# Separating Linear Modalities

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#### Abstract

TODO

### 1 Introduction

TODO [1]

## 2 Categorical Models

## 2.1 Lambek Categories

**TODO: Define Lambek Categories** 

### 2.2 Lambek Categories with Weakening and Contraction

**Definition 1.** A Lambek category with weakening,  $(\mathcal{L}, w, weak)$ , is specified by

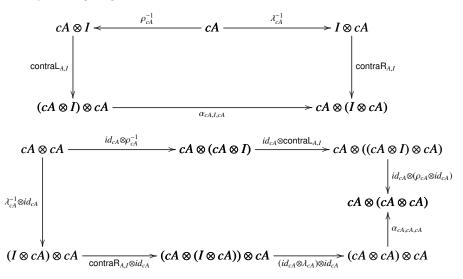
- a monoidal category  $(\mathcal{L}, I, \otimes)$ ,
- a monoidal comonad  $(w, \varepsilon, \delta)$  on  $\mathcal{L}$  with  $q_{A,B} : wA \otimes wB \longrightarrow w(A \otimes B)$  and  $q_I : I \longrightarrow wI$ , and
- a monoidal natural transformation weak on  $\mathcal{L}$  with components weak<sub>A</sub>:  $wA \longrightarrow I$  s.t. weak<sub>A</sub> is a coalgebra morphism, i.e. the following digram commutes:



**Definition 2.** A Lambek category with contraction,  $(\mathcal{L}, c, \text{contraL}, \text{contraR})$ , is specified by

• a monoidal category  $(\mathcal{L}, I, \otimes)$ ,

- a monoidal comonad  $(c, \varepsilon, \delta)$  on  $\mathcal{L}$  with  $q_{A,B} : cA \otimes cB \longrightarrow c(A \otimes B)$  and  $q_I : I \longrightarrow cI$ , and
- monoidal natural transformations contraL and contraR on  $\mathcal{L}$  with components contraL<sub>A,B</sub>:  $cA \otimes B \longrightarrow (cA \otimes B) \otimes cA$  and contraR<sub>A,B</sub>:  $B \otimes cA \longrightarrow cA \otimes (B \otimes cA)$ , s.t. the following diagrams commutes:



### 2.3 Lambek Categories with Exchange

**Definition 3.** A Lambek category with exchange,  $(\mathcal{L}, e, ex)$ , is specified by

- a monoidal category  $(\mathcal{L}, I, \otimes)$ ,
- a monoidal comonad  $(e, \varepsilon, \delta)$  on  $\mathcal{L}$  with  $q_{A,B} : eA \otimes eB \longrightarrow e(A \otimes B)$  and  $q_I : I \longrightarrow eI$ , and
- a monoidal natural transformation ex on  $\mathcal{L}$  with components  $ex_{A,B}: e(A \otimes B) \longrightarrow eB \otimes eA$

subject to the following coherence condition:

and  $ex_{A,B}$  is a coalgebra morphism, i.e. the following diagram commutes:

$$e(A \otimes B) \xrightarrow{\delta_{A \otimes B}} e^{2}(A \otimes B)$$

$$e_{A,B} \downarrow \qquad \qquad \downarrow e_{A,B}$$

$$eB \otimes eA \xrightarrow{\delta_{B} \otimes \delta_{A}} e^{2}B \otimes e^{2}A \xrightarrow{q_{eB,eA}} e(eB \otimes eA)$$

**Lemma 4.** Let  $(e, \varepsilon, \delta)$  be a monoidal comonad over a monoidal category  $(\mathcal{L}, I, \otimes)$  such that  $(\mathcal{L}, e, ex)$  is a Lambek category with exchange. Then in the co-Kleisli category of  $\mathcal{L}$ , there exists a natural isomorphism with components  $\gamma_{A,B} : e(A \otimes B) \longrightarrow B \otimes A$ .

*Proof.* Suppose  $(\mathcal{L}, e, ex)$  is a Lambek category with exchange. We define a natural transformation  $\gamma_{A,B}: e(A \otimes B) \longrightarrow B \otimes A$  as

$$\gamma_{A,B} = (\varepsilon_B \otimes \varepsilon_A) \circ \mathsf{ex}_{A,B}.$$

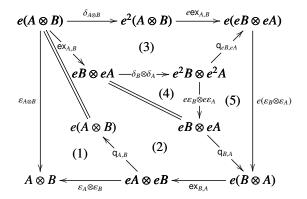
 $\gamma$  is a natural isomorphism if the following diagram commutes:

$$e(A \otimes B) \xrightarrow{\delta_{A \otimes B}} e^{2}(A \otimes B)$$

$$\downarrow e_{\gamma_{A,B}}$$

$$A \otimes B \xleftarrow{\gamma_{B,A}} e(B \otimes A)$$

, which does commute by the diagram chasing below.



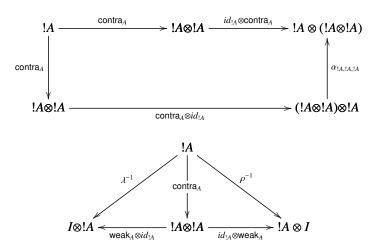
(1) commutes because e is monoidal. (2) and (3) commute by the definition of the Lambek category with exchange. (4) commutes because e is a comonad. And (5) commutes by the naturality of q.

