

# Separating Linear Modalities

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## Abstract

TODO

## 1 Introduction

TODO [1]

### 1.1 Symmetric Monoidal Categories

**Definition 1** A *monoidal category* is a category,  $\mathcal{M}$ , with the following data:

- An object  $\top$  of  $\mathcal{M}$ ,
- A bi-functor  $\otimes : \mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$ ,
- The following natural isomorphisms:

$$\begin{aligned}\lambda_A &: \top \otimes A \longrightarrow A \\ \rho_A &: A \otimes \top \longrightarrow A \\ \alpha_{A,B,C} &: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C)\end{aligned}$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc} ((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\ \downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\ (A \otimes B) \otimes (C \otimes D) & & A \otimes ((B \otimes C) \otimes D) \\ \downarrow \alpha_{A, B, C \otimes D} & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & \\ A \otimes (B \otimes (C \otimes D)) & & A \otimes ((B \otimes C) \otimes D) \end{array}$$
  

$$\begin{array}{ccc} (A \otimes \top) \otimes B & \xrightarrow{\alpha_{A, \top, B}} & A \otimes (\top \otimes B) \\ \searrow \rho_A & & \swarrow \lambda_B \\ & A \otimes B & \end{array}$$

**Definition 2** A *symmetric monoidal category (SMC)* is a category,  $\mathcal{M}$ , with the following data:

- An object  $\top$  of  $\mathcal{M}$ ,
- A bi-functor  $\otimes : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$ ,
- The following natural isomorphisms:

$$\begin{aligned}\lambda_A &: \top \otimes A \rightarrow A \\ \rho_A &: A \otimes \top \rightarrow A \\ \alpha_{A,B,C} &: (A \otimes B) \otimes C \rightarrow A \otimes (B \otimes C)\end{aligned}$$

- A symmetry natural transformation:

$$\beta_{A,B} : A \otimes B \rightarrow B \otimes A$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc} ((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\ \downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\ (A \otimes B) \otimes (C \otimes D) & & \\ \downarrow \alpha_{A, B, C \otimes D} & & \\ A \otimes (B \otimes (C \otimes D)) & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & A \otimes ((B \otimes C) \otimes D) \end{array}$$
  

$$\begin{array}{ccccc} (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\beta_{A, B \otimes C}} & (B \otimes C) \otimes A \\ \downarrow \beta_{A,B} \otimes \text{id}_C & & & & \downarrow \alpha_{B,C,A} \\ (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{\text{id}_B \otimes \beta_{A,C}} & B \otimes (C \otimes A) \end{array}$$
  

$$\begin{array}{ccc} (A \otimes \top) \otimes B & \xrightarrow{\alpha_{A,\top,B}} & A \otimes (\top \otimes B) \\ \searrow \rho_A & & \swarrow \lambda_B \\ & A \otimes B & \end{array}$$
  

$$\begin{array}{ccc} A \otimes B & & \\ \downarrow \beta_{A,B} & \searrow \text{id}_{A \otimes B} & \\ B \otimes A & \xrightarrow{\beta_{B,A}} & A \otimes B \end{array}$$
  

$$\begin{array}{ccc} \top \otimes A & \xrightarrow{\beta_{\top,A}} & A \otimes \top \\ \searrow \lambda_A & & \swarrow \rho_A \\ & A & \end{array}$$

**Definition 3** A *symmetric monoidal closed category (SMCC)* is a symmetric monoidal category,  $(\mathcal{M}, \top, \otimes)$ , such that, for any object  $B$  of  $\mathcal{M}$ , the functor  $- \otimes B : \mathcal{M} \rightarrow \mathcal{M}$  has a specified right adjoint. Hence, for any objects  $A$  and  $C$  of  $\mathcal{M}$  there is an object  $B \multimap C$  of  $\mathcal{M}$  and a natural bijection:

$$\text{Hom}_{\mathcal{M}}(A \otimes B, C) \cong \text{Hom}_{\mathcal{M}}(A, B \multimap C)$$

We call the functor  $\multimap : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$  the internal hom of  $\mathcal{M}$ .

