

Non-Commutative Linear Logic in an Adjoint Model

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Abstract

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1 Introduction

Linear logic is a well-known resource-sensitive logic. It has been used extensively to model attack trees. This paper concerns a non-commutative variant of linear logic and combines the non-commutative variant with Girard's linear logic [1]. We will only focus on the multiplicative (i.e. \otimes , \multimap) part of linear logic for simplicity. We construct the non-commutative variant by using a non-commutative tensor product \triangleright instead of the commutative \otimes , and two implications \leftarrow and \rightarrow for the two directions of \multimap .

We model the non-commutative linear logic categorically using an adjunction between a symmetric monoidal closed category and a Lambek category. Our categorial adjoint model has a similar structure as Benton's adjoint model [2], in which the multiplicative part of intuitionistic linear logic (ILL) is modeled using an adjunction between a cartesian closed category and a symmetric monoidal closed category. On the other hand, Moggi [3] uses monad models to map intuitionistic logic into ILL. As discussed in [4], Benton's adjoint models only gives rise to commutative monad models and the non-commutative part remained as an open problem. Therefore, by combining our adjoint models with Benton's, we would be able to get non-commutative monad models and thus non-commutative ILL.

The rest of the paper is organized as follows. Section 2 discusses existing approaches on constructing non-commutative linear logic. Section 3 contains the basic definitions in category theory that we will be using in our adjoint model. Familiar readers may skip this section. Section 4 contains the definition and essential properties of our adjoint model. Section 5 discusses the sequent calculus and natural deduction rules for our non-commutative linear logic. We prove that our sequent calculus has the property of cut-elimination and the natural deduction is strongly normalizing. Section 6 talks about the preliminary result after combining our non-commutative model with Benton's commutative model. Section 7 briefly mentions how our model could be used in attack trees and other areas. Section 8 concludes this paper with future work.

2 Related Work

Polakow and Pfenning discussed Ordered Linear Logic (OLL) [5], which combines intuitionistic, commutative linear and non-commutative linear logic, OLL contains sequents of the form $\Gamma, \Delta, \Omega \vdash$



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A , where Γ is a multiset of intuitionistic assumptions, Δ is a multiset of commutative linear assumptions, and Ω is a list of non-commutative linear assumptions. OLL contains logical connectives from all three the logics. Therefore, our non-commutative adjoint model is a part of OLL and after combining with Benton's commutative adjoint model, we would get a simplification of OLL.

Greco and Palmigiano [1] also presents a variant of the multiplicative fragment of non-commutative ILL. But they focus on proper display calculi while we use sequent calculi.

de Paiva and Eades [2] also developed categorical models for the non-commutative ILL by adapting the Dialectica categorical models for linear logic.

3 Category Theory Basics

This section contains the basic definitions in category theory that we will be using in our adjoint model. Our model is based on special kinds of monoidal categories: Lambek categories and symmetric monoidal closed categories and Lambek categories, as defined in Definitions 2 and 3.

► **Definition 1.** A **monoidal category** $(\mathcal{M}, \triangleright, I, \alpha, \lambda, \rho)$ is a category \mathcal{M} consists of

- a bifunctor $\triangleright : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$, called the tensor product;
- an object I , called the unit object;
- three natural isomorphisms α , λ , and ρ with components

$$\alpha_{A,B,C} : (A \triangleright B) \triangleright C \rightarrow A \triangleright (B \triangleright C)$$

$$\lambda_A : I \triangleright A \rightarrow A$$

$$\rho_A : A \triangleright I \rightarrow A$$

where α is called associator, λ is left unitor, and ρ is right unitor,

such that the following diagrams commute for any objects A, B, C in \mathcal{M} :

$$\begin{array}{ccc} ((A \triangleright B) \triangleright C) \triangleright D & \xrightarrow{\alpha_{A,B,C} \triangleright id_D} & (A \triangleright (B \triangleright C)) \triangleright D \xrightarrow{\alpha_{A,B \triangleright C,D}} A \triangleright ((B \triangleright C) \triangleright D) \\ \downarrow \alpha_{A \triangleright B,C,D} & & \downarrow id_A \triangleright \alpha_{B,C,D} \\ (A \triangleright B) \triangleright (C \triangleright D) & \xrightarrow{\alpha_{A,B,C \triangleright D}} & A \triangleright (B \triangleright (C \triangleright D)) \end{array}$$

$$\begin{array}{ccc} (A \triangleright I) \triangleright B & \xrightarrow{\alpha_{A,I,B}} & A \triangleright (I \triangleright B) \\ \downarrow \rho_A \triangleright id_B & & \downarrow id_A \triangleright \lambda_B \\ A \triangleright B & & A \triangleright B \end{array}$$

► **Definition 2.** A **Lambek category** (or a **biclosed monoidal category**) is a monoidal category $(\mathcal{M}, \triangleright, I, \alpha, \lambda, \rho)$ equipped with two bifunctors $\multimap : \mathcal{M}^{op} \times \mathcal{M} \rightarrow \mathcal{M}$ and $\multimap : \mathcal{M} \times \mathcal{M}^{op} \rightarrow \mathcal{M}$ that are both right adjoint to the tensor product. That is, the following natural bijections hold:

$$\text{Hom}_{\mathcal{L}}(X \triangleright A, B) \cong \text{Hom}_{\mathcal{L}}(X, A \multimap B)$$

$$\text{Hom}_{\mathcal{L}}(A \triangleright X, B) \cong \text{Hom}_{\mathcal{L}}(X, B \multimap A)$$

► **Definition 3.** A **symmetric monoidal category** (SMCC) is a monoidal category $(\mathcal{M}, \otimes, I, \alpha, \lambda, \rho)$ together with a natural transformation with components $\text{ex}_{A,B} : A \otimes B \rightarrow B \otimes A$, called **exchange**,

such that the following diagrams commute:

$$\begin{array}{ccc}
 A \otimes I & \xrightarrow{\text{ex}_{A,I}} & I \otimes A \\
 \rho_A \searrow & & \nearrow \lambda_A \\
 & A &
 \end{array}
 \qquad
 \begin{array}{ccc}
 A \otimes B & \xrightarrow{id_{A \otimes B}} & A \otimes B \\
 \text{ex}_{A,B} \searrow & & \nearrow \text{ex}_{B,A} \\
 & B \otimes A &
 \end{array}$$

$$\begin{array}{ccccc}
 (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\text{ex}_{A,B \otimes C}} & (B \otimes C) \otimes A \\
 \text{ex}_{A,B} \otimes id_C \downarrow & & & & \downarrow \alpha_{B,A,C} \\
 (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{id_B \otimes \text{ex}_{A,C}} & B \otimes (C \otimes A)
 \end{array}$$

We use \triangleright for non-symmetric monoidal categories while \otimes for symmetric ones.

► **Definition 4.** A **symmetric monoidal closed category** $(\mathcal{M}, \otimes, I, \alpha, \lambda, \rho)$ is a symmetric monoidal category equipped with a bifunctor $\multimap: \mathcal{M}^{op} \times \mathcal{M} \rightarrow \mathcal{M}$ that is right adjoint to the tensor product. That is, the following natural bijection $\text{Hom}_{\mathcal{M}}(X \otimes A, B) \cong \text{Hom}_{\mathcal{M}}(X, A \multimap B)$ holds.

The relation between SMMCs and Lambek categories are demonstrated in Lemma 5 and Corollary 6.

► **Lemma 5.** Let A and B be two objects in a Lambek category with the exchange natural transformation. Then $(A \multimap B) \cong (B \multimap A)$.

Proof. First, notice that for any object C we have

$$\begin{aligned}
 \text{Hom}[C, A \multimap B] &\cong \text{Hom}[C \otimes A, B] && \mathcal{L} \text{ is a Lambek category} \\
 &\cong \text{Hom}[A \otimes C, B] && \text{By the exchange } \text{ex}_{C,A} \\
 &\cong \text{Hom}[C, B \multimap A] && \mathcal{L} \text{ is a Lambek category}
 \end{aligned}$$

Thus, $A \multimap B \cong B \multimap A$ by the Yoneda lemma. ◀

► **Corollary 6.** A Lambek category with exchange is symmetric monoidal closed.

The essential component in our non-commutative adjoint model is a monoidal adjunction, defined in Definitions 7-11.

