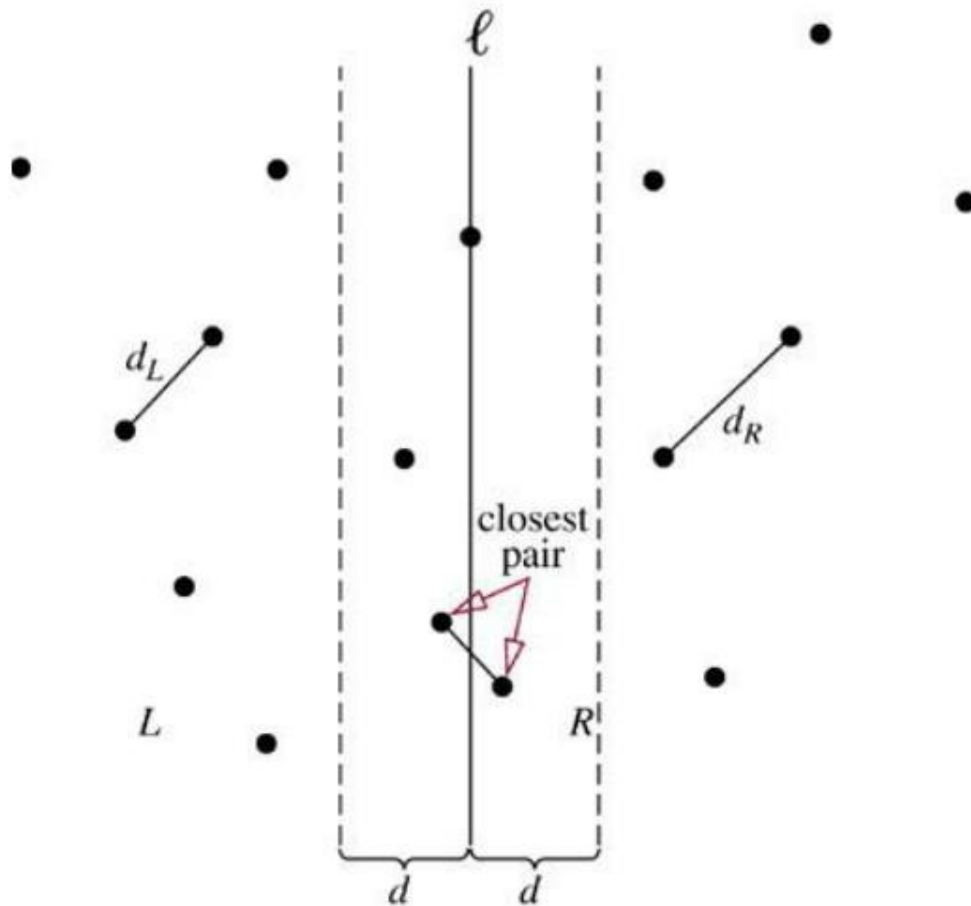
- $n=b^k$
- $f(n)=a^k f(1)+\sum_{j=0..k-1} a^j g(n/b^j)$
- Theorem 1: Let f be an increasing function that satisfies the recurrence relation $f(n)=af(n/b)+c$. Then $f(n)$ is $O(n^{\log b^a})$ if $a>1$, $O(\log n)$ if $a=1$.
When $n=b^k$, $f(n)=C_1 n^{\log b^a}+C_2$, where $C_1=f(1)+c/(a-1)$, $C_2=-c/(a-1)$
 - Proof.
 - Ex. 6: $f(n)=5f(n/2)+3$, $f(1)=7$. Find $f(2^k)$.
 - Ex.7: binary search
 - Ex.8: locate the maximum and minimum elements in a sequence.

- Theorem 2 (**Master Theorem**): Let f be an increasing function that satisfies the recurrence relation

