

Separating Linear Modalities

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Abstract

TODO

1 Introduction

TODO [1]

2 Categorical Models

2.1 Lambek Categories

TODO: Define Lambek Categories

2.2 Lambek Categories with Weakening and Contraction

Definition 1. A *Lambek category with weakening*, $(\mathcal{L}, w, \text{weak})$, is a Lambek category equipped with a monoidal comonad (w, ε, δ) , and a monoidal natural transformation $\text{weak}_A : wA \longrightarrow I$. Furthermore, weak must be a coalgebra morphism. That is, the following digram must commute:

$$\begin{array}{ccc} wA & \xrightarrow{\text{weak}_A} & I \\ \delta_A \downarrow & & \downarrow q_I \\ w^2A & \xrightarrow{w\text{weak}_A} & wI \end{array}$$

Definition 2. A *Lambek category with contraction*, $(\mathcal{L}, c, \text{contraL}, \text{contraR})$, is a Lambek category equipped with a monoidal comonad (c, ε, δ) , and two monoidal natural transformations:

$$\begin{aligned} \text{contraL}_{A,B} &: cA \otimes B \longrightarrow (cA \otimes B) \otimes cA \\ \text{contraR}_{A,B} &: B \otimes cA \longrightarrow cA \otimes (B \otimes cA) \end{aligned}$$

Furthermore, the following diagrams must commute:

$$\begin{array}{ccccc} cA \otimes I & \xleftarrow{\rho_{cA}^{-1}} & cA & \xrightarrow{\lambda_{cA}^{-1}} & I \otimes cA \\ \text{contraL}_{A,I} \downarrow & & & & \downarrow \text{contraR}_{A,I} \\ (cA \otimes I) \otimes cA & \xrightarrow{\alpha_{cA,I,cA}} & & & cA \otimes (I \otimes cA) \end{array}$$

$$\begin{array}{ccccc}
cA \otimes cA & \xrightarrow{id_{cA} \otimes \rho_{cA}^{-1}} & cA \otimes (cA \otimes I) & \xrightarrow{id_{cA} \otimes \text{contraL}_{A,I}} & cA \otimes ((cA \otimes I) \otimes cA) \\
\downarrow \lambda_{cA}^{-1} \otimes id_{cA} & & & & \downarrow id_{cA} \otimes (\rho_{cA} \otimes id_{cA}) \\
(I \otimes cA) \otimes cA & \xrightarrow{\text{contraR}_{A,I} \otimes id_{cA}} & (cA \otimes (I \otimes cA)) \otimes cA & \xrightarrow{(id_{cA} \otimes \lambda_{cA}) \otimes id_{cA}} & (cA \otimes cA) \otimes cA \\
& & & & \uparrow \alpha_{cA, cA, cA}
\end{array}$$

2.3 Lambek Categories with Exchange

Definition 3. A *Lambek category with exchange*, $(\mathcal{L}, e, \text{ex})$, is a Lambek category equipped with a monoidal comonad (e, ε, δ) on \mathcal{L} , and a monoidal natural transformation $\text{ex}_{A,B} : eA \otimes eB \longrightarrow eB \otimes eA$. Furthermore, the following must hold:

- $\text{ex}_{A,B}$ must be an involution:

$$\begin{array}{ccc}
eA \otimes eB & \xrightarrow{\text{ex}_{A,B}} & eB \otimes eA \\
& \searrow & \downarrow \text{ex}_{B,A} \\
& & eA \otimes eB
\end{array}$$

- The following coherence diagrams with respect to the comonad must commute:

$$\begin{array}{ccc}
e^2 A \otimes e^2 B & \xrightarrow{\text{ex}_{eA, eB}} & e^2 B \otimes e^2 A \\
\downarrow \varepsilon_{eA} \otimes \varepsilon_{eB} & & \uparrow \delta_B \otimes \delta_A \\
eA \otimes eB & \xrightarrow{\text{ex}_{A,B}} & eB \otimes eA
\end{array}$$

Definition 4. Suppose $(\mathcal{L}, e, \text{ex})$ is a Lambek category with exchange. Then the *Eilenberg Moore category*, \mathcal{L}^e , of the comonad (e, ε, δ) has as objects all the e -coalgebras $(A, h_A : A \longrightarrow eA)$, and as morphisms all the coalgebra morphisms. We call h_A the action of the coalgebra. Furthermore, the following (action) diagrams must commute:

$$\begin{array}{ccc}
A & \xrightarrow{h_A} & eA \\
h_A \downarrow & & \downarrow eh_A \\
eA & \xrightarrow{\delta_A} & e^2 A
\end{array}
\quad
\begin{array}{ccc}
A & & \\
h_A \downarrow & \searrow & \\
eA & \xrightarrow{\varepsilon_A} & A
\end{array}$$

Lemma 5 (The Eilenberg Moore Category is Monoidal). *Then the category \mathcal{L}^e is monoidal.*

Proof. We must first define the unitors, and then the associator. Then we show that they respect the monoidal coherence diagrams. Throughout this proof we will make use of the coalgebra (A, h_A) , (B, h_B) , and (C, h_C) .

The tensor product of (A, h_A) and (B, h_B) is $(A \otimes B, q_{A,B} \circ (h_A \otimes h_B))$, and the unit of the tensor product is (I, q_I) ; both actions are easily shown to satisfy the action diagrams of the Eilenberg Moore category. The left and right unitors are $\lambda : I \otimes A \longrightarrow A$ and $\rho : A \otimes I \longrightarrow A$, because they are indeed coalgebra morphisms.

The respective diagram for the right unitor is as follows:

$$\begin{array}{ccccccc}
A \otimes I & \xrightarrow{h_A \otimes \text{id}} & eA \otimes I & \xrightarrow{\text{id} \otimes q_I} & eA \otimes eI & \xrightarrow{q_{A,I}} & e(A \otimes I) \\
\downarrow \rho & & & & \searrow \lambda & & \downarrow e\rho \\
A & & & & & & eA
\end{array}$$

h_A

The top-left diagram commutes by naturality of ρ , the top-right diagram commutes by the fact that e is a monoidal functor, and the bottom diagram commutes by the action diagrams for the coalgebra (A, h_A) . Showing the left unitor is a coalgebra morphism is similar.

The unitors are natural and isomorphisms, because they are essentially inherited from the underlying Lambek category.

The associator $\alpha : (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C)$ is also a coalgebra morphism. First, notice that:

$$q_{A \otimes B, C} \circ ((q_{A, B} \circ (h_A \otimes h_B)) \otimes h_C) = q_{A \otimes B, C} \circ (q_{A, B} \otimes \text{id}) \circ ((h_A \otimes h_B) \otimes h_C)$$

where the left-hand side is the action of the coalgebra $(A \otimes B) \otimes C$. Similarly, the following is the action of the coalgebra $A \otimes (B \otimes C)$:

$$q_{A, B \otimes C} \circ (h_A \otimes (q_{B, C} \circ (h_B \otimes h_C))) = q_{A, B \otimes C} \circ (\text{id} \otimes q_{B, C}) \circ (h_A \otimes (h_B \otimes h_C))$$

The following diagram must commute:

$$\begin{array}{ccc}
 (A \otimes B) \otimes C & \xrightarrow{(h_A \otimes h_B) \otimes h_C} & (eA \otimes eB) \otimes eC \\
 \downarrow \alpha & & \downarrow \alpha \\
 A \otimes (B \otimes C) & \xrightarrow{h_A \otimes (h_B \otimes h_C)} & eA \otimes (eB \otimes eC) \\
 \uparrow \alpha & & \uparrow \alpha \\
 (A \otimes B) \otimes eC & \xrightarrow{q \otimes \text{id}} & e(A \otimes B) \otimes eC \\
 \uparrow q & & \uparrow q \\
 e((A \otimes B) \otimes C) & \xrightarrow{e\alpha} & e(A \otimes (B \otimes C))
 \end{array}$$

The left diagram commutes by naturality of α , and the right diagram commutes because e is a monoidal functor.

Composition in \mathcal{L}^e is the same as \mathcal{L} , and thus, the monoidal coherence diagrams hold in \mathcal{L}^e as well. Thus, \mathcal{L}^e is monoidal. We now show that it is symmetric. \square

Lemma 6. *In \mathcal{L}^e there is a natural transformation $\beta_{A, B} : A \otimes B \longrightarrow B \otimes A$.*

Doesn't seem right.