### 2.4 Linear Categories

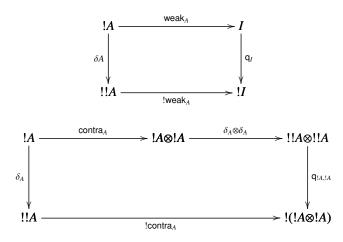
**Definition 5.** A linear category,  $(\mathcal{L}, !, weak, contra)$ , is specified by

- a symmetric monoidal closed category  $(\mathcal{L}, I, \otimes, \multimap)$ ,
- a symmetric monoidal comonad  $(!, \varepsilon, \delta)$  on  $\mathcal{L}$ , with  $q_{A,B} : !A \otimes !B \longrightarrow !(A \otimes B)$  and  $q_I : I \longrightarrow !I$ ;
- monoidal natural transformations on  $\mathcal{L}$  with components weak<sub>A</sub>:! $A \longrightarrow I$  and contra<sub>A</sub>:! $A \longrightarrow !A \otimes !A$ , s.t.

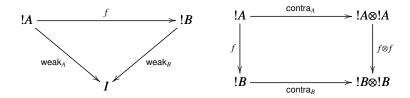
- each (!A, weak<sub>A</sub>, contra<sub>A</sub>) is a commutative comonoid, i.e. the following diagrams commute and  $\beta \circ \text{contra}_A = \text{contra}_A$  where  $\beta_{B,C} : B \otimes C \longrightarrow C \otimes B$  is the symmetry natural transformation of  $\mathcal{L}$ ;



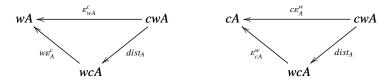
 weak<sub>A</sub> and contra<sub>A</sub> are coalgebra morphisms, i.e. the following diagrams commute;



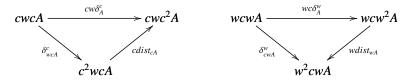
- any coalgebra morphism  $f:(!A,\delta_A) \longrightarrow (!B,\delta_B)$  between free coalgebras preserve the comonoid structure given by weak and contra, i.e. the following diagrams commute.



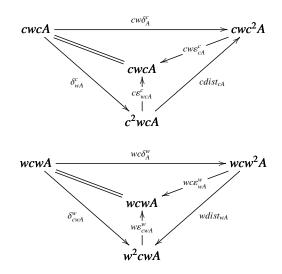
**Definition 6.** Given two comonads  $(c, \varepsilon^c, \delta^c)$  and  $(w, \varepsilon^w, \delta^w)$  on a category  $\mathcal{L}$  such that  $(\mathcal{L}, c, \mathsf{contraL}, \mathsf{contraR})$  is a Lambek category with contraction and  $(\mathcal{L}, w, \mathsf{weak})$  is a Lambek category with weakening, we define a **distributive law** of c over w to be a natural transformation with components  $dist_A : cwA \longrightarrow wcA$ , subject to the following coherence diagrams:



**Lemma 7.** Given two comonads  $(c, \varepsilon^c, \delta^c)$  and  $(w, \varepsilon^w, \delta^w)$  on a category  $\mathcal{L}$  such that  $(\mathcal{L}, c, \text{contraL}, \text{contraR})$  is a Lambek category with contraction and  $(\mathcal{L}, w, \text{weak})$  is a Lambek category with weakening, the following two diagrams commute:



*Proof.* The two diagrams above commute by the diagram chasings below:



, which commute by the distribute law dist and the comonad laws of c and w.

**Lemma 8.** Let  $(c, \varepsilon^c, \delta^c)$  and  $(w, \varepsilon^w, \delta^w)$  be two monoidal comonads on a Lambek category with weakening and contraction  $(\mathcal{L}, I, \otimes, w, \mathsf{weak}^w, c, \mathsf{contraL}, \mathsf{contraR})$ . Then the composition of c and w using the distributive law dist<sub>A</sub>:  $cwA \longrightarrow wcA$  is a monoidal comonad  $(cw, \varepsilon, \delta)$  on  $\mathcal{L}$ .

*Proof.* Suppose  $(c, \varepsilon^c, \delta^c)$  and  $(w, \varepsilon^w, \delta^w)$  are monoidal comonads, and  $(\mathcal{L}, I, \otimes, w, \mathsf{weak}^w, c, \mathsf{contraL}, \mathsf{contraR})$  is a Lambek category with weakening and contraction. Since by definition  $c, w: \mathcal{L} \longrightarrow \mathcal{L}$  are monoidal functors we know that their composition  $cw: \mathcal{L} \longrightarrow \mathcal{L}$  is a monoidal functor:

$$q_{A,B} : cwA \otimes cwB \longrightarrow cw(A \otimes B)$$

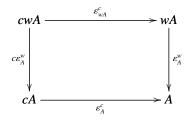
$$q_{A,B} = cq_{A,B}^{w} \circ q_{wA,wB}^{c}$$

$$q_{I} : I \longrightarrow cwI$$

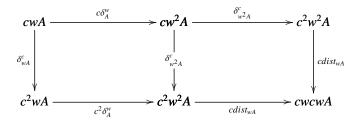
$$q_{I} = cq_{I}^{w} \circ q_{I}^{c}$$

We must now define both  $\varepsilon_A: cwA \longrightarrow A$  and  $\delta_A: cwA \longrightarrow cwcwA$ , and then show that they are monoidal natural transformations subject to the comonad laws. Since we are composing two comonads each of  $\varepsilon$  and  $\delta$  can be given two definitions, but they are equivalent:

•  $\varepsilon_A$ :  $cwA \longrightarrow A$  is defined as in the diagram below, which commutes by the naturality of  $\varepsilon^c$ .



