Computer Systems 1 Lecture 16

Nested Conditionals

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Topics

- Blocks
 - Syntax for blocks
 - Enter at beginning, exit at end
- Nested conditionals
- Case and jump tables
- 4 Real world incident: the SSL Bug

Blocks

- You'll need to learn many programming languages, eventually
- There are concepts that appear in most languages
- It's useful to focus on the general concept
- One important concept is a block
- A block or compound statement is a single statement that contains several statements
- The purpose is to let you have a group of statements in a loop, or controlled by a conditional
- The syntax is the detailed punctuation used to indicate a block, and this varies in different languages
 - ▶ There are lots of small syntax differences between languages
 - ▶ Some languages use := to mean assign, = to mean equals
 - ▶ Some languages use = to mean assign, == to mean equals

Python syntax style for blocks: layout

The layout (the indentation) determines what is inside the if statement

```
a = 1
if x < y:
b = 2
c = 3
d = 4
```

If you change the indentation, you change the meaning of the program

In Python, you write: to mean then, and else to mean else

Algol/Pascal style for blocks: begin—end

You write then to mean then, and else to mean else

Statements must be separated by semicolon;

C style for blocks: braces

```
a = 1;
if (x < y) b = 2;
else
    { d = 4;
        e = 5; }
f = 6;</pre>
```

In C you don't write then at all, but this means the condition $\mathbf{x} < \mathbf{y}$ must be enclosed in parentheses so the compiler can tell where the condition ends and the then-statement begins

Block structured style for blocks: matching keywords

```
a := 1;
if x<v
  then b := 2;
       while i<n do
           sum := sum + x[i]:
           i := i+1
       endwhile
  else d := 4
       e := 5
endif;
f := 6;
```

This style introduces a lot of keywords (endif, endwhile, endfor, endrepeat) but you don't need braces around a block. It makes code more readable and enables the compiler to produce better error messages

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Enter at beginning, exit at end

- A common programming style is to require
 - Each block enters only at the beginning of the block
 - ► Each block exits only at the end of the block
- This style is helpful in some programming languages, but in some languages it sometimes makes code less readable
- In assembly language, and for compilation patterns, it is necessary to follow this style
- In high level languages this style is sometimes helpful, but not always

Single entrance/exit for compilation patterns

- It is straightforward to translate high level control constructs using the compilation patterns
- These require that the blocks of code always start at the beginning and finish at the end
- It's bad to jump into the middle of a block, or exit out of the middle because
 - The compilation patterns won't work correctly
 - You'll have to duplicate a lot of code
 - ► Example: returning from a procedure requires restoring the registers, resetting the stack pointer, loading the return address
 - ▶ That code should not be duplicated in several places in a procedure

Systematic approach to programming?

- Start by understanding what your program should do
- Express the algorithm using high level language notation (and it's ok to mix in some English too)
- Translate the high level to the low level
 - Assignment statements: x := expression
 - ▶ I/O statements: Write string
 - goto L
 - if boolean then goto L
- Translate the low level to assembly language
- Retain the high and low level code as comments
- Do hand execution at every level

Why use this systematic approach to programming?

- It enables you to write correct code at the outset, and minimise debugging
- If there is a bug, it helps you to catch it early (e.g. in translation to goto form)
- If there's a bug in an instruction, the comments enable you to find it quickly
 - ▶ A common error is to use a wrong register number: add R9,R3,R4
 - ▶ Poor comments don't help: add R9,R3,R4 ; R9 := R3+R4
 - ► Good comments help a lot: add R9,R3,R4; x := alpha + (a[i]*b[i])
 - ► Look at the register usage comments (oops, x is in R8, not R9, now I know how to fix it and I don't have to read the entire program)
- Professional software needs to be maintained; the comments make the software easier to read and more valuable

Nested if-then-else

Conditional statements can be nested deeply

```
if b1
    then S1
        if b2
        then S3
        else S4
        S5
    else S6
        if b3
            then S7
        S8
```

Special case of nested if-then-else

Often the nesting isn't random, but has this specific structure:

```
if b1
  then S1
  else if b2
         then S2
         else if b3
                 then S3
                 else if b4
                        then S4
                        else if b5
                                then S5
                                else S6
```

(Some languages require punctuation to avoid ambiguity)

Another way to write it

To avoid running off the right side of the window, it's usually indented like this:

```
if b1
  then S1
else if b2
  then S2
else if b3
  then S3
else if b4
  then S4
else if b5
  then S5
else S6
```

Some programming languages have elsif or elif

```
if b1 then S1
elif b2 then S2
elif b3 then S3
elif b4 then S4
eleif b5 then S5
else S6
```

- It avoids ambiguity
- It signals to the compiler and to a human programmer that this specific construct is being used
- It allows good indentation layout without violating the basic principle of indentation
- Some languages have this, some don't

A common application: numeric code

Nested if-then-else but the boolean conditions are not arbitrary: they are checking the value of a code:

```
if code = 0
  then S1
else if code = 1
  then S2
else if code = 2
  then S3
else if code = 3
  then S4
else if code = 4
  then S5
```

The case statement

```
case n of
    0 -> Stmt
    1 -> Stmt
    2 -> Stmt
    3 -> Stmt
    4 -> Stmt
    5 -> Stmt
else -> Stmt // handle error
```

This means: execute the statement corresponding to the value of n

Many programming languages have this; the syntax varies a lot but that isn't what's important

Example: numeric code specifies a command

```
The input data is an array of records, each specifying an operation
   Command : record
     code : Int
                    ; specify which operation to perform
          : Int
                    ; index into array of lists
     i
          : Int
                    : value of list element
      X
The meaning of a command depends on the code:
   0 terminate the program
     insert x into set[i]
   2 delete x from set[i]
   3 return 1 if set[i] contains x, otherwise 0
   4 print the elements of set[i]
```

Selecting the command with a case statement

```
Initialize
   BuildHeap ()
Execute the commands in the input data
   finished := 0
   while InputPtr <= InputEnd && not finished
      CurrentCode := (*InputPtr).code
      p := set[*InputPtr] ; linked list
      x := (*InputPtr).x ; value to insert, delete, search
      case CurrentCode of
         0 : <CmdTerminate>
         1 : <CmdInsert>
         2 : <CmdDelete>
         3 : <CmdSearch>
         4 : <CmdPrint>
         else : <>
      InputPtr := InputPtr + sizeof(Command)
Terminate the program
   halt.
```

Finding a numeric code

- It's tedious and inefficient to go through the possible values of a numeric code in sequence
 - If you're looking up Dr Zhivago in the phone book, do you look at Arnold Aardvark, and Anne Anderson, on on and on?
 - You go straight to the end of the book
- We want to find the statement corresponding to a numeric code directly, without checking all the other values

A problem with efficiency

- The problem
 - There are many applications of case statements
 - ► They are executed frequently
 - ▶ The number of cases can be large (not just 4 or 5; can be hundreds)
 - ► The implementation of the if-then-else requires a separate compare and jump for each condition
- The solution
 - A technique called jump tables

Jump tables: the basic idea

- For each target statement (S1, S2, S3, etc) in the conditional, introduce a jump to it: jump S1[R0], jump S2[R0], etc
- Make an array "jt" of these jump instructions

```
jt[0] = jump S0[R0]
jt[1] = jump S1[R0]
jt[2] = jump S2[R0]
jt[3] = jump S3[R0]
jt[4] = jump S4[R0]
```

- Given the *code*, Jump to jt[code]
- Each element of the array is an instruction that requires two words
- So jump to &jt $+ 2 \times code$

