Module 04: Graph Algorithms Analysis and Design of Algorithms

Ammar Sherif

Nile University



- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- 3 Weighted Graphs
- 4 Path Construction

- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- Weighted Graphs
- 4 Path Construction

Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

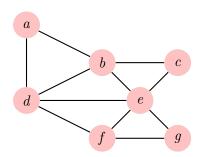
Graphs & Graph Algorithms

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Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

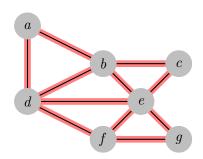
Nodes/Vertices



Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

- Nodes/Vertices
- Edges

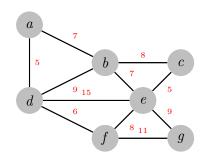




Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

- Nodes/Vertices
- Edges
- Weighted Graph





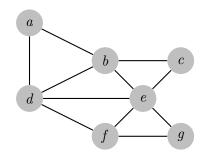
Graphs & Graph Algorithms

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Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

- Nodes/Vertices
- Edges
- Weighted Graph
- **Unweighted Graph**

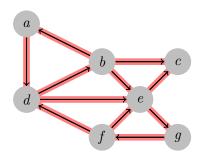




Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

- Nodes/Vertices
- Edges
- Weighted Graph
- Unweighted Graph
- Directed Graph

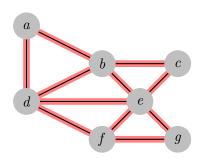




Graphs

From a mathematical perspective, consist of a nodes/vertices and edges.

- Nodes/Vertices
- Edges
- Weighted Graph
- Unweighted Graph
- Directed Graph
- Undirected Graph





Why Graphs & Graph Algorithms?

 Shortest Path and Route planning



- Shortest Path and Route planning
 - Robotics

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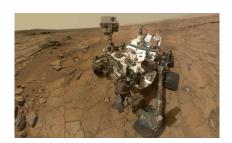
- Shortest Path and Route planning
- Robotics
 - Warehouses



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Why Graphs & Graph Algorithms?

- Shortest Path and Route planning
- Robotics
 - Warehouses
 - Space Robots



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Why Graphs & Graph Algorithms?

- Shortest Path and Route planning
- Robotics
 - Warehouses
 - Space Robots
 - Rescue Robots

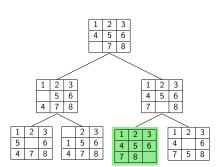


- Shortest Path and Route planning
- Robotics

Graphs & Graph Algorithms

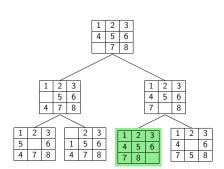
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- Warehouses
- Space Robots
- Rescue Robots
- Games



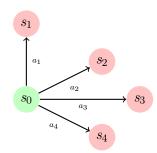
Why Graphs & Graph Algorithms?

- Shortest Path and Route planning
- Robotics
 - Warehouses
 - Space Robots
 - Rescue Robots
- Games
- Optimization **Problems**



Why Graphs & Graph Algorithms?

- Shortest Path and Route planning
- Robotics
 - Warehouses
 - Space Robots
 - Rescue Robots
- Games
- Optimization **Problems**
- Any decision-based problem



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Algorithms are implemented via programming, so how to represent graphs?

What our representation should provide?

Algorithms are implemented via programming, so how to represent graphs?

What our representation should provide?

Know the neighbors

Algorithms are implemented via programming, so how to represent graphs?

What our representation should provide?

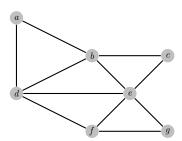
- Know the neighbors
- Weights of links

Algorithms are implemented via programming, so how to represent graphs?

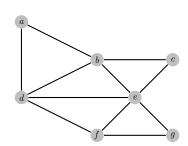
What our representation should provide?

- Know the neighbors
- Weights of links
- list of nodes/edges

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Adjacency Matrix

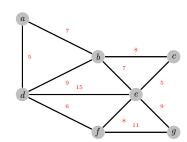


	a	b	c	d	e	\int	g
a	_	1		1			
\overline{b}	1	_	1	1	1		
\overline{c}		1	_		1		
\overline{d}	1	1		_	1	1	
\overline{e}		1	1	1	_	1	1
\overline{f}				1	1	_	1
g					1	1	_

Graphs & Graph Algorithms

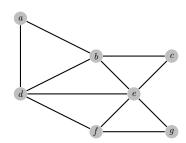
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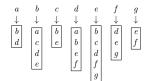
Adjacency Matrix



	a	b	c	d	e	f	g
\overline{a}	-	7		5			
\overline{b}	7	_	8	9	7		
c		8	_		5		
\overline{d}	5	9		_	15	8	
e		7	5	15	_	8	9
\overline{f}				8	8	_	11
\overline{g}					9	11	_

