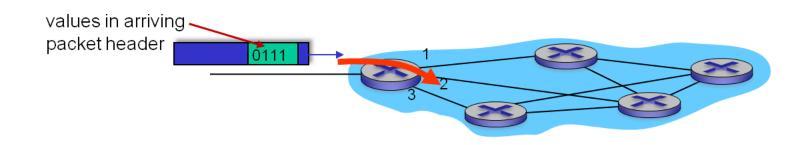
Network layer: "data plane" roadmap

- Network layer: overview
 - data plane
 - control plane
- What's inside a router
 - input ports, switching, output ports
 - buffer management, scheduling
- IP: the Internet Protocol
 - datagram format
 - addressing
 - network address translation
 - IPv6



- Generalized Forwarding, SDN
 - Match+action
 - OpenFlow: match+action in action
- Middleboxes

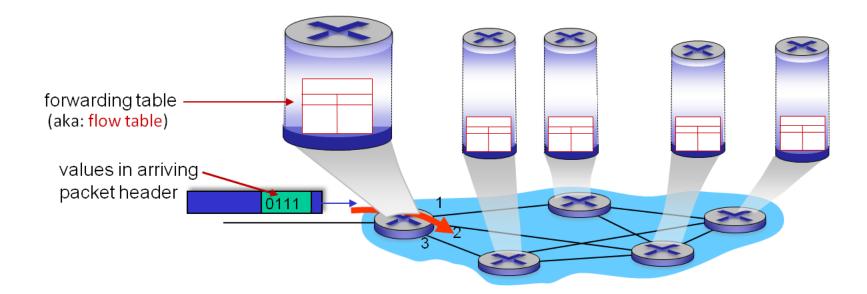
Generalized forwarding: match plus action



Generalized forwarding: match plus action

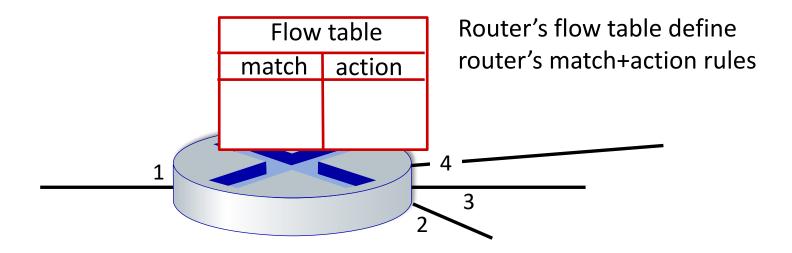
Review: each router contains a forwarding table (aka: flow table)

- "match plus action" abstraction: match bits in arriving packet, take action
 - destination-based forwarding: forward based on dest. IP address
 - generalized forwarding:
 - many header fields can determine action
 - many action possible: drop/copy/modify/log packet



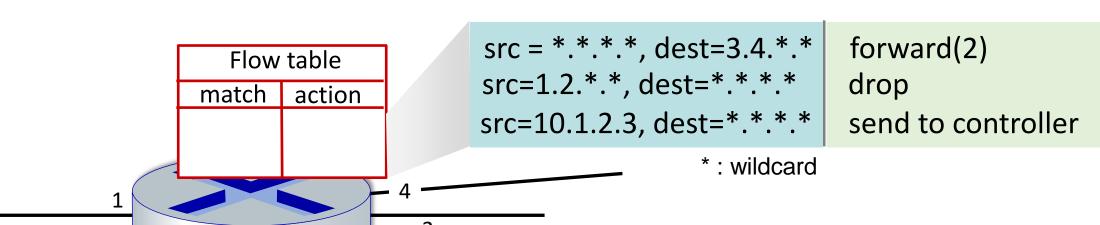
Flow table abstraction

- flow: defined by header fields values (in link-, network-, transport-layer fields)
- generalized forwarding: simple packet-handling rules
 - match: pattern values in packet header fields
 - actions: for matched packet: drop, forward, modify, matched packet or send matched packet to controller
 - priority: disambiguate overlapping patterns
 - counters: #bytes and #packets