Jump table

```
Jump to operation specified by code
    add
           R4.R4.R4
                               : code := 2*code
           R5,CmdJumpTable[R0]; R5 := pointer to jump table
    lea
    add
           R4.R5.R4
                               ; address of jump to operation
    jump
           0 [R4]
                               ; jump to jump to operation
CmdJumpTable
    jump
           CmdTerminate[RO]
                               ; code 0: terminate the program
           CmdInsert[R0]
                               : code 1: insert
    jump
          CmdDelete[R0]
                               : code 2: delete
    jump
           CmdSearch[RO]
                               : code 3: search
    jump
           CmdPrint[R0]
    jump
                               ; code 4: print
CmdDone
    load
           R5, InputPtr[R0]
           R6,3[R0]
   lea
    bha
           R5,R5,R6
           R5, InputPtr[R0]
    store
           CommandLoop[R0]
    jump
```

We have to be careful!

- What if code is negative, or larger than the number of cases?
- The jump to the jump table could go anywhere!
- It might not even go to an instruction
- But whatever is in memory at the place it goes to, will be interpreted as an instruction that will be executed
- The program will go haywire
- Debugging it will be hard: the only way to catch the error is to read the code and/or to single step
- Solution: before jumping into the jump table, check to see if code is invalid (too big or too small)

Checking whether the code is invalid

```
CommandLoop
    load
           R5, InputPtr [R0]
                                ; R1 := InputPtr
           R4,0[R5]
    load
                                ; R4 := *InputPtr.code
: Check for invalid code
           R4,R0
                                ; compare (*InputPtr).code, 0
    cmp
    jumplt CmdDone[R0]
                                  skip invalid code (negative)
           R6,4[R0]
    lea
                                : maximum valid code
                                  compare code with max valid code
           R4,R6
    cmp
    jumpgt CmdDone[R0]
                                ; skip invalid code (too large)
. . .
CmdDone
    load
           R5, InputPtr[R0]
           R6,3[R0]
    lea
    add
           R5,R5,R6
           R5, InputPtr[R0]
    store
           CommandLoop[R0]
```

qmuj

A real world incident!

The moral of the following true story: be meticulous!

- On 21 February 2014, Apple announced a bug in core security software
- It affected iPhone, iPad, Macintosh, Safari
- The bug is serious: it skips a crucial security check in SSL
- A security update was made available quickly
 - That's why you should always apply your security updates!
- It was fixed quickly but we can learn some important lessons from it

What is SSL?

- Secure Sockets Layer
- It arranges that communication between your device and a server is encrypted
- When you see a "padlock icon" on a browser, it means the SSL protocol is being used

Source file: sslKeyExchange.c

- 1,969 lines of code in C
- You can see the code—click on this link, or type it in:
- http://opensource.apple.com/source/Security/
 Security-55471/libsecurity_ssl/lib/sslKeyExchange.c
- The bug made it possible to break open "secure" Internet traffic
- It appears somewhere on the next three slides: can you spot it?

The function SSLVerifySignedServerKeyExchange (1)

```
static OSStatus
SSLVerifvSignedServerKevExchange(SSLContext *ctx, bool isRsa, SSLBuffer signedParams,
                                 uint8 t *signature, UInt16 signatureLen)
{
    OSStatus
                    err;
    SSLBuffer
                    hashOut, hashCtx, clientRandom, serverRandom;
                    hashes[SSL_SHA1_DIGEST_LEN + SSL_MD5_DIGEST_LEN];
   uint8_t
   SSLBuffer
                    signedHashes;
   uint8 t
                                *dataToSign:
                                dataToSignLen;
        size t
        signedHashes.data = 0:
   hashCtx.data = 0:
    clientRandom.data = ctx->clientRandom:
    clientRandom.length = SSL CLIENT SRVR RAND SIZE:
    serverRandom.data = ctx->serverRandom;
    serverRandom.length = SSL_CLIENT_SRVR_RAND_SIZE;
```