- Adjacency Matrix
- Adjacency List

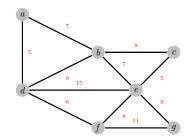


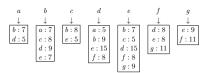


Graphs & Graph Algorithms

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- Adjacency Matrix
- Adjacency List





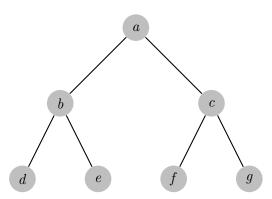
- Graphs & Graph Algorithms
- 2 Unweighted Graphs

Depth First Search (DFS)
Breadth First Search (BFS)

- 3 Weighted Graphs
- 4 Path Construction

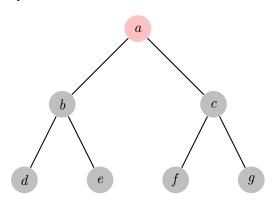
- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs Depth First Search (DFS) Breadth First Search (BFS)
- 3 Weighted Graphs
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Depth has the max priority

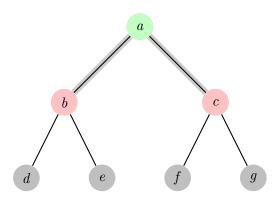


Graphs & Graph Algorithms

Depth has the max priority

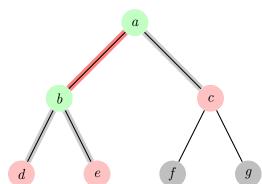


Depth has the max priority



Depth has the max priority Child \leftarrow Parent

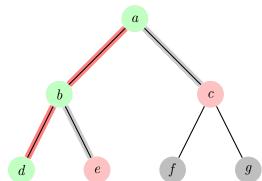




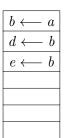
Graphs & Graph Algorithms

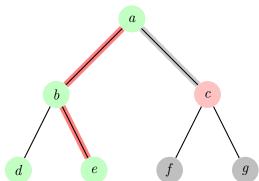
Depth has the max priority Child ← Parent





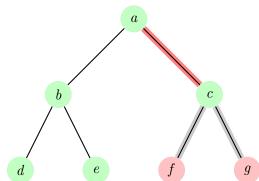
 $\begin{array}{c} \textit{Depth} \text{ has the max priority} \\ \text{Child} \longleftarrow \text{Parent} \end{array}$





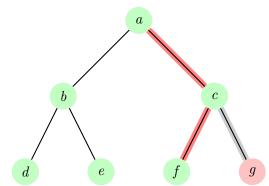
Depth has the max priority Child \leftarrow Parent

$b \longleftarrow$	a
$d \longleftarrow$	b
$e \longleftarrow$	b
$c \longleftarrow$	a



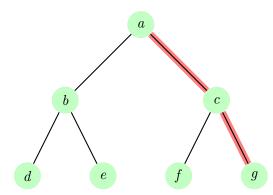
 $\begin{array}{c} \textit{Depth} \text{ has the max priority} \\ \text{Child} \longleftarrow \text{Parent} \end{array}$

$b \longleftarrow$	a
$d \longleftarrow$	b
$e \longleftarrow$	b
$c \longleftarrow$	\overline{a}
$f \longleftarrow$	c

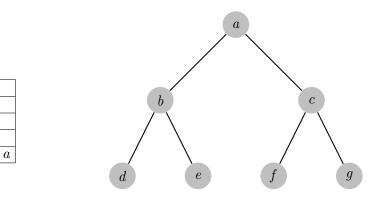


Therefore, the below table summarizes how did we get to any node through our traversal Child — Parent

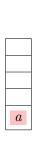
$b \longleftarrow$	a
$d \leftarrow$	b
$e \leftarrow$	b
$c \leftarrow$	\overline{a}
$f \leftarrow$	\overline{c}
$g \leftarrow$	\overline{c}

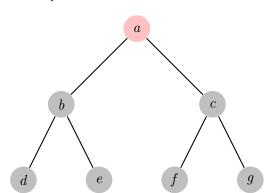


Let's do it again, to notice the pattern of nodes to be visited



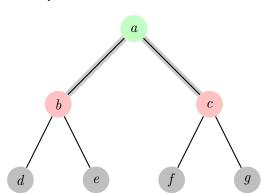
Let's do it again, to notice the pattern of nodes to be visited





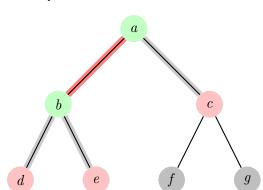
Let's do it again, to notice the pattern of nodes to be visited



