

Separating Linear Modalities

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Abstract

TODO

1 Introduction

TODO [1]

1.1 Symmetric Monoidal Categories

Definition 1. A *monoidal category* is a category, \mathcal{M} , with the following data:

- An object \top of \mathcal{M} ,
- A bi-functor $\otimes : \mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$,
- The following natural isomorphisms:

$$\begin{aligned}\lambda_A &: \top \otimes A \longrightarrow A \\ \rho_A &: A \otimes \top \longrightarrow A \\ \alpha_{A,B,C} &: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C)\end{aligned}$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc} ((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\ \downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\ (A \otimes B) \otimes (C \otimes D) & & A \otimes ((B \otimes C) \otimes D) \\ \downarrow \alpha_{A, B, C \otimes D} & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & \\ A \otimes (B \otimes (C \otimes D)) & & A \otimes ((B \otimes C) \otimes D) \end{array}$$

$$\begin{array}{ccc} (A \otimes \top) \otimes B & \xrightarrow{\alpha_{A, \top, B}} & A \otimes (\top \otimes B) \\ \searrow \rho_A & & \swarrow \lambda_B \\ & A \otimes B & \end{array}$$

Definition 2. A *symmetric monoidal category (SMC)* is a category, \mathcal{M} , with the following data:

- An object \top of \mathcal{M} ,
- A bi-functor $\otimes : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$,
- The following natural isomorphisms:

$$\begin{aligned}\lambda_A &: \top \otimes A \rightarrow A \\ \rho_A &: A \otimes \top \rightarrow A \\ \alpha_{A,B,C} &: (A \otimes B) \otimes C \rightarrow A \otimes (B \otimes C)\end{aligned}$$

- A symmetry natural transformation:

$$\beta_{A,B} : A \otimes B \rightarrow B \otimes A$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc} ((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\ \downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\ (A \otimes B) \otimes (C \otimes D) & & \\ \downarrow \alpha_{A, B, C \otimes D} & & \\ A \otimes (B \otimes (C \otimes D)) & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & A \otimes ((B \otimes C) \otimes D) \end{array}$$

$$\begin{array}{ccccc} (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\beta_{A, B \otimes C}} & (B \otimes C) \otimes A \\ \downarrow \beta_{A,B} \otimes \text{id}_C & & & & \downarrow \alpha_{B,C,A} \\ (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{\text{id}_B \otimes \beta_{A,C}} & B \otimes (C \otimes A) \end{array}$$

$$\begin{array}{ccc} (A \otimes \top) \otimes B & \xrightarrow{\alpha_{A,\top,B}} & A \otimes (\top \otimes B) \\ \downarrow \rho_A \otimes \text{id}_B & & \downarrow \text{id}_A \otimes \lambda_B \\ & A \otimes B & \end{array}$$

$$\begin{array}{ccc} A \otimes B & & \\ \downarrow \beta_{A,B} & \searrow \text{id}_{A \otimes B} & \\ B \otimes A & \xrightarrow{\beta_{B,A}} & A \otimes B \end{array}$$

$$\begin{array}{ccc} \top \otimes A & \xrightarrow{\beta_{\top,A}} & A \otimes \top \\ \downarrow \lambda_A & & \downarrow \rho_A \\ & A & \end{array}$$

Definition 3. A *symmetric monoidal closed category (SMCC)* is a symmetric monoidal category, $(\mathcal{M}, \top, \otimes)$, such that, for any object B of \mathcal{M} , the functor $- \otimes B : \mathcal{M} \rightarrow \mathcal{M}$ has a specified right adjoint. Hence, for any objects A and C of \mathcal{M} there is an object $B \multimap C$ of \mathcal{M} and a natural bijection:

$$\text{Hom}_{\mathcal{M}}(A \otimes B, C) \cong \text{Hom}_{\mathcal{M}}(A, B \multimap C)$$

We call the functor $\multimap : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$ the internal hom of \mathcal{M} .

Definition 4. Suppose we are given two monoidal categories $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$ and $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$. Then a **monoidal functor** is a functor $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$, a map $m_{\top_1} : \top_2 \rightarrow F\top_1$ and a natural transformation $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$ subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A, B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$

$$\begin{array}{ccc} \top_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\ \downarrow m_{\top_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\ F\top_1 \otimes_2 FA & \xrightarrow{m_{\top_1, A}} & F(\top_1 \otimes_1 A) \end{array} \quad \begin{array}{ccc} FA \otimes_2 \top_2 & \xrightarrow{\rho_{2FA}} & FA \\ \downarrow \text{id}_{FA} \otimes m_{\top_1} & & \uparrow F\rho_{1A} \\ FA \otimes_2 F\top_1 & \xrightarrow{m_{A, \top_1}} & F(A \otimes_1 \top_1) \end{array}$$

Need to notice that the composition of monoidal functors is also monoidal, subject to the above coherence conditions.