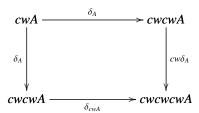
•  $\delta_A : cwA \longrightarrow cwcwA$  is defined as in the diagram:



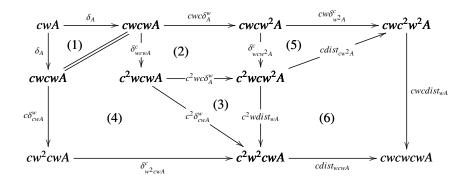
The left part of the diagram commutes by the naturality of  $\delta^c$  and the right part commutes trivially.

The remainder of the proof shows that the comonad laws hold.

#### Case 1:

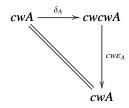


The previous diagram commutes because the following one does.

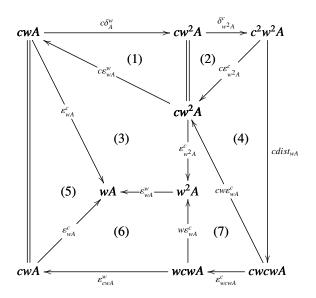


(1) commutes by equality and we will not expand  $\delta_A$  for simplicity. (2) and (4) commutes by the naturality of  $\delta^c$ . (3), (5) commutes by the conditions of *dist*. (6) commutes by the naturality of *dist*.

#### Case 2:

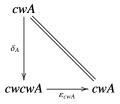


The triangle commutes because of the following diagram chasing.

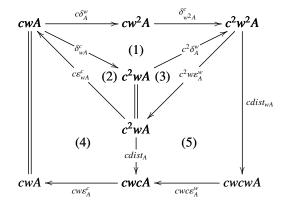


(1) commutes by the comonad law for w with components  $\delta_A^w$  and  $\varepsilon_{wA}^w$ . (2) commutes by the comonad law for c with components  $\delta_{w^2A}^c$  and  $\varepsilon_{w^2A}^c$ . (3) and (7) commute by the naturality of  $\varepsilon^c$ . (4) commutes by the condition of dist. (5) commutes trivially. And (6) commutes by the naturality of  $\varepsilon^w$ .

## Case 3:



The previous triangle commutes because the following diagram chasing does.

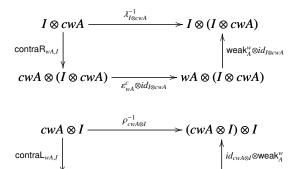


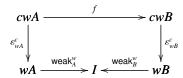
(1) commutes by the naturality of  $\delta^c$ . (2) is the comonad law for c with components  $\delta^c_{wA}$  and  $\varepsilon^c_{wA}$ . (3) is the comonad law for w with components  $\delta^w_A$  and  $\varepsilon^w_A$ . (4) commutes by the condition of dist. And (5) commute by the naturality of dist.

**Definition 9.** A Lambek category with cw,  $(\mathcal{L}, cw, weak^w, contral, contral)$ , is specified by

- a monoidal category  $(\mathcal{L}, I, \otimes)$ ;
- a monoidal comonad  $(c, \varepsilon^c, \delta^c)$  with monoidal natural transformations contral and contraR on  $\mathcal{L}$  s.t.  $(\mathcal{L}, c, \text{contraL}, \text{contraR})$  is a Lambek category with contraction;
- a monoidal comonad  $(w, \varepsilon^w, \delta^w)$  with a monoidal natural transformation weak on  $\mathcal{L}$  s.t.  $(\mathcal{L}, w, \text{weak}^w)$  is a Lambek category with weakening;
- a natural transformation with components  $dist_A : cwA \longrightarrow wcA$ ;

where cw is the composite monoidal comonad cw defined in the proof of Lemma 8, s.t. the following additional coherence diagrams commute:



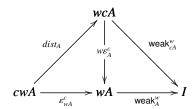


for any coalgebra morphism  $f:(cwA, \delta_A) \longrightarrow (cwB, \delta_B)$  between free coalgebras.

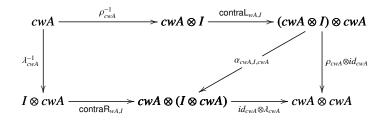
**Lemma 10.** Let  $(\mathcal{L}, cw, weak^w, contral, contral)$  be a Lambek category with cw. Then the following three conditions are satisfied:

- 1. There exist two natural transformations weak<sub>A</sub>:  $cwA \longrightarrow I$  and  $contra_A$ :  $cwA \longrightarrow cwA \otimes cwA$ .
- 2. Each (cwA, weak<sub>A</sub>, contra<sub>A</sub>) is a comonoid.
- 3. weak<sub>A</sub> and contra<sub>A</sub> are coalgebra morphisms.
- 4. Any coalgebra morphism  $f:(cwA, \delta_A) \longrightarrow (cwB, \delta_B)$  between free coalgebras preserves the comonoid structure given by weak and contra.

