## HOMEWORK I (Part I)

Due day: 9:00am Oct.11 (Thursday), 2012

This homework is to let you familiar with tools and Verilog language. It includes A) basic exercises and B) the first part of a simplified multi-cycle (**SMILE**) CPU with only 17 instructions shown in the last page. Part B is to let the first-time CPU designer be familiar with the instruction set architecture, dataflow modeling and controller design. Some problems have reference code. They are for your reference only. Most likely, you need to modify them to fit the problem specification.

# General rules for deliverables

- This homework needs to be completed by INDIVIDUAL student.
- Compress all files described in the problem statements into one zip or rar.
- Submit the compress file to the course website before the due day. Warning!

  AVOID submit in the last minute. Late submission is not accepted.

# **Grading Notes**

- Important! DO remember to include your Verilog code. NO code, NO grades. Also, if your code can not be recompiled by TA successfully <u>using</u> tools in SoC Lab, you receive NO credit.
- Write your report seriously and professionally. Incomplete description and information will reduce your chances to get more credits.
- If extra works (like synthesis, post-simulation or additional instructions) are done, please describe them in your final report clearly for bonus points.
- Please follow course policy.

# **Deliverables**

- 1. All CPU Verilog codes including components, testbenches and machine code for each lab exercise. NOTE: Please DO NOT include source code in the report!
- 2. A homework report that includes
  - a. A summary in the beginning to state what has been done (such as SMILE CPU, synthesis, post-synthesis simulation, additional branch instruction with verification);
  - b. Simulated waveforms with proper explanation;
  - c. Report requirements in each assignment;

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- d. Learned lessons (skills learned, where do you stuck most and what do you find, exciting things etc.)
- 3. Please write in MS word using the Homework report template and follow the convention for the file name of your report: n26984795\_ 李 大 鳴 hw1p1 report.doc

# 1. (10/100) Write and verify Verilog a 16-bit performance counter by using two universal 8-bit counter/timer with the following specification:

Universal 8-bit counter/timer Design Requirements

- a. Number format: unsigned
- b. Inputs: clk (1 bit), \_areset (1 bit), \_load (1 bit), preld-val (8 bit), \_updown (1 bit), \_wrapstop (1 bit)
- c. Outputs: dcout(8 bit), overflow(1 bit),
- d. The counter is triggered by the clock signal (*clk*).
- e. \_areset =1 makes the counter set to zero asynchronously, while \_aset =1 let the counter set to its maximum value asynchronously. If none of them are high, the counter is operated regularly.
- f. If the signal *load* is active (high), the preload value (*preld-val*) will be loaded, otherwise, no effect to the counter.
- g. The priority of the control signal is \_areset > \_aset > \_load. This is to avoid un-intentionally activate these three signals at the same time.
- h. \_updown indicates whether the counter is counted upward (when high) or downward (when low).
- i. \_wrapstop is used to set the counting behavior. If \_wrapstop = 1, the counter will wrap around. (Ex. If the counting is upward, the counter will go back to zero and then recount again.) If \_wrapstop = 0, the counter will stop and issue overflow.
- j. If the result is out of range, overflow will be flagged to logic 1, otherwise, 0.
- k. You may add other signal additional pins and function as long as you clearly specify and explain.
- 1. Hierarchy design is allowed; however, they must be specified.
- m. The design shall be implemented by using structural style and be able to identified major components.

# Report Requirements

a. Proper explanation of your design is required for full credits.

# HOMEWORK I (Part I)

- b. A figure (block diagram with logic gates) shall be draw to depict your design in the end.
- c. Verify your code with a testbench. This testbench needs to show all major situations. Show your snapshot of waveform for different cases in your reports and illustrate the correctness of your results.
- d. Synthesize this counter and re-run testbenches on the synthesized code and observe your waveforms. Provide a snapshot of the waveform and pin point the location of dubious errors, the synthesized block diagram, gates that are actually used by the synthesizer and report 5 most frequent warning/errors during synthesis.
- e. Use nLint to analyze your code, report the final results and 3~5 most frequent warning/errors in your code.

## HOMEWORK I (Part I)

- 2. (20/100) Write and verify a 32-bit ALU that could perform the following arithmetic and logic operations:
  - Logic operations: AND, OR, XOR, NOP
  - Arithmetic operations: ADD, SUB
  - Shift & Rotate operations: SLLI (Shift Left Immediate), SRLI (Shift Right Immediate) and ROTRI (Rotate Right Immediate)
  - Instruction set architecture could be found on Page 8-9.