$f(n) = af(n/b) + cn^d$. When $n = b^k$, Then

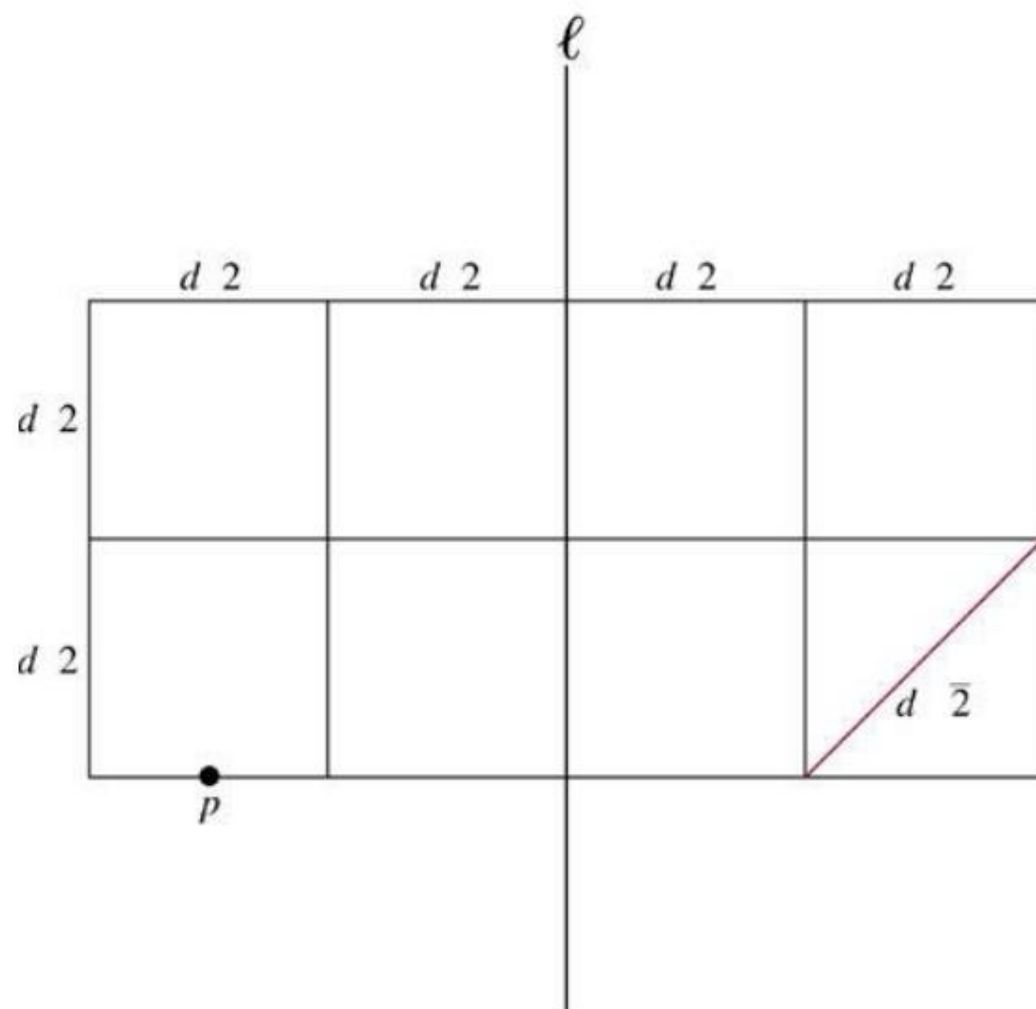
$f(n)$ is $O(n^d)$ if $a < b^d$, $O(n^d \log n)$ if $a = b^d$, $O(n^{\log b^a})$ if $a > b^d$.

- Ex.9: merge sort
- Ex.10: fast multiplication algorithm
- Ex.11: fast matrix multiplication
- Ex.12: The Closest-Pair Problem



In this illustration the problem of finding the closest pair in a set of 16 points is reduced to two problems of finding the closest pair in a set of eight points *and* the problem of determining whether there are points closer than $d = \min(d_L, d_R)$ within the strip of width $2d$ centered at ℓ .

FIGURE 1 The Recursive Step of the Algorithm for Solving the Closest-Pair Problem.



At most eight points, including p , can lie in or on the $2d \times d$ rectangle centered at p because at most one point can lie in or on each of the eight $(d/2) \times (d/2)$ squares.

FIGURE 2 Showing That There Are at Most Seven Other Points to Consider for Each Point in the Strip.

7.4 Generating Functions

- To represent sequences efficiently by coding the terms of a sequence as coefficients of powers of a variable x in a formal power series
 - To solve recurrence relations
 - To prove combinatorial identities
 - To study properties of sequences

- Definition 1: The generating function for the sequence $a_0, a_1, \dots, a_k, \dots$ of real numbers is the infinite series

$$G(x) = a_0 + a_1x + \dots + a_kx^k + \dots = \sum_{k=0.. \infty} a_kx^k.$$

- Ordinary generating functions
- Ex.1
- Ex.2
- Ex.3

Useful Facts about Power Series

- Ex.4: $f(x)=1/(1-x)$
- Ex.5: $f(x)=1/(1-ax)$
- Theorem 1: Let $f(x)=\sum_{k=0..∞}a_kx^k$ and $g(x)=\sum_{k=0..∞}b_kx^k$. Then
 $f(x)+g(x)=\sum_{k=0..∞}(a_k+b_k)x^k$ and $f(x)g(x)=\sum_{k=0..∞}(\sum_{j=0..k}a_jb_{k-j})x^k$ and.
 - It's valid only for power series that converge in an interval.
 - Ex.6: $f(x)=1/(1-x)^2$.