Proof. We define β as follows:

$$\beta_{A,B} := A \otimes B \xrightarrow{h_A \otimes h_B} eA \otimes eB \xrightarrow{ex_{A,B}} eB \otimes eA \xrightarrow{\varepsilon_B \otimes \varepsilon_A} B \otimes A$$

Suppose $f : A \rightarrow A'$ and $g : B \rightarrow B'$ are two coalgebra morphisms. Then the following diagram shows that $\beta_{A,B}$ is a natural transformation:

$$\begin{array}{ccccccc} A \otimes B & \xrightarrow{h_A \otimes h_B} & eA \otimes eB & \xrightarrow{ex_{A,B}} & eB \otimes eA & \xrightarrow{\varepsilon_B \otimes \varepsilon_A} & B \otimes A \\ \downarrow f \otimes g & & \downarrow ef \otimes eg & & \downarrow eg \otimes ef & & \downarrow g \otimes f \\ A' \otimes B' & \xrightarrow{h_{A'} \otimes h_{B'}} & eA' \otimes eB' & \xrightarrow{ex_{A',B'}} & eB' \otimes eA' & \xrightarrow{\varepsilon_{B'} \otimes \varepsilon_{A'}} & B' \otimes A' \end{array}$$

The left diagram commutes because f and g are both coalgebra morphisms, the middle diagram commutes because $ex_{A,B}$ is a natural transformation, and the right diagram commutes by naturality of ε . \square

Corollary 7. For any coalgebras (A, h_A) and (B, h_B) the following commutes:

$$\begin{array}{ccccccc} A \otimes B & \xrightarrow{h_A \otimes h_B} & eA \otimes eB & \xrightarrow{ex_{A,B}} & eB \otimes eA & \xrightarrow{\varepsilon_B \otimes \varepsilon_A} & B \otimes A \\ \parallel & & & & & & \downarrow h_B \otimes h_A \\ A \otimes B & \xrightarrow{h_A \otimes h_B} & eA \otimes eB & \xrightarrow{ex_{A,B}} & eB \otimes eA & = & eB \otimes eA \end{array}$$

Proof. This proof follows by the fact that the following diagram commutes:

$$\begin{array}{ccccc} A \otimes B & = & A \otimes B & = & A \otimes B \\ \downarrow h_A \otimes h_B & & \downarrow h_A \otimes h_B & & \downarrow h_A \otimes h_B \\ eA \otimes eB & = & eA \otimes eB & & eA \otimes eB \\ & \swarrow \varepsilon_{eA} \otimes \varepsilon_{eB} & \downarrow h_{eA} \otimes h_{eB} & & \\ & & e^2A \otimes e^2B & & \\ & & \downarrow ex_{eA,eB} & & \\ & & e^2B \otimes e^2A & & \\ & \swarrow \varepsilon_{eB} \otimes \varepsilon_{eA} & & & \\ & & B \otimes A & \xleftarrow{\varepsilon_B \otimes \varepsilon_A} & eB \otimes eA \\ & \swarrow h_B \otimes h_A & & & \downarrow ex_{A,B} \\ & & & & B \otimes A \end{array}$$

The diagram on the right commutes because $\beta_{A,B}$ is a natural transformation, and the other diagrams commute by either naturality of $ex_{A,B}$ or the action diagrams. \square

define coforks.

Lemma 8. In \mathcal{L}^e , β is a coalgebra morphism.

Proof. β is a coalgebra morphism means the following diagram commutes:

$$\begin{array}{ccc} A \otimes B & \xrightarrow{h_{A \otimes B}} & e(A \otimes B) \\ \beta_{A,B} \downarrow & & \downarrow e\beta_{A,B} \\ B \otimes A & \xrightarrow{h_{B \otimes A}} & e(B \otimes A) \end{array}$$

□

Lemma 9 (The Eilenberg-Moore Category is Symmetric Monoidal). *The category \mathcal{L}^e is symmetric monoidal.*

Proof. The following diagram shows that $\beta_{B,A} \circ \beta_{A,B} = \text{id}_{A \otimes B}$:

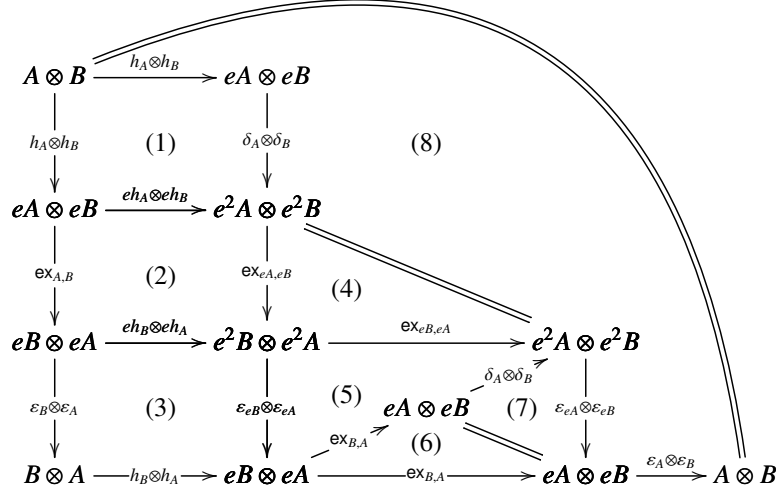


Diagram one commutes by the action diagrams of the Eilenberg Moore category, diagram two commutes by naturality of ex , diagram three commutes by naturality of ε , diagram four and five commute by the coherence diagrams of ex , diagram six clearly commutes, diagram seven commutes because (e, ε, δ) is a comonad, and diagram eight commutes by both the action diagrams of the Eilenberg Moore category and the fact that (e, ε, δ) is a comonad.

At this point we must verify that β respects the coherence diagrams of a symmetric monoidal category; see Definition 22. Thus, we must show that each of the following diagrams hold:

Case

$$\begin{array}{ccccc}
 (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\beta_{A,B \otimes C}} & (B \otimes C) \otimes A \\
 \downarrow \beta_{A,B} \otimes \text{id}_C & & & & \downarrow \alpha_{B,C,A} \\
 (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{\text{id}_B \otimes \beta_{A,C}} & B \otimes (C \otimes A)
 \end{array}$$

Case

$$\begin{array}{ccc}
 A \otimes B & & \\
 \downarrow \beta_{A,B} & \searrow \text{id}_{A \otimes B} & \\
 B \otimes A & \xrightarrow{\beta_{B,A}} & A \otimes B
 \end{array}$$

Case

$$\begin{array}{ccc}
 \top \otimes A & \xrightarrow{\beta_{\top,A}} & A \otimes \top \\
 \searrow \lambda_A & & \swarrow \rho_A \\
 & A &
 \end{array}$$

□

Definition 10. Let $(\mathcal{L}, e, \text{ex})$ be a Lambek category with exchange. The **coKleisli Category** of e , \mathcal{L}_e , is a category with the same objects as \mathcal{L} . There is an arrow $\hat{f} : A \rightarrow B$ in \mathcal{L}_e if there is an arrow $f : eA \rightarrow B$ in \mathcal{L} . The identity arrow $\hat{id}_A : A \rightarrow A$ is the arrow $\varepsilon_A : eA \rightarrow A$ in \mathcal{L} . Given $\hat{f} : A \rightarrow B$ and $\hat{g} : B \rightarrow C$ in \mathcal{L}_e , which are arrows $f : eA \rightarrow B$ and $g : eB \rightarrow C$ in \mathcal{L} , the composition $\hat{g} \circ \hat{f} : A \rightarrow C$ is defined as $g \circ ef \circ \delta_A$.