*Proof.* 1. Each of weak and contracan be given two equivalent definitions. weak<sub>A</sub>:  $cwA \longrightarrow I$  is defined as in the diagram below. The left triangle commutes by the definition of *dist* and the right triangle commutes by the definition of weak<sup>w</sup>.



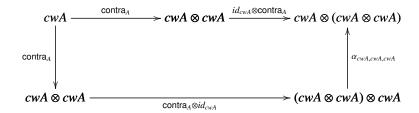
contra<sub>A</sub>:  $cwA \longrightarrow cwA \otimes cwA$  is defined as below. The left part of the diagram commutes by the definitions of contraL and of contraR, and the right part commutes because  $\mathcal{L}$  is monoidal.



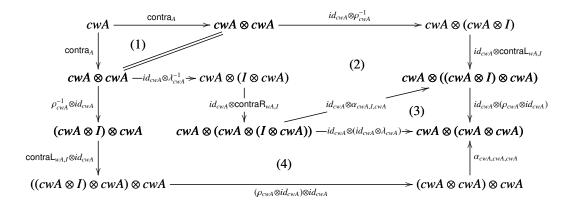
2. Each (cwA, weak<sub>A</sub>, contra<sub>A</sub>) is a comonoid.

I modified this diagram by decomposing weak as weak  $^{w} \circ \varepsilon$  because we haven't defined weak yet.

#### Case 1:

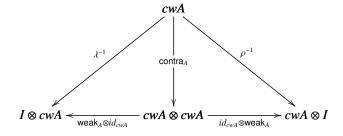


The previous diagram commutes by the following diagram chasing.

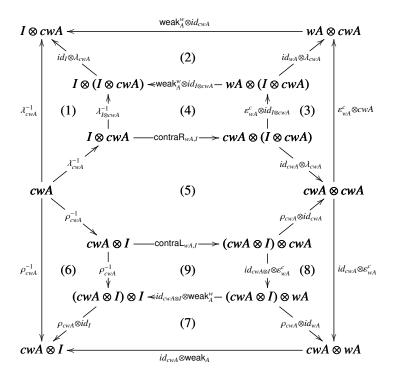


(1) commutes trivially and we would not expand contra for simplicity. (2) and (4) commute because  $(\mathcal{L}, c, \text{contraL}, \text{contraR})$  is a Lambek category with contraction. (3) commutes because  $\mathcal{L}$  is monoidal.

## Case 2:



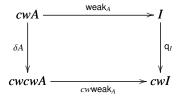
The diagram above commutes by the following diagram chasing.



(1), (2) and (3) commute by the functionality of  $\lambda$ . (6), (7) and (8) commute by the functionality of  $\rho$ . (4) and (9) are conditions of the Lambek category with cw. And (5) is the definition of contra.

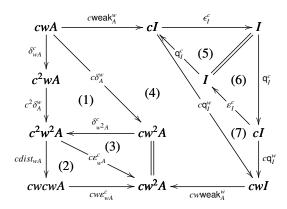
#### 3. weak and contra are coalgebra morphisms.

### Case 1:

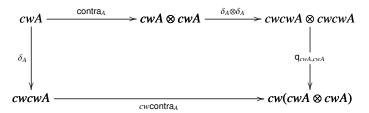


The previous diagram commutes by the diagram below. (1) commutes by the naturality of  $\delta^c$ . (2) commutes by the condition of  $dist_{wA}$ . (3), (5) and (6) commute because c is a monoidal comonad. (4) commutes because  $(\mathcal{L}, w, \mathsf{weak}^w)$  is a Lambek category with weakening. (7) commutes be-

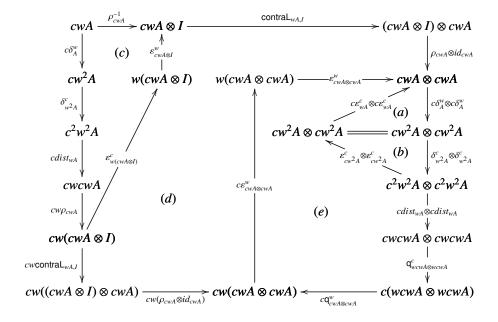
cause c and w are monoidal comonads.



#### Case 2:

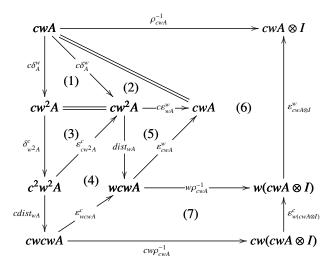


To prove the previous diagram commute, we first expand it, Then we divide it into five parts as shown belovee, and prove each part commutes.

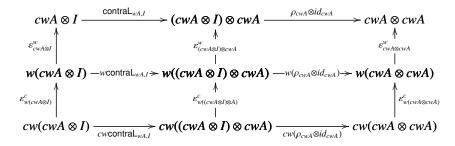


Part (a) and (b) are comonad laws.

Part (c) commutes by the following diagram chase. (1) is equality. (2) is the comonad law for w. (3) is the comonad law for c. (4) commutes by the naturality of  $\varepsilon^c$ . (5) is one of the conditions for  $dist_{wA}$ . (6) commutes by the naturality of  $\varepsilon^w$ . And (7) commutes by the naturality of  $\varepsilon^c$ .