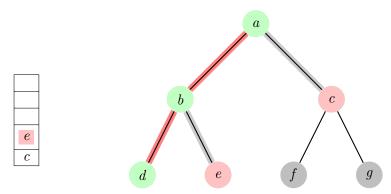


Let's do it again, to notice the pattern of nodes to be visited





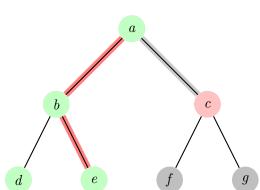
Let's do it again, to notice the pattern of nodes to be visited



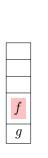
Graphs & Graph Algorithms

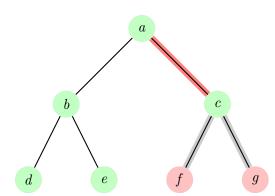
Did you get the pattern of nodes to be visited?



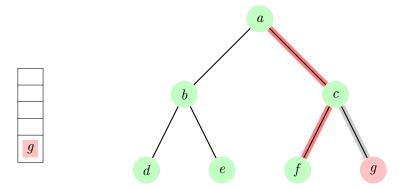


Did you get the pattern of nodes to be visited? **Last** inserted element is **first** to explore.





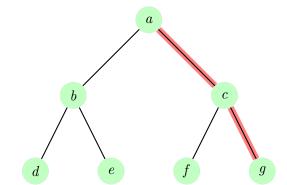
Did you notice what might be this structure? Hint: Last-in First-out



Did you notice what might be this structure? Hint: Last-in First-out

Yes, it is Stack

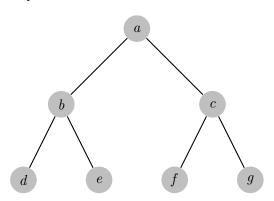


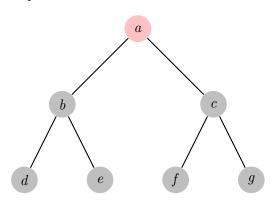


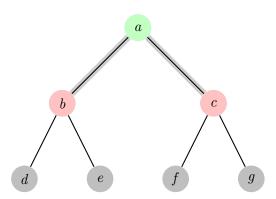
Algorithm 1: Depth-First(root)

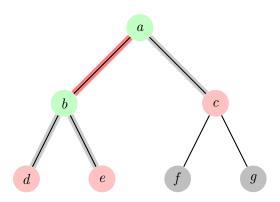
```
def S to be Stack:
visited \leftarrow \{\};
S.\mathtt{push}(root);
while S \neq \phi do
    node \leftarrow S.pop();
    if node \notin visited then
         visited \leftarrow visited \cup \{node\};
         for n \in adjacent(node) do
             S.\mathtt{push}(n);
         end
    end
end
```

- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs Depth First Search (DFS) Breadth First Search (BFS)
- 3 Weighted Graphs
- 4 Path Construction

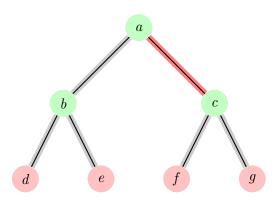


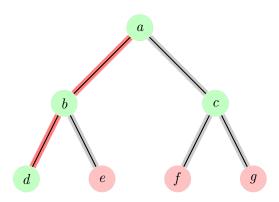




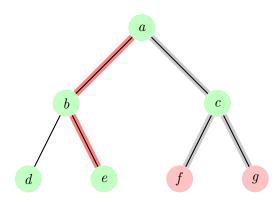


Breadth has the max priority

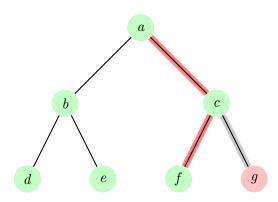


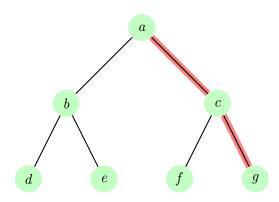


Breadth has the max priority



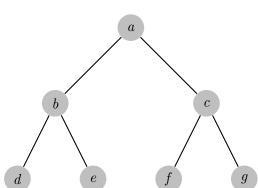
Breadth has the max priority



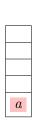


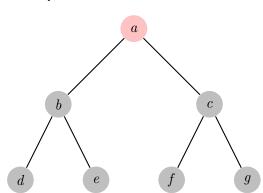
Let's do it again, to notice the pattern of nodes to be visited



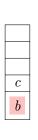


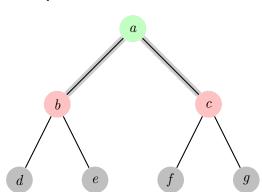
Let's do it again, to notice the pattern of nodes to be visited



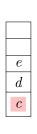


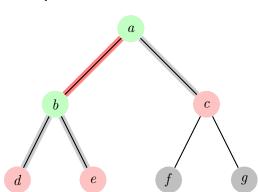
Let's do it again, to notice the pattern of nodes to be visited