► **Definition 7.** Let $(\mathcal{M}, \triangleright, I, \alpha, \lambda, \rho)$ and $(\mathcal{M}', \triangleright', I', \alpha', \lambda', \rho')$ be monoidal categories. A **monoidal functor** (F, m) from \mathcal{M} to \mathcal{M}' is a functor $F: \mathcal{M} \rightarrow \mathcal{M}'$ together with a morphism $m_I: I' \rightarrow F(I)$ and a natural transformation $m_{A,B}: FA' \triangleright' FB' \rightarrow F(A \triangleright B)$, such that the following diagrams commute for any objects A, B , and C in \mathcal{M} :

$$\begin{array}{ccccc}
 (FA \triangleright' FB) \triangleright' FC & \xrightarrow{\alpha'_{FA,FB,FC}} & FA \triangleright' (FB \triangleright' FC) & \xrightarrow{id_{FA \triangleright'} m_{A,B}} & FA \triangleright' F(B \triangleright C) \\
 m_{A,B \triangleright'} id_{FC} \downarrow & & & & \downarrow m_{A,B \triangleright C} \\
 F(A \triangleright B) \triangleright' FC & \xrightarrow{m_{A \triangleright B,C}} & F((A \triangleright B) \triangleright C) & \xrightarrow{F\alpha_{A,B,C}} & F(A \triangleright (B \triangleright C))
 \end{array}$$

$$\begin{array}{ccc}
 I' \triangleright' FA & \xrightarrow{\lambda'_{FA}} & FA \\
 m_I \triangleright id_{FA} \downarrow & & \uparrow F\lambda_A \\
 FI \triangleright' FA & \xrightarrow{m_{I,A}} & F(I \triangleright A)
 \end{array}
 \qquad
 \begin{array}{ccc}
 FA \triangleright' I' & \xrightarrow{\rho'_{FA}} & FA \\
 id_{FA \triangleright'} m_I \downarrow & & \uparrow F\rho_A \\
 FA \triangleright' FI & \xrightarrow{m_{A,I}} & F(A \triangleright I)
 \end{array}$$

► **Definition 8.** Let $(\mathcal{M}, \otimes, I, \alpha, \lambda, \rho)$ and $(\mathcal{M}', \otimes', I', \alpha', \lambda', \rho')$ be symmetric monoidal categories. A **symmetric monoidal functor** $F : \mathcal{M} \rightarrow \mathcal{M}'$ is a monoidal functor (F, m) that satisfies the following coherence diagram:

$$\begin{array}{ccc} FA \otimes' FB & \xrightarrow{\text{ex}_{FA, FB}} & FB \otimes' FA \\ \downarrow m_{A, B} & & \downarrow m_{B, A} \\ F(A \otimes B) & \xrightarrow{F\text{ex}_{A, B}} & F(B \otimes A) \end{array}$$

► **Definition 9.** An **adjunction** between categories \mathcal{C} and \mathcal{D} consists of two functors $F : \mathcal{D} \rightarrow \mathcal{C}$, called the **left adjoint**, and $G : \mathcal{C} \rightarrow \mathcal{D}$, called the **right adjoint**, and two natural transformations $\eta : id_{\mathcal{D}} \rightarrow GF$, called the **unit**, and $\varepsilon : FG \rightarrow id_{\mathcal{C}}$, called the **counit**, such that the following diagrams commute for any object A in \mathcal{C} and B in \mathcal{D} :

$$\begin{array}{ccc} FB & \xrightarrow{F\eta_B} & FGFB \\ \parallel & \searrow \varepsilon_{FB} & \\ FB & & \end{array} \quad \begin{array}{ccc} GA & \xrightarrow{\eta_{GA}} & GFGA \\ \parallel & \searrow G\varepsilon_A & \\ GA & & \end{array}$$

► **Definition 10.** Let (F, m) and (G, n) be monoidal functors from a monoidal category $(\mathcal{M}, \otimes, I, \alpha, \lambda, \rho)$ to a monoidal category $(\mathcal{M}', \otimes', I', \alpha', \lambda', \rho')$. A **monoidal natural transformation** from (F, m) to (G, n) is a natural transformation $\theta : (F, m) \rightarrow (G, n)$ such that the following diagrams commute for any objects A and B in \mathcal{M} :

$$\begin{array}{ccc} FA \triangleright' FB & \xrightarrow{m_{A, B}} & F(A \triangleright B) \\ \downarrow \theta_A \triangleright' \theta_B & & \downarrow \theta_{A \triangleright B} \\ GA \triangleright' GB & \xrightarrow{n_{A, B}} & G(A \triangleright B) \end{array} \quad \begin{array}{ccc} FI & \xrightarrow{\theta_I} & GI \\ \downarrow m_I & \searrow n_I & \\ I' & & \end{array}$$

► **Definition 11.** Let $(\mathcal{M}, \triangleright, I, \alpha, \lambda, \rho)$ and $(\mathcal{M}', \triangleright', I', \alpha', \lambda', \rho')$ be monoidal categories, $F : \mathcal{M} \rightarrow \mathcal{M}'$ and $G : \mathcal{M}' \rightarrow \mathcal{M}$ be functors. The adjunction $F : \mathcal{M} \dashv \mathcal{M}' : G$ is a **monoidal adjunction** if F and G are monoidal functors, and the unit η and the counit ε are monoidal natural transformations.

In Moggi's monad model [], the monad is required to be strong, as defined in Definitions 12 and 13.

► **Definition 12.** Let \mathcal{C} be a category. A **monad** on \mathcal{C} consists of an endofunctor $T : \mathcal{C} \rightarrow \mathcal{C}$ together with two natural transformations $\eta : id_{\mathcal{C}} \rightarrow T$ and $\mu : T^2 \rightarrow T$, where $id_{\mathcal{C}}$ is the identity functor on \mathcal{C} , such that the following diagrams commute:

$$\begin{array}{ccc} T^3 & \xrightarrow{T\mu} & T^2 \\ \downarrow \mu_T & & \downarrow \mu \\ T^2 & \xrightarrow{\mu} & T \end{array} \quad \begin{array}{ccc} T & \xrightarrow{\eta_T} & T^2 \\ \downarrow T\eta & \searrow \mu & \downarrow \mu \\ T^2 & \xrightarrow{\mu} & T \end{array}$$

► **Definition 13.** Let $(\mathcal{M}, \triangleright, I, \alpha, \lambda, \rho)$ be a monoidal category and (T, η, μ) be a monad on \mathcal{M} . T is a **strong monad** if there is natural transformation τ , called the **tensorial strength**, with components

$\tau_{A,B} : A \triangleright TB \rightarrow T(A \triangleright B)$ such that the following diagrams commute:

$$\begin{array}{ccc}
 I \triangleright TA & \xrightarrow{\tau_{I,A}} & T(I \triangleright A) \\
 \searrow \lambda_{TA} & & \swarrow T\lambda_A \\
 & TA &
 \end{array}
 \qquad
 \begin{array}{ccc}
 A \triangleright B & \xrightarrow{id_A \triangleright \eta_B} & A \triangleright TB \\
 \searrow \eta_{A \triangleright B} & & \swarrow \tau_{A,B} \\
 & T(A \triangleright B) &
 \end{array}$$

$$\begin{array}{ccc}
 (A \triangleright B) \triangleright TC & \xrightarrow{\tau_{A \triangleright B, C}} & T((A \triangleright B) \triangleright C) \\
 \downarrow \alpha_{A,B,TC} & & \downarrow T\alpha_{A,B,C} \\
 A \triangleright (B \triangleright TC) & \xrightarrow{id_A \triangleright \tau_{B,C}} A \triangleright T(B \triangleright C) \xrightarrow{\tau_{A,B \triangleright C}} & T(A \triangleright (B \triangleright C))
 \end{array}$$

$$\begin{array}{ccc}
 A \triangleright T^2 B & \xrightarrow{\tau_{A, TB}} T(A \triangleright TB) \xrightarrow{T\tau_{A,B}} & T^2(A \triangleright B) \\
 \downarrow id_A \triangleright \mu_B & & \downarrow \mu_{A \triangleright B} \\
 A \triangleright TB & \xrightarrow{\tau_{A,B}} & T(A \triangleright B)
 \end{array}$$

► **Definition 14.** Let $(\mathcal{M}, \otimes, I, \alpha, \lambda, \rho)$ be a symmetric monoidal category with exchange ex , and (T, η, μ) be a strong monad on \mathcal{M} . Then there is a “**twisted**” **tensorial strength** $\tau'_{A,B} : TA \otimes B \rightarrow T(A \otimes B)$ defined as $\tau'_{A,B} = T\text{ex} \circ \tau_{B,A} \circ \text{ex}$. We can construct a pair of natural transformations Φ, Φ' with components $\Phi_{A,B}, \Phi'_{A,B} : TA \otimes TB \rightarrow T(A \otimes B)$ defined as $\Phi_{A,B} = \mu_{A \otimes B} \circ T\tau'_{A,B} \circ \tau_{TA,B}$ and $\Phi'_{A,B} = \mu_{A \otimes B} \circ T\tau_{A,B} \circ \tau'_{A,TB}$. If $\Phi = \Phi'$, then the monad T is **commutative**.

► **Definition 15.** Let \mathcal{L} be a category. A **comonad** on \mathcal{L} consists of an endofunctor $S : \mathcal{L} \rightarrow \mathcal{L}$ together with two natural transformations $\varepsilon : S \rightarrow id_{\mathcal{L}}$ and $\delta : S^2 \rightarrow S$ such that the following diagrams commute:

$$\begin{array}{ccc}
 S & \xrightarrow{\delta} & S^2 \\
 \delta \downarrow & & \downarrow S\delta \\
 S^2 & \xrightarrow{\delta_S} & S^3
 \end{array}
 \qquad
 \begin{array}{ccc}
 S^2 & \xrightarrow{S\varepsilon} & S \\
 \varepsilon_S \downarrow & & \downarrow \delta \\
 S & \xrightarrow{\delta} & S^2
 \end{array}$$

4 Lambek Adjoint Model

Our adjoint model, Lambek Adjoint Model (LAM), has a similar structure as Benton’s LNL model []. Benton’s LNL model consists of a symmetric monoidal adjunction $F : \mathcal{C} \dashv \mathcal{L} : G$ between a cartesian closed category \mathcal{C} and a symmetric monoidal closed category \mathcal{L} . LAM consists of a monoidal adjunction between a symmetric monoidal closed category and a Lambek category.

► **Definition 16.** A **Lambek Adjoint Model (LAM)** $(\mathcal{C}, \mathcal{L}, F, G, \eta, \varepsilon)$ consists of

- a symmetric monoidal closed category $(\mathcal{C}, \otimes, I, \alpha, \lambda, \rho)$;
- a Lambek category $(\mathcal{L}, \triangleright, I', \alpha', \lambda', \rho')$;
- a monoidal adjunction $F : \mathcal{C} \dashv \mathcal{L} : G$ with unit $\eta_X : X \rightarrow GFX$ and counit $\varepsilon : FG \rightarrow id_{\mathcal{L}}$, where $(F : \mathcal{C} \rightarrow \mathcal{L}, m)$ and $(G : \mathcal{L} \rightarrow \mathcal{C}, n)$ are monoidal functors.

Thus, in LAM, the following four diagrams commute because η and ε are monoidal natural transformations:

$$\begin{array}{ccc}
 X \otimes Y & \xrightarrow{id_{X \otimes Y}} & X \otimes Y \\
 \eta_X \otimes \eta_Y \downarrow & & \downarrow \eta_{X \otimes Y} \\
 GFX \otimes GFY & \xrightarrow{n_{FX,FY}} G(FX \otimes FY) \xrightarrow{Gm_{X,Y}} & GF(X \otimes Y)
 \end{array}
 \qquad
 \begin{array}{ccc}
 I & \xrightarrow{\eta_I} & GFI \\
 \parallel & & \uparrow Gm_I \\
 I & \xrightarrow{n_{I'}} & GI'
 \end{array}$$

$$\begin{array}{ccc}
 FGA \otimes FGB & \xrightarrow{m_{GA,GB}} F(GA \otimes GB) \xrightarrow{Fn_{A,B}} & FG(A \otimes B) \\
 \varepsilon_A \otimes \varepsilon_B \downarrow & & \downarrow \varepsilon_{A \otimes B} \\
 A \otimes B & \xrightarrow{id_{A \otimes B}} & A \otimes B
 \end{array}
 \qquad
 \begin{array}{ccc}
 FGI' & \xrightarrow{\varepsilon_{I'}} & I' \\
 F n_{I'} \uparrow & & \parallel \\
 FI & \xleftarrow{m_I} & I'
 \end{array}$$

And the following two triangles commute because of the adjunction:

$$\begin{array}{ccc}
 FX & \xrightarrow{F\eta_X} & FGFX \\
 \parallel & \searrow \varepsilon_{FX} & \\
 FX & &
 \end{array}
 \qquad
 \begin{array}{ccc}
 GA & \xrightarrow{\eta_{GA}} & GFGA \\
 \parallel & \searrow G\varepsilon_A & \\
 GA & &
 \end{array}$$

Following the tradition, we use letters X, Y, Z for objects in \mathcal{C} and A, B, C for objects in \mathcal{L} . The rest of this section proves essential properties of a LAM.

4.1 An Isomorphism

Let $(\mathcal{C}, \mathcal{L}, F, G, \eta, \varepsilon)$ be a LAM, where (F, m) and (G, n) are monoidal functors. Similarly as in Benton's LNL model, $m_{X,Y}$ are components of a natural isomorphism and m_I is an isomorphism. This is essential for deriving certain rules of our non-commutative linear logic, such as tensor elimination in natural deduction.

We define the inverses of $m_{X,Y} : FX \triangleright FY \rightarrow F(X \otimes Y)$ and $m_I : I' \rightarrow FI$ as:

$$\begin{aligned}
 p_{X,Y} : F(X \otimes Y) &\xrightarrow{F(\eta_X \otimes \eta_Y)} F(GFX \otimes GFY) \xrightarrow{Fn_{FX,FY}} FG(FX \triangleright FY) \xrightarrow{\varepsilon_{FX \triangleright FY}} FX \triangleright FY \\
 p_I : FI &\xrightarrow{Fn_{I'}} FGI' \xrightarrow{\varepsilon_{I'}} I'
 \end{aligned}$$

► **Theorem 17.** $m_{X,Y}$ are components of a natural isomorphism and their inverses are $p_{X,Y}$.

Proof. We need to show that $m_{X,Y} \circ p_{X,Y} = id_{F(X \otimes Y)}$ and $p_{X,Y} \circ m_{X,Y} = id_{FX \triangleright FY}$. The two equations hold because the following diagrams commute: (1)-adjunction; (2)- η is a monoidal natural transformation; (3)-naturality of ε ; (4)-adjunction; (5)-naturality of m ; (6)- ε is a monoidal natural

transformation.

$$\begin{array}{ccc}
 F(X \otimes Y) & \xrightarrow{F(\eta_X \otimes \eta_Y)} & F(GFX \otimes GFY) \\
 \parallel & \searrow F\eta_{X \otimes Y} & \downarrow F\eta_{FX, FY} \\
 & FGF(X \otimes Y) & \\
 (1) \swarrow \varepsilon_{F(X \otimes Y)} & & \searrow FGm_{X, Y} \\
 F(X \otimes Y) & \xleftarrow{m_{X, Y}} FX \triangleright FY \xleftarrow{\varepsilon_{FX \triangleright FY}} & FG(FX \triangleright FY)
 \end{array}
 \quad (2)$$

$$\begin{array}{ccc}
 FX \triangleright FY & \xrightarrow{m_{X, Y}} & F(X \otimes Y) \\
 \parallel & \searrow F\eta_X \triangleright F\eta_Y & \downarrow F(\eta_X \otimes \eta_Y) \\
 & FGFX \triangleright FGFY & \\
 (4) \swarrow \varepsilon_{FX \triangleright FY} & & \searrow m_{GFX, GFY} \\
 FX \triangleright FY & \xleftarrow{\varepsilon_{FX \triangleright FY}} FG(FX \triangleright FY) \xleftarrow{F\eta_{FX, FY}} & F(GFX \otimes GFY)
 \end{array}
 \quad (5)$$

► **Theorem 18.** m_I is an isomorphism and its inverse is p_I .

Proof. This is equivalent to equations $m_I \circ p_I = id_{FI}$ and $p_I \circ m_I = id_{I'}$, equivalent to the following diagrams, which commute because ε is a monoidal natural transformation.

$$\begin{array}{ccc}
 FI & \xrightarrow{F\eta_{I'}} & FGI' \\
 \parallel & & \downarrow \varepsilon_{I'} \\
 FI & \xleftarrow{m_I} & I'
 \end{array}
 \quad
 \begin{array}{ccc}
 I' & \xrightarrow{m_I} & FI \\
 \parallel & & \downarrow F\eta_{I'} \\
 I' & \xleftarrow{\varepsilon_{I'}} & FGI'
 \end{array}$$

4.2 Monad on \mathcal{C}

We first show that the monad on \mathcal{C} in LAM is strong but non-commutative. In Benton's LNL model, the monad on the cartesian closed category is commutative.

► **Lemma 19.** *The monad on the symmetric monoidal closed category \mathcal{C} in LAM is monoidal.*

Proof. Let $(\mathcal{C}, \mathcal{L}, F, G, \eta, \varepsilon)$ be a LAM. We define the monad $(T, \eta : id_{\mathcal{C}} \rightarrow T, \mu : T^2 \rightarrow T)$ on \mathcal{C} as

$$T = GF \qquad \eta_X : X \rightarrow GFX \qquad \mu_X = G\varepsilon_{FX} : GFGFX \rightarrow GFX$$

Since (F, m) and (G, n) are monoidal functors, we have

$$t_{X, Y} = Gm_{X, Y} \circ n_{FX, FY} : TX \otimes TY \rightarrow T(X \otimes Y) \qquad t_I = Gm_I \circ n_{I'} : I \rightarrow TI$$

The monad T being monoidal means:

1. T is a monoidal functor, i.e. the following diagrams commute:

$$\begin{array}{ccc}
 (TX \otimes TY) \otimes TZ & \xrightarrow{\alpha_{TX,TY,TZ}} & TX \otimes (TY \otimes TZ) \xrightarrow{id_{TX} \otimes t_{Y,Z}} TX \otimes T(Y \otimes Z) \\
 \downarrow t_{X,Y} \otimes id_{TZ} & & \downarrow t_{X,Y \otimes Z} \\
 T(X \otimes Y) \otimes TZ & \xrightarrow{t_{X \otimes Y, Z}} & T((X \otimes Y) \otimes Z) \xrightarrow{T\alpha_{X,Y,Z}} T(X \otimes (Y \otimes Z))
 \end{array} \quad (1)$$

$$\begin{array}{ccc}
 I \otimes TX & \xrightarrow{\lambda_{TX}} & TX \\
 \downarrow t_I \otimes id_{TX} & & \uparrow T\lambda_X \\
 TI \otimes TX & \xrightarrow{t_{I,X}} & T(I \otimes X)
 \end{array} \quad (2)$$

$$\begin{array}{ccc}
 TX \otimes I & \xrightarrow{\rho_{TX}} & TX \\
 \downarrow id_{TX} \otimes t_I & & \uparrow T\rho_X \\
 TX \otimes TI & \xrightarrow{t_{X,I}} & T(X \otimes I)
 \end{array} \quad (3)$$

We write GF instead of T in the proof for clarity.

By replacing $t_{X,Y}$ with its definition, diagram (1) above commutes by the following commutative diagram, in which the two hexagons commute because G and F are monoidal functors, and the two quadrilaterals commute by the naturality of n .

$$\begin{array}{ccccc}
 (GFX \otimes GFY) \otimes GFZ & \xrightarrow{\alpha_{GFX,GFY,GFZ}} & GFX \otimes (GFY \otimes GFZ) & \xrightarrow{id_{GFX} \otimes n_{FY,FZ}} & GFX \otimes G(FY \triangleright FZ) \\
 \downarrow n_{FX,FY} \otimes id_{GFZ} & & \downarrow n_{FX,FY \otimes FZ} & & \downarrow id_{GFX} \otimes Gm_{YZ} \\
 G(FX \triangleright FY) \otimes GFZ & & G(FX \triangleright (FY \triangleright FZ)) & & GFX \otimes GF(Y \otimes Z) \\
 \downarrow Gm_{X,Y} \otimes id_{GFZ} & \swarrow n_{FX \triangleright FY, FZ} & \uparrow G\alpha'_{FX,FY,FZ} & \searrow G(id_{FX} \triangleright m_{YZ}) & \downarrow n_{FX, F(Y \otimes Z)} \\
 GF(X \otimes Y) \otimes GFZ & & G((FX \triangleright FY) \triangleright FZ) & & G(FX \triangleright F(Y \otimes Z)) \\
 \downarrow n_{F(X \otimes Y), FZ} & \swarrow G(m_{X,Y} \otimes id_{FZ}) & & & \downarrow Gm_{X,Y \otimes Z} \\
 G(F(X \otimes Y) \triangleright FZ) & \xrightarrow{Gm_{X \otimes Y, Z}} & GF((X \otimes Y) \otimes Z) & \xrightarrow{GF\alpha_{X,Y,Z}} & GF(X \otimes (Y \otimes Z))
 \end{array}$$

Diagram (2) commutes by the following commutative diagrams, in which the top quadrilateral commutes because G is monoidal, the right quadrilateral commutes because F is monoidal, and the left square commutes by the naturality of n .

$$\begin{array}{ccc}
 I \otimes GFX & \xrightarrow{\lambda_{GFX}} & GFX \\
 \downarrow n_{I'} \otimes id_{GFX} & & \downarrow G\lambda'_{FX} \\
 G(I' \otimes GFX) & \xrightarrow{n_{I',FX}} & G(I' \triangleright FX) \\
 \downarrow Gm_{I'} \otimes id_{GFX} & & \downarrow G(m_{I'} \triangleright id_{FX}) \\
 GF(I \otimes X) & \xrightarrow{n_{FI,FX}} & GF(I \triangleright FX) \xrightarrow{Gm_{I,X}} GF(I \otimes X)
 \end{array}$$

Similarly, diagram (3) commutes as follows:

$$\begin{array}{ccccc}
 GFX \otimes I & \xrightarrow{\rho_{GFX}} & GFX & & \\
 \downarrow id_{GFX} \otimes \eta_{I'} & & \uparrow G\rho'_{FX} & & \\
 GFX \otimes GI' & \xrightarrow{\eta_{FX,I'}} & G(FX \triangleright I') & & \\
 \downarrow id_{GFX} \otimes Gm_I & & \downarrow G(id_{FX} \otimes m_I) & & \\
 GFX \otimes GF I & \xrightarrow{\eta_{FX,FI}} & G(FX \triangleright FI) & \xrightarrow{Gm_{X,I}} & GF(X \otimes I)
 \end{array}$$

2. η is a monoidal natural transformation. In fact, since η is the unit of the monoidal adjunction, η is monoidal by definition and thus the following two diagrams commute.

$$\begin{array}{ccc}
 X \otimes Y & \xlongequal{\quad} & X \otimes Y \\
 \eta_X \otimes \eta_Y \downarrow & & \downarrow \eta_{X \otimes Y} \\
 TX \otimes TY & \xrightarrow{t_{X,Y}} & T(X \otimes Y)
 \end{array}
 \qquad
 \begin{array}{ccc}
 I & \xrightarrow{\eta_I} & TI \\
 \parallel & & \uparrow t_I \\
 I & & I
 \end{array}$$

3. μ is a monoidal natural transformation. It is obvious that since $\mu = G\varepsilon_{FA}$ and ε is monoidal, so is μ . Thus the following diagrams commute.

$$\begin{array}{ccc}
 T^2X \otimes T^2Y & \xrightarrow{t_{TX,TY}} & T(TX \otimes TY) \xrightarrow{Tt_{X,Y}} T^2(X \otimes Y) \\
 \mu_X \otimes \mu_Y \downarrow & & \downarrow \mu_{X \otimes Y} \\
 TX \otimes TY & \xrightarrow{t_{X,Y}} & T(X \otimes Y)
 \end{array}
 \qquad
 \begin{array}{ccc}
 T^2I & \xrightarrow{\mu_I} & TI \\
 Tt_I \uparrow & & \uparrow t_I \\
 TI & \xleftarrow{t_I} & I
 \end{array}$$

However, the monad is not symmetric because the following diagram does not commute, for the Lambek category \mathcal{L} is not symmetric.

$$\begin{array}{ccccc}
 GFX \otimes GFY & \xrightarrow{\Theta_{XGFX,GFY}} & GFY \otimes GFX & \xrightarrow{\eta_{FY,FX}} & G(FY \triangleright FX) \\
 \eta_{FX,FY} \downarrow & & & & \downarrow Gm_{Y,X} \\
 G(FX \triangleright FY) & \xrightarrow{Gm_{X,Y}} & GF(X \otimes Y) & \xrightarrow{GF\Theta_{X,Y}} & GF(Y \otimes X)
 \end{array}$$

► **Lemma 20.** *The monad on the symmetric monoidal closed category in LAM is strong.*

Proof. Let $(C, \mathcal{L}, F, G, \eta, \varepsilon)$ be a LAM, where $(C, \otimes, I, \alpha, \lambda, \rho)$ is symmetric monoidal closed, $(\mathcal{L}, \triangleright, I', \alpha', \lambda', \rho')$ is Lambek. In Lemma 19, we have proved that the monad $(T = GF, \eta, \mu)$ is monoidal with the natural transformation $t_{X,Y} : TX \otimes TY \rightarrow T(X \otimes Y)$ and the morphism $t_I : I \rightarrow TI$. We define the tensorial strength $\tau_{X,Y} : X \otimes TY \rightarrow T(X \otimes Y)$ as $\tau_{X,Y} = t_{X,Y} \circ (\eta_X \otimes id_{TY})$. Since η is a monoidal natural transformation, we have $\eta_I = Gm_I \circ \eta_{I'}$. Therefore $\eta_I = t_I$. Thus the following diagram commutes because T is monoidal, where the composition $t_{I,X} \circ (t_I \otimes id_{TX})$ is the definition of $\tau_{I,X}$. So the first triangle in Definition 13 commutes.

$$\begin{array}{ccc}
 I \otimes TX & \xrightarrow{t_I \otimes id_{TX}} & TI \otimes TX \\
 \lambda_{TX} \downarrow & & \downarrow t_{I,X} \\
 TX & \xleftarrow{T\lambda_X} & T(I \otimes X)
 \end{array}$$

Similarly, by using the definition of τ , the the second triangle in the definition is equivalent to the following diagram, which commutes because η is a monoidal natural transformation:

$$\begin{array}{ccc}
 X \otimes Y & \xrightarrow{id_X \otimes \eta_Y} & X \otimes TY \\
 \eta_X \otimes Y \downarrow & \searrow \eta_X \otimes \eta_Y & \downarrow \eta_X \otimes id_{TY} \\
 T(X \otimes Y) & \xleftarrow{t_{X,Y}} & TX \otimes TY
 \end{array}$$

The first pentagon in the definition commutes by the following commutative diagrams, because η and α are natural transformations and T is monoidal:

$$\begin{array}{ccccc}
 (X \otimes Y) \otimes TZ & \xrightarrow{\eta_X \otimes Y \otimes id_{TZ}} & T(X \otimes Y) \otimes TZ & \xrightarrow{t_{X \otimes Y, Z}} & T((X \otimes Y) \otimes Z) \\
 \alpha_{X,Y,TZ} \downarrow & \searrow (\eta_X \otimes \eta_Y) \otimes id_{TZ} & \uparrow t_{X,Y} \otimes id_{TZ} & & \downarrow T\alpha_{X,Y,Z} \\
 X \otimes (Y \otimes TZ) & & (TX \otimes TY) \otimes TZ & & T(X \otimes (Y \otimes Z)) \\
 \downarrow id_X \otimes (\eta_Y \otimes id_{TZ}) & \searrow \eta_X \otimes (\eta_Y \otimes id_{TZ}) & \downarrow \alpha_{TX,TY,TZ} & & \uparrow t_{X,Y \otimes Z} \\
 X \otimes (TY \otimes TZ) & & TX \otimes (TY \otimes TZ) & & \\
 \downarrow id_X \otimes t_{Y,Z} & \searrow \eta_X \otimes id_{TY \otimes TZ} & \downarrow id_{TX} \otimes t_{Y,Z} & & \\
 X \otimes T(Y \otimes Z) & \xrightarrow{id_X \otimes t_{Y,Z}} & X \otimes T(Y \otimes Z) & \xrightarrow{\eta_X \otimes id_{T(Y \otimes Z)}} & TX \otimes T(Y \otimes Z)
 \end{array}$$

The last diagram in the definition commutes by the following commutative diagram, because T is a monad, t is a natural transformation, and μ is a monoidal natural transformation:

$$\begin{array}{ccccc}
 X \otimes T^2 Y & \xrightarrow{\eta_X \otimes id_{T^2 Y}} & TX \otimes T^2 Y & \xrightarrow{t_{X, TY}} & T(X \otimes TY) \\
 id_X \otimes \mu_Y \downarrow & \searrow id_{TX} \otimes \mu_Y & \parallel & \searrow T\eta_X \otimes id_{T^2 Y} & \downarrow T(\eta_X \otimes id_{TY}) \\
 X \otimes TY & \xrightarrow{id_{TX} \otimes \mu_Y} & TX \otimes T^2 Y & \xleftarrow{\mu_X \otimes id_{T^2 Y}} & T^2 X \otimes T^2 Y \xrightarrow{t_{TX, TY}} T(TX \otimes TY) \\
 \eta_X \otimes id_{TY} \downarrow & \searrow id_{TX} \otimes \mu_Y & \downarrow \mu_X \otimes \mu_Y & & \downarrow T t_{X,Y} \\
 TX \otimes TY & \xrightarrow{t_{X,Y}} & T(X \otimes Y) & \xleftarrow{\mu_{X \otimes Y}} & T^2(X \otimes Y)
 \end{array}$$

The following lemma is adopted from [].

► **Lemma 21.** *Let \mathcal{M} be a symmetric monoidal category and T be a strong monad on \mathcal{M} . Then T is commutative iff it is symmetric.*

► **Theorem 22.** *The monad on the SMCC in LAM is strong but non-commutative.*

Proof. The proof is obvious. Based on Lemma 20 and Lemma 21, the monad is non-commutative. ◀

4.3 Comonad on \mathcal{L}

► **Lemma 23.** *The comonad on the Lambek category in a LAM is monoidal.*

Proof. We define the comonad $(S, \varepsilon : S \rightarrow id_{\mathcal{L}}, \delta : S \rightarrow S^2)$ on the Lambek category \mathcal{L} as:

$$S = FG \quad \varepsilon_A : SA \rightarrow A \quad \delta_A = F\eta_{GA} : SA \rightarrow S^2A$$

Thus, we have natural transformation s and morphism s_I defined as:

$$s_{A,B} = F\eta_{A,B} \circ m_{GA,GB} : SA \triangleright SB \rightarrow SA \triangleright SB \quad s_I = F\eta_{I'} \circ m_I : I' \rightarrow SI'$$

The comonad S being monoidal means

1. S is a monoidal functor, i.e. the following diagrams commute:

$$\begin{array}{ccc} (SA \triangleright SB) \triangleright SC & \xrightarrow{\alpha'_{SA,SB,SC}} & SA \triangleright (SB \triangleright SC) \xrightarrow{id_{SA} \triangleright s_{B,C}} SA \triangleright S(B \triangleright C) \\ \downarrow s_{A,B} \triangleright id_{SC} & & \downarrow s_{A,B \triangleright C} \\ S(A \triangleright B) \triangleright SC & \xrightarrow{s_{A \triangleright B, C}} & S((A \triangleright B) \triangleright C) \xrightarrow{S\alpha'_{A,B,C}} S(A \triangleright (B \triangleright C)) \end{array}$$

$$\begin{array}{ccc} I' \triangleright SA & \xrightarrow{\lambda'_{SA}} & SA \\ \downarrow s_{I'} \triangleright id_{SA} & & \uparrow S\lambda'_A \\ SI' \triangleright SA & \xrightarrow{s_{I',A}} & S(I' \triangleright A) \end{array} \quad \begin{array}{ccc} SA \triangleright I' & \xrightarrow{\rho'_{SA}} & SA \\ \downarrow id'_{SA} \triangleright s_{I'} & & \uparrow S\rho'_A \\ SA \triangleright SI' & \xrightarrow{s_{A,I'}} & S(A \triangleright I') \end{array}$$

2. ε is a monoidal natural transformation:

$$\begin{array}{ccc} SA \triangleright SB & \xrightarrow{s_{A,B}} & S(A \triangleright B) \\ \downarrow \varepsilon_A \triangleright \varepsilon_B & & \downarrow \varepsilon_{A \triangleright B} \\ A \triangleright B & \xlongequal{\quad} & A \triangleright B \end{array} \quad \begin{array}{ccc} SI' & \xrightarrow{\varepsilon_{I'}} & I' \\ \swarrow s_{I'} & & \nearrow \\ & I' & \end{array}$$

3. δ is a monoidal natural transformation:

$$\begin{array}{ccc} SA \triangleright SA & \xrightarrow{s_{A,B}} & S(A \triangleright B) \\ \downarrow \delta_A \triangleright \delta_B & & \downarrow \delta_{A \triangleright B} \\ S^2A \triangleright S^2B & \xrightarrow{s_{SA,SB}} S(SA \triangleright SB) \xrightarrow{Ss_{A,B}} S^2(A \triangleright B) \end{array} \quad \begin{array}{ccc} SI' & \xrightarrow{\delta_{I'}} & S^2I' \\ \uparrow s_{I'} & & \uparrow Ss_{I'} \\ I' & \xrightarrow{s_{I'}} & SI' \end{array}$$

The proof for the commutativity of the diagrams are similar as the proof in Lemma 19. We do not include the proof here for simplicity. \blacktriangleleft

We then show that the co-Eilenberg-Moore category of the comonad S is symmetric monoidal closed.

► **Definition 24.** Let (S, ε, δ) be a comonad on a category \mathcal{L} . Then the **co-Eilenberg-Moore category** \mathcal{L}^S of the comonad has

- as objects the S -coalgebras $(A, h_A : A \rightarrow SA)$, where A is an object in \mathcal{L} , s.t. the following diagrams commute:

$$\begin{array}{ccc} A & \xrightarrow{h_A} & SA \\ h_A \downarrow & & \downarrow \delta_A \\ SA & \xrightarrow{Sh_A} & S^2A \end{array} \quad \begin{array}{ccc} & SA & \\ h_A \nearrow & & \searrow \varepsilon_A \\ A & \xlongequal{\quad} & A \end{array}$$

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- as morphisms the coalgebra morphisms, i.e. morphisms $f : (A, h_A) \rightarrow (B, h_B)$ between coalgebras s.t. the diagram commutes:

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ h_A \downarrow & & \downarrow h_B \\ SA & \xrightarrow{Sf} & SB \end{array}$$

► **Lemma 25.** *Given a LAM $(C, \mathcal{L}, F, G, \eta, \varepsilon)$ and the comonad S on \mathcal{L} , the co-Eilenberg-Moore category \mathcal{L}^S has an exchange natural transformation $\text{ex}_{A,B}^S : A \triangleright B \rightarrow B \triangleright A$.*

Proof. We define the exchange $\text{ex}_{A,B}^S : A \triangleright B \rightarrow B \triangleright A$ as

$$A \triangleright B \xrightarrow{h_A \triangleright h_B} FGA \triangleright FGB \xrightarrow{m_{GA,GB}} F(GA \otimes GB) \xrightarrow{F\text{ex}_{GA,GB}} F(GB \otimes GA) \xrightarrow{F\eta_{B,A}} FG(B \triangleright A) \xrightarrow{\varepsilon_{B \triangleright A}} B \triangleright A$$

in which (F, m) and (G, n) are monoidal functors, and ex is the exchange for C . Then ex^S is a natural transformation because the following diagrams commute for morphisms $f : A \rightarrow A'$ and $g : B \rightarrow B'$:

$$\begin{array}{ccccccc} A \triangleright B & \xrightarrow{h_A \triangleright h_B} & FGA \triangleright FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes GB) & \xrightarrow{F\text{ex}_{A,B}} & F(GB \otimes GA) \xrightarrow{F\eta_{B,A}} FG(B \triangleright A) \xrightarrow{\varepsilon_{B \triangleright A}} B \triangleright A \\ f \triangleright g \downarrow & & \downarrow FGf \triangleright FGg & & \downarrow F(Gf \otimes Gg) & & \downarrow F(Gg \otimes Gf) \downarrow FG(g \triangleright f) \downarrow g \triangleright f \\ A' \triangleright B' & \xrightarrow{h_{A'} \triangleright h_{B'}} & FGA' \triangleright FGB' & \xrightarrow{m_{GA',GB'}} & F(GA' \otimes GB') & \xrightarrow{F\text{ex}_{A',B'}} & F(GB' \otimes GA') \xrightarrow{F\eta_{B',A'}} FG(B' \triangleright A') \xrightarrow{\varepsilon_{B' \triangleright A'}} B' \triangleright A' \end{array}$$

