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PHYSICS + COMPUTER SCIENCE

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quantum computation

(density op, quantum op, entanglement, circuit,
quantum algorithms)

Level II : quantum error correction (soon.)

(Planck)

(Turing)

(Shannon)

quantum theory + computer science + info-theory

= quantum info science

QUANTUM INFORMATION SCIENCE

- i) quantum sensing
- ii) quantum cryptography
- iii) quantum networking
- iv) quantum simulation
- v) quantum computing
- vi) quantum information concepts

Two fundamental ideas

(1) quantum complexity

(why we think quantum computing is powerful?)

(2) quantum error correction

(why we think quantum computing is scalable?)

Quantum Entanglement

classical book and Quantum Book

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The information does not reside in the pages. It is in the correlations between the pages.
You can't access the information if you read the book, one page at a time.

A complete description of a typical quantum state of just 300 qubits requires more bits than the number of atoms in the visible universe.

Why we think Quantum Computing is powerful?

- (1) Problems believed to be hard classically, which are easy for quantum computers. e.g. factoring
- (2) Complexity theory arguments indicating that quantum computers are hard to simulate classically.
- (3) We don't know how to simulate a quantum computer efficiently using a digital ("classical") computer. The cost of the best known simulation algorithm rises exponentially with the number of qubits.

But... the power of quantum computing is limited: we don't believe that quantum computers can efficiently solve worst-case instances of NP-hard optimization problem.

Classically Easy

Quantum Easy

Quantum Hard
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The Theory of Everything
The theory of everything is not even remotely a theory of every thing.

We know this equation is correct. However it cannot be solved accurately when the number of particles exceed about 10. No computer existing or that will ever exist, can break the barrier, because it will be a catastrophe of dimension.

R.B. Laughlin and D. Pines

Why quantum computing is hard?

- We want qubits to interact strongly with one another.
- We don't want qubits to interact with the environment
- Except when we control or measure them.

Decoherence

$$\frac{1}{\sqrt{2}} (\text{wake} + \text{sleep})$$

Environment

Decoherence explains why quantum phenomena, though observable in the microscopic system, studied in physics lab, are not manifested in the macroscopic physical systems that we encounter in our ordinary experience.

Quantum computer $\xrightarrow{\text{decoherence}}$ Environment
 \rightarrow ERROR.

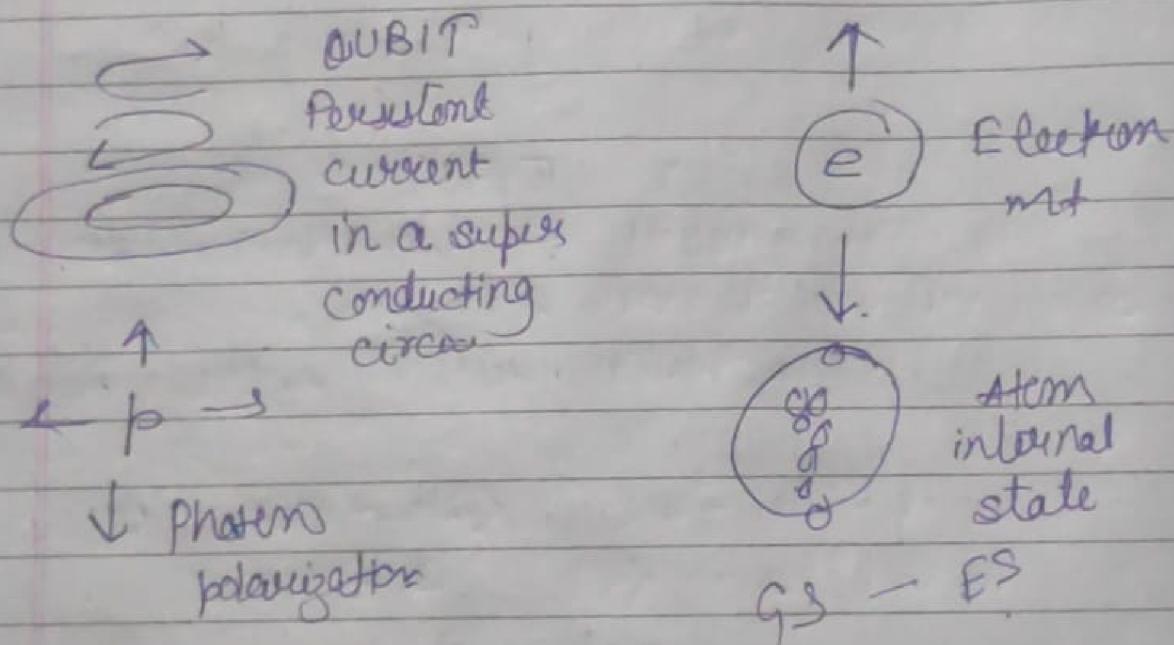
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To resist decoherence, we must prevent the environment from "learning" about the state of quantum computer during the computation

Quantum Error Correction

The protected "logical" quantum information is encoded in a highly entangled state of many qubits.