**Definition 4** Suppose we are given two monoidal categories  $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$ . Then a **monoidal functor** is a functor  $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$ , a map  $m_{\top_1} : \top_2 \rightarrow F\top_1$  and a natural transformation  $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$  subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A, B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$
  

$$\begin{array}{ccc} \top_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\ \downarrow m_{\top_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\ F\top_1 \otimes_2 FA & \xrightarrow{m_{\top_1, A}} & F(\top_1 \otimes_1 A) \end{array} \quad \begin{array}{ccc} FA \otimes_2 \top_2 & \xrightarrow{\rho_{2FA}} & FA \\ \downarrow \text{id}_{FA} \otimes m_{\top_1} & & \uparrow F\rho_{1A} \\ FA \otimes_2 F\top_1 & \xrightarrow{m_{A, \top_1}} & F(A \otimes_1 \top_1) \end{array}$$

**Definition 5** Suppose we are given two symmetric monoidal closed categories  $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$ . Then a **symmetric monoidal functor** is a functor  $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$ , a map  $m_{\top_1} : \top_2 \rightarrow F\top_1$  and a natural transformation  $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$  subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A, B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$

$$\begin{array}{ccc}
\top_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\
\downarrow m_{\top_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\
F\top_1 \otimes_2 FA & \xrightarrow{m_{\top_1, A}} & F(\top_1 \otimes_1 A)
\end{array}
\quad
\begin{array}{ccc}
FA \otimes_2 \top_2 & \xrightarrow{\rho_{2FA}} & FA \\
\downarrow \text{id}_{FA} \otimes m_{\top_1} & & \uparrow F\rho_{1A} \\
FA \otimes_2 F\top_1 & \xrightarrow{m_{A, \top_1}} & F(A \otimes_1 \top_1)
\end{array}$$
  

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{\beta_{2FA, FB}} & FB \otimes_2 FA \\
\downarrow m_{A, B} & & \downarrow m_{B, A} \\
F(A \otimes_1 B) & \xrightarrow{F\beta_{1A, B}} & F(B \otimes_1 A)
\end{array}$$

**Definition 6** Suppose  $(\mathcal{M}_1, \top_1, \otimes_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2)$  are monoidal categories, and  $(F, m)$  and  $(G, n)$  are monoidal functors between  $\mathcal{M}_1$  and  $\mathcal{M}_2$ . Then a **monoidal natural transformation** is a natural transformation,  $f : F \rightarrow G$ , subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A, B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A, B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\top_1 & \xrightarrow{f_{\top_1}} & G\top_1 \\
\swarrow m_{\top_1} & & \searrow n_{\top_1} \\
& \top_2 &
\end{array}$$

**Definition 7** Suppose  $(\mathcal{M}_1, \top_1, \otimes_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2)$  are SMCs, and  $(F, m)$  and  $(G, n)$  are symmetric monoidal functors between  $\mathcal{M}_1$  and  $\mathcal{M}_2$ . Then a **symmetric monoidal natural transformation** is a natural transformation,  $f : F \rightarrow G$ , subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A, B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A, B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\top_1 & \xrightarrow{f_{\top_1}} & G\top_1 \\
\swarrow m_{\top_1} & & \searrow n_{\top_1} \\
& \top_2 &
\end{array}$$

**Definition 8** Suppose  $(\mathcal{M}_1, \top_1, \otimes_1)$  and  $(\mathcal{M}_2, \top_2, \otimes_2)$  are monoidal categories, and  $(F, m)$  is a monoidal functor between  $\mathcal{M}_1$  and  $\mathcal{M}_2$  and  $(G, n)$  is a monoidal functor between  $\mathcal{M}_2$  and  $\mathcal{M}_1$ . Then a **monoidal adjunction** is an ordinary adjunction  $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$  such that the unit,  $\eta_A : A \rightarrow GFA$ , and the counit,  $\varepsilon_A : FGA \rightarrow A$ ,

are monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow Fn_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_{A \otimes_1 B}} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{Fn_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{Gm_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