Definition 5. Suppose we are given two symmetric monoidal closed categories $(\mathcal{M}_1, \top_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$ and $(\mathcal{M}_2, \top_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$. Then a **symmetric monoidal functor** is a functor $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$, a map $m_{\top_1} : \top_2 \rightarrow F\top_1$ and a natural transformation $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$ subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A, B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$

$$\begin{array}{ccc}
\top_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\
\downarrow m_{\top_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\
F\top_1 \otimes_2 FA & \xrightarrow{m_{\top_1, A}} & F(\top_1 \otimes_1 A)
\end{array}
\quad
\begin{array}{ccc}
FA \otimes_2 \top_2 & \xrightarrow{\rho_{2FA}} & FA \\
\downarrow \text{id}_{FA} \otimes m_{\top_1} & & \uparrow F\rho_{1A} \\
FA \otimes_2 F\top_1 & \xrightarrow{m_{A, \top_1}} & F(A \otimes_1 \top_1)
\end{array}$$

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{\beta_{2FA, FB}} & FB \otimes_2 FA \\
\downarrow m_{A, B} & & \downarrow m_{B, A} \\
F(A \otimes_1 B) & \xrightarrow{F\beta_{1A, B}} & F(B \otimes_1 A)
\end{array}$$

Definition 6. Suppose $(\mathcal{M}_1, \top_1, \otimes_1)$ and $(\mathcal{M}_2, \top_2, \otimes_2)$ are monoidal categories, and (F, m) and (G, n) are monoidal functors between \mathcal{M}_1 and \mathcal{M}_2 . Then a **monoidal natural transformation** is a natural transformation, $f : F \rightarrow G$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A, B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A, B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\top_1 & \xrightarrow{f_{\top_1}} & G\top_1 \\
\swarrow m_{\top_1} & & \searrow n_{\top_1} \\
& \top_2 &
\end{array}$$

Definition 7. Suppose $(\mathcal{M}_1, \top_1, \otimes_1)$ and $(\mathcal{M}_2, \top_2, \otimes_2)$ are SMCs, and (F, m) and (G, n) are symmetric monoidal functors between \mathcal{M}_1 and \mathcal{M}_2 . Then a **symmetric monoidal natural transformation** is a natural transformation, $f : F \rightarrow G$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A, B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A, B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\top_1 & \xrightarrow{f_{\top_1}} & G\top_1 \\
\swarrow m_{\top_1} & & \searrow n_{\top_1} \\
& \top_2 &
\end{array}$$

Definition 8. Suppose $(\mathcal{M}_1, \top_1, \otimes_1)$ and $(\mathcal{M}_2, \top_2, \otimes_2)$ are monoidal categories, and (F, m) is a monoidal functor between \mathcal{M}_1 and \mathcal{M}_2 and (G, n) is a monoidal functor between \mathcal{M}_2 and \mathcal{M}_1 . Then a **monoidal adjunction** is an ordinary adjunction $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$ such that the unit, $\eta_A : A \rightarrow GFA$, and the counit, $\varepsilon_A : FGA \rightarrow A$, are

monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow Fn_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_{A \otimes_1 B}} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{Fn_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{Gm_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

Definition 9. Suppose $(\mathcal{M}_1, \tau_1, \otimes_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2)$ are SMCs, and (F, m) is a symmetric monoidal functor between \mathcal{M}_1 and \mathcal{M}_2 and (G, n) is a symmetric monoidal functor between \mathcal{M}_2 and \mathcal{M}_1 . Then a **symmetric monoidal adjunction** is an ordinary adjunction $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$ such that the unit, $\eta_A : A \rightarrow GFA$, and the counit, $\varepsilon_A : FGA \rightarrow A$, are symmetric monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow Fn_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_{A \otimes_1 B}} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{Fn_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{Gm_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

Definition 10. A **monoidal comonad** on a monoidal category \mathcal{C} is a triple (T, ε, δ) , where (T, m) is a monoidal endofunctor on \mathcal{C} , $\varepsilon_A : TA \rightarrow A$ and $\delta_A : TA \rightarrow T^2A$ are monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\quad
\begin{array}{ccc}
& TA & \\
\swarrow & \downarrow \delta_A & \searrow \\
TA & \xleftarrow{\varepsilon_{TA}} T^2A \xrightarrow{T\varepsilon_A} & TA
\end{array}$$

The assumption that ε and δ are monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\quad
\begin{array}{ccc}
T\top & \xleftarrow{m_\top} & \top \\
& \searrow \varepsilon_\top & \downarrow \\
& & \top
\end{array}$$

$$\begin{array}{ccccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) & & \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} & & \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} & T(TA \otimes TB) & \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}$$

$$\begin{array}{ccc}
\top & \xrightarrow{m_\top} & T\top \\
\downarrow m_\top & & \downarrow \delta_\top \\
T\top & \xrightarrow{Tm_\top} & T^2\top
\end{array}$$