◀

► **Lemma 26.** *The following diagrams commute in the co-Eilenberg-Moore category \mathcal{L}^S :*

$$\begin{array}{ccc} F((GA \otimes GB) \otimes GC) & \xrightarrow{F(\eta_{A,B} \otimes id_{GC})} & F(G(A \triangleright B) \otimes GC) \xrightarrow{F(\text{ex}_{A,B} \otimes id_{GC})} FG((A \triangleright B) \triangleright C) \\ \downarrow F(\text{ex}_{A,B} \otimes id_{GC}) & & \downarrow \varepsilon_{(A \triangleright B) \triangleright C} \\ F(G(B \triangleright A) \otimes GC) & & (A \triangleright B) \triangleright C \\ \downarrow F(\eta_{B,A} \otimes id_{GC}) & & \downarrow \text{ex}_{A,B}^S \triangleright id_C \\ F(G(B \triangleright A) \otimes GC) & \xrightarrow{F\eta_{B \triangleright A, C}} & FG((B \triangleright A) \triangleright C) \xrightarrow{\varepsilon_{(B \triangleright A) \triangleright C}} (B \triangleright A) \triangleright C \end{array}$$

$$\begin{array}{ccc} F(GB \otimes (GC \otimes GA)) & \xrightarrow{F(id_{GB} \otimes \eta_{C,A})} & F(GB \otimes G(C \triangleright A)) \xrightarrow{F\eta_{B, C \triangleright A}} FG(B \triangleright (C \triangleright A)) \\ \downarrow F(id_{GB} \otimes \text{ex}_{C,A}) & & \downarrow \varepsilon_{B \triangleright (C \triangleright A)} \\ F(GB \otimes (GA \otimes GC)) & & B \triangleright (C \triangleright A) \\ \downarrow F(id_{GB} \otimes \eta_{A,C}) & & \downarrow id_A \triangleright \text{ex}_{C,A}^S \\ F(GB \otimes G(A \triangleright C)) & \xrightarrow{F\eta_{B, A \triangleright C}} & FG(B \triangleright (A \triangleright C)) \xrightarrow{\varepsilon_{B \triangleright (A \triangleright C)}} B \triangleright (A \triangleright C) \end{array}$$

Proof. We only write the proof for the first diagram. The proof for the second one is similar. (1), (2), (3)–naturality of m ; (4)– F is monoidal; (5), (12)– ε is monoidal; (6), (7), (8), (9), (10)–obvious;

(11)–coalgebra.

(5)

► **Theorem 27.** *The co-Eilenberg-Moore category \mathcal{L}^S of S is symmetric monoidal closed.*

Proof. Let $(\mathcal{L}, \triangleright, I', \alpha', \lambda, \rho')$ be the Lambek category in a SMCC-Lambek model and S be the comonad on \mathcal{L} . Since \mathcal{L} is a Lambek category, it is obvious that \mathcal{L} is also Lambek. By Corollary 6, we only need to prove the exchange defined in Lemma 25 satisfies the three commutative diagrams in Definition 3.

The first triangle in Definition 3 commutes as follows: (1)–coalgebra; (2)– ε is monoidal; (3)–naturality of ρ ; (4)–naturality of ε ; (5)–naturality of m ; (6)– F is monoidal; (7)– C is symmetric; (8)–naturality of ex ; (9)– G is monoidal.

The diagram illustrates the coherence of the monoidal product \otimes with the functor F and the monoidal product \triangleright . It shows how the natural transformations ε and η interact with the functor F and the monoidal product \otimes . The diagram is organized into several rows and columns, with arrows labeled with natural transformations and their components. Key nodes include A , I' , FGA , FGI' , $F(GA \otimes GI')$, $FGA \triangleright I'$, $FGA \triangleright FI$, $F(GA \otimes I)$, $FGA \triangleright I'$, $F(I \otimes GA)$, $FG(I' \triangleright A)$, and $F(GI' \otimes GA)$. Arrows are labeled with transformations like $h_A \triangleright h_{I'}$, $m_{GA,GI'}$, $id_{FGA} \triangleright \varepsilon_{I'}$, $id_{FGA} \triangleright \eta_{I'}$, $\varepsilon_{FGA} \triangleright id_{I'}$, $id_{FGA} \triangleright m_I$, ρ_{FGA} , $F\rho_{GA}$, $F\varepsilon_{GA,I}$, $F\lambda_{GA}$, $FG\lambda'_A$, $F(n_{I'} \otimes id_{GA})$, $F\eta'_{I',A}$, ε_A , and $\varepsilon_{I' \triangleright A}$. Some arrows are double-lined, indicating specific properties or identities. The diagram is divided into regions by these arrows, with some regions labeled with numbers in parentheses.

The second triangle in the proof commutes as follows: (1) and (5)–coalgebra; (2) and (4)– ε is

monoidal; (3)– C is symmetric.

$$\begin{array}{ccccccc}
 A \triangleright B & \xrightarrow{h_A \triangleright h_B} & FGA \triangleright FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes GB) & \xrightarrow{F\epsilon_{A,B}} & F(GB \otimes GA) \xrightarrow{F\eta_{B,A}} FG(B \triangleright A) \\
 \parallel & \swarrow \epsilon_A \triangleright \epsilon_B & & \parallel & \parallel & & \parallel \\
 A \triangleright B & \xleftarrow{\epsilon_{A \triangleright B}} & FG(A \triangleright B) & \xleftarrow{F\eta_{A,B}} & F(GA \otimes GB) & \xleftarrow{F\epsilon_{B,A}} & F(GB \otimes GA) \xleftarrow{m_{GB,GA}} FGB \triangleright FGA \xleftarrow{h_A \triangleright h_A} B \triangleright A
 \end{array}$$

(1) (2) (3) (4) (5)

The third diagram commutes as follows: (1) and (7)–Lemma 26; (2)–naturality of α' ; (3) and (8)–naturality of ϵ ; (4), (6) and (12)– G is a monoidal functor; (5)– C is symmetrical monoidal closed; (9)–coalgebra; (10)– ϵ is a monoidal natural transformation; (11)–naturality of ϵ .

$$\begin{array}{c}
 (A \triangleright B) \triangleright C \xrightarrow{\alpha'_{A,B,C}} A \triangleright (B \triangleright C) \xrightarrow{h_A \triangleright h_{B \triangleright C}} FGA \triangleright FG(B \triangleright C) \\
 \downarrow \epsilon_{(A \triangleright B) \triangleright C} \quad \quad \quad \downarrow \epsilon_{A \triangleright (B \triangleright C)} \quad \quad \quad \downarrow \epsilon_{A \triangleright B \triangleright C} \\
 \begin{array}{ccc}
 FG((A \triangleright B) \triangleright C) & \xrightarrow{F\eta_{A \triangleright B, C}} & F(G(A \triangleright B) \otimes GC) \\
 \uparrow F\eta_{A \triangleright B, C} & & \uparrow F\eta_{A \triangleright B, C} \\
 F((GA \otimes GB) \otimes GC) & \xrightarrow{F\alpha_{GA,GB,GC}} & F(GA \otimes (GB \otimes GC)) \\
 \downarrow F(\epsilon_{A,B} \otimes id_{GC}) & & \downarrow F(\epsilon_{A,B} \otimes id_{GC}) \\
 F((GB \otimes GA) \otimes GC) & \xrightarrow{F\alpha_{GB,GA,GC}} & F(GB \otimes (GA \otimes GC)) \\
 \downarrow F(\eta_{B,A} \otimes id_{GC}) & & \downarrow F(\eta_{B,A} \otimes id_{GC}) \\
 F(G(B \triangleright A) \otimes GC) & \xrightarrow{F\alpha} & F(GB \otimes (GA \otimes GC)) \\
 \downarrow F\eta_{B \triangleright A, C} & & \downarrow F\eta_{B \triangleright A, C} \\
 FG((B \triangleright A) \triangleright C) & \xrightarrow{FG\alpha'_{B,A,C}} & F(GB \otimes G(A \triangleright C)) \\
 \downarrow \epsilon_{(B \triangleright A) \triangleright C} & & \downarrow \epsilon_{B \triangleright A \triangleright C} \\
 B \triangleright (A \triangleright C) & \xleftarrow{\epsilon_{B \triangleright (A \triangleright C)}} & FG(B \triangleright (A \triangleright C)) \\
 \downarrow id_B \triangleright \epsilon_{A,C}^S & & \downarrow id_B \triangleright \epsilon_{A,C}^S \\
 B \triangleright (C \triangleright A) & \xleftarrow{\alpha'_{B,C,A}} & (B \triangleright C) \triangleright A \xleftarrow{\epsilon_{(B \triangleright C) \triangleright A}} FG((B \triangleright C) \triangleright A)
 \end{array}
 \end{array}$$

(1) (2) (3) (4) (5) (6) (7) (8) (9) (10) (11) (12)

5 Non-Commutative Linear Logic

In a LAM, the SMCC C models the commutative linear logic and the Lambek category \mathcal{L} models the non-commutative variant. In Section 5.1, we will present the term assignment for sequent calculus of both sides and prove the cut elimination theorem. In Section 5.2, we present the term assignment for natural deduction of both sides and prove the logic is strongly normalizing.

A sequent in the commutative side is of the form $\Psi, \Phi \vdash_C t : X$. The types must be X, Y, Z , etc., which are objects in the SMCC C . Tye typing contexts are multisets. Suppose Ψ is the set $x_1 : X_1, x_2 : X_2, \dots, x_m : X_m$ and Φ is the set $y_1 : Y_1, y_2 : Y_2, \dots, y_n : Y_n$, then the categorical interpretation of the sequent is the morphism $(X_1 \otimes X_2 \otimes \dots \otimes X_m) \otimes (Y_1 \otimes Y_2 \otimes \dots \otimes Y_n) \rightarrow X$.

A sequent in the non-commutative side is of the form $\Gamma, \Delta \vdash_{\mathcal{L}} s : A$. The types must be A, B, C , etc., which are objects in the Lambek category \mathcal{L} . Tye typing contexts are lists instead of multisets. The typing contexts are mixed in the sense that they could include contexts from the commutative side. When a commutative context $\Psi = \{x_1 : X_1, \dots, x_m : X_m\}$ is included, it is interpreted as the object $F(X_1 \otimes \dots \otimes X_m)$. Therefore, the interpretation of the sequent $\Psi, \Gamma \vdash_{\mathcal{L}} s : A$, where Ψ is defined as above and Γ is the list $y_1 : A_1, \dots, y_n : A_n$, is the morphism $F(X_1 \otimes \dots \otimes X_m) \triangleright (A_1 \triangleright \dots \triangleright A_n) \rightarrow A$.

For the commutative side, since the contexts are multisets, the following exchange rule is implicit in both sequent calculus and natural deduction:

$$\frac{\Phi, x : X, y : Y, \Psi \vdash_C t : Z}{\Phi, z : Y, w : X, \Psi \vdash_C \text{ex } w, z \text{ with } x, y \text{ in } t : Z} \quad \text{T_BETA}$$

5.1 Sequent Calculus

The term assignment for sequent calculus of the commutative part of the model, i.e. the SMCC of the adjunction, is defined in Figure 1. And the term assignme for the non-commutative part, i.e. the Lambek category of the adjunction, is defined in Figure 2. We do not have the structural rules except for exchange because the calculus is for linear logic.

$\frac{}{x : X \vdash_C x : X} \quad \text{T_VAR}$	$\frac{\Phi, \Psi \vdash_C t : X}{\Phi, x : \text{UnitT}, \Psi \vdash_C \text{let } x : \text{UnitT be trivT in } t : X} \quad \text{T_UNITL}$	$\frac{}{\cdot \vdash_C \text{trivT} : \text{UnitT}} \quad \text{T_UNITR}$
$\frac{\Phi, x : X, y : Y, \Psi \vdash_C t : Z}{\Phi, z : X \otimes Y, \Psi \vdash_C \text{let } z : X \otimes Y \text{ be } x \otimes y \text{ in } t : Z} \quad \text{T_TENL}$	$\frac{\Phi \vdash_C t_1 : X \quad \Psi \vdash_C t_2 : Y}{\Phi, \Psi \vdash_C t_1 \otimes t_2 : X \otimes Y} \quad \text{T_TENR}$	
$\frac{\Phi \vdash_C t_1 : X \quad \Psi_1, x : Y, \Psi_2 \vdash_C t_2 : Z}{\Psi_1, \Phi, y : X \multimap Y, \Psi_2 \vdash_C [\text{app } y t_1 / x] t_2 : Z} \quad \text{T_IMPL}$	$\frac{\Phi, x : X, \Psi \vdash_C t : Y}{\Phi, \Psi \vdash_C \lambda x : X. t : X \multimap Y} \quad \text{T_IMPR}$	
$\frac{\Phi \vdash_{\mathcal{L}} s : A}{\Phi \vdash_C \text{Gs} : \text{GA}} \quad \text{T_Gr}$	$\frac{\Phi \vdash_C t_1 : X \quad \Psi_1, x : X, \Psi_2 \vdash_C t_2 : Y}{\Psi_1, \Phi, \Psi_2 \vdash_C [t_1 / x] t_2 : Y} \quad \text{T_CUT}$	

■ **Figure 1** Sequent Calculus: Commutative Part

We prove cut elimination for the sequent calculus using a similar method as Benton's []. First, we replace the cut rules S_GL by the following n-ary forms respectively:

5.2 Natural Deduction

The term assignment for natural deduction of the commutative part of the model, i.e. the SMCC of the adjunction, is defined in Figure 3. And the term assignme for the non-commutative part, i.e. the Lambek category of the adjunction, is defined in Figure 4.

We could derive exchange comonadically as follows:

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$$\begin{array}{c}
\frac{}{x : A \vdash_{\mathcal{L}} x : A} \text{S_AX} \quad \frac{}{\cdot \vdash_{\mathcal{L}} \text{trivS} : \text{UnitS}} \text{S_UNITR} \quad \frac{\Gamma, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, x : \text{UnitT}, \Delta \vdash_{\mathcal{L}} \text{let } x : \text{UnitT} \text{ be } \text{trivT} \text{ in } s : A} \text{S_UNITL1} \\
\frac{\Gamma, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, x : \text{UnitS}, \Delta \vdash_{\mathcal{L}} \text{let } x : \text{UnitS} \text{ be } \text{trivS} \text{ in } s : A} \text{S_UNITL2} \quad \frac{\Gamma, x : X, y : Y, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, z : Y, w : X, \Delta \vdash_{\mathcal{L}} \text{ex } w, z \text{ with } x, y \text{ in } s : A} \text{S_BETA} \\
\frac{\Gamma, x : X, y : Y, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, z : X \otimes Y, \Delta \vdash_{\mathcal{L}} \text{let } z : X \otimes Y \text{ be } x \otimes y \text{ in } s : A} \text{S_TENL1} \quad \frac{\Gamma, x : A, y : B, \Delta \vdash_{\mathcal{L}} s : C}{\Gamma, z : A \multimap B, \Delta \vdash_{\mathcal{L}} \text{let } z : A \multimap B \text{ be } x \multimap y \text{ in } s : C} \text{S_TENL2} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad \Delta \vdash_{\mathcal{L}} s_2 : B}{\Gamma, \Delta \vdash_{\mathcal{L}} s_1 \multimap s_2 : A \multimap B} \text{S_TENR} \quad \frac{\Phi \vdash_C t : X \quad \Gamma, x : Y, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, \Phi, y : X \multimap Y, \Delta \vdash_{\mathcal{L}} [\text{app } y \text{ t/x}] s : A} \text{S_IMPL} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad \Delta, x : B \vdash_{\mathcal{L}} s_2 : C}{\Delta, \Gamma, y : A \multimap B \vdash_{\mathcal{L}} [\text{app } y s_1/x] s_2 : C} \text{S_IMPLR} \quad \frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad x : B, \Delta \vdash_{\mathcal{L}} s_2 : C}{y : B \leftarrow A, \Gamma, \Delta \vdash_{\mathcal{L}} [\text{app } y s_1/x] s_2 : C} \text{S_IMPLL} \\
\frac{\Gamma, x : A \vdash_{\mathcal{L}} s : B}{\Gamma \vdash_{\mathcal{L}} \lambda_{r,x} : A.s : A \multimap B} \text{S_IMPRR} \quad \frac{x : A, \Gamma \vdash_{\mathcal{L}} s : B}{\Gamma \vdash_{\mathcal{L}} \lambda_{l,x} : A.s : B \leftarrow A} \text{S_IMPLR} \quad \frac{\Phi \vdash_C t : X}{\Phi \vdash_{\mathcal{L}} \text{Ft} : \text{FX}} \text{S_FR} \\
\frac{\Gamma, x : X, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, y : \text{FX}, \Delta \vdash_{\mathcal{L}} \text{let } y : \text{FX} \text{ be } \text{F } x \text{ in } s : A} \text{S_FL} \quad \frac{\Gamma, x : A, \Delta \vdash_{\mathcal{L}} s : B}{\Gamma, y : \text{GA}, \Delta \vdash_{\mathcal{L}} \text{let } y : \text{GA} \text{ be } \text{G } x \text{ in } s : B} \text{S_GL} \\
\frac{\Phi \vdash_C t : X \quad \Gamma_1, x : X, \Gamma_2 \vdash_{\mathcal{L}} s : A}{\Gamma_1, \Phi, \Gamma_2 \vdash_{\mathcal{L}} [t/x] s : A} \text{S_CUT1} \quad \frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad \Delta_1, x : A, \Delta_2 \vdash_{\mathcal{L}} s_2 : B}{\Delta_1, \Gamma, \Delta_2 \vdash_{\mathcal{L}} [s_1/x] s_2 : B} \text{S_CUT2}
\end{array}$$

■ Figure 2 Sequent Calculus: Non-Commutative Part

$$\begin{array}{c}
\frac{}{x : X \vdash_C x : X} \text{T_ID} \quad \frac{}{\cdot \vdash_C \text{trivT} : \text{UnitT}} \text{T_UNITI} \quad \frac{\Phi \vdash_C t_1 : \text{UnitT} \quad \Psi \vdash_C t_2 : Y}{\Phi, \Psi \vdash_C \text{let } t_1 : \text{UnitT} \text{ be } \text{trivT} \text{ in } t_2 : Y} \text{T_UNITE} \\
\frac{\Phi \vdash_C t_1 : X \quad \Psi \vdash_C t_2 : Y}{\Phi, \Psi \vdash_C t_1 \otimes t_2 : X \otimes Y} \text{T_TENI} \quad \frac{\Phi \vdash_C t_1 : X \otimes Y \quad \Psi_1, x : X, y : Y, \Psi_2 \vdash_C t_2 : Z}{\Psi_1, \Phi, \Psi_2 \vdash_C \text{let } t_1 : X \otimes Y \text{ be } x \otimes y \text{ in } t_2 : Z} \text{T_TENE} \\
\frac{\Phi, x : X \vdash_C t : Y}{\Phi \vdash_C \lambda x : X. t : X \multimap Y} \text{T_IMPI} \quad \frac{\Phi \vdash_C t_1 : X \multimap Y \quad \Psi \vdash_C t_2 : X}{\Phi, \Psi \vdash_C \text{app } t_1 t_2 : Y} \text{T_IMPE} \quad \frac{\Phi \vdash_{\mathcal{L}} s : A}{\Phi \vdash_C \text{Gs} : \text{GA}} \text{T_GI} \\
\frac{\Gamma, x : X, y : Y, \Delta \vdash_{\mathcal{L}} s : A}{\Gamma, z : Y, w : X, \Delta \vdash_{\mathcal{L}} \text{ex } w, z \text{ with } x, y \text{ in } s : A} \text{S_BETA}
\end{array}$$