**Definition 9** Suppose  $(\mathcal{M}_1, \tau_1, \otimes_1)$  and  $(\mathcal{M}_2, \tau_2, \otimes_2)$  are SMCs, and  $(F, m)$  is a symmetric monoidal functor between  $\mathcal{M}_1$  and  $\mathcal{M}_2$  and  $(G, n)$  is a symmetric monoidal functor between  $\mathcal{M}_2$  and  $\mathcal{M}_1$ . Then a **symmetric monoidal adjunction** is an ordinary adjunction  $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$  such that the unit,  $\eta_A : A \rightarrow GFA$ , and the counit,  $\varepsilon_A : FGA \rightarrow A$ , are symmetric monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow Fn_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_{A \otimes_1 B}} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{Fn_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{Gm_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

**Definition 10** A **monoidal comonad** on a monoidal category  $\mathcal{C}$  is a triple  $(T, \varepsilon, \delta)$ , where  $(T, m)$  is a monoidal endofunctor on  $\mathcal{C}$ ,  $\varepsilon_A : TA \rightarrow A$  and  $\delta_A : TA \rightarrow T^2A$  are monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\quad
\begin{array}{ccccc}
& & TA & & \\
& \swarrow & \downarrow \delta_A & \searrow & \\
TA & \xleftarrow{\varepsilon_{TA}} & T^2A & \xrightarrow{T\varepsilon_A} & TA
\end{array}$$

The assumption that  $\varepsilon$  and  $\delta$  are monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\qquad
\begin{array}{ccc}
T\top & \xrightarrow{m_\top} & \top \\
& \searrow \varepsilon_\top & \downarrow \\
& & \top
\end{array}$$

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} T(TA \otimes TB) \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}$$

$$\begin{array}{ccc}
\top & \xrightarrow{m_\top} & T\top \\
\downarrow m_\top & & \downarrow \delta_\top \\
T\top & \xrightarrow{Tm_\top} & T^2\top
\end{array}$$

**Definition 11** A *symmetric monoidal comonad* on a symmetric monoidal category  $\mathcal{C}$  is a triple  $(T, \varepsilon, \delta)$ , where  $(T, m)$  is a symmetric monoidal endofunctor on  $\mathcal{C}$ ,  $\varepsilon_A : TA \rightarrow A$  and  $\delta_A : TA \rightarrow T^2A$  are symmetric monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\qquad
\begin{array}{ccccc}
& & TA & & \\
& \swarrow & \downarrow \delta_A & \searrow & \\
TA & \xleftarrow{\varepsilon_{TA}} & T^2A & \xrightarrow{T\varepsilon_A} & TA
\end{array}$$

The assumption that  $\varepsilon$  and  $\delta$  are symmetric monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\qquad
\begin{array}{ccc}
T\top & \xrightarrow{m_\top} & \top \\
& \searrow \varepsilon_\top & \downarrow \\
& & \top
\end{array}$$

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} T(TA \otimes TB) \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}$$
  

$$\begin{array}{ccc}
\mathbb{T} & \xrightarrow{m_{\mathbb{T}}} & T\mathbb{T} \\
\downarrow m_{\mathbb{T}} & & \downarrow \delta_{\mathbb{T}} \\
T\mathbb{T} & \xrightarrow{Tm_{\mathbb{T}}} & T^2\mathbb{T}
\end{array}$$

