×

Week 3 - Problem Set

Back to Week 3

8/10 points earned (80%)

Quiz passed!

1 / 1
points

1.
Suppose a MAC system (S, V) is used to protect files in a file system

by appending a MAC tag to each file. The MAC signing algorithm S

is applied to the file contents and nothing else. What tampering attacks

are not prevented by this system?

- ○ Changing the last modification time of a file.

Correct
The MAC signing algorithm is only applied to the file contents and

does not protect the file meta data.

- ○ Replacing the contents of a file with the concatenation of two files

  on the file system.
- ○ Changing the first byte of the file contents.
- ○ Replacing the tag and contents of one file with the tag and contents of a file

  from another computer protected by the same MAC system, but a different key.

1 / 1
points

2.
Let (S,V) be a secure MAC defined over (K,M,T) where $M=\{0,1\}^n$ and $T=\{0,1\}^{128}$.
That is, the key space is K, message space is $\{0,1\}^n$, and tag space is $\{0,1\}^{128}$.

Which of the following is a secure MAC: (as usual, we use $\|$ to denote string concatenation)

- [ ] $S'(k,\ m) = [t \leftarrow S(k,m),\ \text{output}\ (t,t)\,)$ and

  $$V'(k,m,(t_1,t_2)) = \begin{cases} V(k,m,t_1) & \text{if } t_1 = t_2 \\ \text{"0"} & \text{otherwise} \end{cases}$$

  (i.e., $V'(k,m,(t_1,t_2))$ only outputs "1"

  if t_1 and t_2 are equal and valid)

  Correct
  a forger for (S',V') gives a forger for (S,V).

- [ ] $S'(k,m) = S(k, m \oplus m)$ and

  $$V'(k,m,t) = V(k,\ m \oplus m,\ t)$$

  Un-selected is correct

- [ ] $S'(k,m) = (S(k,m), S(k,0^n))$ and

  $$V'(k,m,(t_1,t_2)) = [V(k,m,t_1) \text{ and } V(k,0^n,t_2)]$$

  (i.e., $V'(k,m,(t_1,t_2))$ outputs ``1'' if both t_1 and t_2 are valid tags)

  Un-selected is correct

- [ ] $S'(k,m) = S(k,\ \ m[0,...,n-2] \,\|\, 0)$ and

  $$V'(k,m,t) = V(k,\ \ m[0,...,n-2] \,\|\, 0,\ \ t)$$

  Un-selected is correct

- [ ] $S'(k,m) = S(k,\ m \,\|\, m)$ and

  $$V'(k,m,t) = V(k,\ m \,\|\, m,\ t).$$

  Correct
  a forger for (S',V') gives a forger for (S,V).

- [ ] $S'(k,m) = S(k, m \oplus 1^n)$ and

1 / 1
points

3.
Recall that the ECBC-MAC uses a fixed IV (in the lecture we simply set the IV to 0). Suppose instead we chose a random IV for every message being signed and include the IV in the tag. In other words, $S(k,m) := \left(r,\ \ \mathrm{ECBC}_r(k,m)\right)$

where $\mathrm{ECBC}_r(k,m)$ refers to the ECBC function using r as

the IV. The verification algorithm V given key k, message m,

and tag (r,t) outputs ``1'' if $t = \mathrm{ECBC}_r(k,m)$ and outputs

``0'' otherwise.

The resulting MAC system is insecure.

An attacker can query for the tag of the 1-block message m

and obtain the tag (r,t). He can then generate the following

existential forgery: (we assume that the underlying block cipher

operates on n-bit blocks)

- The tag $(r \oplus 1^n,\ t)$ is a valid tag for the 1-block

  message $m \oplus 1^n$.

Correct
The CBC chain initiated with the IV $r \oplus m$ and applied

to the message 0^n will produce exactly the same output

as the CBC chain initiated with the IV r and applied to the

message m. Therefore, the tag $(r \oplus 1^n,\ t)$ is a valid

existential forgery for the message $m \oplus 1^n$.

- The tag $(r,\ t \oplus r)$ is a valid tag for the 1-block message 0^n.
- The tag $(r \oplus t,\ m)$ is a valid tag for the 1-block message 0^n.
- The tag $(m \oplus t,\ t)$ is a valid tag for the 1-block message 0^n.

0 / 1
points

4.
Suppose Alice is broadcasting packets to 6 recipients

$B_1, ..., B_6$. Privacy is not important but integrity is.