Definition 11. A *symmetric monoidal comonad* on a symmetric monoidal category \mathcal{C} is a triple (T, ε, δ) , where (T, m) is a symmetric monoidal endofunctor on \mathcal{C} , $\varepsilon_A : TA \rightarrow A$ and $\delta_A : TA \rightarrow T^2A$ are symmetric monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\quad
\begin{array}{ccccc}
& & TA & & \\
& \swarrow & \downarrow \delta_A & \searrow & \\
TA & \xleftarrow{\varepsilon_{TA}} & T^2A & \xrightarrow{T\varepsilon_A} & TA
\end{array}$$

The assumption that ε and δ are symmetric monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\quad
\begin{array}{ccc}
T\top & \xleftarrow{m_\top} & \top \\
& \searrow \varepsilon_\top & \downarrow \\
& & \top
\end{array}$$

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} T(TA \otimes TB) \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}$$

$$\begin{array}{ccc}
T & \xrightarrow{m_T} & TT \\
\downarrow m_T & & \downarrow \delta_T \\
TT & \xrightarrow{Tm_T} & T^2T
\end{array}$$

1.2 Linear Category

Definition 12. A *linear category*, $(\mathcal{L}, !, \text{weak}, \text{contra})$, is specified by

- a symmetric monoidal closed category $(\mathcal{L}, I, \otimes, \multimap)$,
- a symmetric monoidal comonad $(!, \varepsilon, \delta)$ on \mathcal{L} , with $q_{A,B} : !A \otimes !B \longrightarrow !(A \otimes B)$ and $q_I : I \longrightarrow !I$;
- monoidal natural transformations on \mathcal{L} with components $\text{weak}_A : !A \longrightarrow I$ and $\text{contra}_A : !A \longrightarrow !A \otimes !A$, s.t.
 - each $(!A, \text{weak}_A, \text{contra}_A)$ is a commutative comonoid, i.e. the following diagrams commute and $\beta \circ \text{contra}_A = \text{contra}_A$ where $\beta_{B,C} : B \otimes C \longrightarrow C \otimes B$ is the symmetry natural transformation of \mathcal{L} ;

$$\begin{array}{ccccc}
!A & \xrightarrow{\text{contra}_A} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \text{contra}_A} & !A \otimes (!A \otimes !A) \\
\downarrow \text{contra}_A & & & & \uparrow \alpha_{!A, !A, !A} \\
!A \otimes !A & \xrightarrow{\text{contra}_A \otimes id_{!A}} & (!A \otimes !A) \otimes !A & & \\
\\
& & !A & & \\
& \swarrow \lambda & \downarrow \text{contra}_A & \searrow \rho & \\
I \otimes !A & \xleftarrow{\text{weak}_A \otimes id_{!A}} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \text{weak}_A} & !A \otimes I
\end{array}$$

- weak_A and contra_A are coalgebra morphisms, i.e. the following diagrams commute;

$$\begin{array}{ccc}
 !A & \xrightarrow{\text{weak}_A} & I \\
 \delta_A \downarrow & & \downarrow q_I \\
 !!A & \xrightarrow{! \text{weak}_A} & !I
 \end{array}$$

$$\begin{array}{ccccc}
 !A & \xrightarrow{\text{contra}_A} & !A \otimes !A & \xrightarrow{\delta_A \otimes \delta_A} & !!A \otimes !!A \\
 \delta_A \downarrow & & & & \downarrow q_{!!A, !A} \\
 !!A & \xrightarrow{! \text{contra}_A} & & & !(A \otimes A)
 \end{array}$$

- any coalgebra morphism $f : (!A, \delta_A) \longrightarrow (!B, \delta_B)$ between free coalgebras preserve the comonoid structure given by weak and contra , i.e. the following diagrams commute.

$$\begin{array}{ccc}
 !A & \xrightarrow{!f} & !B \\
 \text{weak}_A \searrow & & \swarrow \text{weak}_B \\
 & I &
 \end{array}$$

$$\begin{array}{ccc}
 !A & \xrightarrow{\text{contra}_A} & !A \otimes !A \\
 f \downarrow & & \downarrow f \otimes f \\
 !B & \xrightarrow{\text{contra}_B} & !B \otimes !B
 \end{array}$$

Definition 13. A linear category with weakening, $(\mathcal{L}, w, \text{weak})$, is specified by

- a monoidal category $(\mathcal{L}, I, \otimes)$,
- a monoidal comonad (w, ε, δ) on \mathcal{L} with $q_{A,B} : wA \otimes wB \longrightarrow w(A \otimes B)$ and $q_I : I \longrightarrow wI$, and
- a monoidal natural transformation weak on \mathcal{L} with components $\text{weak}_A : wA \longrightarrow I$ s.t. the following diagrams commute:

$$\begin{array}{ccc}
 wA & \xrightarrow{\text{weak}_A} & I \\
 \delta_A \downarrow & & \downarrow q_I \\
 wwA & \xrightarrow{w\text{weak}_A} & wI
 \end{array}$$

$$\begin{array}{ccc}
 wA & \xrightarrow{wf} & wB \\
 \text{weak}_A \searrow & & \swarrow \text{weak}_B \\
 & I &
 \end{array}$$