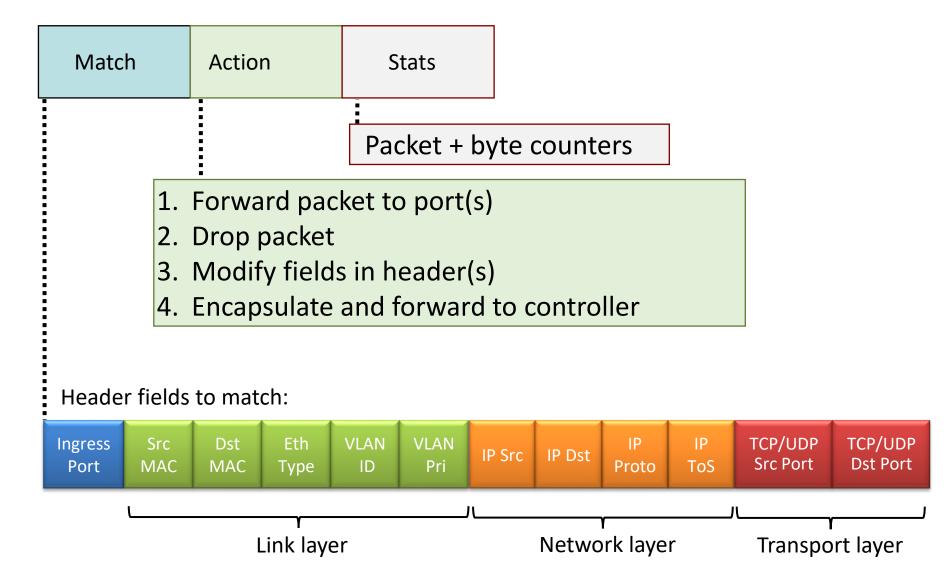


Flow table abstraction

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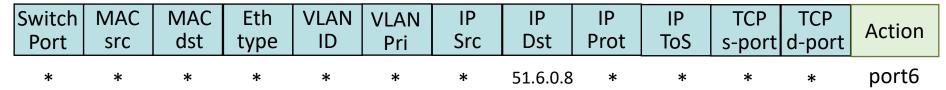


OpenFlow: flow table entries



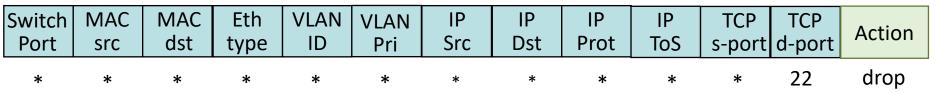
OpenFlow: examples

Destination-based forwarding:

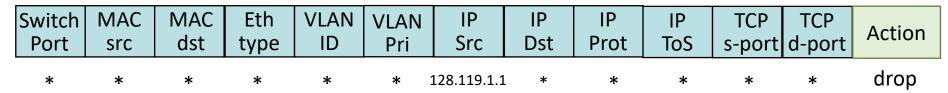


IP datagrams destined to IP address 51.6.0.8 should be forwarded to router output port 6

Firewall:



Block (do not forward) all datagrams destined to TCP port 22 (ssh port #)



Block (do not forward) all datagrams sent by host 128.119.1.1

OpenFlow: examples

Layer 2 destination-based forwarding:

Switch	MAC	MAC	Eth	VLAN	VLAN	IP	IP	IP	IP	TCP	TCP	Action
Port	src	dst	type	ID	Pri	Src	Dst	Prot	ToS	s-port	d-port	
*	*	22:A7:23: 11:E1:02	*	*	*	*	*	*	*	*	*	port3

layer 2 frames with destination MAC address 22:A7:23:11:E1:02 should be forwarded to output port 3

