

Project Specification

*Project due Thursday, 23 May 2019, 5:00PM
Worth 15%*

The objective of this project is to practice and assess your understanding of logic programming and Prolog. You will write code to solve maths puzzles.

Maths Puzzles

A maths puzzle is a square grid of squares, each to be filled in with a single digit 1–9 (zero is not permitted) satisfying these constraints:

- each row and each column contains no repeated digits;
- all squares on the diagonal line from upper left to lower right contain the same value; and
- the heading of each row and column (leftmost square in a row and topmost square in a column) holds either the sum or the product of all the digits in that row or column

Note that the row and column headings are not considered to be part of the row or column, and so may be filled with a number larger than a single digit. The upper left corner of the puzzle is not meaningful.

When the puzzle is originally posed, most or all of the squares will be empty, with the headings filled in. The goal of the puzzle is to fill in all the squares according to the rules. A proper maths puzzle will have at most one solution.

Here is an example puzzle as posed (left) and solved (right):

	14	10	35
14			
15			
28		1	

	14	10	35
14	7	2	1
15	3	7	5
28	4	1	7

The Program

You will write Prolog code to solve maths puzzles. Your program should supply a predicate `puzzle_solution(Puzzle)` that holds when *Puzzle* is the representation of a solved maths puzzle.

A maths puzzle will be represented as a list of lists, each of the same length, representing a single row of the puzzle. The first element of each list is considered to be the header for that

row. Each element but the first of the first list in the puzzle is considered to be the header of the corresponding column of the puzzle. The first element of the first element of the list is the corner square of the puzzle, and thus is ignored.

You can assume that when your `puzzle_solution/1` predicate is called, its argument will be a proper list of proper lists, and all the header squares of the puzzle (plus the ignored corner square) are bound to integers. Some of the other squares in the puzzle may also be bound to integers, but the others will be unbound. When `puzzle_solution/1` succeeds, its argument must be ground. You may assume your code will only be tested with proper puzzles, which have at most one solution. Of course, if the puzzle is not solvable, the predicate should fail, and it should never succeed with a puzzle argument that is not a valid solution. For example, your program would solve the above puzzle as below:

```
?- Puzzle=[[0,14,10,35],[14,_,_,_],[15,_,_,_],[28,_,1,_],
| puzzle_solution(Puzzle).
Puzzle = [[0, 14, 10, 35], [14, 7, 2, 1], [15, 3, 7, 5], [28, 4, 1, 7]] ;
false.
```

Your `puzzle_solution/1` predicate should be written in the file `proj2.pl`, but you may submit other prolog files if you like. If you do use multiple files, make sure `proj2.pl` loads the other files. You may also use Prolog library modules supported by SWI Prolog as installed on the servers, which is version 7.6.4. You may, but need not, use SWI Prolog's Constraint Logic Programming facilities to solve the problem.

Hints

1. Begin by unifying all the squares on the diagonal. This only needs to be done once to ensure that the puzzle satisfies the first constraint.
2. It is fairly simple to check if a row of the puzzle (one of the inner lists) is a valid solution. Simply check that the first list element is equal to the sum or product of the other elements. Checking the columns is a bit more complicated. However, the library provides a `transpose/2` predicate. If you transpose the puzzle, the columns become rows, so you can check the columns by checking the rows of the transposed puzzle. Load the library module containing the `transpose/2` predicate with the directive

```
:- ensure_loaded(library(clpfd)).
```

3. A very simple strategy for solving puzzles is to backtrack over all possible values for each puzzle square, and test that it is a valid puzzle solution. For a 2×2 puzzle, there are $9^4 = 6561$ filled puzzles. This is very tractable, as long as the code to check that a puzzle is valid is reasonably efficient. This is a good way to get started.
4. However, even for a 3×3 puzzle, there are $9^9 = 387,420,489$ filled puzzles, and even for very efficient code to check if a solution is valid, the strategy of Hint 3 will not be tractable.

But note that if the first row of the filled puzzle is not valid, the whole puzzle will be invalid, so there is no point exploring those possibilities. For example, if the heading of the first row is **23**, only 6 of the $9^3 = 729$ possible ways to fill it could possibly be

valid. More importantly, if you check that each row and column is valid as soon as it is generated, rather than generating the whole grid before checking, you can cut the search space enormously. If all three rows of a 3×3 puzzle have only 6 solutions, then there are only $6^3 = 216$ candidate puzzles to check (since once all the rows have been filled in, then so have all the columns).

5. For some headings, there are far more than 6 possible rows or columns. For example, a three square row or column with a heading of **16** has 54 valid ways to fill it, and for a four square row or column, there are 192 possibilities. Thus for some puzzles, it will be necessary to further cut the number of possibilities to explore.

One way to do this is to first fill the row or column with the fewest possibilities. Consider the example puzzle: the last row has only two possible ways to fill it: 7—1—4 and 4—1—7, while the last column has six possible ways.

Once one row or column is filled, that will reduce the number of possibilities in other rows and columns. In the example above, if the last row is filled with 7—1—4, There are, in fact, no ways to fill the last column, thus the choice of 7—1—4 for the last row can be immediately dismissed (since no selection for any of the other squares can get around the fact that there are no two digits that can be added to or multiplied by 4 to get 35). This leaves 4—1—7 as the only possible way to fill in the last row. By repeatedly choosing the row or column with the fewest alternatives and (nondeterministically) filling it with valid values, you can achieve adequate performance for this project.

6. The easiest way to do all this in Prolog is to reason about which variables are bound and which are unbound as a way to determine which squares have already been filled. For example, you can use the built-in predicate `ground(Term)`, which holds when *Term* is ground at the time the test is performed. When this succeeds for the puzzle, that means the entire puzzle has been filled in. At this point you can check that it is bound to a valid solution, and if so, succeed. If the puzzle is not ground, you will need to select a row or column and fill it in. Doing this will fill in one square in all perpendicular columns or rows, substantially cutting down on the number of valid ways they can be filled. In fact, if you have taken Hint 1, then filling in one row will fill in one square in all the other rows, and one or two squares in all of the columns, and the reverse if you fill a column.

The easiest way to decide which row or column has the fewest solutions is to use the `bagof/3` predicate, and take the length of the resulting list of solutions.

Assessment

Your project will be assessed on the following criteria:

- 30%** Quality of your code and documentation;
- 20%** The ability of your program to correctly solve 2×2 maths puzzles;
- 30%** The ability of your program to correctly solve 3×3 maths puzzles; and
- 20%** The ability of your program to correctly solve 4×4 maths puzzles.

Note that timeouts will be imposed on all tests. You will have at least 20 seconds to solve each puzzle, regardless of difficulty. Executions taking longer than that will be unceremoniously terminated, leading to that test being assessed as failing. Twenty seconds should be ample with careful implementation.

See the Project Coding Guidelines on the LMS for detailed suggestions for coding style. These guidelines will form the basis of the quality assessment of your code and documentation.

Submission

You **must** submit your project from either of the unix servers `dimefox.eng.unimelb.edu.au`, `nutmeg.eng.unimelb.edu.au`, or `nutmeg2.eng.unimelb.edu.au`. As `nutmeg2` is a faster machine than `dimefox` and `nutmeg`, we recommend that you test and submit your project on `nutmeg2`. Make sure the version of your program source files you wish to submit is on this host, then `cd` to the directory holding your source code and issue the command:

```
submit COMP90048 proj2 proj2.pl
```

If your code spans multiple source files, add the extra ones to the end of that command line.

Important: you must wait a minute or two (or more if the servers are busy) after submitting, and then issue the command

```
verify COMP90048 proj2 | less
```

This will show you the test results from your submission, as well as the file(s) you submitted. If the test results show any problems, correct them and submit again. You may submit as often as you like; only your final submission will be assessed.

If you wish to (re-)submit after the project deadline, you may do so by adding `“.late”` to the end of the project name (*i.e.*, `proj2.late`) in the `submit` and `verify` commands. But note that a penalty, described below, will apply to late submissions, so you should weigh the points you will lose for a late submission against the points you expect to gain by revising your program and submitting again.

It is your responsibility to verify your submission.

Your submission will be tested on the server `nutmeg2.eng.unimelb.edu.au` (note the “2”). We will run your submission using SWI Prolog version 7.6.4, which is probably older than the version you will develop on. On `nutmeg2`, this version of SWI Prolog can be accessed via the command:

```
/usr/local/lib/swipl-7.6.4/bin/x86_64-linux/swipl
```

You are advised to test your program on this server before submitting; in the unlikely case that your program uses some Prolog features or libraries not supported by SWI Prolog 7.6.4, it will be much easier to discover this.

Note that these hosts are only available through the university’s network. If you wish to use these machines from off campus, you will need to use the university’s Virtual Private Network. The LMS Resources list gives instructions.

Windows users should see the LMS Resources list for instructions for downloading the (free) MobaXterm or Putty and Winscp programs to allow you to use and copy files to the department servers from windows computers. Mac OS X and Linux users can use the `ssh`, `scp`, and `sftp` programs that come with your operating system.

Late Penalties

Late submissions will incur a penalty of 0.5% of the possible value of that submission per hour late, including evening and weekend hours. This means that a perfect project that is much more than 4 days late will receive less than half the marks for the project. If you have a medical or similar compelling reason for being late, you should contact the lecturer as early as possible to ask for an extension (preferably before the due date).

Note Well:

This project is part of your final assessment, so cheating is not acceptable. Any form of material exchange between teams, whether written, electronic or any other medium, is considered cheating, and so is the soliciting of help from electronic newsgroups. Providing undue assistance is considered as serious as receiving it, and in the case of similarities that indicate exchange of more than basic ideas, formal disciplinary action will be taken for all involved parties. If you have questions regarding these rules, please ask the lecturer.