

Preamble

```
#include <bits/stdc++.h>

using namespace std;

#define int long long
#define REP(i, a, b) for (int i = a; i < (b); ++i)

signed main() {
    cin.tie(NULL)->sync_with_stdio(false);
}
```

Debug Memory Usage

```
long long get_memory_usage() {
    struct rusage usage;
    getrusage(RUSAGE_SELF, &usage);
    return usage.ru_maxrss; // Maximum resident set size (in kilobytes
                             on Linux, bytes on macOS)
}
```

Output

```
// Fixed precision.
cout << fixed << setprecision(6) << lf << '\n';
// Binary output
cout << format("{:06b}", b) << "fixed length binary";
cout << format("{:b}", b) << "variable length binary";
```

Linear Algebra

Gauss-Jordan

Partial Pivot RREF - Rectangular

```
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
tuple<int,double> rref(VVT &a) {
    int n = a.size();
    int m = a[0].size();
    int r = 0;
    double det = 1.;
    for (int c = 0; c < m && r < n; c++) {
        int j = r;
        for (int i = r + 1; i < n; i++)
            if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
        if (fabs(a[j][c]) < EPSILON) continue;
        swap(a[j], a[r]);
        if (j != r) det *= -1.;
        det *= a[r][c];
        T s = 1.0 / a[r][c];
        for (int j = 0; j < m; j++) a[r][j] *= s;
        for (int i = 0; i < n; i++) if (i != r) {
            T t = a[i][c];
            for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
        }
        r++;
    }
    return {r,det};
}
```

Full Pivot - Inverse, Square, Solving  $(n \times n) \cdot (n \times m) = (n. \times m)$

- Solving systems of linear equations ( $AX = B$ )
- Inverting matrices ( $AX = I$ )
- Computing determinants of square matrices

Runs in  $\mathcal{O}(n^3)$

Output:

- $X$  stored in b
- $A^{-1}$  stored in a

```
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
```

```
T GaussJordan(VVT &a, VVT &b) {
    const int n = a.size();
    const int m = b[0].size();
    VI irow(n), icol(n), ipiv(n);
    T det = 1;
    for (int i = 0; i < n; i++) {
        int pj = -1, pk = -1;
        for (int j = 0; j < n; j++) if (!ipiv[j])
            for (int k = 0; k < n; k++) if (!ipiv[k])
                if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
        if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det *= -1;
        irow[i] = pj;
        icol[i] = pk;
        T c = 1.0 / a[pk][pk];
        det *= a[pk][pk];
        a[pk][pk] = 1.0;
        for (int p = 0; p < n; p++) a[pk][p] *= c;
        for (int p = 0; p < m; p++) b[pk][p] *= c;
        for (int p = 0; p < n; p++) if (p != pk) {
            c = a[p][pk];
            a[p][pk] = 0;
            for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
            for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
        }
    }
    for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
        for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
    }
    return det;
}
```

XOR Basis

Small vectors

```
vector<int> basis;
void add(int x) {
    for (int i = 0; i < basis.size(); i++) { // reduce x using the
        current basis vectors
        x = min(x, x ^ basis[i]);
    }
    if (x != 0) { basis.push_back(x); }
}
```

Arbitrarily large vectors

```
bool non_zero(const vector<uint64_t>& x) {
    bool non_zero = false;
    for(const auto& a : x) {
        non_zero |= (a != (uint64_t) 0);
    }
    return non_zero;
}
struct Basis {
    vector<vector<uint64_t>> basis;
    vector<uint64_t> reduce(vector<uint64_t> x) {
        for(int i = 0; i < basis.size(); i++) {
            int state = 0;
            for(int j = 0; j < x.size(); j++) {
                int cur = basis[i][j] ^ x[j];
                if (state == 0 and cur < x[j]) state = -1;
                if (state == 0 and cur > x[j]) state = 1;
                if (state <= 0) x[j] = cur;
            }
        }
        return x;
    }
    void add(vector<uint64_t> x) {
        x = reduce(x);
        if (non_zero(x)) basis.push_back(x);
    }
}
```

```

}
bool equal(const Basis& other) {
    if (other.basis.size() != basis.size()) return false;
    bool ans = true;
    for(const auto & v : other.basis) {
        ans &= !non_zero(reduce(v));
    }
    return ans;
}
};

```

# Number Theory

## Extended Euclidean Algorithm

Finds  $x$  and  $y$  for which  $ax + by = \gcd(a, b)$ .