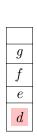


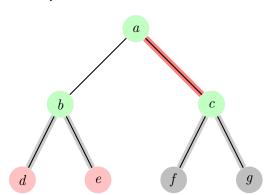
Let's do it again, to notice the pattern of nodes to be visited



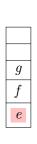


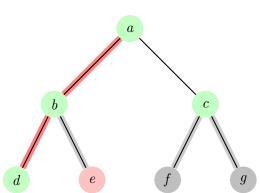
Let's do it again, to notice the pattern of nodes to be visited





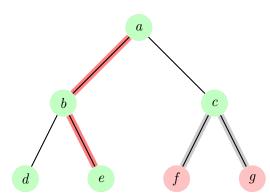
Did you get the pattern of nodes to be visited?



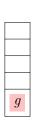


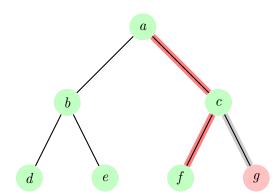
Did you get the pattern of nodes to be visited? **First** inserted element is **first** to explore.





Did you notice what might be this structure? Hint: First-in First-out

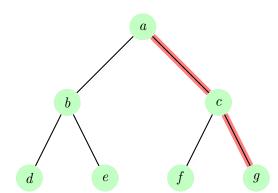




Did you notice what might be this structure? Hint: First-in First-out

Yes, it is Queue





Algorithm 2: Breadth-First(root)

```
def S to be Queue:
visited \leftarrow \{\};
S.\mathtt{enqueue}(root);
while S \neq \phi do
    node \leftarrow S.dequeue();
    if node \notin visited then
        visited \leftarrow visited \cup \{node\};
        for n \in adjacent(node) do
             S.\mathtt{enqueue}(n);
        end
    end
end
```

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- 2 Unweighted Graphs
- 3 Weighted Graphs

Branch and Bound
Search Space Pruning
Branch and Bound + Visited List
Branch and Bound + Heuristics
A* Algorithm
Heuristic Design

4 Path Construction



- 1 Graphs & Graph Algorithms
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- Weighted Graphs
 Branch and Bound
 Search Space Pruning
 Branch and Bound + Visited List
 Branch and Bound + Heuristics
 - A* Algorithm
 Heuristic Design
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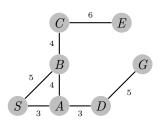
Branch and Bound

Right now we have weights, now what should we prioritize?

Branch and Bound

Graphs & Graph Algorithms

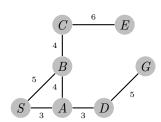
Right now we have weights, now what should we prioritize? Min/Max weights



Right now we have weights, now what should we prioritize? Min/Max weights







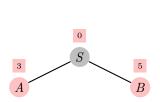
S:0

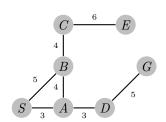


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Graphs & Graph Algorithms

Right now we have weights, now what should we prioritize? Min/Max weights





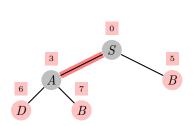
iterations/visits: 1

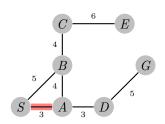
 $A:3 \; B:5$

14 / 26

Graphs & Graph Algorithms

Right now we have weights, now what should we prioritize? Min/Max weights





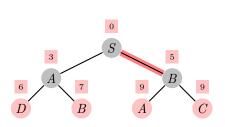
iterations/visits: 2

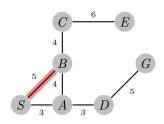
$$B:5 \ D:6 \ B:7$$

14 / 26

Graphs & Graph Algorithms

Can you notice the pattern of the structure?

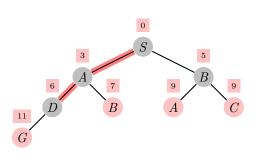


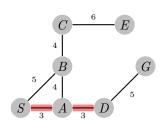


iterations/visits: 3

D:6 B:7 A:9 C:9

Can you notice the pattern of the structure? Also, we found G, so shall we stop?

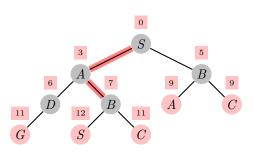


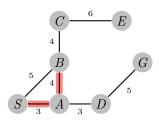


iterations/visits: 4

 $B:7 \mid A:9 \mid C:9 \mid G:11$

Have you noticed what happened?

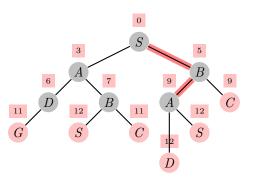


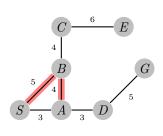


iterations/visits: 5

A:9 C:9 C:11 G:11 S:12

Now, what is such structure?