The environment can't access this information if it interacts locally with the protected system.



Intrinsic resistance to decoherence.

classical systems cannot simulate quantum systems efficiently

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About Sycamore Quantum Device vs Classical solver

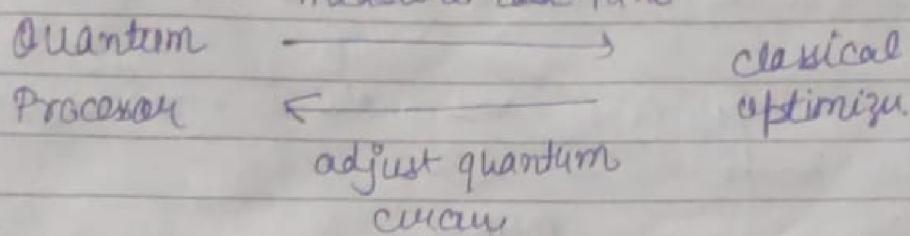
A fully programmable circuit-based qc ($n=53$)
working qubits in 2-d array with coupling of
nearest neighbors

circuit with 20 layers of 2-qubit gates
executed million times

out $\times \times \times$
coupler

HYBRID QUANTUM / CLASSICAL

measure, cost func



The era of quantum heuristics

Sometimes algorithms are effective in practice even though theorists are not able to validate their performance in advance. Eg: Deep Learning
possible quantum examples:

Quantum Annealers, approximate opt.
qml.

dialog b/w quantum algo experts and
application users.

OQC carries a high overhead cost in number
of qubits and gates

20 million physical qubits to break RSA 2048,
for gate error rate 10^{-3}

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Quantum Error
Correction

HYBRID

Quantum Info vs classical info.

1) Randomness: Clicks in Geiger counter are intrinsically random, not pseudo random.

2) Uncertainty: operators 'A' and 'B' commute means that measuring A influences B

do not

3) Entanglement: The whole is more definite than the parts.

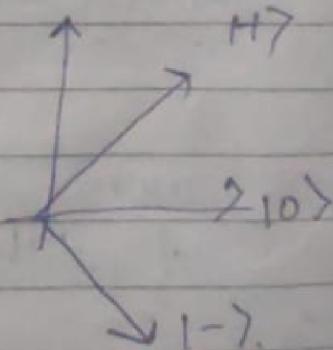
QUBIT

A vector (actually a 'ray' because the normalization is 1, and overall phase does not matter) in a two dimensional complex Hilbert space.

$$|\Psi\rangle = a|0\rangle + b|1\rangle \quad |a|^2 + |b|^2 = 1$$
$$a, b \in \mathbb{C} \quad |\Psi\rangle \sim e^{i\alpha}|\Psi\rangle$$

Two orthogonal states $|0\rangle$ and $|1\rangle$ are perfectly distinguishable

$$|\pm\rangle = \frac{1}{\sqrt{2}}(|0\rangle \pm |1\rangle)$$



Prove it:

But chance of winning: $\cos^2\left(\frac{\pi}{8}\right) = 0.853$

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If Alice sends either $|+\rangle$ or $|-\rangle$

Quantum key distribution

Non-orthogonal states cannot be perfectly distinguishable

Information vs. disturbance

Suppose Alice prepares either $|\psi\rangle$ or $|\phi\rangle$. To distinguish them, Eve performs a unitary transformation that rotates her probe, while leaving Alice's state intact.

$$U: |\psi\rangle_A \otimes |0\rangle_E \rightarrow |\psi\rangle_A \otimes |e\rangle_E$$
$$|\phi\rangle_A \otimes |0\rangle_E \rightarrow |\phi\rangle_A \otimes |f\rangle_E$$

U unitary, hence preserves inner product

$$\langle \psi | \psi \rangle \cdot \langle f | e \rangle = \langle \psi | \psi \rangle$$

and if $|\psi\rangle$ and $|\phi\rangle$ are non-orthogonal
then $\langle f | e \rangle = 0$

states are same

Eve's measurement of the probe cannot reveal any info. about the state.