**Time:**  $\mathcal{O}(\log n)$

```

// Returns {x,y,gcd} where xa + yb = gcd
array<int,3> gcd_ext(int a,int b) {
    auto oa=a,ob=b;
    int x=0,y=1,u=1,v=0;
    while(a!=0) {
        auto q=b/a,r=b%a;
        auto m=x-u*q,n=y-v*q;
        b=a, a=r, x=u,y=v,u=m,v=n;
    }
    assert(oa*x+ob*y==b);
    return {x,y,b};
}

```

## Modular Inverse

Finds  $x$  such that  $ax = 1 \bmod m$ .

**Time:**  $\mathcal{O}(\log n)$

```

int inv(int a, int m) {
    auto [x,y,g] = gcd_ext(a, m);
    if (g != 1) {
        // No solution!!!
        return -1;
    }
    else {
        // Inverse
        return (x % m + m) % m;
    }
}

```

TODO: All modular inverses in  $\mathcal{O}(m)$ : <https://cp-algorithms.com/algebra/module-inverse.html>

## Linear Congruence Equation

**Time:**  $\mathcal{O}(\log n)$

```

// Returns {solution, modulo}
pair<int,int> linear_congruence(int a, int b, int n) {
    int d;
    if ((d = gcd(a,n)) != 1) {
        // No solution
        if (b % d != 0) return {-1, -1};
        a /= d; b /= d; n /= d;
    }
    int i = inv(a, n);
    return {(b * i) % n, n};
}

```

## Linear Prime Sieve

This calculates the minimum prime factor  $pr[j]$  for all  $j$  up to  $n$ . From this, we can calculate the prime factorisation of all these numbers.

**Time:**  $\mathcal{O}(n)$

```

const int N = 10000000;
vector<int> lp(N+1);
vector<int> pr;
for (int i=2; i <= N; ++i) {
    if (lp[i] == 0) { lp[i] = i; pr.push_back(i); }
    for (int j = 0; i * pr[j] <= N; ++j) {
        lp[i * pr[j]] = pr[j];
    }
}

```

```

    if (pr[j] == lp[i]) break;
}
}

```

## Extended Chinese Remainder Theorem

Works for non-coprime moduli

```

struct ChineseRemainder {
    int a=0,m=0;
    void add(int b, int n) {
        b=(b%n+n)%n;
        if(m==-1) return;
        if(m==0) { a=b; m=n; return; }
        auto [u,v,g] = gcd_ext(m,n);
        if((a-b)%g!=0) { m=-1;return; }
        int lam = (a-b)/g;
        m=m/g*n;
        a = b + (lam*v)%m*n;
        a = (a%m+m)%m;
    }
    int get(int x) {return a+m*x;}
};

```

## Fast Fourier Transform

Useful for multiplying polynomials, or computing convolutions.  $c[k] = \sum_i a[i]b[k-i]$ . For sliding element-wise multiplication, reverse one of the arrays. Rounding is safe if  $(\sum a_i^2 + \sum b_i^2) \log_2 N < 9 \cdot 10^{14}$ . ( $N = |A| + |B|$ ). In practice, with random inputs, bound is  $10^{16}$ .

**Time:**  $\mathcal{O}(N \log N)$

```

#define SZ(x) (int)(x).size()
#define ALL(x) begin(x), end(x)

typedef vector<int> vi;
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C> &a) {
    int n = SZ(a), L = 31 - __builtin_clz(n);
    static vector<complex<long double>> R(2, 1);
    static vector<C> rt(2, 1); // (^ 10% faster i f double )
    for (static int k = 2; k < n; k *= 2) {
        R.resize(n);
        rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k);
        REP(i, k, 2 * k) rt[i] = R[i] = i & 1 ? R[i / 2] * x : R[i / 2];
    }
    vi rev(n);
    REP(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
    REP(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int k = 1; k < n; k *= 2)
        for (int i = 0; i < n; i += 2 * k) REP(j, 0, k) {
            C z = rt[j + k] *
                a[i + j + k]; // (25% faster i f hand-r o l l e d )
            a[i + j + k] = a[i + j] - z;
            a[i + j] += z;
        }
}
vd conv(const vd &a, const vd &b) {
    if (a.empty() || b.empty()) return {};
    vd res(SZ(a) + SZ(b) - 1);
    int L = 32 - __builtin_clz(SZ(res)), n = 1 << L;
    vector<C> in(n), out(n);
    copy(ALL(a), begin(in));
    REP(i, 0, SZ(b)) in[i].imag(b[i]);
    fft(in);
    for (C &x : in) x *= x;
    REP(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
    fft(out);
    REP(i, 0, SZ(res)) res[i] = imag(out[i]) / (4 * n);
    return res;
}

```

## Geometry

### Preamble

This gives us vector addition, scalar and complex multiplication, angle `arg()`, and polar form initialisation `cis()`.

```
typedef complex<double> C;
```

### Dot Product

```
double dotp(C a , C b){return (conj(a)*b).real();}
double dist2(C a, C b){return dotp(a-b, a-b);}
```

$$a_0b_0 + a_1b_1 = |a||b|\cos(\theta)$$

### Cross Product

```
double crossp(C a , C b){return (conj(a)*b).imag();}
double orient(C a, C b, C c){return crossp(b-c,b-a);}
```

$$a_0b_1 - a_1b_0 = |a||b|\cos(\theta)$$

### Ordering By Orientation

```
bool topHalf(C a) {
    return (a.imag() > 0) || (a.imag() == 0 && a.real() >= 0);
}
bool cmp(const C &a, const C &b) {
    bool ha = topHalf(a);
    bool hb = topHalf(b);
    if (ha != hb) return ha;
    return orient(a, {0,0}, b) > 0;
}
```

## String Matching

### Z-Algorithm

```
vector<int> z_algo(const string& s) {
    int n = s.size();
    vector<int> z(n);
    int l = 0, r = 0;
    for(int i = 1; i < n; i++) {
        if(i < r) z[i] = min(r - i, z[i - l]);
        while(i + z[i] < n && s[z[i]] == s[i + z[i]]) z[i]++;
        if(i + z[i] > r) { l = i; r = i + z[i]; }
    }
    return z;
}
```

### Aho-Curasick

Creates a string automaton for matching a dictionary of patterns. We hit a success state for each match of a pattern. Linear time on the total length of all patterns.

```
struct Node {
    int par;
    char c;
    map<char, int> next;
    int link = -1;
    bool terminal = false;
    Node(int par, char c) : par(par), c(c) {}
};
vector<Node> nodes;
int new_node(int par, char c) {
    Node node = Node(par, c);
    nodes.push_back(node);
    return nodes.size() - 1;
}
int aho_curasick(const vector<string>& words) {
    // Root
    new_node(-1, '!');
    // Trie construction
    for (const auto& word : words) {
        int cur = 0;
        REP(i, 0, word.size()) {
            char c = word[i];
            if (nodes[cur].next.find(c) == nodes[cur].next.end()) {
                int nw = new_node(cur, c);
                nodes[cur].next[c] = nw;
            }
        }
    }
}
```

```
cur = nodes[cur].next[c];
}
nodes[cur].terminal = true;
}
// Initialize root.
deque<int> q;
nodes[0].link = 0;
for (char c = 'a'; c <= 'z'; c++) {
    if (nodes[0].next.find(c) == nodes[0].next.end()) {
        nodes[0].next[c] = 0;
    } else {
        q.push_back(nodes[0].next[c]);
    }
}
// BFS - initialise suffix links and failiure states.
while (!q.empty()) {
    int i = q.front();
    q.pop_front();
    if (nodes[i].par == 0) {
        nodes[i].link = 0;
    } else {
        nodes[i].link =
            nodes[nodes[nodes[i].par].link].next[nodes[i].c];
    }
    for (char c = 'a'; c <= 'z'; c++) {
        if (nodes[i].next.find(c) == nodes[i].next.end()) {
            nodes[i].next[c] = nodes[nodes[i].link].next[c];
        } else {
            q.push_back(nodes[i].next[c]);
        }
    }
}
return 0;
}
```

### Ukkonen's

Linear time suffix tree construction. Useful for string matching.

```
const int MAXN = 8000005;
string s;
int n;
struct Node {
    int l, r, par, link;
    vector<pair<char, int>> next;
    Node(int l = 0, int r = 0, int par = -1) : l(l), r(r), par(par), link(-1) {}
    int len() { return r - l; }
    // More space efficient than map, can use alternatively.
    int& get(char c) {
        for (auto& [a, b] : next)
            if (a == c) return b;
        next.push_back({c, -1});
        return next.back().second;
    }
};
Node t[MAXN];
int sz;
struct State {
    int v, pos;
    State(int v, int pos) : v(v), pos(pos) {}
};
State ptr(0, 0);
State go(State st, int l, int r) {
    while (l < r)
        if (st.pos == t[st.v].len()) {
            st = State(t[st.v].get(s[l]), 0);
            if (st.v == -1) return st;
        } else {
            if (s[t[st.v].l + st.pos] != s[l]) return State(-1, -1);
            if (r - l < t[st.v].len() - st.pos)
                return State(st.v, st.pos + r - l);
            l += t[st.v].len() - st.pos;
            st.pos = t[st.v].len();
        }
}
```

```

    return st;
}

int split(State st) {
    if (st.pos == t[st.v].len()) return st.v;
    if (st.pos == 0) return t[st.v].par;
    Node v = t[st.v];
    int id = sz++;
    t[id] = Node(v.l, v.l + st.pos, v.par);
    t[v.par].get(s[v.l]) = id;
    t[id].get(s[v.l + st.pos]) = st.v;
    t[st.v].par = id;
    t[st.v].l += st.pos;
    return id;
}

int get_link(int v) {
    if (t[v].link != -1) return t[v].link;
    if (t[v].par == -1) return 0;
    int to = get_link(t[v].par);
    return t[v].link = split(go(State(to, t[to].len()), t[v].l +
(t[v].par == 0), t[v].r));
}

void tree_extend(int pos) {
    for (;;) {
        State nptr = go(ptr, pos, pos + 1);
        if (nptr.v != -1) {
            ptr = nptr;
            return;
        }
        int mid = split(ptr);
        int leaf = sz++;
        t[leaf] = Node(pos, n, mid);
        t[mid].get(s[pos]) = leaf;

        ptr.v = get_link(mid);
        ptr.pos = t[ptr.v].len();
        if (!mid) break;
    }
}

void build_tree() {
    sz = 1;
    for (int i = 0; i < n; ++i) tree_extend(i);
}

```

## Segment Trees!!!

### Basic

```

struct BasicSegmentTree {
    using Value = int;
    Value identity = INT_MAX;
    Value binop(Value a, Value b) {return min(a, b);}
    vector<Value> arr;
    int size;
    BasicSegmentTree(int n) : arr(4*n + 2,identity), size(n) {};
    void update(int cur, int i, Value v, int l, int r) {
        if (l == r) {arr[cur] = v; return; }
        int mid = midpoint(l, r);
        if (i <= mid) update(2*cur, i, v, l, mid);
        else update(2*cur + 1, i, v, mid + 1, r);
        arr[cur] = binop(arr[2*cur],arr[2*cur + 1]);
    }
    void update(int i, int v) {update(1,i,v,0,size - 1);}
    Value query(int cur, int ql, int qr, int l, int r) {
        if (l == ql and r == qr) return arr[cur];
        int mid = midpoint(l,r);
        Value val = identity;
        if (ql <= mid) val = binop(val,
query(2*cur,ql,min(mid,qr),l,mid));
        if (qr > mid) val = binop(val,query(2*cur + 1,max(mid +
1,ql),qr,mid+1,r));
        return val;
    }
    Value query(int ql, int qr) {return query(1,ql,qr,0,size - 1);}
};

```

## Lazy Update

```

struct LazyUpdateTree {
    using Value = int;
    using Update = int;
    Value identity = LLONG_MIN;
    Value def = 0;
    Update idUpdate = 0;
    Value binop(Value a, Value b) { return max(a, b); }
    Value applyUpdate(Update a, Value u, int l, int r) { return u +
a; }
    Update mergeUpdate(Update old, Update nw) { return old + nw; }
    vector<Value> arr;
    vector<Update> lazy;
    int size;
    LazyUpdateTree(int n) : arr(4 * n + 2, def), lazy(4 * n + 2,
idUpdate), size(n) {};
    void push(int cur, int l, int r) {
        if (l != r) {
            int mid = midpoint(l, r);
            update(2 * cur, l, mid,lazy[cur],l,mid);
            update(2 * cur + 1, mid + 1, r, lazy[cur],mid + 1, r);
        }
        lazy[cur] = idUpdate;
    }
    void update(int cur, int ql, int qr, Update u, int l, int r) {
        if (l == ql and r == qr) {
            lazy[cur] = mergeUpdate(lazy[cur], u);
            arr[cur] = applyUpdate(u, arr[cur], l, r);
            return;
        }
        push(cur, l, r);
        int mid = midpoint(l, r);
        if (ql <= mid) update(2 * cur, ql, min(mid, qr), u, l, mid);
        if (qr > mid) update(2 * cur + 1, max(mid + 1, ql), qr, u, mid +
1, r);
        arr[cur] = binop(arr[2 * cur], arr[2 * cur + 1]);
    }
    void update(int ql, int qr, Update u) { update(1, ql, qr, u, 0,
size - 1); }
    Value query(int cur, int ql, int qr, int l, int r) {
        if (l == ql and r == qr) return arr[cur];
        push(cur, l, r);
        int mid = midpoint(l, r);
        Value val = identity;
        if (ql <= mid) val = binop(val, query(2 * cur, ql, min(mid, qr),
l, mid));
        if (qr > mid) val = binop(val, query(2 * cur + 1, max(mid + 1,
ql), qr, mid + 1, r));
        return val;
    }
    Value query(int ql, int qr) { return query(1, ql, qr, 0, size -
1); }
};

```

## Lazy Create Segment Tree

```

struct LazyCreateNode {
    using Value = int;
    static const Value identity = 0;
    Value binop(Value a, Value b) { return a + b; }
    Value val = identity;
    LazyCreateNode *lt = NULL, *rt = NULL;
    LazyCreateNode() {};
    void create_children() {
        if (lt != NULL) return;
        lt = new LazyCreateNode();
        rt = new LazyCreateNode();
    }
    void update(int i, Value v, int l, int r) {
        if (l == r) { val = v; return; }
        create_children();
        int mid = midpoint(l, r);
        if (i <= mid) lt->update(i,v,l,mid);
        else rt->update(i,v,mid+1,r);
        val = binop(lt->val,rt->val);
    }
};

```

```

Value query(int ql, int qr, int l, int r) {
    if (l == ql and r == qr) return val;
    create_children();
    int mid = midpoint(l, r);
    Value val = identity;
    if (ql <= mid) val = lt->query(ql, min(mid, qr), l, mid);
    if (qr > mid) val = binop(val, rt->query(max(mid + 1, ql), qr, mid + 1, r));
    return val;
}
};

```

## DP Optimisations

### LineContainer

Container where you can add lines of the form  $kx+m$ , and query maximum values at points  $x$ . Useful for dynamic programming (convex hull trick").  $\mathcal{O}(\log n)$ .

```

struct Line {
    mutable ll k, m, p;
    bool operator<(const Line& o) const { return k < o.k; }
    bool operator<(ll x) const { return p < x; }
};

struct LineContainer : multiset<Line, less<>> {
    // (for doubles, use inf = 1/.0, div(a,b) = a/b)
    static const ll inf = LLONG_MAX;
    ll div(ll a, ll b) { // floored division
        return a / b - ((a ^ b) < 0 && a % b); }
    bool isect(iterator x, iterator y) {
        if (y == end()) return x->p = inf, 0;
        if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
        else x->p = div(y->m - x->m, x->k - y->k);
        return x->p >= y->p;
    }
    void add(ll k, ll m) {
        auto z = insert({k, m, 0}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
        if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y->p)
            isect(x, erase(y));
    }
    ll query(ll x) {
        assert(!empty());
        auto l = *lower_bound(x);
        return l.k * x + l.m;
    }
};

```

### Convex Hull Trick

From a set of linear functions, finds the minimum value at a point.