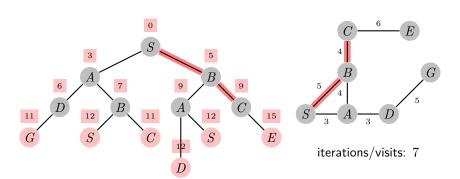




iterations/visits: 6

C:9 C:11 G:11 S:12 D:12 S:12

Now, what is such structure?

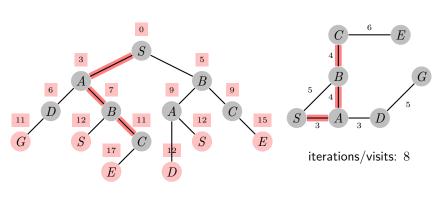


 $C: 11 \mid G: 11 \mid S: 12 \mid D: 12 \mid S: 12 \mid E: 15$



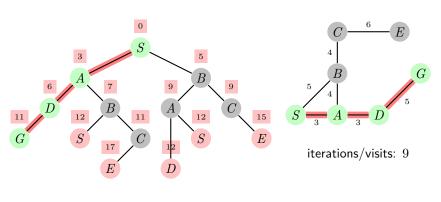
Graphs & Graph Algorithms

Now, what is such structure?



 $G: 11 \mid S: 12 \mid D: 12 \mid S: 12 \mid E: 15 \mid E: 17$

Yes, it is a **Priority Queue**, where the priority is the overall path cost.



 $S: 12 \mid D: 12 \mid S: 12 \mid E: 15 \mid E: 17$

Algorithm 3: Branch-Bound(root, goal)

- **1** Graphs & Graph Algorithms
- 2 Unweighted Graphs

Graphs & Graph Algorithms

3 Weighted Graphs

Search Space Pruning

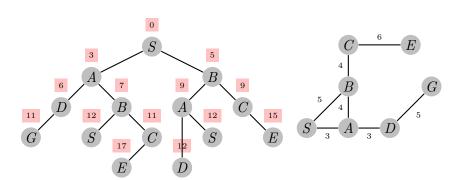
Branch and Bound + Visited List A* Algorithm

Path Construction

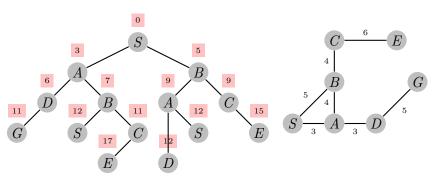
Graphs & Graph Algorithms

One question over the previous algorithm is: cannot we do better?

To answer this, we need to check the produced **search space**, on the left.

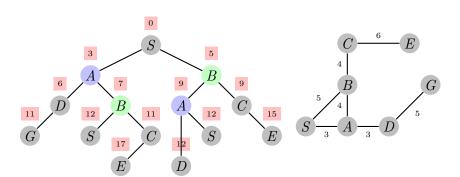


We know the **denser** the search space, more nodes to visit, the more time our algorithms takes. Hence, **pruning** it, eliminating some of the nodes, should improve it, how can we do this?

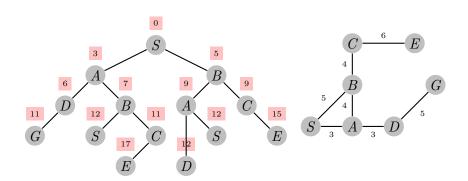


Ammar Sherif

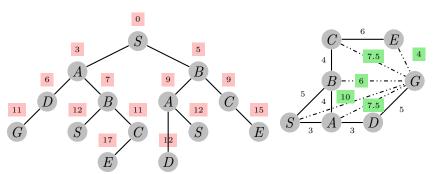
One way is to remove duplicates; do you think we need to check ${\cal A}$ or ${\cal B}$ twice? We can do this using a simple **visited list**



Could you think of other techniques to prune the tree?

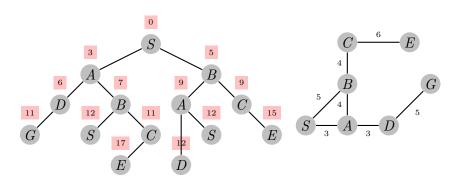


One solution is if we could have an estimate of where the goal is; this may make us more oriented towards searching in particular regions than others; this is briefly, the **heuristics**.



Ammar Sherif

We are covering both the **visited list** and **heuristic** techniques in the next subsections, applying them to branch and bound.



- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- 3 Weighted Graphs

Branch and Bound
Search Space Pruning

 $Branch\ and\ Bound\ +\ Visited\ List$

Branch and Bound + Heuristics

A* Algorithm

Heuristic Design

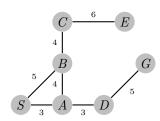
Path Construction



Let's follow on the same graph, but while using a **visited list**.