Definition 14. A linear category with contraction, $(\mathcal{L}, c, \text{contraL}, \text{contraR})$, is specified by

- a monoidal category $(\mathcal{L}, I, \otimes)$,

- a monoidal comonad (c, ε, δ) on \mathcal{L} with $q_{A,B} : cA \otimes cB \rightarrow c(A \otimes B)$ and $q_I : I \rightarrow cI$, and
- monoidal natural transformations contraL and contraR on \mathcal{L} with components $\text{contraL}_{A,B} : cA \otimes B \rightarrow (cA \otimes B) \otimes cA$ and $\text{contraR}_{A,B} : B \otimes cA \rightarrow cA \otimes (B \otimes cA)$, s.t. the following diagram commutes:

$$\begin{array}{ccc}
cA \otimes I & \xleftarrow{\rho_{cA}^{-1}} & cA & \xrightarrow{\lambda_{cA}^{-1}} & I \otimes cA \\
\downarrow \text{contraL}_{A,I} & & & & \downarrow \text{contraR}_{A,I} \\
(cA \otimes I) \otimes cA & \xrightarrow{\alpha_{cA,I,cA}} & cA \otimes (I \otimes cA)
\end{array}$$

Definition 15. Given two comonads $(c, \varepsilon^c, \delta^c)$ and $(w, \varepsilon^w, \delta^w)$ on a category \mathcal{L} , a **distributive law** of c over w is a natural transformation with components $\text{dist}_A : cwA \rightarrow wcA$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
cwA & \xrightarrow{cw(f)} & cwB \\
\downarrow \text{dist}_A & & \downarrow \text{dist}_B \\
wcA & \xrightarrow{wc(f)} & wcB
\end{array}$$

$$\begin{array}{ccc}
wA & \xleftarrow{\varepsilon_{wA}^c} & cwA \\
\swarrow w\varepsilon_A^c & & \searrow \text{dist}_A \\
& wcA &
\end{array}
\quad
\begin{array}{ccc}
cA & \xleftarrow{c\varepsilon_A^w} & cwA \\
\swarrow \varepsilon_{cA}^w & & \searrow \text{dist}_A \\
& wcA &
\end{array}$$

$$\begin{array}{ccc}
cwA & \xrightarrow{cw\delta_A^c} & cwc^2A \\
\swarrow \delta_{wcA}^c & & \searrow c\text{dist}_{cA} \\
& c^2wcA &
\end{array}
\quad
\begin{array}{ccc}
wcwA & \xrightarrow{cw\delta_A^w} & wcw^2A \\
\swarrow \delta_{wcA}^w & & \searrow w\text{dist}_{wA} \\
& w^2cwA &
\end{array}$$

Lemma 16. Let $(c, \varepsilon^c, \delta^c)$ and $(w, \varepsilon^w, \delta^w)$ be two monoidal comonads on a linear category with weakening and contraction $(\mathcal{L}, I, \otimes, w, \text{weak}^w, c, \text{contraL}^c, \text{contraR}^c)$. Then the composition of c and w using the distributive law $\text{dist}_A : cwA \rightarrow wcA$ is a monoidal comonad $(cw, \varepsilon, \delta)$ on \mathcal{L} .

Proof. Suppose $(c, \varepsilon^c, \delta^c)$ and $(w, \varepsilon^w, \delta^w)$ are monoidal comonads, and $(\mathcal{L}, I, \otimes, w, \text{weak}^w, c, \text{contraL}^c, \text{contraR}^c)$ is a linear category with weakening and contraction. Since by definition $w, c : \mathcal{L} \rightarrow \mathcal{L}$ are monoidal functors we know that their composition $wc : \mathcal{L} \rightarrow \mathcal{L}$ is a monoidal functor:

$q_{A,B} : cwA \otimes cwB \longrightarrow cw(A \otimes B)$ is defined as: $q_{A,B} = cq_{A,B}^w \circ q_{wA,wB}^c$,
and $q_I : I \longrightarrow cwI$ is defined as: $q_I = cq_I^w \circ q_I^c$

We must now define both $\varepsilon_A : cwA \longrightarrow A$ and $\delta_A : cwA \longrightarrow cwcwA$, and then show that they are monoidal natural transformations subject to the comonad laws. Since we are composing two comonads each of ε and δ can be given two definitions, but they are equivalent:

- $\varepsilon_A : cwA \longrightarrow A$ is defined as in the diagram below, which commutes by the naturality of ε^c .

$$\begin{array}{ccc}
 cwA & \xrightarrow{\varepsilon_{wA}^c} & wA \\
 \downarrow c\varepsilon_A^w & & \downarrow \varepsilon_A^w \\
 cA & \xrightarrow{\varepsilon_A^c} & A
 \end{array}$$

- $\delta_A : cwA \longrightarrow cwcwA$ is defined as in the diagram:

$$\begin{array}{ccccc}
 cwA & \xrightarrow{c\delta_A^w} & cw^2A & \xrightarrow{\delta_{w^2A}^c} & c^2w^2A \\
 \downarrow \delta_{wA}^c & & \downarrow \delta_{w^2A}^c & & \downarrow cdist_{wA} \\
 c^2wA & \xrightarrow{c^2\delta_A^w} & c^2w^2A & \xrightarrow{cdist_{wA}} & cwcwA
 \end{array}$$

The left part of the diagram commutes by the naturality of δ^c and the right part commutes trivially.