## 1.2 LNL Model

**Definition 12** A *linear/non-linear (LNL) model*  $(\mathcal{C}, \mathcal{L}, F, G)$  consists of

- a cartesian closed category  $(\mathcal{C}, 1, \times, \longrightarrow)$
- a symmetric monoidal closed category  $(\mathcal{L}, I, \otimes, \multimap)$ ,
- a pair of symmetric monoidal functors  $(G, n) : \mathcal{L} \longrightarrow \mathcal{C}$  and  $(F, m) : \mathcal{C} \longrightarrow \mathcal{L}$  that form a symmetric monoidal adjunction  $\mathcal{C} : F \dashv G : \mathcal{L}$ , subject to the following coherence conditions, where  $\eta$  and  $\varepsilon$  are the unit and counit of the adjunction and  $A, B \in \text{Ob}(\mathcal{L})$ ,  $X, Y \in \text{Ob}(\mathcal{C})$ .

$$\begin{array}{ccc}
GA & \xrightarrow{\eta_{GA}} & GFGA \\
\parallel & \searrow G\varepsilon_A & \\
GA & & 
\end{array}
\quad
\begin{array}{ccc}
FX & \xrightarrow{F\eta_X} & FGFX \\
\parallel & \searrow \varepsilon_{FX} & \\
FX & & 
\end{array}$$
  

$$\begin{array}{ccc}
X \times Y & \xrightarrow{\eta_{X \times Y}} & GF(X \times Y) \\
\downarrow \eta_{X \times Y} & & \downarrow Gm_{X,Y} \\
GF(X \times Y) & \xleftarrow{Gm_{X,Y}} & G(FX \otimes GY)
\end{array}
\quad
\begin{array}{ccc}
FGA \otimes FGB & \xrightarrow{m_{GA,GB}} & F(GA \times GB) \\
\downarrow \varepsilon_A \otimes \varepsilon_B & & \downarrow F\eta_{A,B} \\
A \otimes B & \xleftarrow{\eta_{A \otimes B}} & FG(A \otimes B)
\end{array}$$
  

$$\begin{array}{ccc}
1 & \xrightarrow{\eta_1} & GF(1) \\
\parallel & \searrow Gm_{\mathbb{T},1} \circ \eta_I & \\
1 & & 
\end{array}
\quad
\begin{array}{ccc}
FG(I) & \xrightarrow{\varepsilon_I} & I \\
\downarrow F\eta_I \circ m_1 & & \parallel \\
I & & 
\end{array}$$

**Lemma 13** Given a LNL model  $(C, \mathcal{L}, F, G)$ , define  $\mathfrak{p}_{X,Y} : F(X \times Y) \longrightarrow FX \otimes FY$  as the composition  $\varepsilon_{FX \otimes FY} \circ F\eta_{FX, FY} \circ F(\eta_X \times \eta_Y)$ , and  $\mathfrak{p}_I : F1 \longrightarrow I$  as the composition  $\varepsilon_I \circ F\eta_I$ . Then for  $F \dashv G$ ,  $\mathfrak{m}_{X,Y}$  are components of a natural isomorphism with inverses  $\mathfrak{p}_{X,Y}$ , and  $\mathfrak{m}_I$  is an isomorphism with inverse  $\mathfrak{p}_I$ , i.e.  $F(X) \otimes F(Y) \cong F(X \times Y)$ , and  $I \cong F(1)$ .

**Lemma 14** Given a LNL model  $(C, \mathcal{L}, F, G)$ , the adjunction  $F \dashv G$  induces a symmetric monoidal comonad  $(!, \varepsilon, \delta)$  on  $\mathcal{L}$ , where  $!$  represents  $FG$ ,  $\varepsilon : FG \longrightarrow 1$  is the counit of the adjunction, and  $\delta : FG \longrightarrow FGFG$ , i.e.  $!$  is a symmetric monoidal functor and  $\varepsilon, \delta$  are monoidal natural transformations.

**Lemma 15** Given a LNL model  $(C, \mathcal{L}, F, G)$ ,  $F$  is a strong functor, i.e.  $F$  preserves the monoidal structure up to an isomorphism. And a strong functor induces a unique monoidal structure.