S

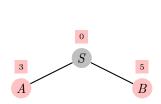


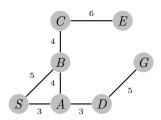
$$PQ$$
: $S:0$

wisited:

Graphs & Graph Algorithms

Let's follow on the same graph, but while using a **visited list**.



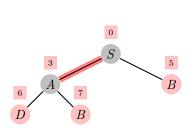


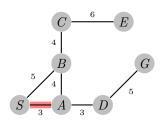
iterations/visits: 1

$$PQ$$
: $A:3$ $B:5$ $visited$: S

Graphs & Graph Algorithms

Let's follow on the same graph, but while using a visited list.



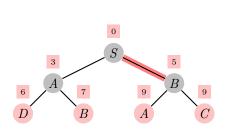


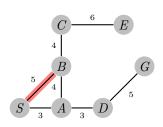
iterations/visits: 2

$$PQ$$
: $B:5$ $D:6$ $B:7$ $visited$: S A

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Let's follow on the same graph, but while using a visited list.



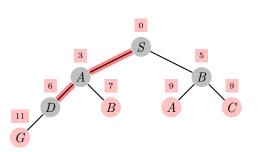


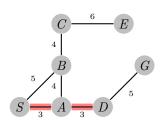
iterations/visits: 3

$$PQ:$$
 $D:6$ $B:7$ $A:9$ $C:9$ $visited:$ S A B

18 / 26

Now, the queue has B, should we visit it next?

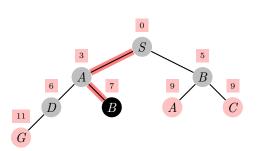


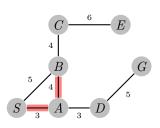


iterations/visits: 4

$$PQ:$$
 $B:7 \mid A:9 \mid C:9 \mid G:11$ $visited:$ $\mid S \mid A \mid B \mid D \mid$

Well, no: B is already visited, so there is already a shorter path to it.



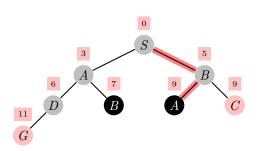


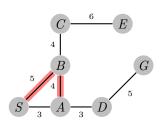
iterations/visits: 5

PQ: A:9 C:9 G:11 visited: S A B D

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Again, the next was A, so it is blocked as well.





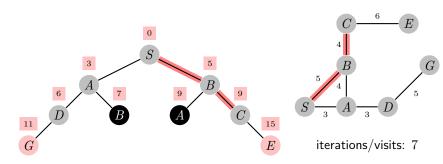
iterations/visits: 6

PQ: C:9 G:11 visited: S A B D

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Graphs & Graph Algorithms

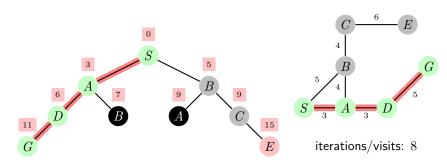
We, then, proceed



PQ: $G: 11 \mid E: 15 \mid$ $S \mid A \mid B \mid D \mid C$

18 / 26

Finally, we got our solution. Could you see the pruning effect?



PQ: E:15visited:

18 / 26

Graphs & Graph Algorithms

Algorithm 4: Branch-Bound-Visited(root, goal)

```
def PQ to be Priority Queue
visited \leftarrow \{\}
PQ.\mathtt{enqueue}(root, 0)
node \leftarrow root
while PQ \neq \phi \land node \neq qoal do
    node, path cost \leftarrow PQ.dequeue()
    if node∉ visited then
         for n \in adjacent(node) do
             // loop over the links and their costs
              PQ.\mathtt{enqueue}(n, path\_cost + \mathtt{cost}(n))
         end
    end
end
```

- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- 3 Weighted Graphs

Search Space Pruning

Branch and Bound + Visited List

Branch and Bound + Heuristics

A* Algorithm Heuristic Design

Path Construction



Now, we start with showing what heuristics are, first. Heuristics are **optimistic estimation** to the future cost.

Graphs & Graph Algorithms

Now, we start with showing what heuristics are, first. Heuristics are **optimistic estimation** to the future cost. Although they are not accurate, but because they capture some of the information, we use them within our algorithms to guide our search to more promising areas than others, leading to faster findings.

We will be using the same map, where our heuristic values, in dashed lines, represent the direct Euclidean distance between the nodes. They are not accurate because shorter distances does not necessarily mean we are close to our goal.

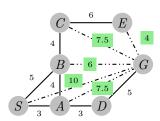
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Graphs & Graph Algorithms

Our cost, priority, = actual cost to current node + the estimate to the goal. Let's do it without a visited list.