In other words, each of $B_1, ..., B_6$ should be assured that the

packets he is receiving were sent by Alice.

Alice decides to use a MAC. Suppose Alice and $B_1, ..., B_6$ all

share a secret key k. Alice computes a tag for every packet she

sends using key k. Each user B_i verifies the tag when

receiving the packet and drops the packet if the tag is invalid.

Alice notices that this scheme is insecure because user B_1 can

use the key k to send packets with a valid tag to

users $B_2, ..., B_6$ and they will all be fooled into thinking

that these packets are from Alice.

Instead, Alice sets up a set of 4 secret keys $S = \{k_1, ..., k_4\}$.

She gives each user B_i some subset $S_i \subseteq S$

of the keys. When Alice transmits a packet she appends 4 tags to it

by computing the tag with each of her 4 keys. When user B_i receives

a packet he accepts it as valid only if all tags corresponding

to his keys in S_i are valid. For example, if user B_1 is given keys $\{k_1, k_2\}$ he will accept an incoming packet only if the first and second tags are valid. Note that B_1 cannot validate the 3rd and 4th tags because he does not have k_3 or k_4.

How should Alice assign keys to the 6 users so that no single user

can forge packets on behalf of Alice and fool some other user?

- [ ] $S_1 = \{k_1, k_2\},\ S_2 = \{k_1, k_3\},\ S_3 = \{k_1, k_4\},\ S_4 = \{k_2, k_3, k_4\},\ S_5 = \{k_2, k_3\},\ S_6 = \{k_3, k_4\}$

Un-selected is correct

- [ ] $S_1 = \{k_1, k_2\},\ S_2 = \{k_1, k_3\},\ S_3 = \{k_1, k_4\},\ S_4 = \{k_2, k_3\},\ S_5 = \{k_2, k_4\},\ S_6 = \{k_3, k_4\}$

This should be selected

- [ ] $S_1 = \{k_1, k_2\},\ S_2 = \{k_2, k_3\},\ S_3 = \{k_3, k_4\},\ S_4 = \{k_1, k_3\},\ S_5 = \{k_1, k_2\},\ S_6 = \{k_1, k_4\}$

1 / 1
points

5.
Consider the encrypted CBC MAC built from AES. Suppose we

compute the tag for a long message m comprising of n AES blocks.

Let m' be the n-block message obtained from m by flipping the

last bit of m (i.e. if the last bit of m is b then the last bit

of m' is $b \oplus 1$). How many calls to AES would it take

to compute the tag for m' from the tag for m and the MAC key? (in this question please ignore message padding and simply assume that the message length is always a multiple of the AES block size)

- 3
- 4

Correct
You would decrypt the final CBC MAC encryption step done using k_2,

the decrypt the last CBC MAC encryption step done using k_1,

flip the last bit of the result, and re-apply the two encryptions.

- $n + 1$
- 2

1 / 1
points

6.
Let $H: M \to T$ be a collision resistant hash function.

Which of the following is collision resistant:

(as usual, we use ‖ to denote string concatenation)

- [ ] $H'(m) = H(m \parallel m)$

Correct
a collision finder for H' gives a collision finder for H.

- [ ] $H'(m) = H(m) \oplus H(m)$

Un-selected is correct

- [ ] $H'(m) = H(0)$

Un-selected is correct

- [ ] $H'(m) = H(m) \oplus H(m \oplus 1^{|m|})$

(where $m \oplus 1^{|m|}$ is the complement of m)

Un-selected is correct

- [ ] $H'(m) = H(m) \parallel H(m)$

Correct
a collision finder for H' gives a collision finder for H.

- [ ] $H'(m) = H(H(m))$

Correct
a collision finder for H' gives a collision finder for H.

- [ ] $H'(m) = H(|m|)$

(i.e. hash the length of m)

Un-selected is correct

1 / 1
points

7.
Suppose H_1 and H_2 are collision resistant

hash functions mapping inputs in a set M to $\{0,1\}^{256}$.

Our goal is to show that the function $H_2(H_1(m))$ is also

collision resistant. We prove the contra-positive:

suppose $H_2(H_1(\cdot))$ is not collision resistant, that is, we are

given $x \neq y$ such that $H_2(H_1(x)) = H_2(H_1(y))$.

We build a collision for either H_1 or for H_2.