**Definition 16** A *linear category* is specified by

- a symmetric monoidal closed category  $(\mathcal{L}, I, \otimes, \multimap)$ ,
- a symmetric monoidal comonad  $(!, \varepsilon, \delta)$  on  $\mathcal{L}$ , with  $\mathfrak{q}_{A,B} : !A \otimes !B \longrightarrow !(A \otimes B)$  and  $\mathfrak{q}_I : I \longrightarrow !I$ ;
- monoidal natural transformations on  $\mathcal{L}$  with components  $\mathfrak{e}_A : !A \longrightarrow I$  and  $\mathfrak{d}_A : !A \longrightarrow !A \otimes !A$ , s.t.
  - each  $(!A, \mathfrak{e}_A, \mathfrak{d}_A)$  is a commutative comonoid, i.e. the following diagrams commute and  $\beta \circ \mathfrak{d}_A = \mathfrak{d}_A$  where  $\beta_{B,C} : B \otimes C \longrightarrow C \otimes B$  is the symmetry natural transformation of  $\mathcal{L}$ ;

$$\begin{array}{ccccc}
!A & \xrightarrow{\mathfrak{d}_A} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \mathfrak{d}_A} & !A \otimes (!A \otimes !A) \\
\downarrow \mathfrak{d}_A & & & & \downarrow \alpha_{!A, !A, !A} \\
!A \otimes !A & \xrightarrow{\mathfrak{d}_A \otimes id_{!A}} & & & (!A \otimes !A) \otimes !A
\end{array}$$

$$\begin{array}{ccccc}
& & !A & & \\
& \nearrow \lambda & \downarrow \mathfrak{d}_A & \nwarrow \rho & \\
I \otimes !A & \xleftarrow{\mathfrak{e}_A \otimes id_{!A}} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \mathfrak{e}_A} & !A \otimes I
\end{array}$$

- $\mathfrak{e}_A$  and  $\mathfrak{d}_A$  are coalgebra morphisms, i.e. the following diagrams commute;

$$\begin{array}{ccc}
!A & \xrightarrow{\mathfrak{e}_A} & I \\
\downarrow \delta A & & \downarrow \mathfrak{q}_I \\
!!A & \xleftarrow{!e_A} & !I
\end{array}$$



$$\begin{array}{ccccc}
!A & \xrightarrow{d_A} & !A \otimes !A & \xrightarrow{\delta_! A \otimes \delta_A} & !!A \otimes !!A \\
\downarrow \delta_A & & & & \downarrow q_{!A, !A} \\
!!A & \xrightarrow{!d_A} & & & !(A \otimes A)
\end{array}$$

- any coalgebra morphism  $f : (!A, \delta_A) \longrightarrow (!B, \delta_B)$  between free coalgebras preserve the comonoid structure given by  $e$  and  $d$ , i.e. the following diagrams commute.

$$\begin{array}{ccc}
!A & \xrightarrow{f} & !B \\
& \searrow e_A & \nearrow e_B \\
& I &
\end{array}
\qquad
\begin{array}{ccc}
!A & \xrightarrow{d_A} & !A \otimes !A \\
\downarrow f & & \downarrow f \otimes f \\
!B & \xrightarrow{d_B} & !B \otimes !B
\end{array}$$

**Theorem 17** Any LNL model is a linear category.

**Theorem 18** Any linear category gives rise to an LNL model, though it is not in general unique.

## 2 Related Work

TODO

## 3 Conclusion

TODO

## References

- [1] P. N. Benton. A mixed linear and non-linear logic: Proofs, terms and models (preliminary report). Technical Report UCAM-CL-TR-352, University of Cambridge Computer Laboratory, 1994. Accessible online at <http://research.microsoft.com/en-us/um/people/nick/mixed3.ps>.