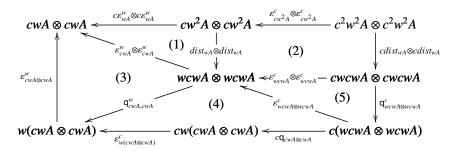


Part (d) commutes by the following diagram chase. The upper two squares both commute by the naturality of  $\varepsilon^w$ , and the lower two squares commute by the naturality of  $\varepsilon^c$ .



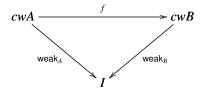
Part (e) commutes by the following diagram. (1) commutes by the condition of  $dist_{wA}$ . (2) and (4) commute by the naturality of  $\varepsilon^c$ . (3) and (5)

commute because w and c are monoidal comonads.



4. Any coalgebra morphism  $f:(cwA, \delta_A) \longrightarrow (cwB, \delta_B)$  between free coalgebras preserves the comonoid structure given by weak and contra.

Case 1: This coherence diagram is given in the definition of the Lambek category with cw.



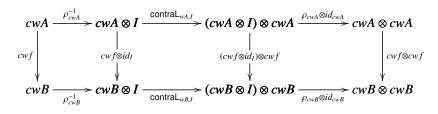
#### Case 2:

$$cwA \xrightarrow{\text{contra}_A} cwA \otimes cwA$$

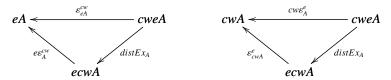
$$\downarrow f \qquad \qquad \downarrow f \otimes f$$

$$cwB \xrightarrow{\text{contra}_B} cwB \otimes cwB$$

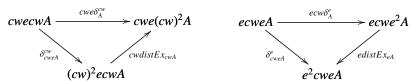
The square commutes by the diagram chasing below, which commutes by the naturality of  $\rho$  and contral.



**Definition 11.** Given two comonads  $(cw, \varepsilon^{cw}, \delta^{cw})$  and  $(e, \varepsilon^e, \delta^e)$  on a category  $\mathcal{L}$  such that  $(\mathcal{L}, cw, weak, contra)$  is a Lambek category with cw and  $(\mathcal{L}, e, ex)$  is a Lambek category with exchange, we define a **distributive law for exchange** of cw over e to be a natural isomorphism with components  $distEx_A : cweA \longrightarrow ecwA$ , subject to the following coherence diagrams:



**Lemma 12.** Given two comonads  $(cw, \varepsilon^{cw}, \delta^{cw})$  and  $(e, \varepsilon^{e}, \delta^{e})$  on a category  $\mathcal{L}$  such that  $(\mathcal{L}, cw, weak, contra)$  is a Lambek category with cw and  $(\mathcal{L}, e, ex)$  is a Lambek category with exchange, the following two digrams also commute:



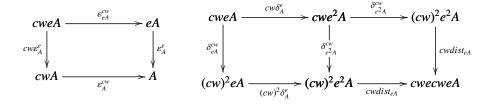
The proof is similar with the proof of Lemma 7 and we will not elaborate it here. Also, notice the difference between dist of c over w and distEx of cw over e. While dist is a natural transformation, distEx is a natural isomorphism.

**Lemma 13.** let  $(cw, \varepsilon^{cw}, \delta^{cw})$  and  $(e, \varepsilon^e, \delta^e)$  be two monoidal comonads on a Lambek category with cw and exchange  $(\mathcal{L}, I, \otimes, cw, weak, contra, e, ex)$ . Then the composition of cw and e using the distributive law for exchange dist $Ex_A$ :  $cweA \longrightarrow ecwA$  is a monoidal comonad  $(cwe, \varepsilon, \delta)$  on  $\mathcal{L}$ .

*Proof.* Suppose  $(cw, \varepsilon^{cw}, \delta^{cw})$  and  $(e, \varepsilon^e, \delta^e)$  are monoidal comonads, and  $(\mathcal{L}, I, \otimes, cw, \text{weak}, \text{contra}, e, \text{ex})$  is a Lambek category with cw and exchange. Since by definition  $cw, e: \mathcal{L} \longrightarrow \mathcal{L}$  are monoidal functors, we know that their composition  $cwe: \mathcal{L} \longrightarrow \mathcal{L}$  is a monoidal functor:

$$\begin{aligned} \mathsf{q}_{A,B} &: cweA \otimes cweB \longrightarrow cwe(A \otimes B) \\ \mathsf{q}_{A,B} &= cw\mathsf{q}_{A,B}^e \circ \mathsf{q}_{eA,eB}^{cw} \\ \mathsf{q}_I &: I \longrightarrow cweI \\ \mathsf{q}_I &= cw\mathsf{q}_I^e \circ \mathsf{q}_I^{cw} \end{aligned}$$

Analogous to the proof of Lemma 8, each of  $\varepsilon$  and  $\delta$  can be given two equivalent definitions:



And the comonad laws can be proved similarly, which we will not elaborate for simplicity.

**Lemma 14.** Let  $(cwe, \varepsilon, \delta)$  be a monoidal comonad over a monoidal category  $(\mathcal{L}, I, \otimes)$  such that  $(\mathcal{L}, I, \otimes, cw, weak, contra, e, ex)$  is a Lambek category with cw and exchange. Then the co-Kleisli category of  $\mathcal{L}$ ,  $\mathcal{L}_{cwe}$ , is a symmetric monoidal category.

*Proof.* The identity object of  $\mathcal{L}_{cwe}$  is still *I*.