The function continued (2)

```
if(isRsa) {
            /* skip this if signing with DSA */
            dataToSign = hashes:
            dataToSignLen = SSL_SHA1_DIGEST_LEN + SSL_MD5_DIGEST_LEN;
            hashOut.data = hashes;
            hashOut.length = SSL MD5 DIGEST LEN:
            if ((err = ReadyHash(&SSLHashMD5, &hashCtx)) != 0)
                    goto fail:
            if ((err = SSLHashMD5.update(&hashCtx, &clientRandom)) != 0)
                    goto fail;
            if ((err = SSLHashMD5.update(&hashCtx, &serverRandom)) != 0)
                    goto fail;
            if ((err = SSLHashMD5.update(&hashCtx, &signedParams)) != 0)
                    goto fail;
            if ((err = SSLHashMD5.final(&hashCtx, &hashOut)) != 0)
                    goto fail;
    else {
            /* DSA, ECDSA - just use the SHA1 hash */
            dataToSign = &hashes[SSL_MD5_DIGEST_LEN];
            dataToSignLen = SSL SHA1 DIGEST LEN:
    hashOut.data = hashes + SSL_MD5_DIGEST_LEN;
hashOut.length = SSL SHA1 DIGEST LEN:
```

The function continued (3)

```
if ((err = SSLFreeBuffer(&hashCtx)) != 0)
        goto fail:
    if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
        goto fail:
    if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
        goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
        goto fail:
    if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
        goto fail:
        goto fail:
    if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
        goto fail;
        err = sslRawVerify(ctx,
                       ctx->peerPubKey,
                       dataToSign.
                                                                 /* plaintext */
                       dataToSignLen,
                                                         /* plaintext length */
                       signature,
                       signatureLen):
        if(err) {
                sslErrorLog("SSLDecodeSignedServerKeyExchange: sslRawVerify "
                    "returned %d\n", (int)err);
                goto fail:
        }
fail:
    SSLFreeBuffer(&signedHashes):
   SSLFreeBuffer(&hashCtx);
    return err:
```

A bit of C

```
Syntax of the if-statement
if (condition) statement;
Examples:
if (x < y) n=n+1;
if (n==5)
  {p = 3;}
    q = 4;
    printf ("hello p = %3d\n", p);
```

A closer look

```
if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
   goto fail;
if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
   goto fail;
   goto fail;
if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
   goto fail;
```

The second goto fail is unconditional so the check in the last if statement is never executed

What can we learn from this?

- To be a good programmer you need to
 - Know your programming language
 - Be meticulous
 - Read your code carefully
 - Do thorough testing
- It was very helpful that Apple published the source code for this bug

A mystery: The Attempted Linux Trapdoor Hack of 2003

- A trapdoor is a flaw (bug? hack?) in software that allows a hostile user to gain abilities they should not have
 - It's like giving the key to the safe in a bank to a burglar
- The master copy of Linux was maintained on Bitkeeper (no longer—now git is used)
- A secondary copy was on cvs
- Someone—unknown to this day—made a small edit to an obscure piece of code on the cvs version

The offending code: can you spot the trapdoor?

- A change was made to a function that a user could call to wait for an event
- System functions typically take options, and they need to check to see if the options are valid
- If not, a return code retval is set to a non-zero value

So, here's the code:

```
if ((options == (__WCLONE|__WALL)) && (current->uid = 0))
    retval = -EINVAL;
```

What's going on?

- In C, == is the comparison operator
- is the assignment operator (bad notation, but we're stuck with it)
- This code looks like it is checking:
 - If either the WCLONE or WALL option flag is set and the current user is root, then return an error condition
- What it actually says is
 - Check the irrelevant flags, set the current user to be root, get the irrelevant result 0 and do nothing else

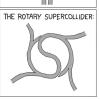
What's the significance of root?

- Many processes are running in the computer
- Each has its own set of privileges
- There is one special user "root" which has all privileges: root can do anything at all
- The effect of this change to the source code (a change of one character! is
 - ▶ If a user program calls this obscure function with just the right set of obscure options, it suddenly gains full control over the machine
- Don't worry this faulty code never made it into the master copy of Linux

HIGHWAY ENGINEER PRANKS:



THE ZERO-CHOICE INTERCHANGE:



https://xkcd.com/253/