2.4 Linear Categories

Definition 11. A **linear category**, $(\mathcal{L}, !, \text{weak}, \text{contra})$, is a symmetric monoidal closed category $(\mathcal{L}, I, \otimes, \multimap)$ equipped with a symmetric monoidal comonad $(!, \varepsilon, \delta)$ with $q_{A,B} : !A \otimes !B \rightarrow !(A \otimes B)$ and $q_I : I \rightarrow !I$, and two monoidal natural transformations with components $\text{weak}_A : !A \rightarrow I$ and $\text{contra}_A : !A \rightarrow !A \otimes !A$, satisfying the following conditions:

- each $(!A, \text{weak}_A, \text{contra}_A)$ is a commutative comonoid, i.e. the following diagrams commute and $\beta \circ \text{contra}_A = \text{contra}_A$ where $\beta_{B,C} : B \otimes C \rightarrow C \otimes B$ is the symmetry natural transformation of \mathcal{L} :

$$\begin{array}{ccccc}
 !A & \xrightarrow{\text{contra}_A} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \text{contra}_A} & !A \otimes (!A \otimes !A) \\
 \text{contra}_A \downarrow & & & & \uparrow \alpha_{!A, !A, !A} \\
 !A \otimes !A & \xrightarrow{\text{contra}_A \otimes id_{!A}} & (!A \otimes !A) \otimes !A & & \\
 & & \downarrow \text{contra}_A & & \\
 & & !A & & \\
 & \swarrow \lambda^{-1} & \downarrow \text{contra}_A & \searrow \rho^{-1} & \\
 I \otimes !A & \xleftarrow{\text{weak}_A \otimes id_{!A}} & !A \otimes !A & \xrightarrow{id_{!A} \otimes \text{weak}_A} & !A \otimes I
 \end{array}$$

- weak_A and contra_A are coalgebra morphisms, i.e. the following diagrams commute;

$$\begin{array}{ccc}
 !A & \xrightarrow{\text{weak}_A} & I \\
 \delta_A \downarrow & & \downarrow q_I \\
 !!A & \xrightarrow{! \text{weak}_A} & !I
 \end{array}$$

$$\begin{array}{ccccc}
 !A & \xrightarrow{\text{contra}_A} & !A \otimes !A & \xrightarrow{\delta_A \otimes \delta_A} & !!A \otimes !!A \\
 \delta_A \downarrow & & & & \downarrow q_{!A, !A} \\
 !!A & \xrightarrow{! \text{contra}_A} & !(A \otimes A) & &
 \end{array}$$

- any coalgebra morphism $f : (!A, \delta_A) \rightarrow (!B, \delta_B)$ between free coalgebras preserve the comonoid structure given by weak and contra , i.e. the following diagrams commute.

$$\begin{array}{ccc}
 !A & \xrightarrow{f} & !B \\
 \text{weak}_A \searrow & & \swarrow \text{weak}_B \\
 & I &
 \end{array}$$

$$\begin{array}{ccc}
 !A & \xrightarrow{\text{contra}_A} & !A \otimes !A \\
 f \downarrow & & \downarrow f \otimes f \\
 !B & \xrightarrow{\text{contra}_B} & !B \otimes !B
 \end{array}$$

Definition 12. Given two comonads $(c, \varepsilon^c, \delta^c)$ and $(w, \varepsilon^w, \delta^w)$ on a category \mathcal{L} such that $(\mathcal{L}, c, \text{contraL}, \text{contraR})$ is a Lambek category with contraction and $(\mathcal{L}, w, \text{weak})$ is a Lambek category with weakening, we define a **distributive law**

of c over w to be a natural transformation with components $\text{dist}_A : cwA \rightarrow wcA$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
 wA & \xleftarrow{\varepsilon_{wA}^c} & cwA \\
 & \nwarrow w\varepsilon_A^c & \nearrow \text{dist}_A \\
 & wcA &
 \end{array}
 \qquad
 \begin{array}{ccc}
 cA & \xleftarrow{c\varepsilon_A^w} & cwA \\
 & \nwarrow \varepsilon_{cA}^w & \nearrow \text{dist}_A \\
 & wcA &
 \end{array}$$

Lemma 13. Given two comonads $(c, \varepsilon^c, \delta^c)$ and $(w, \varepsilon^w, \delta^w)$ on a category \mathcal{L} such that $(\mathcal{L}, c, \text{contraL}, \text{contraR})$ is a Lambek category with contraction and $(\mathcal{L}, w, \text{weak})$ is a Lambek category with weakening, the following two diagrams commute:

$$\begin{array}{ccc}
 cwA & \xrightarrow{cw\delta_A^c} & cwc^2A \\
 \searrow \delta_{wcA}^c & & \nearrow c\text{dist}_{cA} \\
 & c^2wcA &
 \end{array}
 \qquad
 \begin{array}{ccc}
 wcwA & \xrightarrow{wc\delta_A^w} & wcw^2A \\
 \searrow \delta_{wcA}^w & & \nearrow w\text{dist}_{wA} \\
 & w^2cwA &
 \end{array}$$

Proof. The two diagrams above commute because the following ones commute by the distributive law and the comonad laws for c and w .

$$\begin{array}{ccc}
 cwA & \xrightarrow{cw\delta_A^c} & cwc^2A \\
 \searrow \delta_{wcA}^c & \swarrow cw\varepsilon_{cA}^c & \nearrow c\text{dist}_{cA} \\
 & cwA & \\
 \uparrow c\varepsilon_{wcA}^c & & \\
 c^2wcA & &
 \end{array}
 \qquad
 \begin{array}{ccc}
 wcwA & \xrightarrow{wc\delta_A^w} & wcw^2A \\
 \searrow \delta_{wcA}^w & \swarrow wc\varepsilon_{wA}^w & \nearrow w\text{dist}_{wA} \\
 & wcwA & \\
 \uparrow w\varepsilon_{wcA}^w & & \\
 w^2cwA & &
 \end{array}$$

□

Lemma 14 (Composition of Weakening and Contraction). Suppose

$(\mathcal{L}, I, \otimes, w, \text{weak}^w, c, \text{contraL}, \text{contraR})$ is a Lambek category with weakening and contraction, where $(w, \varepsilon^w, \delta^w)$ and $(c, \varepsilon^c, \delta^c)$ are the respective monoidal comonads. Then the composition of c and w using the distributive law $\text{dist}_A : cwA \rightarrow wcA$ is a monoidal comonad on \mathcal{L} .

Proof. For the complete proof see Appendix B.1. □

Definition 15. A *Lambek category with cw* , $(\mathcal{L}, cw, \text{weak}^w, \text{contraL}, \text{contraR}, \text{dist})$, is a Lambek category with weakening and contraction, and a distributive law. Furthermore, the following coherence diagrams commute:

$$\begin{array}{ccc}
 I \otimes cwA & \xrightarrow{\lambda_{I \otimes cwA}^{-1}} & I \otimes (I \otimes cwA) \\
 \text{contraR}_{wA, I} \downarrow & & \uparrow \text{weak}_A^w \text{id}_{I \otimes cwA} \\
 cwA \otimes (I \otimes cwA) & \xrightarrow{\varepsilon_{wA}^c \otimes \text{id}_{I \otimes cwA}} & wA \otimes (I \otimes cwA)
 \end{array}
 \qquad
 \begin{array}{ccc}
 cwA \otimes I & \xrightarrow{\rho_{cwA \otimes I}^{-1}} & (cwA \otimes I) \otimes I \\
 \text{contraL}_{wA, I} \downarrow & & \uparrow \text{id}_{cwA \otimes I} \text{weak}_A^w \\
 (cwA \otimes I) \otimes cwA & \xrightarrow{\text{id}_{cwA \otimes I} \otimes \varepsilon_{wA}^c} & (cwA \otimes I) \otimes wA
 \end{array}$$

$$\begin{array}{ccc}
 cwA & \xrightarrow{f} & cwB \\
 \varepsilon_{wA}^c \downarrow & & \downarrow \varepsilon_{wB}^c \\
 wA & \xrightarrow{\text{weak}_A^w} I & \xleftarrow{\text{weak}_B^w} wB
 \end{array}$$

where $f : (cwA, \delta_A) \rightarrow (cwB, \delta_B)$ is any coalgebra morphism between free coalgebras.

Lemma 16. Let $(\mathcal{L}, cw, \text{weak}^w, \text{contraL}, \text{contraR})$ be a Lambek category with cw . Then the following conditions are satisfied:

1. There exist two natural transformations $\text{weak}_A : cwA \longrightarrow I$ and $\text{contra}_A : cwA \longrightarrow cwA \otimes cwA$.
2. Each $(cwA, \text{weak}_A, \text{contra}_A)$ is a comonoid.
3. weak_A and contra_A are coalgebra morphisms.
4. Any coalgebra morphism $f : (cwA, \delta_A) \longrightarrow (cwB, \delta_B)$ between free coalgebras preserves the comonoid structure given by weak and contra .