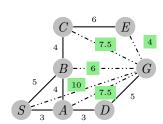
$$PQ$$
: $S:10$



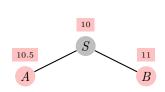
Graphs & Graph Algorithms

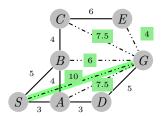
Note our initial cost = 10, which is 0 actual cost +10 as estimate





Again, the cost = true cost + heuristic



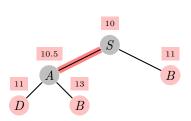


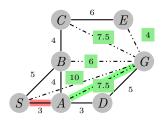
iterations/visits: 1

 $A:10.5 \ B:11$ PQ:

Graphs & Graph Algorithms

Again, the cost = true cost + heuristic



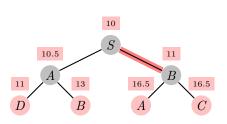


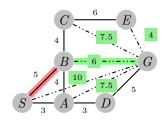
iterations/visits: 2

 $PQ: \quad B: 11 \mid D: 11 \mid B: 13$

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Again, the cost = true cost + heuristic



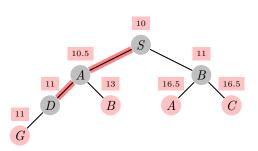


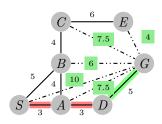
iterations/visits: 3

PQ: $D: 11 \mid B: 13 \mid A: 16.5 \mid C: 16.5$

Graphs & Graph Algorithms

Again, the cost = true cost + heuristic



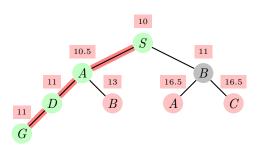


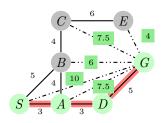
iterations/visits: 4

PQ: $G: 11 \mid B: 13 \mid A: 16.5 \mid C: 16.5 \mid$

Graphs & Graph Algorithms

and we got it in just 5 iterations.





iterations/visits: 5

B:13 A:16.5 C:16.5PQ:

Algorithm 5: Branch-Bound-Heuristics (root, goal)

```
def PQ to be Priority Queue
PQ.enqueue(root, h(root))
node \leftarrow root
while PQ \neq \phi \land node \neq qoal do
    node, node_cost \leftarrow PQ.dequeue()
    path cost \leftarrow node cost - h(node)
    for n \in adjacent(node) do
        // h(n) is the heuristic estimate from n to goal
        total cost \leftarrow path cost + cost(n) + h(n)
        PQ.enqueue(n, total cost)
    end
```

end

- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- 3 Weighted Graphs

Branch and Bound
Search Space Pruning
Branch and Bound + Visited List
Branch and Bound + Heuristics

A* Algorithm

Heuristic Design

Path Construction

A* Algorithm

We studied two cool search pruning techniques:

- visited list
- heuristics

Now, you might ask what happens if we used both of these cool techniques.

Graphs & Graph Algorithms

We studied two cool search pruning techniques:

- visited list
- heuristics

Now, you might ask what happens if we used both of these cool techniques. We get an even more awesome algorithm, which is the A*. It is typically using both techniques with the Branch-and-Bound algorithm. Let's check the pseudo-code

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A* Algorithm

Algorithm 6: A-STAR(root, goal)

```
visited \leftarrow \{\}, node \leftarrow root
PQ.enqueue(root, h(root))
while PQ \neq \phi \land node \neq qoal do
    node, node_cost \leftarrow PQ.dequeue()
   if node∉ visited then
        visited \leftarrow visited \cup \{node\}
        path cost \leftarrow node cost - h(node)
        for n \in adjacent(node) do
            PQ.enqueue(n, path\_cost + cost(n) + h(n))
        end
   end
end
```

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- 1 Graphs & Graph Algorithms
- 2 Unweighted Graphs
- 3 Weighted Graphs

Search and Bound
Search Space Pruning
Branch and Bound + Visited List
Branch and Bound + Heuristics
A* Algorithm

Heuristic Design

4 Path Construction

- **1** Graphs & Graph Algorithms
- 2 Unweighted Graphs

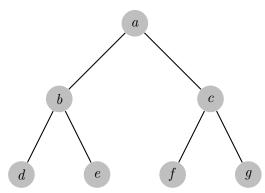
Graphs & Graph Algorithms

- Weighted Graphs
- Path Construction

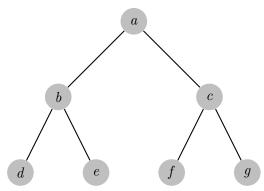
Cool; now, we have shown how many traversing algorithms, *which visits the nodes*, and checked the pseudo-code, but and important question is

how to construct a path from a traversing algorithm? (DFS for example)

First, to differentiate between traversing and path, we revisit the same DFS tree as before; the ordering of the visited nodes is a, b, d, e, c, f, g.