On the other hand if $|\psi\rangle$ and $|\phi\rangle$ are orthogonal, the probe state can also be orthogonal and Eve can copy that info.

"We can't distinguish non-orthogonal states;
without disturbing them".



Tensor product.
System divided into two

\boxed{A}

$\{\lvert i \rangle_A, i=1, 0, \dots d_A\}$

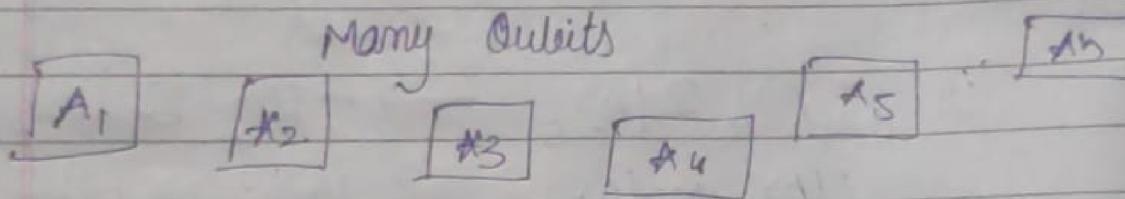
\boxed{B}

$\{\lvert a \rangle_B, a=1, 2, \dots d_B\}$

Basis states of the composite system are distinguishable if they can be distinguished on either Alice's or Bob's.

$$(\langle j_1 \otimes b_1 | (| i \rangle \otimes | a \rangle) = \delta_{ij} \delta_{ab}$$

$\{\lvert 100 \rangle, \lvert 101 \rangle, \lvert 110 \rangle, \lvert 111 \rangle\}$ are all distinguishable



$$\mathcal{C}^{2^n} = \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \dots \otimes \mathbb{C}^2$$

spanned by

$$|\alpha\rangle = |\alpha_{n-1}\rangle \otimes |\alpha_{n-2}\rangle \otimes \dots \otimes |\alpha_1\rangle \otimes |\alpha_0\rangle$$

$$x \in \mathbb{S}^{2n-1}$$

$$\text{where } \langle x | y \rangle = \delta_{xy}$$

For 300 qubits, vector in a space with dimension $2^{300} \sim 10^{90}$ more than the number of atoms in visible universe.

which decomposition into subsystem?
→ Typically dictated by spatial locality.

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+ product state

$|\Psi\rangle = |\Psi_1\rangle_{A_1} \otimes |\Psi_2\rangle_{A_2} \otimes \dots \otimes |\Psi_n\rangle_{A_n}$
has sufficient description; only n real params.

Entanglement

If a pure state is not a product state, it is entangled.

$$\boxed{A_K} \leftrightarrow \boxed{A_{K+1}}$$

Quantum Computer: the circuit model

(1) Hilbert space of 'n' qubits

$$\mathcal{H} = (\mathbb{C}^{2^n})$$

spanned by

$$|\alpha\rangle = |\alpha_{n-1}\rangle \otimes |\alpha_{n-2}\rangle \otimes \dots \otimes |\alpha_1\rangle \otimes |\alpha_0\rangle$$

$\alpha \in \{0, 1\}^n$

The Hilbert Space, is equipped with a natural binary product decomposition into subsystems

$$\mathbb{C}^{2^n} = \underbrace{\mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \dots \otimes \mathbb{C}^2}_{n \text{ times}}$$

or terms

Physically, this decomposition arises from spatial locality.

(2) Initial State $|000..00\rangle = |0\rangle^{\otimes n}$

(3) A finite set of fundamental quantum gates
 $\{U_1, U_2, \dots, U_n\}$

each gate is a unitary transformation acting on a bounded number of qubits.

$$\begin{array}{c} \text{---} \\ | \\ \text{---} \end{array} \quad \left| \begin{array}{c} U \\ \text{---} \end{array} \right| \quad = \quad \begin{array}{c} \text{---} \\ | \\ \text{---} \end{array} \quad \left| \begin{array}{c} U_1 \quad U_2 \\ \text{---} \end{array} \right| \quad \begin{array}{c} \text{---} \\ | \\ \text{---} \end{array}$$

(4) Classical control

The construction of a quantum circuit is directed by a classical computer i.e. using M

(5) Readout

At the end, we read the result by measuring σ_z i.e. projecting onto $|0\rangle, |1\rangle$