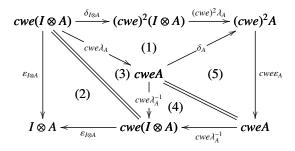
The left and right unitors,  $\hat{\lambda}_A: I \otimes A \longrightarrow A$  and  $\hat{\rho}_A: A \otimes I \longrightarrow A$ , in  $\mathcal{L}_{cwe}$  are morphisms  $cwe(I \otimes A) \longrightarrow A$  and  $cwe(A \otimes I) \longrightarrow A$  in  $\mathcal{L}$ , respectively. Then we define  $\hat{\lambda}$  and  $\hat{\rho}$  as:

$$\hat{\lambda}_A = \varepsilon_A \circ cwe\lambda_A$$
$$\hat{\rho}_A = \varepsilon_A \circ cwe\rho_A,$$

where  $\lambda$  and  $\rho$  are the left and right unitors in  $\mathcal{L}$ , respectively. And we define their inverses as:

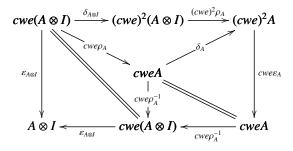
$$\hat{\lambda}_{A}^{-1} = \varepsilon_{I \otimes A} \circ cwe \lambda_{A}^{-1}$$
$$\hat{\rho}_{A}^{-1} = \varepsilon_{A \otimes I} \circ cwe \rho_{A}^{-1}$$

 $\hat{\lambda}$  is a nautral isomorphism with inverse  $\hat{\lambda}^{-1}$  because the following diagram chasing commutes:



(1) commutes by the naturality of  $\delta$ . (2), (3) and (4) commute trivially. And (5) commutes because *cwe* is a comonad.

Similarly,  $\hat{\rho}$  is a natural isomorphism with inverse  $\hat{\rho}^{-1}$  by the following diagram chasing:



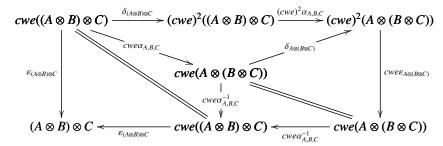
The associator  $\hat{\alpha}_A$ :  $(A \otimes B) \otimes C) \longrightarrow A \otimes (B \otimes C)$  in  $\mathcal{L}_{cwe}$  is the morphism  $cwe((A \otimes B) \otimes C) \longrightarrow A \otimes (B \otimes C)$  in  $\mathcal{L}$ . We define  $\hat{\alpha}$  as:

$$\hat{\alpha}_{A,B,C} = \varepsilon_{A\otimes (B\otimes C)} \circ cwe\alpha_{A,B,C},$$

where  $\alpha$  is the associator of  $\mathcal{L}$ . And its inverse is

$$\hat{\alpha}_{A.B.C}^{-1} = \varepsilon_{(A \otimes B) \otimes C} \circ cwe\alpha_{A.B.C}^{-1}$$

 $\hat{\alpha}$  is a natural isomorphism with inverse  $\hat{\alpha}^{-1}$  because the following diagram chasing commutes:



Therefore,  $\mathcal{L}_{cwe}$  is a monoidal category.

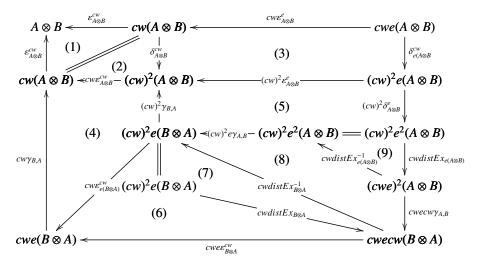
The symmetry,  $\hat{\beta}_{A,B}: A \otimes B \longrightarrow B \otimes A$ , in  $\mathcal{L}_{cwe}$  is the morphism  $cwe(A \otimes B) \longrightarrow B \otimes A$  in  $\mathcal{L}$ , which is defined as:

$$\hat{\beta}_{A,B} = \varepsilon_{B\otimes A}^{cw} \circ cw\gamma_{A,B},$$

where  $\varepsilon_A^{cw}: cwA \longrightarrow A$  is a natural transformation associated with the comonad cw, and  $\gamma$  is the natural isomorphism defined in Lemma 4. Then its inverse is

$$\hat{\beta}_{A,B}^{-1} = \varepsilon_{A\otimes B}^{cw} \circ cw\gamma_{B,A}$$

 $\hat{\beta}$  is a natural isomorphism with inverse  $\hat{\beta}^{-1}$  because the following diagram chasing commutes:



(1), (7) and (9) commute trivially. (2) is the comonad law for cw. (3) commutes by the naturality of  $\delta^{cw}$ . (4) commutes by the naturality of  $\varepsilon^{cw}$ . (5) commutes because  $\gamma$  is a natural isomorphism (Lemma 4). (6) is the definition of distEx. (8) is the naturality of distEx.

In conclusion,  $\mathcal{L}_{cwe}$  is a symmetric monoidal category.

## 3 Related Work

**TODO** 

### 4 Conclusion

**TODO** 

### References

[1] P. N. Benton. A mixed linear and non-linear logic: Proofs, terms and models (preliminary report). Technical Report UCAM-CL-TR-352, University of Cambridge Computer Laboratory, 1994. Accessible online at http://research.microsoft.com/en-us/um/people/nick/mixed3.ps.