■ Figure 3 Natural Deduction: Commutative Part

$$\begin{array}{c}
\frac{}{y_0 : \text{GB} \vdash_C y_0 : \text{GB}} \text{ID} \quad \frac{}{x_0 : \text{GA} \vdash_C x_0 : \text{GA}} \text{ID} \\
\frac{y_0 : \text{GB} \vdash_{\mathcal{L}} F y_0 : \text{FGB}}{y_0 : \text{GB}, x_0 : \text{GA} \vdash_{\mathcal{L}} F y_0 \multimap F x_0 : \text{FGB} \multimap \text{FGA}} \text{FI} \quad \frac{x_0 : \text{GA} \vdash_{\mathcal{L}} F x_0 : \text{FGA}}{x_0 : \text{GA}, y_0 : \text{GB} \vdash_{\mathcal{L}} \text{ex } y_0, x_0 \text{ in } (F y_0 \multimap F x_0) : \text{FGB} \multimap \text{FGA}} \text{FE} \\
\frac{y_0 : \text{GB}, x_0 : \text{GA} \vdash_{\mathcal{L}} F y_0 \multimap F x_0 : \text{FGB} \multimap \text{FGA}}{x_1 : \text{GA}, y_1 : \text{GB} \vdash_{\mathcal{L}} \text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0) : \text{FGB} \multimap \text{FGA}} \text{BETA} \\
\frac{x_2 : \text{FGA} \vdash_{\mathcal{L}} x_2 : \text{FGA}}{x_2 : \text{FGA}, y_2 : \text{FGB} \vdash_{\mathcal{L}} \text{let } F x_1 : \text{FGB} \text{ be } y_2 \text{ in } (\text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0)) : \text{FGB} \multimap \text{FGA}} \text{FI} \\
\frac{x_2 : \text{FGA}, y_2 : \text{FGB} \vdash_{\mathcal{L}} \text{let } F x_1 : \text{FGB} \text{ be } y_2 \text{ in } (\text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0)) : \text{FGB} \multimap \text{FGA}}{z : \text{FGA} \multimap \text{FGB} \vdash_{\mathcal{L}} \text{let } z : \text{FGA} \multimap \text{FGB} \text{ be } x_2 \multimap y_2 \text{ in } (\text{let } F x_1 : \text{FGB} \text{ be } y_2 \text{ in } (\text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0)))) : \text{FGB} \multimap \text{FGA}} \text{TENE2} \\
\frac{\cdot \vdash_{\mathcal{L}} \lambda_{r,z} : \text{FGA} \multimap \text{FGB} \text{ let } z : \text{FGA} \multimap \text{FGB} \text{ be } x_2 \multimap y_2 \text{ in } (\text{let } F x_1 : \text{FGB} \text{ be } y_2 \text{ in } (\text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0)))) : (\text{FGA} \multimap \text{FGB}) \multimap (\text{FGB} \multimap \text{FGA})}{\cdot \vdash_{\mathcal{L}} \lambda_{r,z} : \text{FGA} \multimap \text{FGB} \text{ let } z : \text{FGA} \multimap \text{FGB} \text{ be } x_2 \multimap y_2 \text{ in } (\text{let } F x_1 : \text{FGB} \text{ be } y_2 \text{ in } (\text{ex } y_1, x_1 \text{ with } y_0, x_0 \text{ in } (F y_0 \multimap F x_0)))) : (\text{FGA} \multimap \text{FGB}) \multimap (\text{FGB} \multimap \text{FGA})} \text{IMPL}
\end{array}$$

We also have the three cut rules derivable in the natural deduction:

$$\frac{\Phi \vdash_C t : X}{\Phi \vdash_{\mathcal{L}} \text{Ft} : \text{FX}} \text{FI}$$

$$\begin{array}{c}
\frac{}{x : A \vdash_{\mathcal{L}} x : A} \text{S_ID} \quad \frac{}{\cdot \vdash_{\mathcal{L}} \text{trivS} : \text{UnitS}} \text{S_UNITI} \quad \frac{\Gamma \vdash_{\mathcal{L}} s_1 : \text{UnitS} \quad \Delta \vdash_{\mathcal{L}} s_2 : A}{\Gamma, \Delta \vdash_{\mathcal{L}} \text{let } s_1 : \text{UnitS} \text{ be } \text{trivS} \text{ in } s_2 : A} \text{S_UNITE1} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : \text{UnitS} \quad \Delta \vdash_{\mathcal{L}} s_2 : A}{\Gamma, \Delta \vdash_{\mathcal{L}} \text{let } s_1 : \text{UnitS} \text{ be } \text{trivS} \text{ in } s_2 : A} \text{S_UNITE1} \quad \frac{\Phi \vdash_C t : \text{UnitT} \quad \Gamma \vdash_{\mathcal{L}} s : A}{\Phi, \Gamma \vdash_{\mathcal{L}} \text{let } t : \text{UnitT} \text{ be } \text{trivT} \text{ in } s : A} \text{S_UNITE2} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad \Delta \vdash_{\mathcal{L}} s_2 : B}{\Gamma, \Delta \vdash_{\mathcal{L}} s_1 \triangleright s_2 : A \triangleright B} \text{S_TENE1} \quad \frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \triangleright B \quad \Delta_1, x : A, y : B, \Delta_2 \vdash_{\mathcal{L}} s_2 : C}{\Delta_1, \Gamma, \Delta_2 \vdash_{\mathcal{L}} \text{let } s_1 : A \triangleright B \text{ be } x \triangleright y \text{ in } s_2 : C} \text{S_TENE1} \\
\frac{\Phi \vdash_C t : X \otimes Y \quad \Gamma_1, x : X, y : Y, \Gamma_2 \vdash_{\mathcal{L}} s : A}{\Gamma_1, \Phi, \Gamma_2 \vdash_{\mathcal{L}} \text{let } t : X \otimes Y \text{ be } x \otimes y \text{ in } s : A} \text{S_TENE2} \quad \frac{\Gamma, x : A \vdash_{\mathcal{L}} s : B}{\Gamma \vdash_{\mathcal{L}} \lambda_r x : A. s : A \rightarrow B} \text{S_IMPRI} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \rightarrow B \quad \Delta \vdash_{\mathcal{L}} s_2 : A}{\Gamma, \Delta \vdash_{\mathcal{L}} \text{app}_r s_1 s_2 : B} \text{S_IMPRE} \quad \frac{x : A, \Gamma \vdash_{\mathcal{L}} s : B}{\Gamma \vdash_{\mathcal{L}} \lambda_l x : A. s : B \leftarrow A} \text{S_IMPLI} \\
\frac{\Gamma \vdash_{\mathcal{L}} s_1 : B \leftarrow A \quad \Delta \vdash_{\mathcal{L}} s_2 : A}{\Gamma, \Delta \vdash_{\mathcal{L}} \text{app}_l s_1 s_2 : B} \text{S_IMPLE} \quad \frac{\Phi \vdash_C t : \text{GA}}{\Phi \vdash_{\mathcal{L}} \text{derelict } t : A} \text{S_GE} \quad \frac{\Phi \vdash_C t : X}{\Phi \vdash_{\mathcal{L}} \text{Ft} : \text{FX}} \text{S_FI} \\
\frac{\Gamma \vdash_{\mathcal{L}} y : \text{FX} \quad \Delta_1, x : X, \Delta_2 \vdash_{\mathcal{L}} s : A}{\Delta_1, \Gamma, \Delta_2 \vdash_{\mathcal{L}} \text{let } \text{F}x : \text{FX} \text{ be } y \text{ in } s : A} \text{S_FE}
\end{array}$$

■ **Figure 4** Natural Deduction: Non-Commutative Part

$$\frac{\Phi \vdash_C t_1 : X \quad \Psi_1, x : X, \Psi_2 \vdash_C t_2 : Y}{\Psi_1, \Phi, \Psi_2 \vdash_C [t_1/x]t_2 : Y} \text{T_CUT} \quad \frac{\Phi \vdash_C t : X \quad \Gamma_1, x : X, \Gamma_2 \vdash_{\mathcal{L}} s : A}{\Gamma_1, \Phi, \Gamma_1 \vdash_{\mathcal{L}} [t/x]s : A} \text{S_CUT1} \quad \frac{\Gamma \vdash_{\mathcal{L}} s_1 : A \quad \Delta_1, x : A, \Delta_2 \vdash_{\mathcal{L}} s_2 : B}{\Delta_1, \Gamma, \Delta_2 \vdash_{\mathcal{L}} [s_1/x]s_2 : B} \text{S_CUT2}$$

6 Combining with Benton's Adjoint Model

7 Applications

8 Conclusion

TODO

A Appendix