OpenFlow abstraction

match+action: abstraction unifies different kinds of devices

Router

- match: longest destination IP prefix
- action: forward out a link

Switch

- match: destination MAC address
- action: forward or flood

Firewall

- match: IP addresses and TCP/UDP port numbers
- action: permit or deny

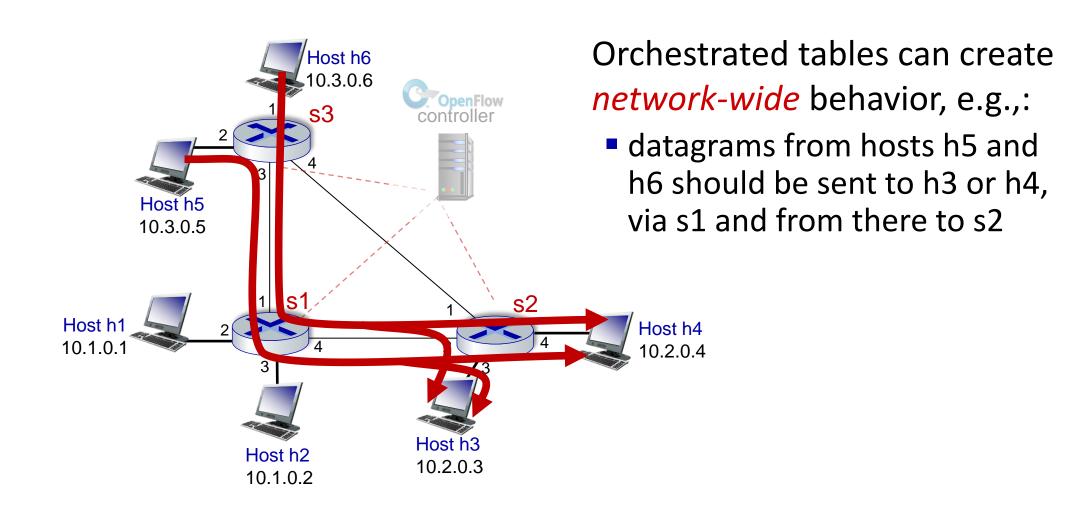
NAT

- match: IP address and port
- action: rewrite address and port

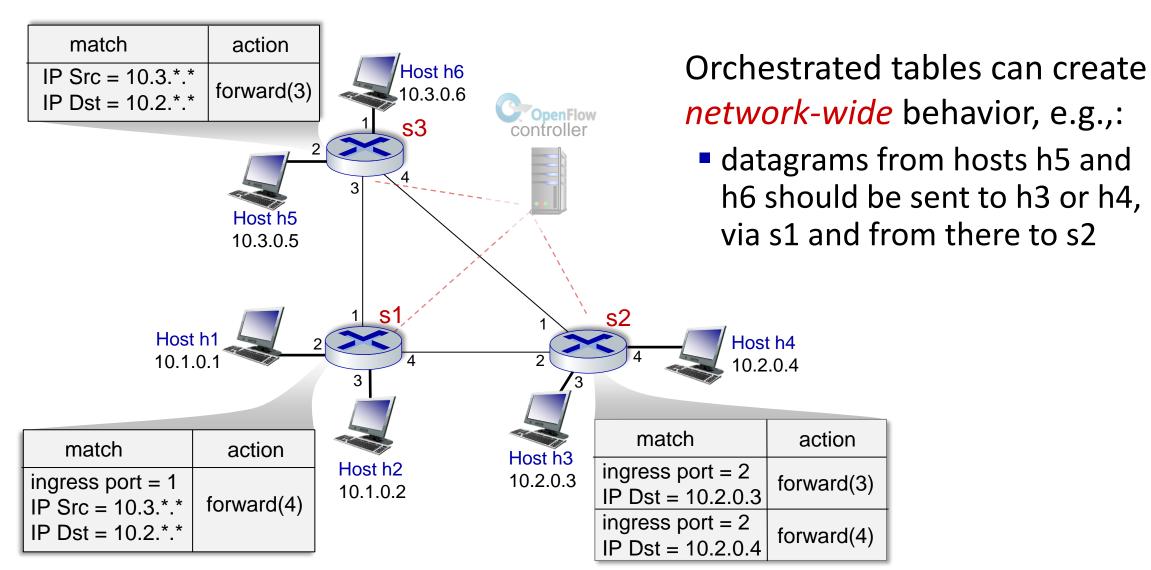
OpenFlow abstraction

- If we can generate the individual forwarding tables in all routers in the network then we can implement any network wide routing we actually want.
- Compute and generate the forwarding tables via program in an SDN controller.
- Hasn't quite been mainstreamed yet!