Proof. We will only prove the first condition by defining weak and contra . For the complete proof see Appendix B.2. Each of weak and contra can be given two equivalent definitions. $\text{weak}_A : cwA \longrightarrow I$ is defined as in the diagram below. The left triangle commutes by the definition of dist and the right triangle commutes by the definition of weak^w .

$$\begin{array}{ccccc}
 & & wcA & & \\
 & \nearrow \text{dist}_A & \downarrow w\varepsilon_A^c & \searrow \text{weak}_{cA}^w & \\
 cwA & \xrightarrow{\varepsilon_{wA}^c} & wA & \xrightarrow{\text{weak}_A^w} & I
 \end{array}$$

$\text{contra}_A : cwA \longrightarrow cwA \otimes cwA$ is defined as below. The left part of the diagram commutes by the definitions of contraL and of contraR , and the right part commutes because \mathcal{L} is monoidal.

$$\begin{array}{ccccc}
 cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I & \xrightarrow{\text{contraL}_{wA,I}} & (cwA \otimes I) \otimes cwA \\
 \downarrow \lambda_{cwA}^{-1} & & & \swarrow \alpha_{cwA,I,cwA} & \downarrow \rho_{cwA} \otimes id_{cwA} \\
 I \otimes cwA & \xrightarrow{\text{contraR}_{wA,I}} & cwA \otimes (I \otimes cwA) & \xrightarrow{id_{cwA} \otimes \lambda_{cwA}} & cwA \otimes cwA
 \end{array}$$

□

Definition 17. Given two comonads $(cw, \varepsilon^{cw}, \delta^{cw})$ and $(e, \varepsilon^e, \delta^e)$ on a category \mathcal{L} such that $(\mathcal{L}, cw, \text{weak}, \text{contra})$ is a Lambek category with cw and $(\mathcal{L}, e, \text{ex})$ is a Lambek category with exchange, we define a **distributive law for exchange** of cw over e to be a natural isomorphism with components $\text{distEx}_A : cweA \longrightarrow ecwA$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
 eA & \xleftarrow{\varepsilon_{eA}^{cw}} & cweA \\
 \swarrow e\varepsilon_A^{cw} & & \searrow \text{distEx}_A \\
 & ecwA &
 \end{array}
 \qquad
 \begin{array}{ccc}
 cwA & \xleftarrow{cwe_A^e} & cweA \\
 \swarrow \varepsilon_{cwA}^e & & \searrow \text{distEx}_A \\
 & ecwA &
 \end{array}$$

Lemma 18. Given two comonads $(cw, \varepsilon^{cw}, \delta^{cw})$ and $(e, \varepsilon^e, \delta^e)$ on a category \mathcal{L} such that $(\mathcal{L}, cw, \text{weak}, \text{contra})$ is a Lambek category with cw and $(\mathcal{L}, e, \text{ex})$ is a Lambek category with exchange, the following two digrams also commute:

$$\begin{array}{ccc}
 cwecwA & \xrightarrow{cwe\delta_A^{cw}} & cwe(cw)^2A \\
 \searrow \delta_{cweA}^{cw} & & \nearrow cwe\text{distEx}_{cwA} \\
 & (cw)^2ecwA &
 \end{array}
 \qquad
 \begin{array}{ccc}
 ecweA & \xrightarrow{ecw\delta_A^e} & ecwe^2A \\
 \searrow \delta_{ecweA}^e & & \nearrow edistEx_{eA} \\
 & e^2cweA &
 \end{array}$$

The proof is similar with the proof of Lemma 13 and we will not elaborate it here. Also, notice the difference between dist of c over w and distEx of cw over e . While dist is a natural transformation, distEx is a natural isomorphism.

Lemma 19. let $(cw, \varepsilon^{cw}, \delta^{cw})$ and $(e, \varepsilon^e, \delta^e)$ be two monoidal comonads on a Lambek category with cw and exchange $(\mathcal{L}, I, \otimes, cw, \text{weak}, \text{contra}, e, \text{ex})$. Then the composition of cw and e using the distributive law for exchange $\text{distEx}_A : cweA \longrightarrow ecwA$ is a monoidal comonad $(cwe, \varepsilon, \delta)$ on \mathcal{L} .

Proof. Suppose $(cw, \varepsilon^{cw}, \delta^{cw})$ and $(e, \varepsilon^e, \delta^e)$ are monoidal comonads, and $(\mathcal{L}, I, \otimes, cw, \text{weak}, \text{contra}, e, \text{ex})$ is a Lambek category with cw and exchange. Since by definition $cw, e : \mathcal{L} \longrightarrow \mathcal{L}$ are monoidal functors, we know that their composition $cwe : \mathcal{L} \longrightarrow \mathcal{L}$ is a monoidal functor:

$$\begin{aligned} q_{A,B} &: cweA \otimes cweB \longrightarrow cwe(A \otimes B) \\ q_{A,B} &= cwq_{A,B}^e \circ q_{eA,eB}^{cw} \\ q_I &: I \longrightarrow cweI \\ q_I &= cwq_I^e \circ q_I^{cw} \end{aligned}$$

Analogous to the proof of Lemma 14, each of ε and δ can be given two equivalent definitions:

$$\begin{array}{ccc} cweA & \xrightarrow{\varepsilon^{cw}_{eA}} & eA \\ \downarrow cw\varepsilon_A^e & & \downarrow \varepsilon_A^e \\ cwA & \xrightarrow{\varepsilon_A^{cw}} & A \end{array} \quad \begin{array}{ccccc} cweA & \xrightarrow{cw\delta_A^e} & cwe^2A & \xrightarrow{\delta_{e^2A}^{cw}} & (cw)^2e^2A \\ \downarrow \delta_{eA}^{cw} & & \downarrow \delta_{e^2A}^{cw} & & \downarrow cwe_{eA} \\ (cw)^2eA & \xrightarrow{(cw)^2\delta_A^e} & (cw)^2e^2A & \xrightarrow{cwe_{eA}} & cwe cweA \end{array}$$

And the comonad laws can be proved similarly, which we will not elaborate for simplicity. \square

Definition 20. Suppose $(\mathcal{L}, cwe, \text{contra}, \text{weak}, \text{ex})$ is a Lambek category with contraction, weakening, and exchange. Then the **Eilenberg Moore category**, \mathcal{L}^{cwe} , of the comonad $(cwe, \varepsilon, \delta)$ has as objects all the cwe -coalgebras $(A, h_A : A \longrightarrow eA)$, and as morphisms all the coalgebras morphisms.

Should the two diagrams in Definition 4 also commute?

3 Related Work

TODO

4 Conclusion

TODO

References

- [1] P. N. Benton. A mixed linear and non-linear logic: Proofs, terms and models (preliminary report). Technical Report UCAM-CL-TR-352, University of Cambridge Computer Laboratory, 1994. Accessible online at <http://research.microsoft.com/en-us/um/people/nick/mixed3.ps>.

A Appendix

A.1 Symmetric Monoidal Categories

Definition 21. A *monoidal category* is a category, \mathcal{M} , with the following data:

- An object \top of \mathcal{M} ,
- A bi-functor $\otimes : \mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$,
- The following natural isomorphisms:

$$\begin{aligned} \lambda_A &: \top \otimes A \longrightarrow A \\ \rho_A &: A \otimes \top \longrightarrow A \\ \alpha_{A,B,C} &: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C) \end{aligned}$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc}
((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\
\downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\
(A \otimes B) \otimes (C \otimes D) & & \\
\downarrow \alpha_{A, B, C \otimes D} & & \\
A \otimes (B \otimes (C \otimes D)) & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & A \otimes ((B \otimes C) \otimes D)
\end{array}$$

$$\begin{array}{ccc}
(A \otimes \top) \otimes B & \xrightarrow{\alpha_{A, \top, B}} & A \otimes (\top \otimes B) \\
\downarrow \rho_A \otimes \text{id}_B & & \downarrow \text{id}_A \otimes \lambda_B \\
& A \otimes B &
\end{array}$$

Definition 22. A *symmetric monoidal category (SMC)* is a category, \mathcal{M} , with the following data:

- An object \top of \mathcal{M} ,
- A bi-functor $\otimes : \mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}$,
- The following natural isomorphisms:

$$\begin{aligned}
\lambda_A &: \top \otimes A \longrightarrow A \\
\rho_A &: A \otimes \top \longrightarrow A \\
\alpha_{A,B,C} &: (A \otimes B) \otimes C \longrightarrow A \otimes (B \otimes C)
\end{aligned}$$