- Adding Equation - Amortized  $\mathcal{O}(1)$
- Finding minimum -  $\mathcal{O}(\log n)$

Requires gradients to be increasing when inserted. Line container allows for online insertion and is preferable unless both queries and insertions are monotonic where we can achieve  $\mathcal{O}(n)$

To find maximum, flip equations on insert, and flip answer.

```

typedef ll F;
typedef complex<F> P;
F dot(P a, P b) {
    return (conj(a) * b).real();
}
F cross(P a, P b) {
    return (conj(a) * b).imag();
}
struct Cht {
    deque<P> hull, vecs;
    // Ascending gradient
    void add_line_back(F k, F b) {
        P nw = {k, b};
        while(!vecs.empty() && dot(vecs.back(), nw - hull.back()) < 0)
            { hull.pop_back(); vecs.pop_back(); }
        if(!hull.empty())

```

```

        vecs.push_back(P(0,1) * (nw - hull.back()));
        hull.push_back(nw);
    }
    // Descending gradient
    void add_line_front(F k, F b) {
        P nw = {k, b};
        while(!vecs.empty() && dot(vecs.front(), nw - hull.front()) < 0)
            { hull.pop_front(); vecs.pop_front(); }
        if(!hull.empty())
            vecs.push_front(P(0,1) * (hull.front() - nw));
    }
    F get(F x) {
        P query = {x, 1};
        auto it = lower_bound(vecs.begin(), vecs.end(), query,
            [](P a, P b) { return cross(a, b) > 0; }
        );
        return dot(query, hull[it - vecs.begin()]);
    }
};

```

## Flow

### Edmund Karp

$\mathcal{O}(VE^2)$

```

int n;
vector<vector<int>> capacity;
vector<vector<int>> adj;

int bfs(int s, int t, vector<int>& parent) {
    fill(parent.begin(), parent.end(), -1);
    parent[s] = -2;
    queue<pair<int, int>> q;
    q.push({s, INF});

    while (!q.empty()) {
        int cur = q.front().first;
        int flow = q.front().second;
        q.pop();

        for (int next : adj[cur]) {
            if (parent[next] == -1 && capacity[cur][next]) {
                parent[next] = cur;
                int new_flow = min(flow, capacity[cur][next]);
                if (next == t) return new_flow;
                q.push({next, new_flow});
            }
        }
    }

    return 0;
}

int maxflow(int s, int t) {
    int flow = 0;
    vector<int> parent(n);
    int new_flow;

    while (new_flow = bfs(s, t, parent)) {
        flow += new_flow;
        int cur = t;
        while (cur != s) {
            int prev = parent[cur];
            capacity[prev][cur] -= new_flow;
            capacity[cur][prev] += new_flow;
            cur = prev;
        }
    }

    return flow;
}

```

### Push Relabel

$\mathcal{O}(V^3)$

```

const int inf = 1000000000;

int n;
vector<vector<int>> capacity, flow;
vector<int> height, excess;

void push(int u, int v)
{
    int d = min(excess[u], capacity[u][v] - flow[u][v]);
    flow[u][v] += d;
    flow[v][u] -= d;
    excess[u] -= d;
    excess[v] += d;
}

void relabel(int u)
{
    int d = inf;
    for (int i = 0; i < n; i++) {
        if (capacity[u][i] - flow[u][i] > 0)
            d = min(d, height[i]);
    }
    if (d < inf) height[u] = d + 1;
}

vector<int> find_max_height_vertices(int s, int t) {
    vector<int> max_height;
    for (int i = 0; i < n; i++) {
        if (i != s && i != t && excess[i] > 0) {
            if (!max_height.empty() && height[i] > height[max_height[0]])
                max_height.clear();
            if (max_height.empty() || height[i] == height[max_height[0]])
                max_height.push_back(i);
        }
    }
    return max_height;
}

int max_flow(int s, int t)
{
    height.assign(n, 0);
    height[s] = n;
    flow.assign(n, vector<int>(n, 0));
    excess.assign(n, 0);
    excess[s] = inf;
    for (int i = 0; i < n; i++) {
        if (i != s) push(s, i);
    }

    vector<int> current;
    while (!(current = find_max_height_vertices(s, t)).empty()) {
        for (int i : current) {
            bool pushed = false;
            for (int j = 0; j < n && excess[i]; j++) {
                if (capacity[i][j] - flow[i][j] > 0 && height[i] ==
                    height[j] + 1) {
                    push(i, j);
                    pushed = true;
                }
            }
        }
        if (!pushed) {
            relabel(i);
            break;
        }
    }

    return excess[t];
}

```