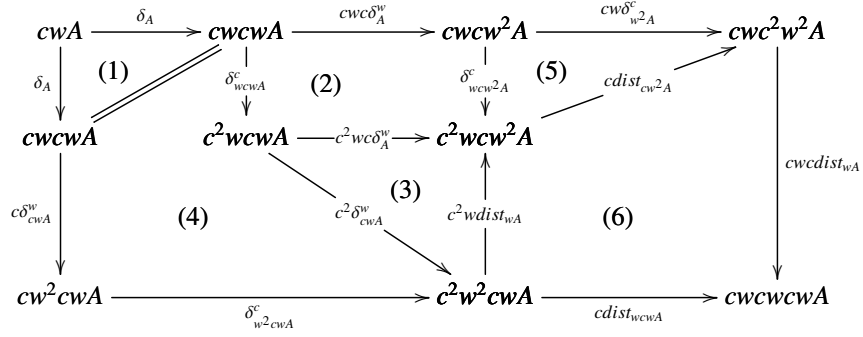
The remainder of the proof shows that the comonad laws hold.

- We show the comonad law $cw\delta_A \circ \delta_A = \delta_{cwA} \circ \delta_A$, expressed in the diagram below,

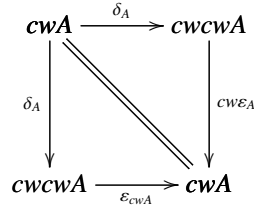
$$\begin{array}{ccc}
 cwA & \xrightarrow{\delta_A} & cwcwA \\
 \downarrow \delta_A & & \downarrow cw\delta_A \\
 cwcwA & \xrightarrow{\delta_{cwA}} & cwcwcwA
 \end{array}$$

is satisfied by proving $cw\delta_A = \delta_{cwA}$, as in the following diagram chasing. (1) commutes by equality and we will not expand δ_A for simplicity. (2) and (4) commutes by the naturality of δ^c . (3), (5) and (6) commute by the conditions of

dist.



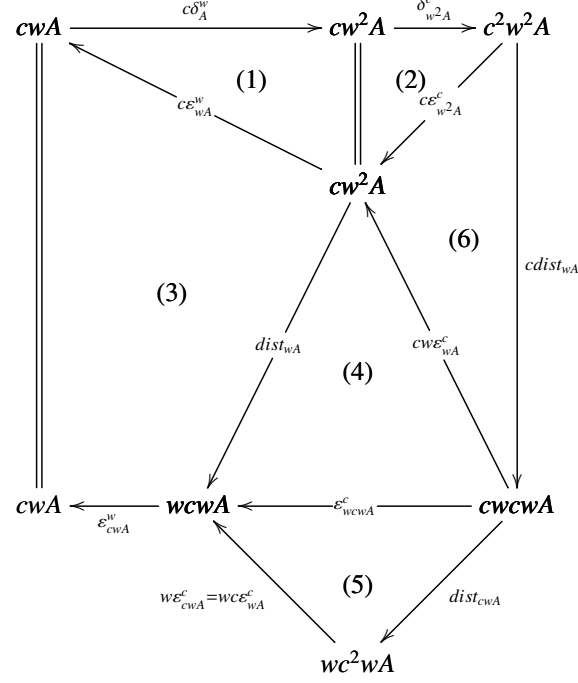
- The comonad law $cw\varepsilon_A \circ \delta_A = \varepsilon_{cwA} \circ \delta_A = id_{cwA}$, expressed in the diagram below,



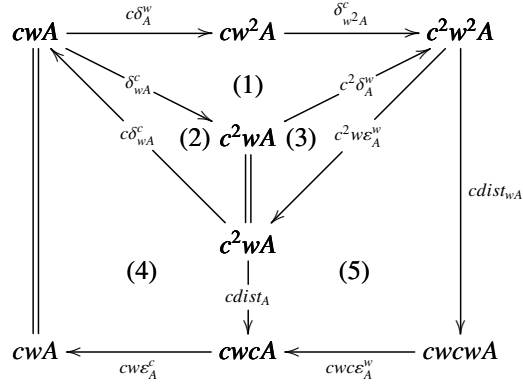
is satisfied by the following two diagram chasings below.

The left triangle is expanded in the following diagram chasing. (1) commutes by the comonad law for w with components δ_A^w and ε_{wA}^w . (2) commutes by the comonad law for c with components $\delta_{w^2A}^c$ and $\varepsilon_{w^2A}^c$. (3), (6) and the combination

of (4) and (5) ommute by the definition of $dist$.