OpenFlow example



OpenFlow example



Generalized forwarding: summary

- "match plus action" abstraction: match bits in arriving packet header(s) in any layers, take action
 - matching over many fields (link-, network-, transport-layer)
 - local actions: drop, forward, modify, or send matched packet to controller
 - "program" network-wide behaviors

Generalized forwarding: summary

- simple form of "network programmability"
 - programmable, per-packet "processing"
 - historical roots: active networking
 - *today:* more generalized programming: P4 (see p4.org).

Network layer: "data plane" roadmap

- Network layer: overview
- What's inside a router
- IP: the Internet Protocol
- Generalized Forwarding
- Middleboxes
 - middlebox functions
 - evolution, architectural principles of the Internet



Middleboxes

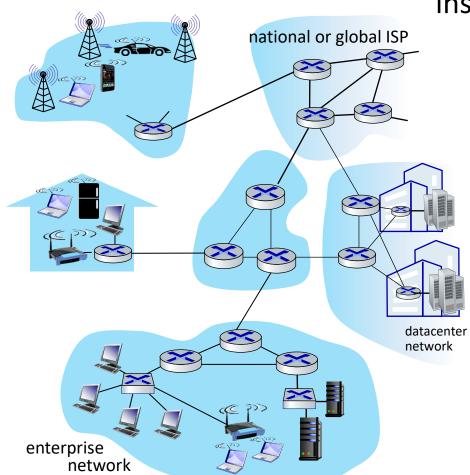
Middlebox (RFC 3234)

"any intermediary box performing functions apart from normal, standard functions of an IP router on the data path between a source host and destination host"

Middleboxes everywhere!

NAT: home, cellular, institutional

Applicationspecific: service
providers,
institutional,
CDN



Firewalls, IDS: corporate, institutional, service providers, ISPs

Load balancers:

corporate, service provider, data center, mobile nets

Caches: service provider, mobile, CDNs

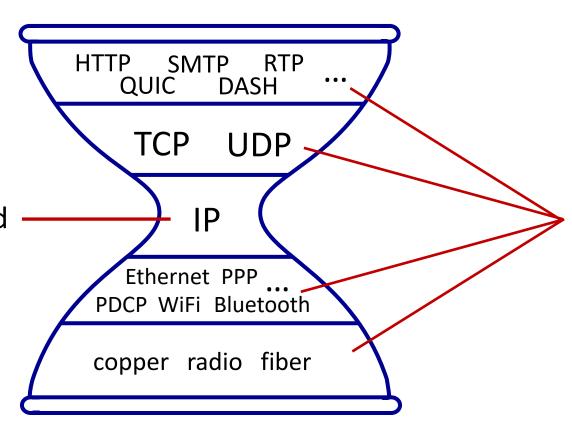
Middleboxes

- initially: proprietary (closed) hardware solutions
- move towards "whitebox" hardware implementing open API
 - move away from proprietary hardware solutions
 - programmable local actions via match+action
 - move towards innovation/differentiation in software
- SDN: (logically) centralized control and configuration management often in private/public cloud
- network functions virtualization (NFV): programmable services over white box networking, computation, storage

The IP hourglass

Internet's "thin waist":

- one network layer protocol: IP
- must be implemented by every (billions) of Internet-connected devices



many protocols in physical, link, transport, and application layers

The IP hourglass, at middle age

RTP HTTP **SMTP** QUIC **DASH** Internet's middle age **TCP** UDP "love handles"? caching VE TAN middleboxes, Firewalls operating inside the Ethernet PPP network PDCP WiFi Bluetooth copper radio fiber

Architectural Principles of the Internet

RFC 1958

"Many members of the Internet community would argue that there is no architecture, but only a tradition, which was not written down for the first 25 years (or at least not by the IAB). However, in very general terms, the community believes that the goal is connectivity, the tool is the Internet