- A symmetry natural isomorphism:

$$\beta_{A,B} : A \otimes B \longrightarrow B \otimes A$$

- Subject to the following coherence diagrams:

$$\begin{array}{ccc}
((A \otimes B) \otimes C) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes \text{id}_D} & (A \otimes (B \otimes C)) \otimes D \\
\downarrow \alpha_{A \otimes B, C, D} & & \downarrow \alpha_{A, B \otimes C, D} \\
(A \otimes B) \otimes (C \otimes D) & & \\
\downarrow \alpha_{A, B, C \otimes D} & & \\
A \otimes (B \otimes (C \otimes D)) & \xleftarrow{\text{id}_A \otimes \alpha_{B,C,D}} & A \otimes ((B \otimes C) \otimes D)
\end{array}$$

$$\begin{array}{ccccc}
(A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\beta_{A, B \otimes C}} & (B \otimes C) \otimes A \\
\downarrow \beta_{A,B} \otimes \text{id}_C & & & & \downarrow \alpha_{B,C,A} \\
(B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{\text{id}_B \otimes \beta_{A,C}} & B \otimes (C \otimes A)
\end{array}$$

$$\begin{array}{ccc}
(A \otimes \top) \otimes B & \xrightarrow{\alpha_{A, \top, B}} & A \otimes (\top \otimes B) \\
\downarrow \rho_A \otimes \text{id}_B & & \downarrow \text{id}_A \otimes \lambda_B \\
& A \otimes B &
\end{array}$$

$$\begin{array}{ccc}
A \otimes B & & \\
\downarrow \beta_{A,B} & \searrow \text{id}_{A \otimes B} & \\
B \otimes A & \xrightarrow{\beta_{B,A}} & A \otimes B
\end{array}$$

$$\begin{array}{ccc}
\top \otimes A & \xrightarrow{\beta_{\top, A}} & A \otimes \top \\
\downarrow \lambda_A & & \downarrow \rho_A \\
& A &
\end{array}$$

Definition 23. A **monoidal biclosed category** is a monoidal category $(\mathcal{M}, \tau, \otimes)$, such that, for any object B of \mathcal{M} , each of the functors $- \otimes B : \mathcal{M} \rightarrow \mathcal{M}$ and $B \otimes - : \mathcal{M} \rightarrow \mathcal{M}$ has a specified right adjoint. Hence, for any object A and C of \mathcal{M} , there are two objects $C \leftarrow B$ and $B \rightarrow C$ of \mathcal{M} and two natural bijections:

$$\begin{aligned}\text{Hom}_{\mathcal{M}}(A \otimes B, C) &\cong \text{Hom}_{\mathcal{M}}(A, C \leftarrow B) \\ \text{Hom}_{\mathcal{M}}(B \otimes A, C) &\cong \text{Hom}_{\mathcal{M}}(A, B \rightarrow C)\end{aligned}$$

Definition 24. A **symmetric monoidal closed category (SMCC)** is a symmetric monoidal category, $(\mathcal{M}, \tau, \otimes)$, such that, for any object B of \mathcal{M} , the functor $- \otimes B : \mathcal{M} \rightarrow \mathcal{M}$ has a specified right adjoint. Hence, for any objects A and C of \mathcal{M} there is an object $B \multimap C$ of \mathcal{M} and a natural bijection:

$$\text{Hom}_{\mathcal{M}}(A \otimes B, C) \cong \text{Hom}_{\mathcal{M}}(A, B \multimap C)$$

We call the functor $\multimap : \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$ the internal hom of \mathcal{M} .

Definition 25. Suppose we are given two monoidal categories $(\mathcal{M}_1, \tau_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$. Then a **monoidal functor** is a functor $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$, a map $m_{\tau_1} : \tau_2 \rightarrow F\tau_1$ and a natural transformation $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$ subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A,B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$

$$\begin{array}{ccc} \tau_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\ \downarrow m_{\tau_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\ F\tau_1 \otimes_2 FA & \xrightarrow{m_{\tau_1, A}} & F(\tau_1 \otimes_1 A) \end{array} \quad \begin{array}{ccc} FA \otimes_2 \tau_2 & \xrightarrow{\rho_{2FA}} & FA \\ \downarrow \text{id}_{FA} \otimes m_{\tau_1} & & \uparrow F\rho_{1A} \\ FA \otimes_2 F\tau_1 & \xrightarrow{m_{A, \tau_1}} & F(A \otimes_1 \tau_1) \end{array}$$

Need to notice that the composition of monoidal functors is also monoidal, subject to the above coherence conditions.

Definition 26. Suppose we are given two symmetric monoidal closed categories $(\mathcal{M}_1, \tau_1, \otimes_1, \alpha_1, \lambda_1, \rho_1, \beta_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2, \alpha_2, \lambda_2, \rho_2, \beta_2)$. Then a **symmetric monoidal functor** is a functor $F : \mathcal{M}_1 \rightarrow \mathcal{M}_2$, a map $m_{\tau_1} : \tau_2 \rightarrow F\tau_1$ and a natural transformation $m_{A,B} : FA \otimes_2 FB \rightarrow F(A \otimes_1 B)$ subject to the following coherence conditions:

$$\begin{array}{ccc} (FA \otimes_2 FB) \otimes_2 FC & \xrightarrow{\alpha_{2FA,FB,FC}} & FA \otimes_2 (FB \otimes_2 FC) \\ \downarrow m_{A,B} \otimes \text{id}_{FC} & & \downarrow \text{id}_{FA} \otimes m_{B,C} \\ F(A \otimes_1 B) \otimes_2 FC & & FA \otimes_2 F(B \otimes_1 C) \\ \downarrow m_{A \otimes_1 B, C} & & \downarrow m_{A,B \otimes_1 C} \\ F((A \otimes_1 B) \otimes_1 C) & \xrightarrow{F\alpha_{1A,B,C}} & F(A \otimes_1 (B \otimes_1 C)) \end{array}$$

$$\begin{array}{ccc} \tau_2 \otimes_2 FA & \xrightarrow{\lambda_{2FA}} & FA \\ \downarrow m_{\tau_1} \otimes \text{id}_{FA} & & \uparrow F\lambda_{1A} \\ F\tau_1 \otimes_2 FA & \xrightarrow{m_{\tau_1, A}} & F(\tau_1 \otimes_1 A) \end{array} \quad \begin{array}{ccc} FA \otimes_2 \tau_2 & \xrightarrow{\rho_{2FA}} & FA \\ \downarrow \text{id}_{FA} \otimes m_{\tau_1} & & \uparrow F\rho_{1A} \\ FA \otimes_2 F\tau_1 & \xrightarrow{m_{A, \tau_1}} & F(A \otimes_1 \tau_1) \end{array}$$

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{\beta_{2FA,FB}} & FB \otimes_2 FA \\
\downarrow m_{A,B} & & \downarrow m_{B,A} \\
F(A \otimes_1 B) & \xrightarrow{F\beta_{1A,B}} & F(B \otimes_1 A)
\end{array}$$

Definition 27. Suppose $(\mathcal{M}_1, \tau_1, \otimes_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2)$ are monoidal categories, and (F, m) and (G, n) are monoidal functors between \mathcal{M}_1 and \mathcal{M}_2 . Then a **monoidal natural transformation** is a natural transformation, $f : F \rightarrow G$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A,B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A,B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{f_{\tau_1}} & G\tau_1 \\
\swarrow m_{\tau_1} & & \searrow n_{\tau_1} \\
& \tau_2 &
\end{array}$$

Definition 28. Suppose $(\mathcal{M}_1, \tau_1, \otimes_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2)$ are SMCs, and (F, m) and (G, n) are symmetric monoidal functors between \mathcal{M}_1 and \mathcal{M}_2 . Then a **symmetric monoidal natural transformation** is a natural transformation, $f : F \rightarrow G$, subject to the following coherence diagrams:

$$\begin{array}{ccc}
FA \otimes_2 FB & \xrightarrow{m_{A,B}} & F(A \otimes_1 B) \\
\downarrow f_A \otimes_2 f_B & & \downarrow f_{A \otimes_1 B} \\
GA \otimes_2 GB & \xrightarrow{n_{A,B}} & G(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{f_{\tau_1}} & G\tau_1 \\
\swarrow m_{\tau_1} & & \searrow n_{\tau_1} \\
& \tau_2 &
\end{array}$$

Definition 29. Suppose $(\mathcal{M}_1, \tau_1, \otimes_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2)$ are monoidal categories, and (F, m) is a monoidal functor between \mathcal{M}_1 and \mathcal{M}_2 and (G, n) is a monoidal functor between \mathcal{M}_2 and \mathcal{M}_1 . Then a **monoidal adjunction** is an ordinary adjunction $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$ such that the unit, $\eta_A : A \rightarrow GFA$, and the counit, $\varepsilon_A : FGA \rightarrow A$, are monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow F n_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_{A \otimes_1 B}} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{F n_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{G m_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