- Definition 2: Let u be a real number and k a nonnegative integer. Then the extended binomial coefficient (u, k) is defined by $(u, k) = u(u-1)\dots(u-k+1)/k!$ if $k > 0$, or 1 if $k = 0$.
 - Ex.7: $(-2, 3), (1/2, 3)$
 - Ex.8: $(-n, r) = (-1)^r C(n+r-1, r)$
- Theorem 2: (The Extended Binomial Theorem)
 Let x be a real number with $|x| < 1$ and let u be a real number. Then $(1+x)^u = \sum_{k=0}^{\infty} (u, k) x^k$.
 - Ex.9: $(1+x)^{-n}, (1-x)^{-n}$.

TABLE 1 Useful Generating Functions.

$G(x)$	a_k
$(1+x)^n = \sum_{k=0}^n C(n, k)x^k$ $= 1 + C(n, 1)x + C(n, 2)x^2 + \cdots + x^n$	$C(n, k)$
$(1+ax)^n = \sum_{k=0}^n C(n, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n, 2)a^2 x^2 + \cdots + a^n x^n$	$C(n, k)a^k$
$(1+x^r)^n = \sum_{k=0}^n C(n, k)x^{rk}$ $= 1 + C(n, 1)x^r + C(n, 2)x^{2r} + \cdots + x^{rn}$	$C(n, k/r)$ if $r \mid k$; 0 otherwise
$\frac{1-x^{n+1}}{1-x} = \sum_{k=0}^n x^k = 1 + x + x^2 + \cdots + x^n$	1 if $k \leq n$; 0 otherwise
$\frac{1}{1-x} = \sum_{k=0}^{\infty} x^k = 1 + x + x^2 + \cdots$	1
$\frac{1}{1-ax} = \sum_{k=0}^{\infty} a^k x^k = 1 + ax + a^2 x^2 + \cdots$	a^k
$\frac{1}{1-x^r} = \sum_{k=0}^{\infty} x^{rk} = 1 + x^r + x^{2r} + \cdots$	1 if $r \mid k$; 0 otherwise
$\frac{1}{(1-x)^2} = \sum_{k=0}^{\infty} (k+1)x^k = 1 + 2x + 3x^2 + \cdots$	$k+1$
$\frac{1}{(1-x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)x^k$ $= 1 + C(n, 1)x + C(n+1, 2)x^2 + \cdots$	$C(n+k-1, k) = C(n+k-1, n-1)$
$\frac{1}{(1+x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)(-1)^k x^k$ $= 1 - C(n, 1)x + C(n+1, 2)x^2 - \cdots$	$(-1)^k C(n+k-1, k) = (-1)^k C(n+k-1, n-1)$
$\frac{1}{(1-ax)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n+1, 2)a^2 x^2 + \cdots$	$C(n+k-1, k)a^k = C(n+k-1, n-1)a^k$
$e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \cdots$	$1/k!$
$\ln(1+x) = \sum_{k=0}^{\infty} \frac{(-1)^{k+1}}{k} x^k = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \cdots$	$(-1)^{k+1}/k$

Note: The series for the last two generating functions can be found in most calculus books when power series are discussed.

Counting Problems and Generating Functions

- Counting the r -combinations from n elements when repetition is allowed
 - $e_1 + e_2 + \dots + e_n = C$
 - Ex.10: $e_1 + e_2 + e_3 = 17$, $e_1:2-5$, $e_2:3-6$, $e_3:4-7$.
 - Ex.11
 - Ex.12
 - Ex.13: k -combinations from n elements.
 - Ex.14: r -combinations from n elements when repetition is allowed.
 - Ex.15: select r objects of n different kinds if we must select at least one object of each kind.

Using Generating Functions to Solve Recurrence Relations

- Ex.16: $a_k = 3a_{k-1}$, $a_0 = 2$.
- Ex.17: $a_n = 8a_{n-1} + 10^{n-1}$, $a_1 = 9$.

Proving Identities via Generating Functions

- Ex.18

7.5 Inclusion-Exclusion

- The principle of inclusion-exclusion
 - $|A \cup B| = |A| + |B| - |A \cap B|$
 - Ex.1: (Fig.1)
 - Ex.2: (Fig.2)
 - Ex.3

$$|A \cup B| = |A| + |B| - |A \cap B| = 25 + 13 - 8 = 30$$

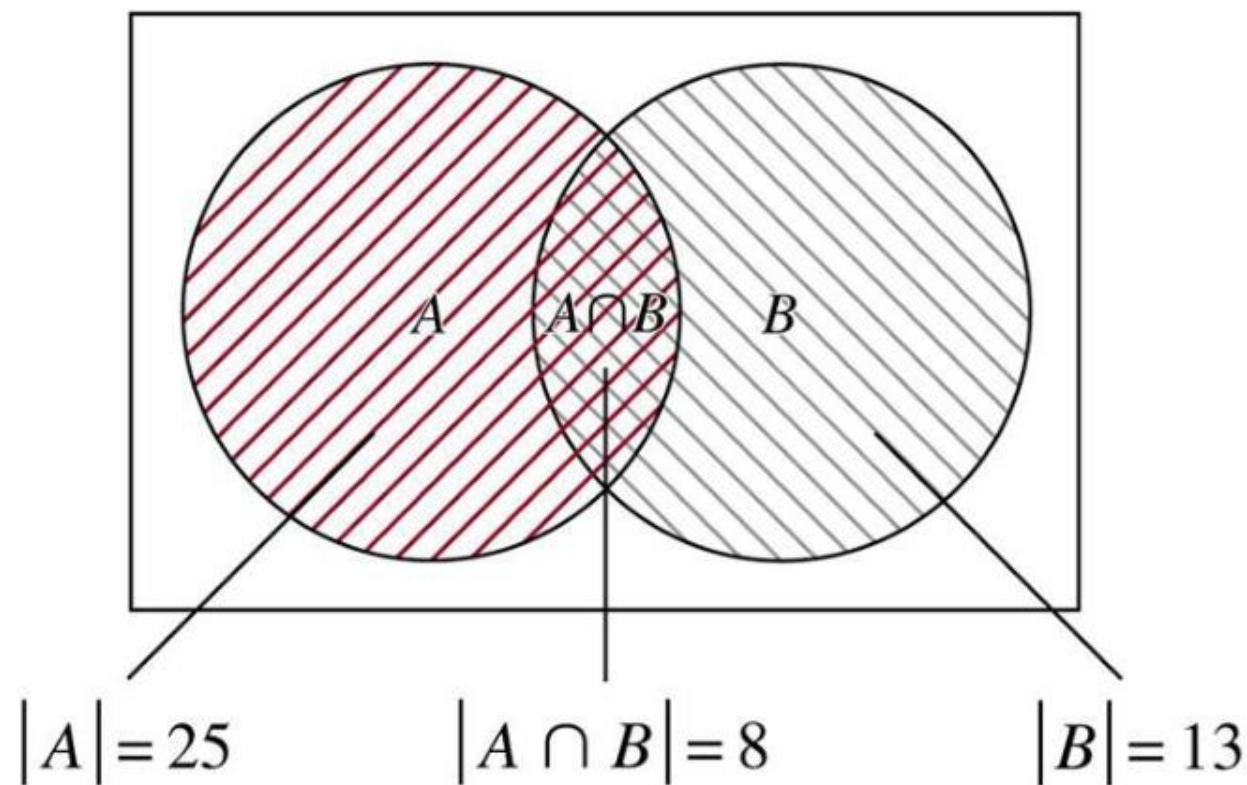


FIGURE 1 The Set of Students in a Discrete Mathematics Class.

$$|A \cup B| = |A| + |B| - |A \cap B| = 142 + 90 - 12 = 220$$

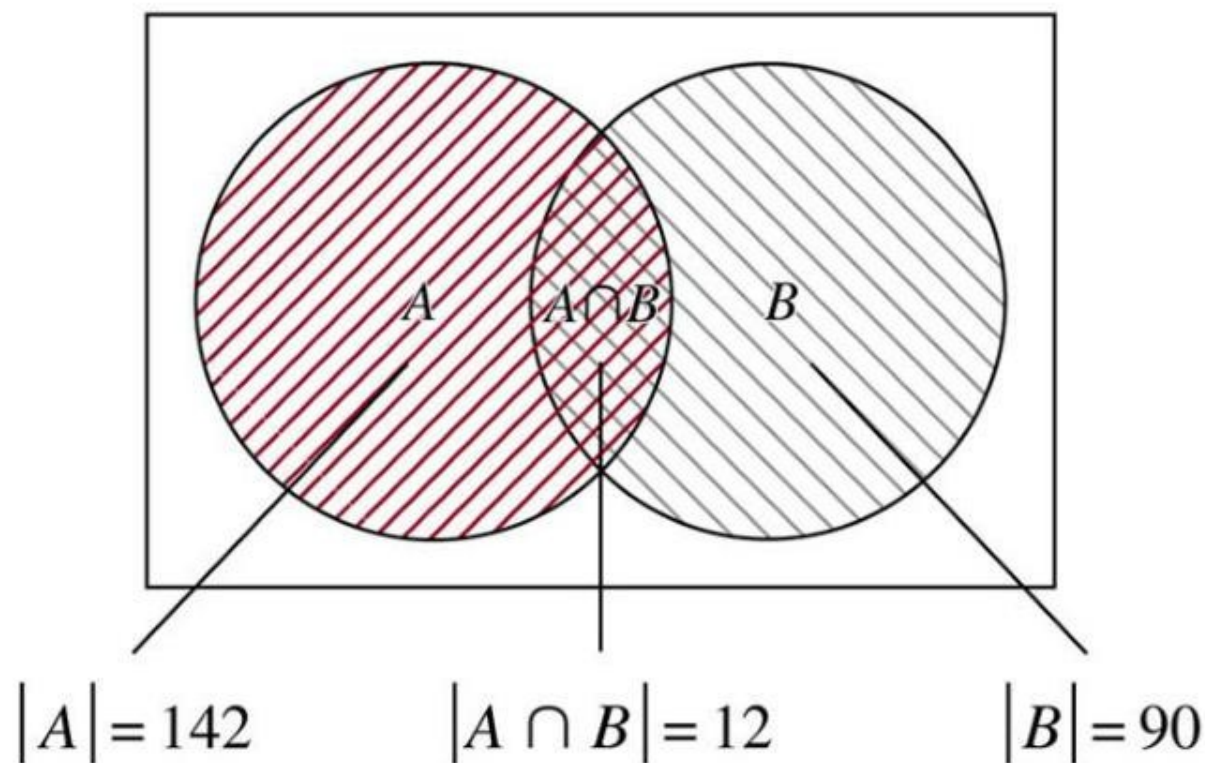
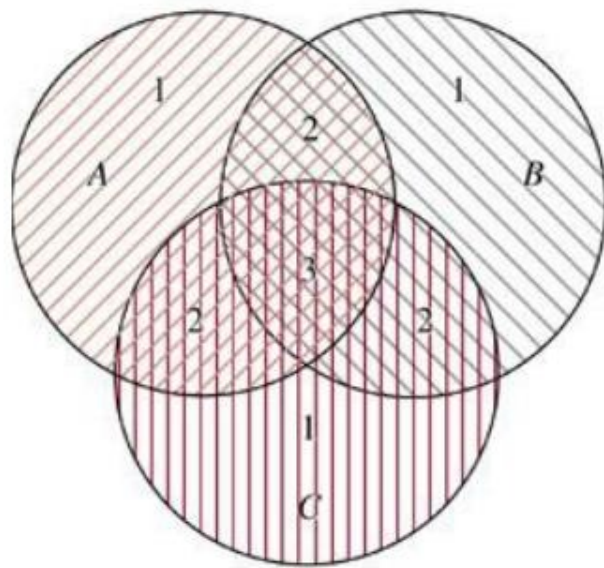
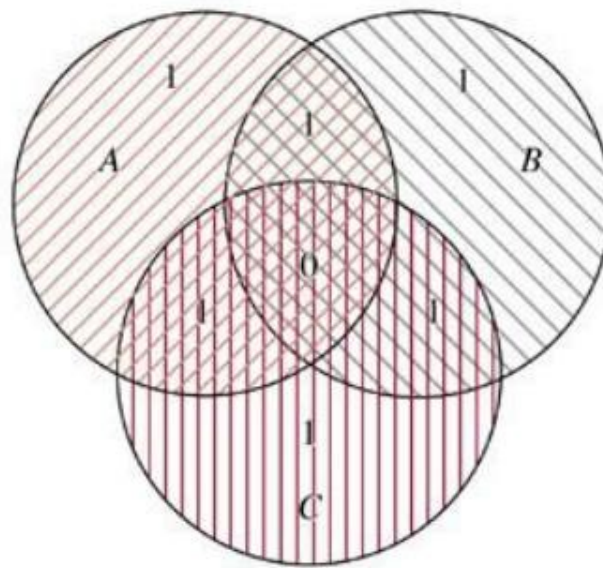


FIGURE 2 The Set of Positive Integers Not Exceeding 1000 Divisible by Either 7 or 11.

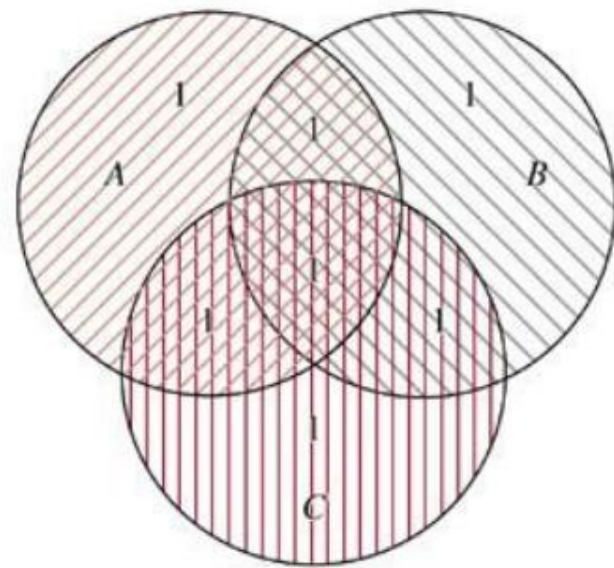
- For three sets: (Fig. 3)
 - $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |C \cap A| + |A \cap B \cap C|$
 - Ex. 4 (Fig. 4)



(a) Count of elements by
 $|A| + |B| + |C|$



(b) Count of elements by
 $|A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C|$



(c) Count of elements by
 $|A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$

FIGURE 3 Finding a Formula for the Number of Elements in the Union of Three Sets.

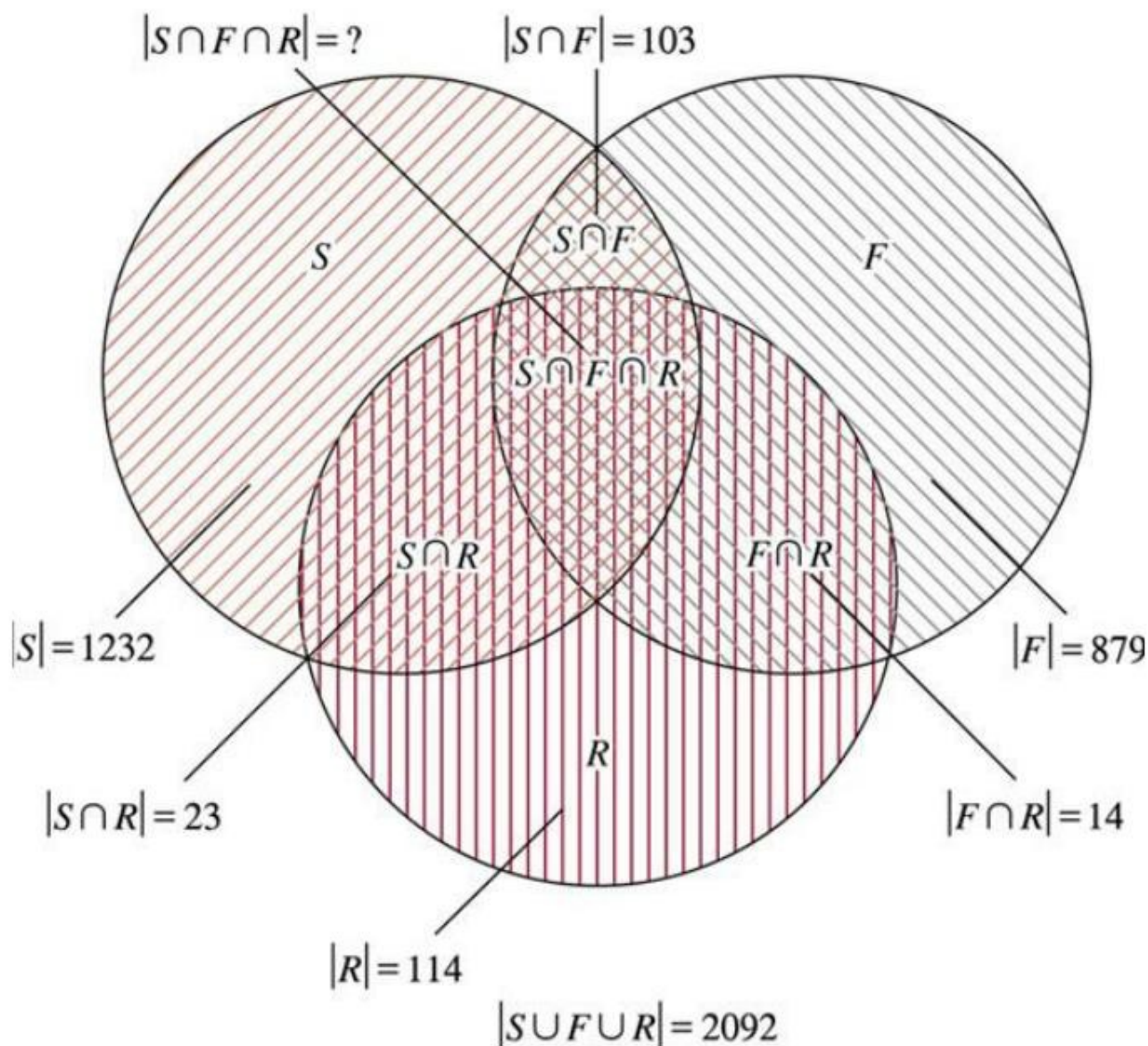


FIGURE 4

The Set of Students Who Have Taken Courses in Spanish, French, and Russian.

- Theorem 1: (The Principle of Inclusion-Exclusion) Let A_1, A_2, \dots, A_n be finite sets. Then

$$|A_1 \cup A_2 \cup \dots \cup A_n| = \sum_{i=1}^n |A_i| - \sum_{i,j=1}^n |A_i \cap A_j| + \sum_{i,j,k=1}^n |A_i \cap A_j \cap A_k| - \dots + (-1)^{n+1} |A_1 \cap A_2 \cap \dots \cap A_n|.$$

– Ex.5

7.6 Applications of Inclusion-Exclusion

- An alternative form of inclusion-exclusion
 - To solve: the number of elements in a set that have none of n properties P_1, P_2, \dots, P_n .
 - Let A_i be the subset that have property P_i .
 - The number of elements with all the properties $P_{i1}, P_{i2}, \dots, P_{ik}$: $N(P_{i1}, P_{i2}, P_{ik})$.
 - $N(P_1', P_2', \dots, P_n') = N - |A_1 \cup A_2 \cup \dots \cup A_n|$
 $= N - \sum_{i=1..n} N(P_i) + \sum_{i,j=1..n} N(P_i P_j) - \sum_{i,j,k=1..n} N(P_i P_j P_k) + \dots$
 $+ (-1)^n N(P_1 P_2 \dots P_n)$.
 - Ex.1

The Sieve of Eratosthenes

- To find all primes not exceeding a specified positive integer
 - For example, the primes not exceeding 100
 - P1: divisible by 2
 - P2: divisible by 3
 - P3: divisible by 5
 - P4: divisible by 7

TABLE 1 The Sieve of Eratosthenes.