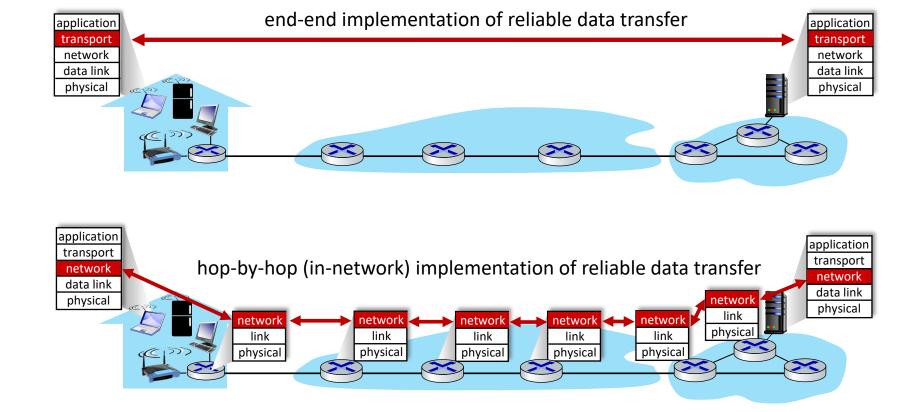
Protocol, and the intelligence is end to end rather than hidden in the network."

Three cornerstone beliefs:

- simple connectivity
- IP protocol: that narrow waist
- intelligence, complexity at network edge

The end-end argument

some network functionality (e.g., reliable data transfer, congestion)
 can be implemented in network, or at network edge



The end-end argument

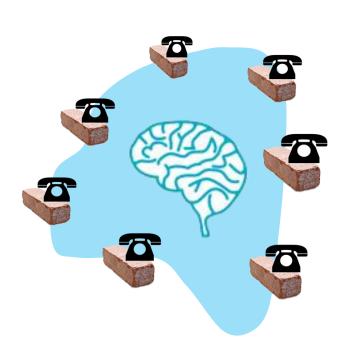
some network functionality (e.g., reliable data transfer, congestion)
 can be implemented in network, or at network edge

"The function in question can completely and correctly be implemented only with the knowledge and help of the application standing at the end points of the communication system. Therefore, providing that questioned function as a feature of the communication system itself is not possible. (Sometimes an incomplete version of the function provided by the communication system may be useful as a performance enhancement.)

We call this line of reasoning against low-level function implementation the "end-to-end argument."

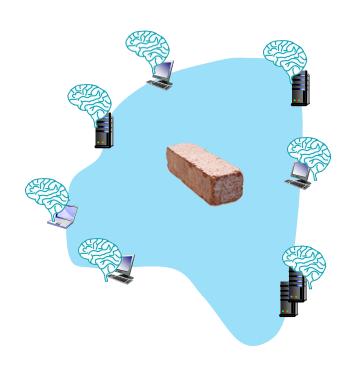
Saltzer, Reed, Clark 1981

Where's the intelligence?



20th century phone net:

intelligence/computing at network switches



Internet (pre-2005)

intelligence, computing at edge



Internet (post-2005)

- programmable network devices
- intelligence, computing, massive application-level infrastructure at edge

Chapter 4: done!

- Network layer: overview
- What's inside a router
- IP: the Internet Protocol
- Generalized Forwarding, SDN
- Middleboxes



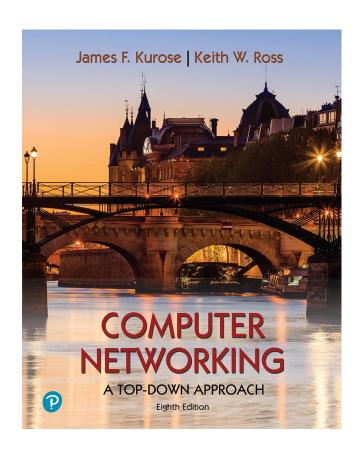
Question: how are forwarding tables (destination-based forwarding) or flow tables (generalized forwarding) computed?

Answer: by the control plane (next chapter)

Important Dates

- 07-12-2024(Friday), Midnight(11:59PM)–
 Homework3 due date
- 07-22-2024 (Monday)- Quiz 2
 From LectureD_11_Chapter_3
 To LectureD_17_Chapter_5
- 07-26-2024 (Friday)- Exam 2
 From LectureD_11_Chapter_3
 To LectureD 19 Chapter 5

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Computer Networking: A Top-Down Approach

8th edition Jim Kurose, Keith Ross Pearson, 2020

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