This will prove that if H_1 and H_2 are collision resistant

then so is $H_2(H_1(\cdot))$. Which of the following must be true:

- Either x, y are a collision for H_2 or

  $H_1(x), H_1(y)$ are a collision for H_1.

- Either x, y are a collision for H_1 or

  $H_1(x), H_1(y)$ are a collision for H_2.

Correct
If $H_2(H_1(x)) = H_2(H_1(y))$ then

either $H_1(x) = H_1(y)$ and $x \neq y$, thereby giving us

a collision on H_1. Or $H_1(x) \neq H_1(y)$ but

$H_2(H_1(x)) = H_2(H_1(y))$ giving us a collision on H_2.

Either way we obtain a collision on H_1 or H_2 as required.

- Either x, y are a collision for H_1 or

  x, y are a collision for H_2.

- Either $H_2(x), H_2(y)$ are a collision for H_1 or

  x, y are a collision for H_2.

0 / 1
points

8.
In this question you are asked to find a collision for the compression function:

$f_1(x, y) = \text{AES}(y, x) \oplus y,$

where $\text{AES}(x, y)$ is the AES-128 encryption of y under key x.

Your goal is to find two distinct pairs (x_1, y_1) and (x_2, y_2) such that $f_1(x_1, y_1) = f_1(x_2, y_2)$.

Which of the following methods finds the required (x_1, y_1) and (x_2, y_2)?

- Choose x_1, y_1, x_2 arbitrarily (with $x_1 \neq x_2$) and let $v := AES(y_1, x_1)$.

  Set $y_2 = AES^{-1}(x_2,\ v \oplus y_1 \oplus x_2)$

This should not be selected
This does not work.

- Choose x_1, y_1, y_2 arbitrarily (with $y_1 \neq y_2$) and let $v := AES(y_1, x_1)$.

  Set $x_2 = AES^{-1}(y_2,\ v \oplus y_1 \oplus y_2)$

- Choose x_1, y_1, y_2 arbitrarily (with $y_1 \neq y_2$) and let $v := AES(y_1, x_1)$.

  Set $x_2 = AES^{-1}(y_2,\ v \oplus y_1)$

- Choose x_1, y_1, y_2 arbitrarily (with $y_1 \neq y_2$) and let $v := AES(y_1, x_1)$.

  Set $x_2 = AES^{-1}(y_2,\ v \oplus y_2)$

1 / 1
points

9.
Repeat the previous question, but now to find a collision for the compression function $f_2(x, y) = \text{AES}(x, x) \oplus y$.

Which of the following methods finds the required (x_1, y_1) and (x_2, y_2)?

- Choose x_1, x_2, y_1 arbitrarily (with $x_1 \neq x_2$) and set

  $y_2 = y_1 \oplus AES(x_1, x_1) \oplus AES(x_2, x_2)$

Correct
Awesome!

- Choose x_1, x_2, y_1 arbitrarily (with $x_1 \neq x_2$) and set

  $y_2 = y_1 \oplus x_1 \oplus AES(x_2, x_2)$

- Choose x_1, x_2, y_1 arbitrarily (with $x_1 \neq x_2$) and set

  $y_2 = y_1 \oplus AES(x_1, x_1)$

- Choose x_1, x_2, y_1 arbitrarily (with $x_1 \neq x_2$) and set

  $y_2 = AES(x_1, x_1) \oplus AES(x_2, x_2)$

1 / 1
points

10.
Let $H: M \to T$ be a random hash function where $|M| \gg |T|$ (i.e. the size of M is much larger than the size of T).

In lecture we showed

that finding a collision on H can be done with $O(|T|^{1/2})$

random samples of H. How many random samples would it take

until we obtain a three way collision, namely distinct strings x, y, z

in M such that $H(x) = H(y) = H(z)$?

- $O(|T|^{2/3})$

Correct
An informal argument for this is as follows: suppose we

collect n random samples. The number of triples among the n

samples is n choose 3 which is $O(n^3)$. For a particular

triple x, y, z to be a 3-way collision we need $H(x) = H(y)$

and $H(x) = H(z)$. Since each one of these two events happens

with probability $1/|T|$ (assuming H behaves like a random

function) the probability that a particular triple is a 3-way

collision is $O(1/|T|^2)$. Using the union bound, the probability

that some triple is a 3-way collision is $O(n^3/|T|^2)$ and since

we want this probability to be close to 1, the bound on n

follows.

- $O(|T|^{3/4})$
- $O(|T|)$
- $O(|T|^{1/4})$