The ALU unit shown in Fig. 2 has two data inputs (SCR1 and SCR2) and one data output (ALU\_result) along with one output flag, ALU\_Overflow. Controlling signals are *opcode*, *sub\_opcode* and Enable\_execute. *opcode* and *sub\_opcode* will be used to select which operation to be executed

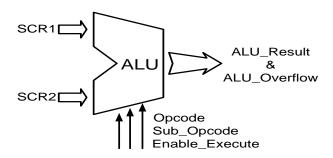


Fig. 2 Conceptual Diagram of the ALU

This ALU consists of three different combinational circuit building blocks:

- Add/Sub to perform addition and subtraction arithmetic operations
- Logic block to perform logical operations
- Barrel Shifter to perform shift and rotate operations

# **Design Requirements**

- a. This ALU shall be designed such that it could be compatible with the given ISA and instruction format.
- b. The ALU will be used in subsequent Homework 1-Part 2. So you shall be aware that more instructions and mechanisms will be added later. When design, please give room for extension.
- c. A skeleton code is given for quick startup. However, it still needs to be modified to fully work.
- d. The ALU needs to be verified with proper testbenches for each instruction.

## HOMEWORK I (Part I)

# Report Requirements

- a. Proper explanation of your design is required for full credits.
- b. A figure (block diagram with logic gates) shall be draw to depict your design in the end.
- c. Verify your code with a testbench before synthesis. This testbench needs to show all major situations. Show your snapshot of waveform for different cases in your reports and illustrate the correctness of your results.
- d. Synthesize this ALU and re-run testbenches on the synthesized code and observe your waveforms. Provide a snapshot of the waveform and pin point the location of dubious errors, the synthesized block diagram, gates that are actually used by the synthesizer and report 5 most frequent warning/errors during synthesis.
- e. Use nLint to analyze your code, report the final results and 3~5 most frequent warning/errors in your code.

Its major portion of the Verilog may look like as follows.

```
Add/Sub block (add sub module in reference code1 for ALU):
    assign result=(control == 0)?(operand2+operand1):(operand2-operand1):
Logic block
               and result=src2&src1;
               or_result=src2|src1;
               xor_result=src2^src1;
Shift block (Barrel Shifter module in reference code1 for ALU):
output [31:0] result; // result operand
reg [95:0] rotate reg;
input [31:0] in1,in2;// input operand; in1: # of shifts; in2: original data
rotate_reg={in2,in2,in2};
     Rotate Operation:
               rotate_reg=(direction)?rotate_reg>>in1:rotate_reg<<in1:
               result=rotate_reg[63:32];
     Shift Operation:
               result=(direction)?in2>>in1:in2<<in1;
```

## HOMEWORK I (Part I)

# Reference code for ALU

```
module alu(alu_result,alu_overflow,scr1,scr2,opcode,sub_opcode,enable_execute,reset);
 parameter NOP=5'b01001,ADD=5'b00000,SUB=5'b00001,AND=5'b00010,
            OR=5'b00100,XOR=5'b00011,SRLI=5'b01001,SLLI=5'b01000,
            ROTRI=5'b01011:
 output reg [31:0]alu_result;
 output reg alu_overflow;
 input [31:0]scr1,scr2;
 input [5:0]opcode;
 input [4:0]sub_opcode;
  input reset;
 input enable_execute;
 reg [63:0]rotate;
 reg a,b;
 always @ ( * )begin
    if(reset)begin
      alu_result=32'b0;
      alu_overflow=1'b0;
   else if(enable_execute)begin
      case(opcode)
        6'b100000 : case (sub_opcode)
                      NOP
                            : begin
                                alu_result=32'b0;
                                alu_overflow=1'b0;
                              end
                      ADD
                            : begin
                                 {a,alu_result[30:0]}=scr1[30:0]+scr2[30:0];
                                 {b,alu_result[31]}=scr1[31]+scr2[31]+a;
                                alu overflow=a ^ b;
                              end
                      SUB
                            : begin
                                 {a,alu_result[30:0]}=scr1[30:0]-scr2[30:0];
                                 {b,alu_result[31]}=scr1[31]-scr2[31]-a;
                                alu_overflow=a ^ b;
                              end
```

## HOMEWORK I (Part I)

```
AND
                        : begin
                             alu_overflow=1'b0;
                             alu_result=scr1&scr2;
                           end
                  OR
                        : begin
                             alu_overflow=1'b0;
                             alu_result=scr1 scr2;
                           end
                  XOR
                        : begin
                             alu_overflow=1'b0;
                             alu_result=scr1^scr2;
                  SRLI : begin
                             alu_overflow=1'b0;
                             alu_result=scr1>>scr2;
                           end
                  SLLI : begin
                             alu_overflow=1'b0;
                             alu_result=scr1<<scr2;</pre>
                           end
                  ROTRI : begin
                             alu_overflow=1'b0;
                             rotate={scr1,scr1}>>scr2;
                             alu_result=rotate[31:0];
                           end
                  default : begin
                               alu_overflow=1'b0;
                               alu_result=32'b0;
                             end
                endcase
    6'b101000 : begin
                  {a,alu_result[30:0]}=scr1[30:0]+scr2[30:0];
                  {b,alu_result[31]}=scr1[31]+scr2[31]+a;
                  alu_overflow=a ^ b;
                end
    6'b101100 : begin
                  alu_overflow=1'b0;
                  alu_result=scr1|scr2;
                end
    6'b101011 : begin
                  alu_overflow=1'b0;
                  alu_result=scr1^scr2;
    default : begin
                alu overflow=1'b0;
                alu_result=32'b0;
              end
  endcase
else begin
  alu_result=32'b0;
  alu_overflow=1'b0;
```

end

end end endmodule

# HOMEWORK I (Part I)

# SMILE Processor Instruction Set Architecture (ISA)

We will follow ISA in this homework and for the reset of homeworks. Please read the ISA carefully. You design needs to comply with this ISA. It is NOTEd that you SHALL NOT revise by yourself. If you have any questions, please talk to TA first.

# **BASICS**

OP code

6bits

\*Data-processing (SE:sign-extened, ZE:zero-extened)

\*\*\* Note that NOP and SRLI has the same OP + sub OP. This is durable. If you exam their DST, SRC1 and SRC2, you will be able to find their differences and find a proper way to implement.

31 30	25 2	24	20 19	15	14	10	9	5	4		0
0 0	OP code	Rt		Ra	Rb or in	nm	0			Sub_OP	
1bits	6bits	5bits		5bits	5bits	S	5bits			5bits	
OP code	Mnemoni	cs DST	SRC1	SRC2	Sub OP		Desc	ript	ion		
100000	NOP	00000	00000	00000	01001		No o	perat	ion		
100000	ADD	\$Rt	\$Ra	\$Rb	00000		Rt = 1	Ra +	Rb		
100000	SUB	\$Rt	\$Ra	\$Rb	00001		Rt =	Ra -	Rb		
100000	AND	\$Rt	\$Ra	\$Rb	00010		Rt = 1	Ra &	Rb		
100000	OR	\$Rt	\$Ra	\$Rb	00100		Rt =	Ra	Rb		
100000	XOR	\$Rt	\$Ra	\$Rb	00011		Rt =	Ra ^	Rb		
100000	SRLI	\$Rt	\$Ra	imm	01001	Sh	ift right : I	Rt = 1	Ra>	>> imm	
100000	SLLI	\$Rt	\$Ra	imm	01000	Sl	nift left : R	t = F	a <	< imm	
100000	ROTRI	\$Rt	\$Ra	imm	01011	Rot	ate right :	Rt =	Ra	>> imm	
31 30	25 2	24	20 19	15	14						0
0 0	OP code	Rt		Ra			immediat	te			
1bits	6bits	5bits		5bits			15bits				
OP code	Mnemoni	cs DST	SRC1	SRC2		l	Descripti	on			
101000	ADDI	\$Rt	\$Ra	imm		Rt =	= Ra + SE	(imm	)		
101100	ORI	\$Rt	\$Ra	imm		Rt	= Ra   ZE(	imm	)		
101011	XORI	\$Rt	\$Ra	imm		Rt =	= Ra ^ ZE(	(imm	)		
31 30	25 2	24	20 19								0
							-				$\overline{}$

5bits

immediate

20bits

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OP code	Mnemonics	DST	SRC1	Description
100010	MOVI	\$Rt	imm	Rt = SE(immediate)

# SMILE CPU has the following instruction format (Andes ISA compatible)

The Andes 32bit instruction formats and the meaning of each filed are described below:

# Type-0 Instruction Format

	0	Opc_6	{sub_1, imm_24}						
>		Type-1 Inst	ruction Form	mat					
	0	Opc 6	rt 5	imm 20					

0	Opc_6	rt_5	imm_20				
0	Opc_6	rt_5	sub_4	imm_16			

Type-2 Instruction Format

0	Opc_6	rt_5	ra_5	imm_15
0	Opc_6	rt_5	ra_5	{sub_1, imm_14}

Type-3 Instruction Format

0	Opc_6	rt_5	ra_5	rb_5	sub_10
0	Opc_6	rt_5	ra_5	imm_5	sub_10

> Type-4 Instruction Format

0	Opc_6	rt_5	ra_5	rb_5	rd_5	sub_5
0	Opc_6	rt_5	ra_5	imm1_5	imm2_5	sub_5

- opc\_6: 6-bit opcode
- ◆ rt\_5: target register in 5-bit index register set
- ◆ ra\_5: source register in 5-bit index register set
- ◆ rb\_5: source register in 5-bit index register set
- ◆ rd\_5: destination register in 5-bit index register set
- ◆ sub\_10: 10-bit sub-opcode
- ◆ sub\_5: 5-bit sub-opcode
- ◆ sub\_4: 4-bit sub-opcode
- ◆ sub\_1: 1-bit sub-opcode
- imm\_24: 24-bit immediate value, for unconditional jump instructions (J, JAL). The
- ◆ PC[31:0] = {current PC[31:25], imm\_24, 1'b0}
- imm\_20: 20-bit immediate value. Sign-extended to 32-bit for MOVI operations.
- imm\_16: signed PC relative address displacement for branch instructions.
- imm\_15: 15-bit immediate value. Zero extended to 32-bit for unsigned operations,
- while sign extended to 32-bit for signed operations.
- ♦ imm\_14: signed PC relative address displacement for branch instructions.
- imm\_5, imm1\_5, imm2\_5: 5-bit unsigned count value or index value