The right triangle is expanded in the following diagram chasing. (1) commutes by the naturality of δ^c . (2) is the comonad law for c with components δ_{wA}^c and ϵ_{wA}^c . (3) is the comonad law for w with components δ_A^w and ϵ_A^w . (3) and (4) commute by the definition of $dist$.



□

Lemma 17. Let $(c, \epsilon^c, \delta^c)$ and $(w, \epsilon^w, \delta^w)$ be two monoidal comonads on a linear category with weakening and contraction $(\mathcal{L}, I, \otimes, w, \text{weak}^w, c, \text{contraL}^c, \text{contraR}^c)$, and

$(cw, \varepsilon, \delta)$ be the monoidal comonad on \mathcal{L} by composing c and w using the distributive law $dist_A : cwA \longrightarrow wA$. The monoidal natural transformations **weak** and **contra** satisfy the following conditions:

- Each $(cwA, \text{weak}_A, \text{contra}_A)$ is a commutative comonoid.
- weak_A and contra_A are coalgebra morphisms.
- Any coalgebra morphism $f : (cwA, \delta_A) \longrightarrow (cwB, \delta_B)$ between (free?) coalgebras preserve the comonoid structure given by **weak** and **contra**.

Proof. We first define **weak** and **contra**. Each of them can also be given two equivalent definitions:

- $\text{weak}_A : cwA \longrightarrow I$ is defined as in the diagram below. The left triangle commutes by the definition of $dist$ and the right triangle commutes by the definition of weak^w .

$$\begin{array}{ccccc}
 & & wcA & & \\
 & \nearrow dist_A & \downarrow w\varepsilon_A^c & \searrow \text{weak}_{cA}^w & \\
 cwA & \xrightarrow{\varepsilon_{wA}^c} & wA & \xrightarrow{\text{weak}_A^w} & I
 \end{array}$$

- $\text{contra}_A : cwA \longrightarrow cwA \otimes cwA$ is defined as below. The left part of the diagram commutes by the definitions of contraL^c and of contraR^c , and the right part commutes because \mathcal{L} is monoidal.

$$\begin{array}{ccccc}
 cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I & \xrightarrow{\text{contraL}_{wA,I}^c} & (cwA \otimes I) \otimes cwA \\
 \downarrow \lambda_{cwA}^{-1} & & & \searrow \alpha_{cwA,I,cwA} & \downarrow \rho_{cwA} \otimes id_{cwA} \\
 I \otimes cwA & \xrightarrow{\text{contraR}_{wA,I}^c} & cwA \otimes (I \otimes cwA) & \xrightarrow{id_{cwA} \otimes \lambda_{cwA}} & cwA \otimes cwA
 \end{array}$$

Then we show each condition is satisfied.

- Each $(cwA, \text{weak}_A, \text{contra}_A)$ is a commutative comonoid, i.e. Diagrams 1 and 2 commute.

$$\begin{array}{ccccc}
cwA & \xrightarrow{\text{contra}_A} & cwA \otimes cwA & \xrightarrow{id_{cwA} \otimes \text{contra}_A} & cwA \otimes (cwA \otimes cwA) \\
\downarrow \text{contra}_A & & & & \uparrow \alpha_{cwA, cwA, cwA} \\
cwA \otimes cwA & \xrightarrow{\text{contra}_A \otimes id_{cwA}} & (cwA \otimes cwA) \otimes cwA & &
\end{array}$$

Diagram 1: Commutative comonoid 1

Diagram 1 commutes by the following diagram chasing. (1) and (3) commute because $(\mathcal{L}, c, \text{contraL}^c, \text{contraR}^c)$ is a linear category with contraction. (2), (4), and (7) commute by the naturality of β . (5) and (6) commute by the coherence conditions of monoidal category.