Definition 30. Suppose $(\mathcal{M}_1, \tau_1, \otimes_1)$ and $(\mathcal{M}_2, \tau_2, \otimes_2)$ are SMCs, and (F, m) is a symmetric monoidal functor between \mathcal{M}_1 and \mathcal{M}_2 and (G, n) is a symmetric monoidal functor between \mathcal{M}_2 and \mathcal{M}_1 . Then a **symmetric monoidal adjunction** is an ordinary adjunction $\mathcal{M}_1 : F \dashv G : \mathcal{M}_2$ such that the unit, $\eta_A : A \rightarrow GFA$, and the counit, $\varepsilon_A : FGA \rightarrow A$, are

symmetric monoidal natural transformations. Thus, the following diagrams must commute:

$$\begin{array}{ccc}
FGA \otimes_2 FGB & \xrightarrow{m_{GA,GB}} & F(GA \otimes_1 GB) \\
\downarrow \varepsilon_A \otimes_1 \varepsilon_B & & \downarrow F\eta_{A,B} \\
A \otimes_2 B & \xleftarrow{\varepsilon_A \otimes_1 B} & FGA \otimes_2 FGB
\end{array}
\quad
\begin{array}{ccc}
F\tau_1 & \xrightarrow{Fn_{\tau_2}} & FG\tau_2 \\
\uparrow m_{\tau_1} & & \downarrow \varepsilon_{\tau_1} \\
\tau_2 & \xlongequal{\quad} & \tau_2
\end{array}$$

$$\begin{array}{ccc}
GFA \otimes_1 GFB & \xleftarrow{\eta_A \otimes_1 \eta_B} & A \otimes_1 B \\
\downarrow n_{FA,FB} & & \downarrow \eta_{A \otimes_1 B} \\
G(FA \otimes_2 FB) & \xrightarrow{m_{A,B}} & GF(A \otimes_1 B)
\end{array}
\quad
\begin{array}{ccc}
G\tau_2 & \xrightarrow{Gm_{\tau_1}} & GF\tau_1 \\
\uparrow n_{\tau_2} & & \uparrow \eta_{\tau_1} \\
\tau_1 & \xlongequal{\quad} & \tau_1
\end{array}$$

Definition 31. A *monoidal comonad* on a monoidal category C is a triple (T, ε, δ) , where (T, m) is a monoidal endofunctor on C , $\varepsilon_A : TA \rightarrow A$ and $\delta_A : TA \rightarrow T^2A$ are monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\quad
\begin{array}{ccc}
& TA & \\
\swarrow \varepsilon_{TA} & \downarrow \delta_A & \searrow T\varepsilon_A \\
TA & T^2A & TA
\end{array}$$

The assumption that ε and δ are monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\quad
\begin{array}{ccc}
T\tau & \xleftarrow{m_\tau} & \tau \\
& \searrow \varepsilon_\tau & \downarrow \\
& & \tau
\end{array}$$

$$\begin{array}{ccccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) & & \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} & & \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} & T(TA \otimes TB) & \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}
\quad
\begin{array}{ccc}
\tau & \xrightarrow{m_\tau} & T\tau \\
\downarrow m_\tau & & \downarrow \delta_\tau \\
T\tau & \xrightarrow{Tm_\tau} & T^2\tau
\end{array}$$

Definition 32. A *symmetric monoidal comonad* on a symmetric monoidal category C is a triple (T, ε, δ) , where (T, m) is a symmetric monoidal endofunctor on C , $\varepsilon_A : TA \rightarrow A$ and $\delta_A : TA \rightarrow T^2A$ are symmetric monoidal natural transformations, which make the following diagrams commute:

$$\begin{array}{ccc}
TA & \xrightarrow{\delta_A} & T^2A \\
\downarrow \delta_A & & \downarrow T\delta_A \\
T^2A & \xrightarrow{\delta_{TA}} & T^3A
\end{array}
\quad
\begin{array}{ccc}
& TA & \\
\swarrow \varepsilon_{TA} & \downarrow \delta_A & \searrow T\varepsilon_A \\
TA & T^2A & TA
\end{array}$$

The assumption that ε and δ are symmetric monoidal natural transformations amount to the following diagrams commuting:

$$\begin{array}{ccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) \\
& \searrow \varepsilon_A \otimes \varepsilon_B & \downarrow \varepsilon_{A \otimes B} \\
& & A \otimes B
\end{array}
\qquad
\begin{array}{ccc}
T\top & \xleftarrow{m_\top} & \top \\
& \searrow \varepsilon_\top & \parallel \\
& & \top
\end{array}$$

$$\begin{array}{ccccc}
TA \otimes TB & \xrightarrow{m_{A,B}} & T(A \otimes B) & & \\
\downarrow \delta_A \otimes \delta_B & & \downarrow \delta_{A \otimes B} & & \\
T^2A \otimes T^2B & \xrightarrow{m_{TA,TB}} & T(TA \otimes TB) & \xrightarrow{Tm_{A,B}} & T^2(A \otimes B)
\end{array}
\qquad
\begin{array}{ccc}
\top & \xrightarrow{m_\top} & T\top \\
\downarrow m_\top & & \downarrow \delta_\top \\
T\top & \xrightarrow{Tm_\top} & T^2\top
\end{array}$$

B Proofs

B.1 Proof of Composition of Weakening and Contraction (Lemma 14)

Since by definition $w : \mathcal{L} \rightarrow \mathcal{L}$ and $c : \mathcal{L} \rightarrow \mathcal{L}$ are monoidal functors we know that their composition $cw : \mathcal{L} \rightarrow \mathcal{L}$ is a monoidal functor:

$$\begin{aligned}
q_{A,B} &: cwA \otimes cwB \rightarrow cw(A \otimes B) \\
q_{A,B} &= cq_{A,B}^w \circ q_{wA,wB}^c \\
q_I &: I \rightarrow cwI \\
q_I &= cq_I^w \circ q_I^c
\end{aligned}$$

We must now define both $\varepsilon_A : cwA \rightarrow A$ and $\delta_A : cwA \rightarrow cwcwA$, and then show that they are monoidal natural transformations subject to the comonad laws. Since we are composing two comonads each of ε and δ can be given two definitions, but they are equivalent:

- $\varepsilon_A : cwA \rightarrow A$ is defined as in the diagram below, which commutes by the naturality of ε^c .

$$\begin{array}{ccc}
cwA & \xrightarrow{\varepsilon_{wA}^c} & wA \\
\downarrow c\varepsilon_A^w & & \downarrow \varepsilon_A^w \\
cA & \xrightarrow{\varepsilon_A^c} & A
\end{array}$$

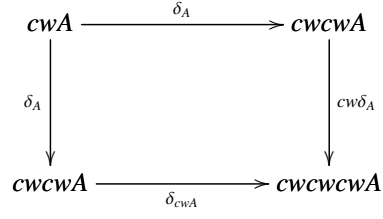
- $\delta_A : cwA \rightarrow cwcwA$ is defined as in the diagram:

$$\begin{array}{ccccc}
cwA & \xrightarrow{c\delta_A^w} & cw^2A & \xrightarrow{\delta_{w^2A}^c} & c^2w^2A \\
\downarrow \delta_{wA}^c & & \downarrow \delta_{w^2A}^c & & \downarrow cdist_{wA} \\
c^2wA & \xrightarrow{c^2\delta_A^w} & c^2w^2A & \xrightarrow{cdist_{wA}} & cwcwA
\end{array}$$

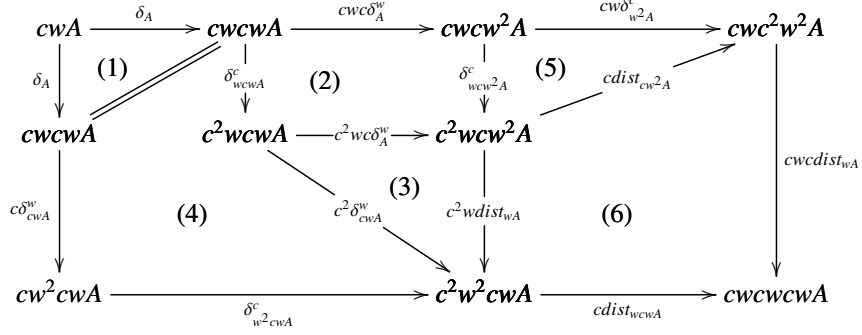
The left part of the diagram commutes by the naturality of δ^c and the right part commutes trivially.

