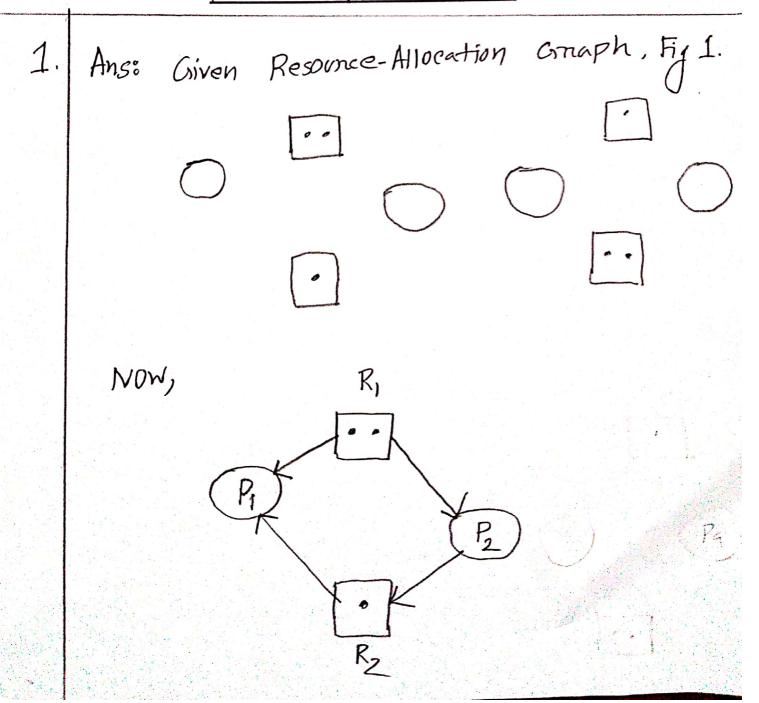


Bangladesh Army University of Engineering & Technology

Qadirabad Cantonment Natore-6431

SL NO: Class Test Script					
Course Code: CSE-3211	Class Test	03	Invigilat	or's Signature & Dat	
Course Title: Openating System			Exam Date: 11-03-21		lass
Batch: 7th Semester: 6th	Department:	CS	E	Section:	
ID NO: 18104050		Obtair	Obtained Mark:		

(Please answer questions from here)



At first here no cycle create, so no possibilites of deadlock. But there is a some possibilites of stanvation here. But when Pa complete work, after free all resources. Then P2 get chance to complete work. So. Safe sequence is

 $(P_1 \rightarrow P_2)$

Again,

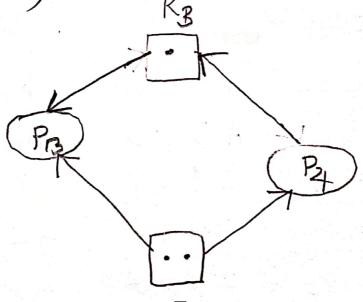


Fig: Safe sequence for graph 2

Hene, also no cycle eneate so no chance for deadlock. Hene At first P2 complete task,
then P1 get reson instance and
after that P1 complete the
Process. So, here some Possibilities
Process. So, here some Possibilities
Of stanvation. Here safe seavence
is, (P3 > P4). So, Now menge
that,
P1 > P2 > P3 > P4
Ans: (i) Hene,

	7.100			
Process	Alloeation	Max	Available	Need
	A B c D	ABCD	ABCD	ABCD
wp,	2111	5234	4992	3123
ve P2	3211	4252	7653	1041
4-P3	2112	2317	81084	0205
re Pg	1431	2543	101195	1112
			1212107	

Safe seavence: P2 > P4 > P1 > P3

9.

In hene no deadlock occurs.

There is possibilities of little

Stanvation. But according to the

Safe seavence no deadlock occur.

(ii) Ars. Hene Our availble is (4,4,4,2). So When a request from process Pz annives for (0,3,0,2) then the negrest be can be granted. Now we need to check after giving Pz, the Other process Will non Safe execute safe. sequence. so. after calculations the need. The safe seavence is 及为月十月一月

So, Pz request can be granted and no deadlock occur.