$$\begin{array}{c}
\begin{array}{c}
\text{cwa} \otimes \text{cwa} \xrightarrow{\text{id}_{\text{cwa}} \otimes \rho_{\text{cwa}}^{-1}} \text{cwa} \otimes (\text{cwa} \otimes I) \xrightarrow{\text{id}_{\text{cwa}} \otimes \text{contral}_{\text{wA}, I}^c} \text{cwa} \otimes ((\text{cwa} \otimes I) \otimes \text{cwa}) \\
\downarrow \rho_{\text{cwa}}^{-1} \otimes \text{id}_{\text{cwa}} \quad \downarrow \text{id}_{\text{cwa}} \otimes \lambda_{\text{cwa}}^{-1} \quad \downarrow \text{id}_{\text{cwa}} \otimes \alpha_{\text{cwa}, I, \text{cwa}} \\
\text{(cwa} \otimes I) \otimes \text{cwa} \quad \text{(2)} \quad \text{cwa} \otimes (I \otimes \text{cwa}) \xrightarrow{\text{id}_{\text{cwa}} \otimes \text{contral}_{\text{wA}, I}^c} \text{cwa} \otimes (\text{cwa} \otimes (I \otimes \text{cwa})) \\
\downarrow \beta_{\text{cwa}, I, \text{cwa}} \quad \downarrow \beta_{\text{cwa}} \otimes \text{id}_{\text{cwa}} \quad \downarrow \beta_{\text{cwa}} \otimes (\text{id}_{\text{cwa}})_{\text{cwa}} \\
\text{(cwa} \otimes I) \otimes \text{cwa} \quad \text{(3)} \quad (I \otimes \text{cwa}) \otimes \text{cwa} \xrightarrow{\text{contral}_{\text{wA}, I}^c \otimes \text{id}_{\text{cwa}}} (\text{cwa} \otimes (I \otimes \text{cwa})) \otimes \text{cwa} \xrightarrow{\text{id}_{\text{cwa}} \otimes (\text{id}_{\text{cwa}} \otimes \lambda_{\text{cwa}})} \text{cwa} \otimes (\text{cwa} \otimes \text{cwa}) \\
\downarrow \text{contral}_{\text{wA}, I}^c \otimes \text{id}_{\text{cwa}} \quad \downarrow \alpha_{\text{cwa}, I, \text{cwa}} \otimes \text{id}_{\text{cwa}} \quad \downarrow (\text{id}_{\text{cwa}} \otimes \lambda_{\text{cwa}}) \otimes \text{id}_{\text{cwa}} \quad \downarrow \text{id}_{\text{cwa}} \otimes (\rho_{\text{cwa}} \otimes \text{id}_{\text{cwa}}) \\
((\text{cwa} \otimes I) \otimes \text{cwa}) \otimes \text{cwa} \xrightarrow{(\rho_{\text{cwa}, I} \otimes \text{id}_{\text{cwa}}) \otimes \text{id}_{\text{cwa}}} (\text{cwa} \otimes \text{cwa}) \otimes \text{cwa} \xrightarrow{\alpha_{\text{cwa}, \text{cwa}, \text{cwa}}} \text{cwa} \otimes (\text{cwa} \otimes \text{cwa})
\end{array}
\end{array}$$

(1) (2) (3) (4) (5) (6) (7)

$$\begin{array}{ccccc}
cwA & \xrightarrow{\text{contra}_A} & cwA \otimes cwA & \xrightarrow{\delta_A \otimes \delta_A} & cwcwA \otimes cwcwA \\
\downarrow \delta_A & & & & \downarrow q_{cwA, cwA} \\
cwcwA & \xrightarrow{cw\text{contra}_A} & & & cw(cwA \otimes cwA)
\end{array}$$

Diagram 4: Coalgebra Morphism 2

parts and prove each part commutes, as shown below.

$$\begin{array}{c}
\begin{array}{ccccccc}
cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I & \xrightarrow{\text{contra}_{cwA, I}^C} & (cwA \otimes I) \otimes cwA & \xrightarrow{\rho_{cwA} \otimes id_{cwA}} & cwA \otimes cwA \\
\downarrow c\delta_A^w & & \downarrow \varepsilon_{cwA \otimes I}^w & & \downarrow \varepsilon_{cwA \otimes cwA}^w & & \downarrow \delta_{wA}^c \delta_{wA}^w \\
cw^2A & & w(cwA \otimes I) & & c^2wA \otimes c^2wA & & c^2wA \otimes c^2wA \\
\downarrow \delta_{w^2A}^c & & \downarrow \varepsilon_{w(cwA \otimes I)}^c & & \downarrow c^2\delta_A^w \otimes c^2\delta_A^w & & \downarrow c^2\delta_A^w \otimes c^2\delta_A^w \\
c^2w^2A & & & & c^2w^2A \otimes c^2w^2A & & c^2w^2A \otimes c^2w^2A \\
\downarrow cdist_{wA} & & & & \downarrow cdist_{wA} \otimes cdist_{wA} & & \downarrow cdist_{wA} \otimes cdist_{wA} \\
cwcwA & \xrightarrow{cwp-1} & cw(cwA \otimes I) & \xrightarrow{cw\text{contra}_{cwA, I}^C} & cw((cwA \otimes I) \otimes cwA) & \xrightarrow{cw(\rho_{cwA} \otimes id_{cwA})} & cw(cwA \otimes cwA) \\
& & \downarrow \varepsilon_{w(cwA \otimes I)}^c & & \downarrow q_{w(cwA \otimes cwA)}^c & & \downarrow cq_{cwA, cwA}^w
\end{array} \\
\begin{array}{c}
(a) \qquad \qquad \qquad (b) \\
\downarrow \varepsilon_{cwA \otimes I}^w \qquad \downarrow \varepsilon_{w(cwA \otimes I)}^c \qquad \downarrow \varepsilon_{cwA \otimes cwA}^w \qquad \downarrow \varepsilon_{w(cwA \otimes cwA)}^c
\end{array}
\end{array}$$

Part (a) commutes by the following diagram chase. (1) is equality. (2) is the comonad law for w . (3) is the comonad law for c . (4) commutes by the naturality of ε^c . (5) is one of the conditions for $dist_{wA}$. (6) commutes by the naturality of ε^w . And (7) commutes by the naturality of ε^c .

$$\begin{array}{ccc}
cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I \\
\downarrow c\delta_A^w & \searrow c\delta_A^w & \nearrow \varepsilon_{cwA \otimes I}^w \\
cw^2A & \xrightarrow{c\varepsilon_{wA}^w} & cwA \\
\downarrow \delta_{w^2A}^c & \nearrow \varepsilon_{cw^2A}^c & \downarrow dist_{wA} \\
c^2w^2A & \xrightarrow{w\rho_{cwA}^{-1}} & w(cwA \otimes I) \\
\downarrow cdist_{wA} & \nearrow \varepsilon_{wcwA}^c & \downarrow \varepsilon_{w(cwA \otimes I)}^c \\
cwcwA & \xrightarrow{cw\rho_{cwA}^{-1}} & cw(cwA \otimes I)
\end{array}$$

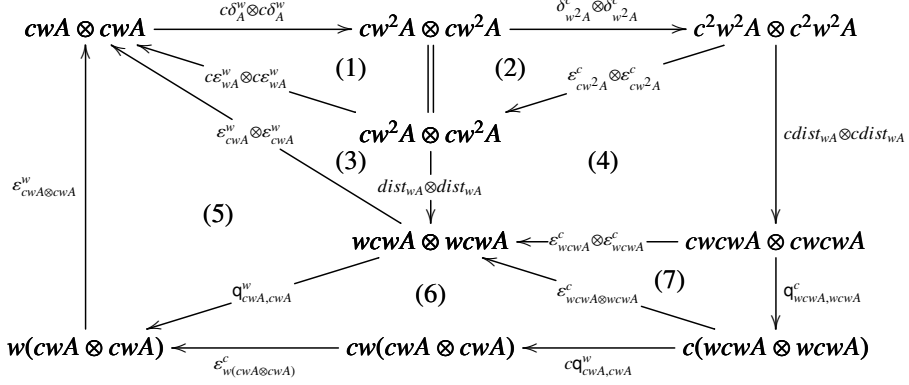
(1) (2) (3) (4) (5) (6) (7)

Part (b) commutes by the following diagram chase. The upper two squares both commute by the naturality of ε^w , and the lower two squares commute by the naturality of ε^c .

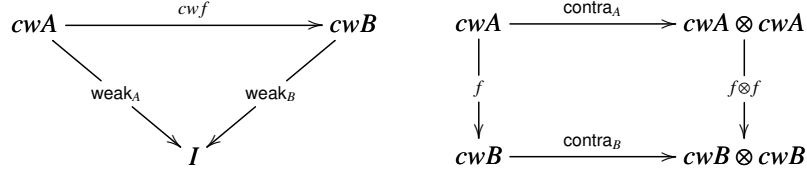
$$\begin{array}{ccccc}
cwA \otimes I & \xrightarrow{contral_{wA,I}^c} & (cwA \otimes I) \otimes cwA & \xrightarrow{\rho_{cwA} \otimes id_{cwA}} & cwA \otimes cwA \\
\uparrow \varepsilon_{cwA \otimes I}^w & & \uparrow \varepsilon_{(cwA \otimes I) \otimes cwA}^w & & \uparrow \varepsilon_{cwA \otimes cwA}^w \\
w(cwA \otimes I) & \xrightarrow{wcontral_{wA,I}^c} & w((cwA \otimes I) \otimes cwA) & \xrightarrow{w(\rho_{cwA} \otimes id_{cwA})} & w(cwA \otimes cwA) \\
\uparrow \varepsilon_{w(cwA \otimes I)}^c & & \uparrow \varepsilon_{w((cwA \otimes I) \otimes cwA)}^c & & \uparrow \varepsilon_{w(cwA \otimes cwA)}^c \\
cw(cwA \otimes I) & \xrightarrow{cwcontral_{wA,I}^c} & cw((cwA \otimes I) \otimes cwA) & \xrightarrow{cw(\rho_{cwA} \otimes id_{cwA})} & cw(cwA \otimes cwA)
\end{array}$$

Part (c) commutes by the following diagram. (1) and (2) are the comonad law. (3) commutes by the condition of $dist_{wA}$. (4) and (6) commute by the naturality

of ε^c . (5) and (7) commute because w and c are monoidal comonads.



- any coalgebra morphism $f : (!A, \delta_A) \longrightarrow (!B, \delta_B)$ between free coalgebras preserve the comonoid structure given by weak and contra, i.e. the following diagrams commute.



□

2 Related Work

TODO

3 Conclusion

TODO

References

- [1] P. N. Benton. A mixed linear and non-linear logic: Proofs, terms and models (preliminary report). Technical Report UCAM-CL-TR-352, University of Cambridge Computer Laboratory, 1994. Accessible online at <http://research.microsoft.com/en-us/um/people/nick/mixed3.ps>.