The remainder of the proof shows that the comonad laws hold.

Case 1:

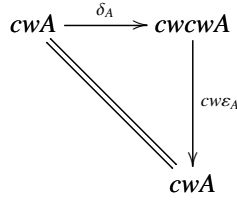


The previous diagram commutes because the following one does.

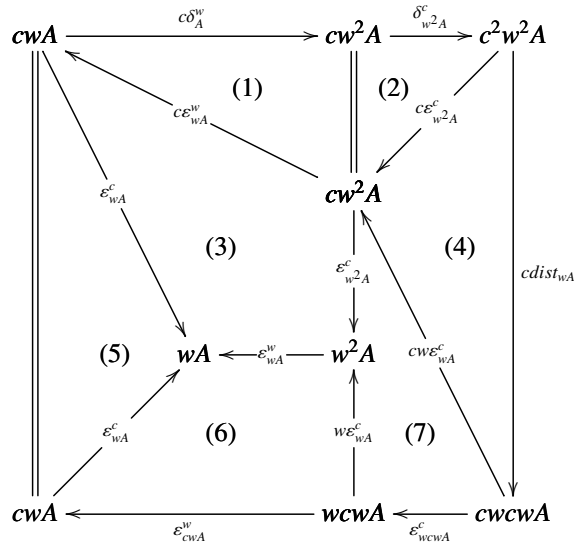


(1) commutes by equality and we will not expand δ_A for simplicity. (2) and (4) commutes by the naturality of δ^c . (3), (5) commutes by the conditions of $dist$. (6) commutes by the naturality of $dist$.

Case 2:

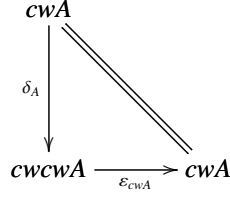


The triangle commutes because of the following diagram chasing.

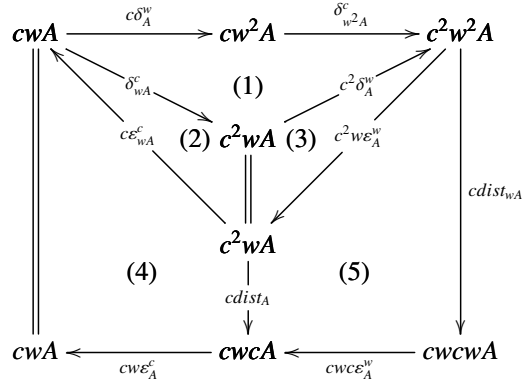


(1) commutes by the comonad law for w with components δ_A^w and ε_{wA}^w . (2) commutes by the comonad law for c with components $\delta_{w^2A}^c$ and $\varepsilon_{w^2A}^c$. (3) and (7) commute by the naturality of ε^c . (4) commutes by the condition of *dist*. (5) commutes trivially. And (6) commutes by the naturality of ε^w .

Case 3:



The previous triangle commutes because the following diagram chasing does.



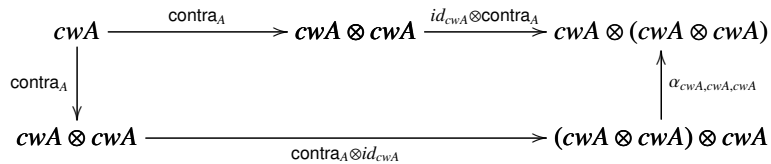
(1) commutes by the naturality of δ^c . (2) is the comonad law for c with components δ_{wA}^c and ε_{wA}^c . (3) is the comonad law for w with components δ_A^w and ε_A^w . (4) commutes by the condition of *dist*. And (5) commute by the naturality of *dist*.

||||| HEAD

B.2 Proof of Conditions of Lambek category with cw (Lemma 16)

1. As shown in the paper.
2. Each $(cwA, \text{weak}_A, \text{contra}_A)$ is a comonoid.

Case 1:



The previous diagram commutes by the following diagram chasing.

$$\begin{array}{ccccc}
cwA & \xrightarrow{\text{contra}_A} & cwA \otimes cwA & \xrightarrow{id_{cwA} \otimes \rho_{cwA}^{-1}} & cwA \otimes (cwA \otimes I) \\
\downarrow \text{contra}_A & \nearrow (1) & \downarrow id_{cwA} \otimes \lambda_{cwA}^{-1} & \nearrow (2) & \downarrow id_{cwA} \otimes \text{contra}_{wA,I} \\
cwA \otimes cwA & \xrightarrow{id_{cwA} \otimes \lambda_{cwA}^{-1}} & cwA \otimes (I \otimes cwA) & & cwA \otimes ((cwA \otimes I) \otimes cwA) \\
\downarrow \rho_{cwA}^{-1} \otimes id_{cwA} & & \downarrow id_{cwA} \otimes \text{contra}_{wA,I} & \nearrow id_{cwA} \otimes \alpha_{cwA,I,cwA} & \downarrow id_{cwA} \otimes (\rho_{cwA} \otimes id_{cwA}) \\
(cwA \otimes I) \otimes cwA & & cwA \otimes (cwA \otimes (I \otimes cwA)) & \xrightarrow{id_{cwA} \otimes (id_{cwA} \otimes \lambda_{cwA})} & cwA \otimes (cwA \otimes cwA) \\
\downarrow \text{contra}_{wA,I} \otimes id_{cwA} & & \downarrow id_{cwA} \otimes \alpha_{cwA,cwA,cwA} & \nearrow (3) & \downarrow \alpha_{cwA,cwA,cwA} \\
((cwA \otimes I) \otimes cwA) \otimes cwA & \xrightarrow{(\rho_{cwA} \otimes id_{cwA}) \otimes id_{cwA}} & (cwA \otimes cwA) \otimes cwA & & \\
& (4) & & &
\end{array}$$

(1) commutes trivially and we would not expand contra for simplicity. (2) and (4) commute because $(\mathcal{L}, c, \text{contraL}, \text{contraR})$ is a Lambek category with contraction. (3) commutes because \mathcal{L} is monoidal.

Case 2:

$$\begin{array}{ccccc}
& & cwA & & \\
& \swarrow \lambda^{-1} & \downarrow \text{contra}_A & \searrow \rho^{-1} & \\
I \otimes cwA & \xleftarrow{\text{weak}_A \otimes id_{cwA}} & cwA \otimes cwA & \xrightarrow{id_{cwA} \otimes \text{weak}_A} & cwA \otimes I
\end{array}$$

The diagram above commutes by the following diagram chasing.

$$\begin{array}{ccccc}
I \otimes cwA & \xleftarrow{\text{weak}_A^w \otimes id_{cwA}} & wA \otimes cwA & & \\
\uparrow id_I \otimes \lambda_{cwA} & & \uparrow id_{wA} \otimes \lambda_{cwA} & & \\
I \otimes (I \otimes cwA) & \xleftarrow{\text{weak}_A^w \otimes id_{I \otimes cwA}} & wA \otimes (I \otimes cwA) & & \\
\uparrow \lambda_{I \otimes cwA}^{-1} & & \uparrow \varepsilon_{wA}^c \otimes id_{I \otimes cwA} & & \\
I \otimes cwA & \xrightarrow{\text{contra}_{wA,I}} & cwA \otimes (I \otimes cwA) & & \\
\uparrow \lambda_{cwA}^{-1} & & \downarrow id_{cwA} \otimes \lambda_{cwA} & & \\
cwA & & cwA \otimes cwA & & \\
\downarrow \rho_{cwA}^{-1} & & \downarrow \rho_{cwA} \otimes id_{cwA} & & \\
cwA \otimes I & \xrightarrow{\text{contra}_{wA,I}} & (cwA \otimes I) \otimes cwA & & \\
\downarrow \rho_{cwA}^{-1} & & \downarrow id_{cwA \otimes I} \otimes \varepsilon_{wA}^c & & \\
(cwA \otimes I) \otimes I & \xleftarrow{id_{cwA \otimes I} \otimes \text{weak}_A^w} & (cwA \otimes I) \otimes wA & & \\
\downarrow \rho_{cwA} \otimes id_I & & \downarrow \rho_{cwA} \otimes id_{wA} & & \\
cwA \otimes I & \xleftarrow{id_{cwA} \otimes \text{weak}_A} & cwA \otimes wA & &
\end{array}$$

(1), (2) and (3) commute by the functionality of λ . (6), (7) and (8) commute by the functionality of ρ . (4) and (9) are conditions of the Lambek category with cw . And (5) is the definition of contra .

3. weak and contra are coalgebra morphisms.