## A Appendix

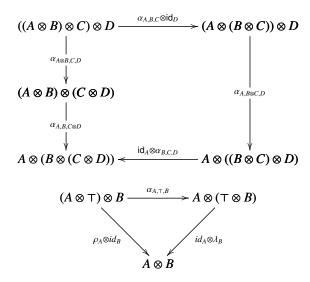
### A.1 Symmetric Monoidal Categories

**Definition 15.** A monoidal category is a category, M, with the following data:

- An object  $\top$  of  $\mathcal{M}$ ,
- A bi-functor  $\otimes$  :  $\mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$ ,
- The following natural isomorphisms:

 $\begin{array}{l} \lambda_A: \top \otimes A \longrightarrow A \\ \rho_A: A \otimes \top \longrightarrow A \\ \alpha_{A,B,C}: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C) \end{array}$ 

• Subject to the following coherence diagrams:



**Definition 16.** A symmetric monoidal category (SMC) is a category, M, with the following data:

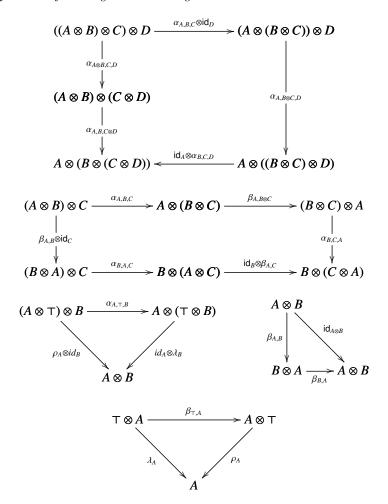
- An object  $\top$  of  $\mathcal{M}$ ,
- A bi-functor  $\otimes$  :  $\mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$ ,
- The following natural isomorphisms:

$$\begin{array}{l} \lambda_A: \top \otimes A \longrightarrow A \\ \rho_A: A \otimes \top \longrightarrow A \\ \alpha_{A,B,C}: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C) \end{array}$$

• A symmetry natural isomorphism:

$$\beta_{A,B}: A \otimes B \longrightarrow B \otimes A$$

• Subject to the following coherence diagrams:



**Definition 17.** A monoidal biclosed category is a monoidal category  $(M, \top, \otimes)$ , such that, for any object B of M, each of the functors  $-\otimes B : M \longrightarrow M$  and  $B \otimes - : M \longrightarrow M$  has a specified right adjoint. Hence, for any object A and C of M, there are two objects  $C \hookrightarrow B$  and  $B \to C$  of M and two natural bijections:

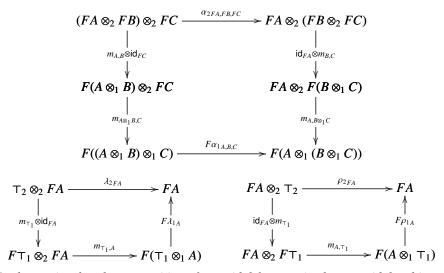
$$\operatorname{\mathsf{Hom}}_{\mathcal{M}}(A\otimes B,C)\cong\operatorname{\mathsf{Hom}}_{\mathcal{M}}(A,C\leftharpoonup B)$$
  
 $\operatorname{\mathsf{Hom}}_{\mathcal{M}}(B\otimes A,C)\cong\operatorname{\mathsf{Hom}}_{\mathcal{M}}(A,B\rightharpoonup C)$ 

**Definition 18.** A symmetric monoidal closed category (SMCC) is a symmetric monoidal category,  $(\mathcal{M}, \top, \otimes)$ , such that, for any object B of  $\mathcal{M}$ , the functor  $-\otimes B : \mathcal{M} \longrightarrow \mathcal{M}$  has a specified right adjoint. Hence, for any objects A and C of  $\mathcal{M}$  there is an object  $B \multimap C$  of  $\mathcal{M}$  and a natural bijection:

$$\operatorname{\mathsf{Hom}}_{\mathcal{M}}(A \otimes B, C) \cong \operatorname{\mathsf{Hom}}_{\mathcal{M}}(A, B \multimap C)$$

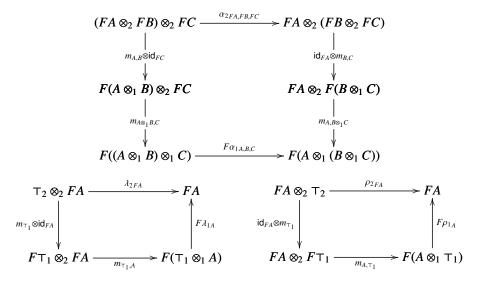
*We call the functor*  $\multimap$ :  $\mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$  *the internal hom of*  $\mathcal{M}$ .

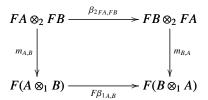
**Definition 19.** Suppose we are given two monoidal categories  $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$ . Then a **monoidal functor** is a functor  $F: \mathcal{M}_1 \longrightarrow \mathcal{M}_2$ , a map  $m_{\top_1} : \top_2 \longrightarrow F \top_1$  and a natural transformation  $m_{A,B} : FA \otimes_2 FB \longrightarrow F(A \otimes_1 B)$  subject to the following coherence conditions:



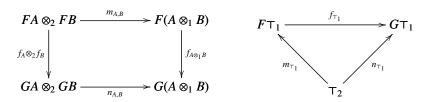
Need to notice that the composition of monoidal functors is also monoidal, subject to the above coherence conditions.