<i>Integers divisible by 2 other than 2 receive an underline.</i>										<i>Integers divisible by 3 other than 3 receive an underline.</i>									
1	2	3	4	5	6	7	8	9	10	1	2	3	4	5	6	7	8	9	10
11	<u>12</u>	13	<u>14</u>	15	<u>16</u>	17	<u>18</u>	19	<u>20</u>	11	<u>12</u>	13	<u>14</u>	<u>15</u>	<u>16</u>	17	<u>18</u>	19	<u>20</u>
21	<u>22</u>	23	<u>24</u>	25	<u>26</u>	27	<u>28</u>	29	<u>30</u>	21	<u>22</u>	23	<u>24</u>	25	<u>26</u>	<u>27</u>	<u>28</u>	29	<u>30</u>
31	<u>32</u>	33	<u>34</u>	35	<u>36</u>	37	<u>38</u>	39	<u>40</u>	31	<u>32</u>	<u>33</u>	<u>34</u>	35	<u>36</u>	37	<u>38</u>	<u>39</u>	<u>40</u>
41	<u>42</u>	43	<u>44</u>	45	<u>46</u>	47	<u>48</u>	49	<u>50</u>	41	<u>42</u>	43	<u>44</u>	<u>45</u>	<u>46</u>	47	<u>48</u>	49	<u>50</u>
51	<u>52</u>	53	<u>54</u>	55	<u>56</u>	57	<u>58</u>	59	<u>60</u>	51	<u>52</u>	53	<u>54</u>	55	<u>56</u>	<u>57</u>	<u>58</u>	59	<u>60</u>
61	<u>62</u>	63	<u>64</u>	65	<u>66</u>	67	<u>68</u>	69	<u>70</u>	61	<u>62</u>	<u>63</u>	<u>64</u>	65	<u>66</u>	67	<u>68</u>	<u>69</u>	<u>70</u>
71	<u>72</u>	73	<u>74</u>	75	<u>76</u>	77	<u>78</u>	79	<u>80</u>	71	<u>72</u>	73	<u>74</u>	<u>75</u>	<u>76</u>	77	<u>78</u>	79	<u>80</u>
81	<u>82</u>	83	<u>84</u>	85	<u>86</u>	87	<u>88</u>	89	<u>90</u>	81	<u>82</u>	83	<u>84</u>	85	<u>86</u>	<u>87</u>	<u>88</u>	89	<u>90</u>
91	<u>92</u>	93	<u>94</u>	95	<u>96</u>	97	<u>98</u>	99	<u>100</u>	91	<u>92</u>	<u>93</u>	<u>94</u>	95	<u>96</u>	97	<u>98</u>	<u>99</u>	<u>100</u>
<i>Integers divisible by 5 other than 5 receive an underline.</i>										<i>Integers divisible by 7 other than 7 receive an underline; integers in color are prime.</i>									
1	2	3	4	5	6	7	8	9	10	1	2	3	4	5	6	7	8	9	10
11	<u>12</u>	13	<u>14</u>	<u>15</u>	<u>16</u>	17	<u>18</u>	19	<u>20</u>	11	<u>12</u>	13	<u>14</u>	<u>15</u>	<u>16</u>	17	<u>18</u>	19	<u>20</u>
21	<u>22</u>	23	<u>24</u>	<u>25</u>	<u>26</u>	<u>27</u>	<u>28</u>	29	<u>30</u>	21	<u>22</u>	23	<u>24</u>	25	<u>26</u>	<u>27</u>	<u>28</u>	29	<u>30</u>
31	<u>32</u>	<u>33</u>	<u>34</u>	<u>35</u>	<u>36</u>	37	<u>38</u>	<u>39</u>	<u>40</u>	31	<u>32</u>	<u>33</u>	<u>34</u>	<u>35</u>	<u>36</u>	37	<u>38</u>	<u>39</u>	<u>40</u>
41	<u>42</u>	43	<u>44</u>	<u>45</u>	<u>46</u>	47	<u>48</u>	49	<u>50</u>	41	<u>42</u>	43	<u>44</u>	<u>45</u>	<u>46</u>	47	<u>48</u>	<u>49</u>	<u>50</u>
51	<u>52</u>	53	<u>54</u>	<u>55</u>	<u>56</u>	<u>57</u>	<u>58</u>	59	<u>60</u>	51	<u>52</u>	53	<u>54</u>	55	<u>56</u>	<u>57</u>	<u>58</u>	59	<u>60</u>
61	<u>62</u>	<u>63</u>	<u>64</u>	<u>65</u>	<u>66</u>	67	<u>68</u>	<u>69</u>	<u>70</u>	61	<u>62</u>	<u>63</u>	<u>64</u>	<u>65</u>	<u>66</u>	67	<u>68</u>	<u>69</u>	<u>70</u>
71	<u>72</u>	73	<u>74</u>	<u>75</u>	<u>76</u>	77	<u>78</u>	79	<u>80</u>	71	<u>72</u>	73	<u>74</u>	<u>75</u>	<u>76</u>	<u>77</u>	<u>78</u>	79	<u>80</u>
81	<u>82</u>	83	<u>84</u>	<u>85</u>	<u>86</u>	<u>87</u>	<u>88</u>	89	<u>90</u>	81	<u>82</u>	83	<u>84</u>	<u>85</u>	<u>86</u>	<u>87</u>	<u>88</u>	89	<u>90</u>
91	<u>92</u>	<u>93</u>	<u>94</u>	<u>95</u>	<u>96</u>	97	<u>98</u>	<u>99</u>	<u>100</u>	91	<u>92</u>	<u>93</u>	<u>94</u>	<u>95</u>	<u>96</u>	97	<u>98</u>	<u>99</u>	<u>100</u>

The Number of Onto Functions

- Ex.2
- Theorem 1: Let m and n be positive integers with $m \geq n$. Then, there are $n^m - C(n,1)(n-1)^m + C(n,2)(n-2)^m - \dots + (-1)^{n-1}C(n,n-1)1^m$ **onto** functions from a set with m elements to a set with n elements.
 - $= n!S(m,n)$, where $S(m,n)$ is a Stirling number of the second kind
 - Ex.3

Derangement

- *Derangement*: a permutation of n objects that leave no objects in their original positions
 - Ex.4: The Hatcheck problem (later)
 - Ex.5
 - Let D_n denote the number of derangements of n objects, $D_3=2$
- Theorem 2: The number of derangements of a set with n elements is
$$D_n = n! \left[1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + (-1)^n \frac{1}{n!} \right]$$
 - Proof.

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TABLE 2 The Probability of a Derangement.

n	2	3	4	5	6	7
$D_n/n!$	0.50000	0.33333	0.37500	0.36667	0.36806	0.36786

Thanks for Your Attention!