Case 1:

$$\begin{array}{ccc}
 cwA & \xrightarrow{\text{weak}_A} & I \\
 \delta_A \downarrow & & \downarrow q_I \\
 cwcwA & \xrightarrow{cw\text{weak}_A} & cwI
 \end{array}$$

The previous diagram commutes by the diagram below. (1) commutes by the naturality of δ^c . (2) commutes by the condition of dist_{wA} . (3), (5) and (6) commute because c is a monoidal comonad. (4) commutes because $(\mathcal{L}, w, \text{weak}^w)$ is a Lambek category with weakening. (7) commutes because c and w are monoidal comonads.

$$\begin{array}{ccccc}
 cwA & \xrightarrow{c\text{weak}_A^w} & cI & \xrightarrow{\epsilon_I^c} & I \\
 \delta_{wA}^c \downarrow & \searrow c\delta_A^w & \swarrow q_I^c & \text{(5)} & \downarrow q_I^c \\
 c^2wA & & & & \\
 c^2\delta_A^w \downarrow & \text{(1)} & \text{(4)} & & \text{(6)} \\
 c^2w^2A & \xleftarrow{\delta_{w^2A}^c} & cw^2A & & cI \\
 c\text{dist}_{wA} \downarrow & \searrow c\epsilon_{wA}^c & \parallel & \text{(7)} & \downarrow cq_I^w \\
 cwcwA & \xrightarrow{cw\epsilon_{wA}^c} & cw^2A & \xleftarrow{cw\text{weak}_A^w} & cwI
 \end{array}$$

Case 2:

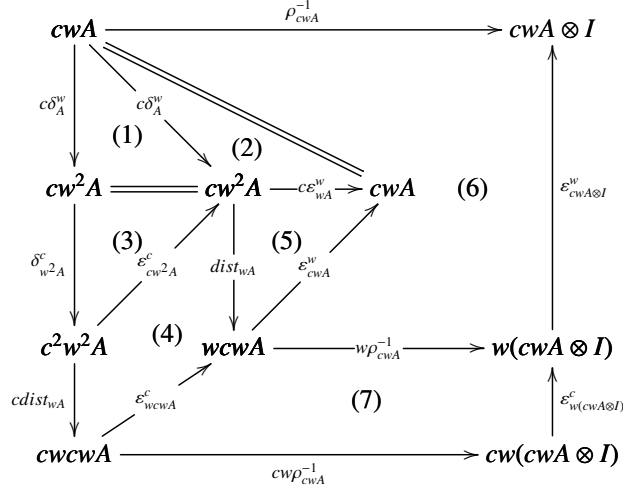
$$\begin{array}{ccccc}
 cwA & \xrightarrow{\text{contra}_A} & cwA \otimes cwA & \xrightarrow{\delta_A \otimes \delta_A} & cwcwA \otimes cwcwA \\
 \delta_A \downarrow & & & & \downarrow q_{cwA, cwA} \\
 cwcwA & \xrightarrow{cw\text{contra}_A} & & & cw(cwA \otimes cwA)
 \end{array}$$

To prove the previous diagram commute, we first expand it, Then we divide it into five parts as shown below, and prove each part commutes.

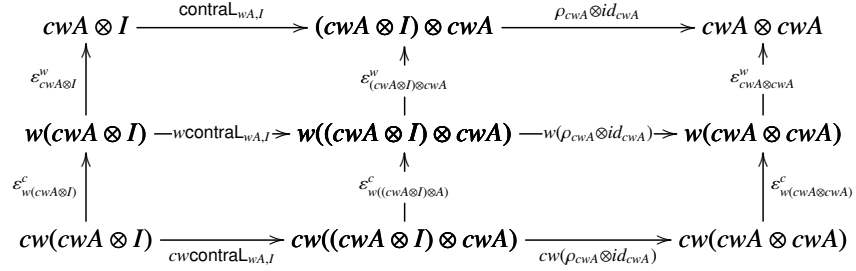
$$\begin{array}{ccccccc}
 cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I & \xrightarrow{\text{contra}_{wA, I}} & (cwA \otimes I) \otimes cwA \\
 \downarrow c\delta_A^w & & \uparrow \epsilon_{cwA \otimes I}^w & & \downarrow \rho_{cwA} \otimes id_{cwA} \\
 cw^2A & & w(cwA \otimes I) & & cwA \otimes cwA \\
 \downarrow \delta_{w^2A}^c & & \uparrow \epsilon_{w(cwA \otimes I)}^c & & \downarrow c\delta_A^w \otimes c\delta_A^w \\
 c^2w^2A & & & & cw^2A \otimes cw^2A \\
 \downarrow c\text{dist}_{wA} & & & & \text{(a)} \\
 cwcwA & & & & cw^2A \otimes cw^2A \\
 \downarrow cw\rho_{cwA} & & & & \text{(b)} \\
 cw(cwA \otimes I) & & & & c^2w^2A \otimes c^2w^2A \\
 \downarrow cw\text{contra}_{wA, I} & & & & \downarrow c\text{dist}_{wA} \otimes c\text{dist}_{wA} \\
 cw((cwA \otimes I) \otimes cwA) & \xrightarrow{cw(\rho_{cwA} \otimes id_{cwA})} & cw(cwA \otimes cwA) & \xleftarrow{cq_{cwA \otimes cwA}^w} & c(wcwA \otimes wcwA)
 \end{array}$$

Part (a) and (b) are comonad laws.

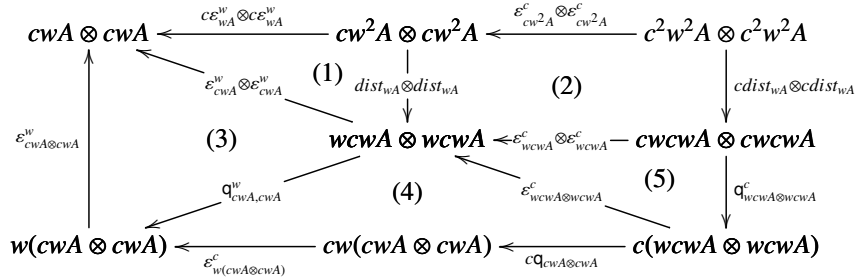
Part (c) commutes by the following diagram chase. (1) is equality. (2) is the comonad law for w . (3) is the comonad law for c . (4) commutes by the naturality of ε^c . (5) is one of the conditions for $dist_{wA}$. (6) commutes by the naturality of ε^w . And (7) commutes by the naturality of ε^c .



Part (d) commutes by the following diagram chase. The upper two squares both commute by the naturality of ε^w , and the lower two squares commute by the naturality of ε^c .

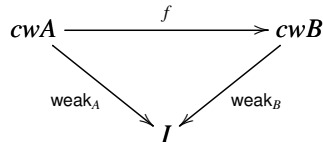


Part (e) commutes by the following diagram. (1) commutes by the condition of $dist_{wA}$. (2) and (4) commute by the naturality of ε^c . (3) and (5) commute because w and c are monoidal comonads.



4. Any coalgebra morphism $f : (cwA, \delta_A) \rightarrow (cwB, \delta_B)$ between free coalgebras preserves the comonoid structure given by weak and contra.

Case 1: This coherence diagram is given in the definition of the Lambek category with cw .



Case 2:

$$\begin{array}{ccc}
 cwA & \xrightarrow{\text{contra}_A} & cwA \otimes cwA \\
 f \downarrow & & \downarrow f \otimes f \\
 cwB & \xrightarrow{\text{contra}_B} & cwB \otimes cwB
 \end{array}$$

The square commutes by the diagram chasing below, which commutes by the naturality of ρ and contraL .

$$\begin{array}{ccccccc}
 cwA & \xrightarrow{\rho_{cwA}^{-1}} & cwA \otimes I & \xrightarrow{\text{contraL}_{wA,I}} & (cwA \otimes I) \otimes cwA & \xrightarrow{\rho_{cwA} \otimes id_{cwA}} & cwA \otimes cwA \\
 \downarrow cw f & & \downarrow cw f \otimes id_I & & \downarrow (cw f \otimes id_I) \otimes cw f & & \downarrow cw f \otimes cw f \\
 cwB & \xrightarrow{\rho_{cwB}^{-1}} & cwB \otimes I & \xrightarrow{\text{contraL}_{wB,I}} & (cwB \otimes I) \otimes cwB & \xrightarrow{\rho_{cwB} \otimes id_{cwB}} & cwB \otimes cwB
 \end{array}$$

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