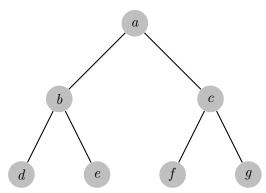


the ordering of the visited nodes is a,b,d,e,c,f,g. Does this mean that the **path** according to our algorithm from $a \to f$ is a,b,d,e,c,f?

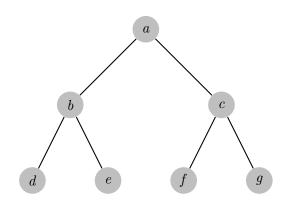


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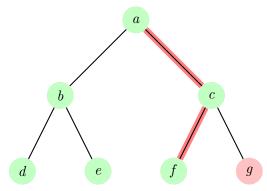
the ordering of the visited nodes is a,b,d,e,c,f,g. Does this mean that the **path** according to our algorithm from $a \to f$ is a,b,d,e,c,f? well, **No**. This is just a traversing order.



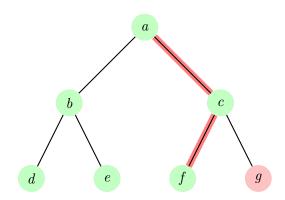
In that case, what is the path constructed from $a \to f$ according to our algorithm?



In that case, what is the path constructed from $a \to f$ according to our algorithm? It is a, c, f as we see from our previous execution.

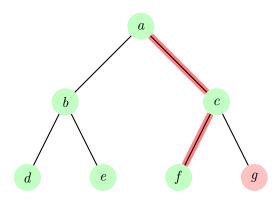


so, how could we construct such path from our algorithm?



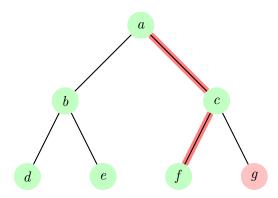
so, how could we construct such path from our algorithm?

• store the whole paths, instead of just nodes



so, how could we construct such path from our algorithm?

• store the whole paths, instead of just nodes; Storage hungry

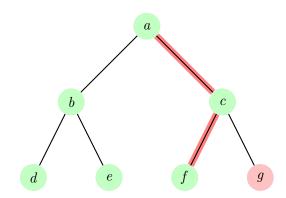


so, how could we construct such path from our algorithm?

using parent-child structure.

Child \leftarrow Parent

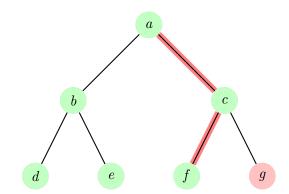
$b \longleftarrow$	a
$d \leftarrow$	b
$e \leftarrow$	b
$c \longleftarrow$	a
$f \leftarrow$	c
$g \longleftarrow$	c



start from your goal, f, and move backward until getting your start.

Child ← Parent

$b \longleftarrow$	a
$d \leftarrow$	b
$e \leftarrow$	b
$c \leftarrow$	a
$f \leftarrow$	c
$g \longleftarrow$	c



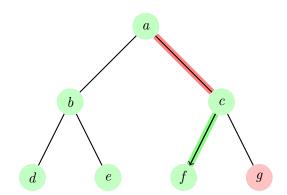
Path: $c \rightarrow f$

Graphs & Graph Algorithms

Child \leftarrow Parent

$$\begin{array}{c}
b \longleftarrow a \\
d \longleftarrow b \\
e \longleftarrow b \\
c \longleftarrow a
\end{array}$$

$$f \longleftarrow c$$

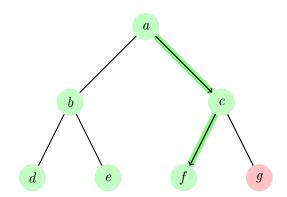


Path:
$$a \to c \to f$$

Child ← Parent

Graphs & Graph Algorithms

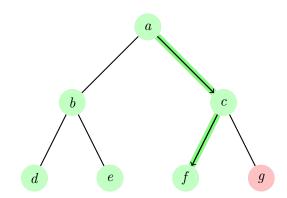
$$\begin{array}{c} b \longleftarrow a \\ d \longleftarrow b \\ e \longleftarrow b \\ \hline c \longleftarrow a \\ f \longleftarrow c \\ g \longleftarrow c \end{array}$$



We can this technique for all the traversing algorithms, mentioned within this module.

Child \leftarrow Parent

$b \longleftarrow$	a
$d \leftarrow$	b
$e \longleftarrow$	b
$c \longleftarrow$	\overline{a}
$f \leftarrow$	c
$g \longleftarrow$	c



Thanks!

Ammar Sherif

Graphs & Graph Algorithms