**Definition 20.** Suppose we are given two symmetric monoidal closed categories  $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$ . Then a **symmetric monoidal** functor is a functor  $F: \mathcal{M}_1 \longrightarrow \mathcal{M}_2$ , a map  $m_{\top_1}: \top_2 \longrightarrow F \top_1$  and a natural transformation  $m_{A,B}: FA \otimes_2 FB \longrightarrow F(A \otimes_1 B)$  subject to the following coherence conditions:

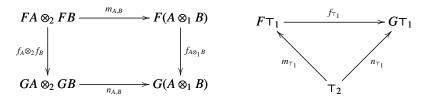




**Definition 21.** Suppose  $(\mathcal{M}_1, \top_1, \otimes_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2)$  are monoidal categories, and (F, m) and (G, n) are monoidal functors between  $\mathcal{M}_1$  and  $\mathcal{M}_2$ . Then a **monoidal natural transformation** is a natural transformation,  $f: F \longrightarrow G$ , subject to the following coherence diagrams:

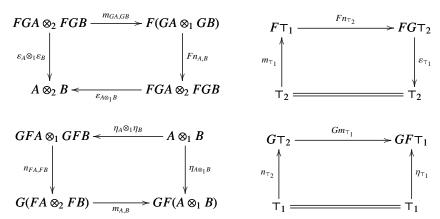


**Definition 22.** Suppose  $(\mathcal{M}_1, \top_1, \otimes_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2)$  are SMCs, and (F, m) and (G, n) are symmetric monoidal functors between  $\mathcal{M}_1$  and  $\mathcal{M}_2$ . Then a symmetric monoidal natural transformation is a natural transformation,  $f: F \longrightarrow G$ , subject to the following coherence diagrams:

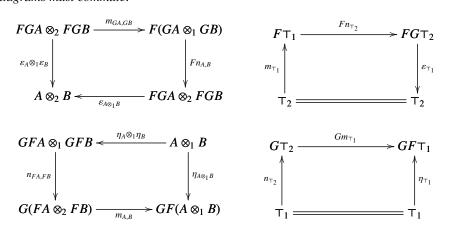


**Definition 23.** Suppose  $(M_1, \top_1, \otimes_1)$  and  $(M_2, \top_2, \otimes_2)$  are monoidal categories, and (F, m) is a monoidal functor between  $M_1$  and  $M_2$  and (G, n) is a monoidal functor between  $M_2$  and  $M_1$ . Then a **monoidal adjunction** is an ordinary adjunction  $M_1$ :  $F \dashv G : M_2$  such that the unit,  $\eta_A : A \to GFA$ , and the counit,  $\varepsilon_A : FGA \to A$ , are

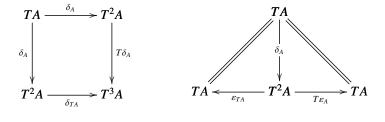
monoidal natural transformations. Thus, the following diagrams must commute:



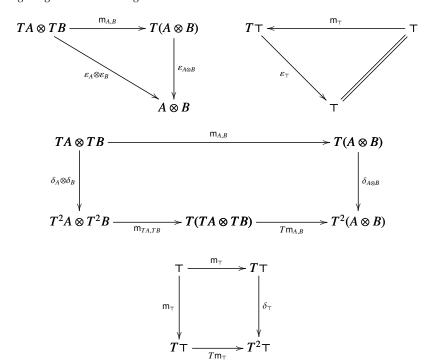
**Definition 24.** Suppose  $(M_1, \top_1, \otimes_1)$  and  $(M_2, \top_2, \otimes_2)$  are SMCs, and (F, m) is a symmetric monoidal functor between  $M_1$  and  $M_2$  and (G, n) is a symmetric monoidal functor between  $M_2$  and  $M_1$ . Then a **symmetric monoidal adjunction** is an ordinary adjunction  $M_1: F \dashv G: M_2$  such that the unit,  $\eta_A: A \to GFA$ , and the counit,  $\varepsilon_A: FGA \to A$ , are symmetric monoidal natural transformations. Thus, the following diagrams must commute:



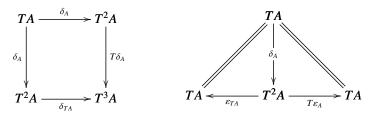
**Definition 25.** A monoidal comonad on a monoidal category C is a triple  $(T, \varepsilon, \delta)$ , where  $(T, \mathsf{m})$  is a monoidal endofunctor on C,  $\varepsilon_A : TA \longrightarrow A$  and  $\delta_A : TA \to T^2A$  are monoidal natural transformations, which make the following diagrams commute:



The assumption that  $\varepsilon$  and  $\delta$  are monoidal natural transformations amount to the following diagrams commuting:



**Definition 26.** A symmetric monoidal comonad on a symmetric monoidal category C is a triple  $(T, \varepsilon, \delta)$ , where (T, m) is a symmetric monoidal endofunctor on C,  $\varepsilon_A$ :  $TA \longrightarrow A$  and  $\delta_A : TA \to T^2A$  are symmetric monoidal natural transformations, which make the following diagrams commute:



The assumption that  $\varepsilon$  and  $\delta$  are symmetric monoidal natural transformations amount to the following diagrams commuting:

