

# Operating Systems & Computer Networks 1. Introduction & Motivation

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### Roadmap

- 1. Introduction and Motivation
- 2. Interrupts and System Calls
- 3. Processes
- 4. Scheduling
- 5. Memory
- 6. I/O and File System
- 7. Booting, Services, and Security

### Lernziele

#### Sie nennen:

- die Aufgaben und Ziele eines Betriebssystems
- welche realen Ressourcen durch welche virtuellen Ressourcen abgebildet werden

#### Sie nennen und beschreiben:

die Schnittstelle zwischen Betriebssystem und Anwendungsschicht

#### Sie beschreiben:

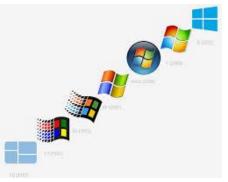
- was unter Protection Rings zu verstehen ist und wie sie mit der Microkernel-vs-Monolithischer-Kernel-Diskussion zusammenhängen
- was unter einem Microkernel BS zu verstehen ist
- was unter einem Monolithischen Kernel BS zu verstehen ist

#### Sie wägen die Vor- und Nachteile ab:

zwischen Microkernel und Monolithischem Kernel

### Motivation







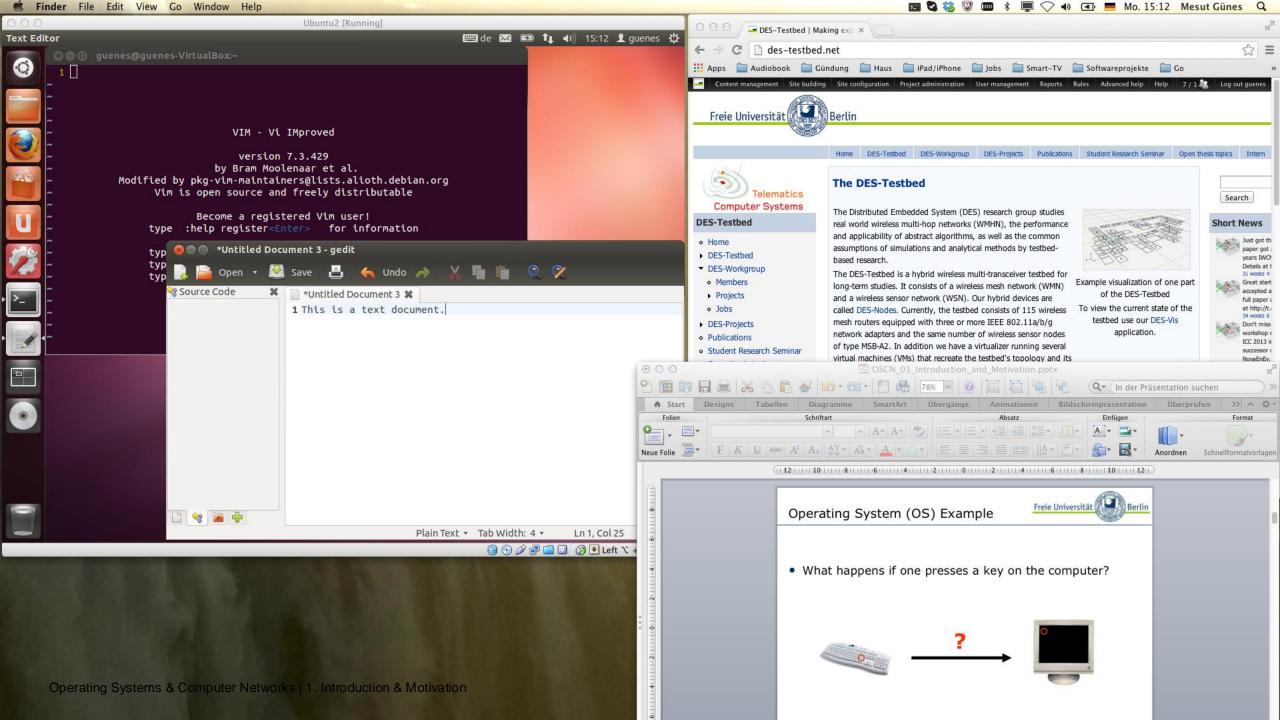












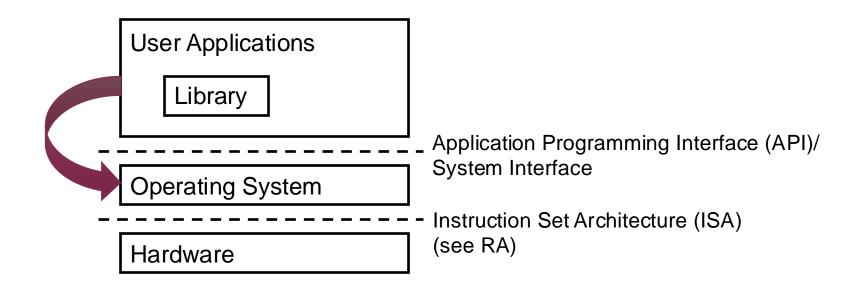
# Layers of Abstraction

User Interface (Shell, GUI, ...) **User Applications** System Interface (system calls, C functions) Operating System / Kernel Hardware Interface (Instruction Set Architecture, I/O Ports, ...) Hardware

# System Interface and System Calls

System interface is the only way for user applications to interact with the operating system.

System interface consists of system calls (supervisor calls) → POSIX.



High-level programming languages hide systems calls in library routines.

### POSIX

Portable Operating System Interface (POSIX)

• E.g. https://standards.ieee.org/standard/1003 1,2013Edition.html

#### **POSIX** defines

- Application programming interface (API)
- Command line shells
- Utilities

**UNIX like Operating Systems** 

POSIX oriented operating systems

- Unix
- Linux
- Windows
- Mac OS X
- ...

# Tasks of an Operating System

Typical services of a **general** purpose OS includes:

- Program execution
- Access to I/O-devices
  - Hardware abstraction
- Controlled access to files
  - Non-volatile memory
- Access control
  - Security / user management
- Error detection and error handling
  - Both hardware and software
- Logging

Special purpose operating systems focus on different services, e.g., real-time or communication requirements.

See: Stallings, W. (2017). Operating Systems. (9th ed.). Pearson International, chapter 2.1, p. 70

# Goals of an Operating System

#### Convenience

→ ease of use for users and programmers

#### **Efficiency**

→ when managing limited resources

#### **Ability to evolve**

- → New hardware standards
- → Changing user requirements

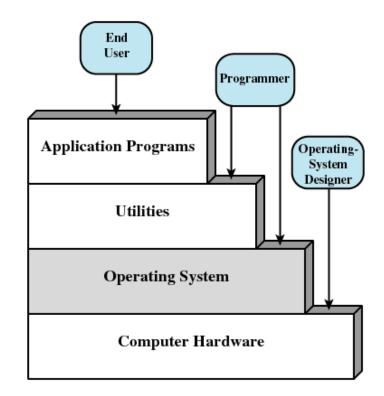


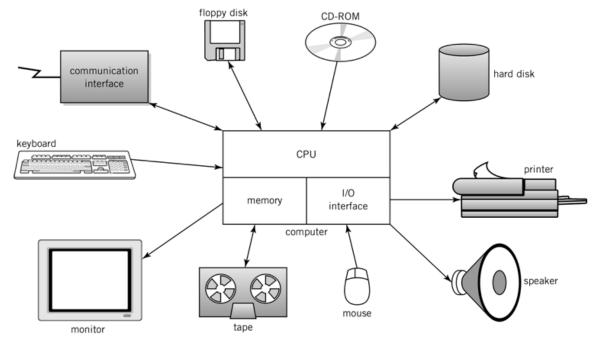
Figure 2.1 Layers and Views of a Computer System

See: Stallings, W. (2017). Operating Systems. (9th ed.). Pearson International, chapter 2.1, p. 69

# Managing Resources

Hardware provides the basic computing resources such as

- Processor(s)
- Memory
- Persistent storage
- Network connection



Englander: The Architecture of Computer Hardware and Systems Software, 2nd edition Chapter 1, Figure 01-06

OS virtualizes resources to permit controlled sharing and isolation

virtual instances of a resource are created
OS provides virtual resources for user applications

### Virtual Resources

#### Virtual resources and corresponding real resources:

Processes processor(s)

Virtual Memory main memory

Files persistent memory

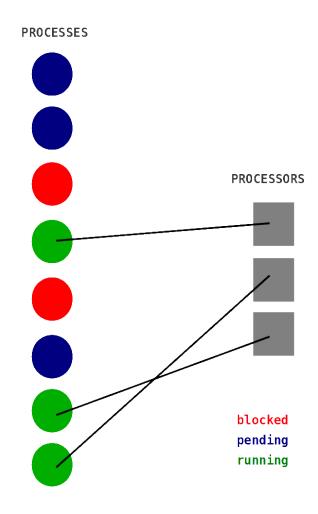
Ports network adapter

#### Advantages:

- Easy to use through procedural interface (system calls)
- Secure against hardware and software errors or manipulation



### Processes



Number of processes is not limited by the number of processors: Multitasking

Processor is used efficiently:

Time is not wasted by processes that are waiting on I/O devices

Reduced latency (=response time)

Different process states, e.g.,

- running executing
- pending ready to execute
- blocked not ready to execute

# Virtual Memory

Managed by the Memory Management Unit (MMU)

#### Transportability:

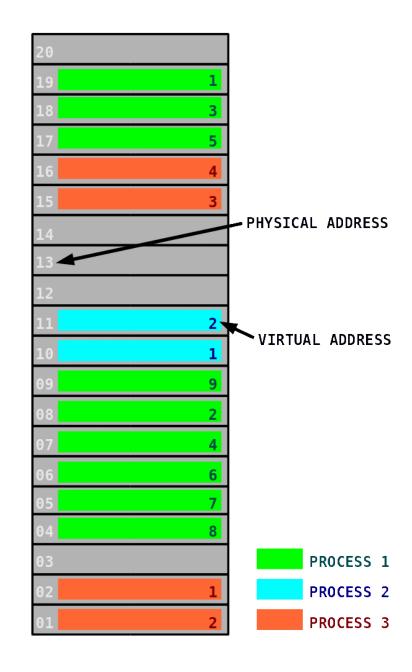
 position independent code – program does not depend on memory architecture

#### Security:

 memory access is restricted to memory units "owned" by a process

#### Efficiency:

external fragmentation is avoided



### **Files**

Managed by a file system

Persistent objects for long-term data storage

Stored in secondary memory (e.g., tape, hard disk, USB flash drive)

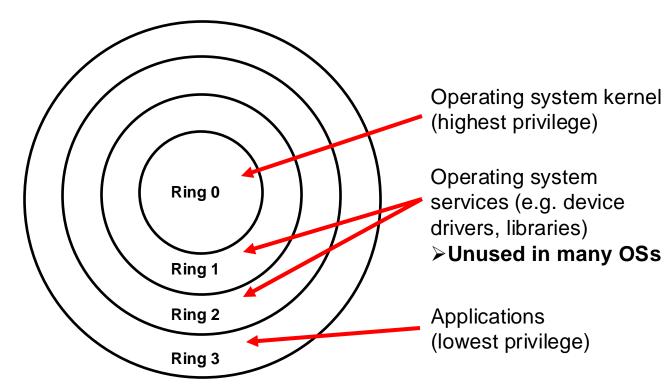
Similar to virtual memory - file name instead of virtual address



### Protection Rings

Hardware provides hierarchical privilege levels

- Inner rings have access to outer rings' resources
- Outer rings may access inner rings through predefined gateways



# Operating System Kernel / Ring 0

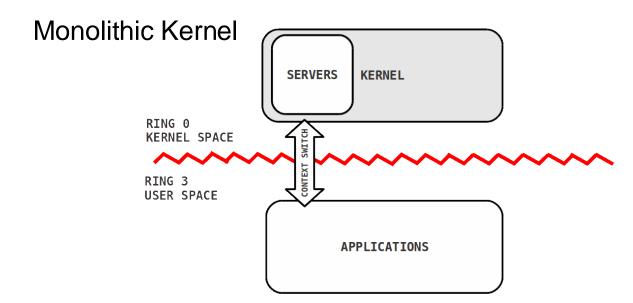
Kernel implements basic layer of abstraction

Runs with full access to hardware (Ring 0)

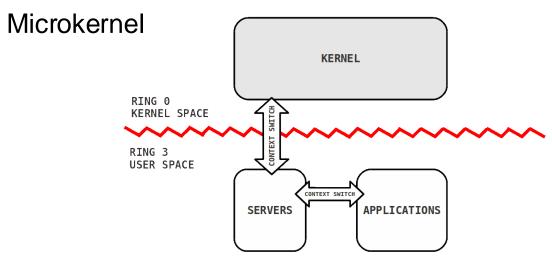
Context Switch: switching from one process to another

 A certain amount of time is required for doing the administration, e.g., saving and loading registers.

### Monolithic versus Microkernel

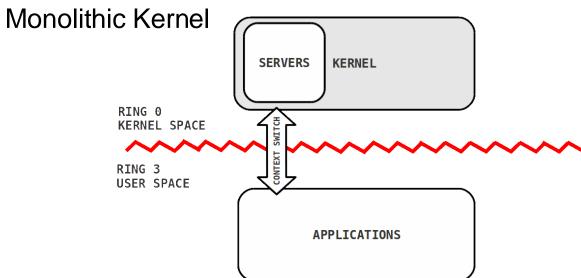


- Implemented as single process
- Sharing same kernel address space



- Very few dedicated functions in kernel:
  - address space management
  - interprocess communication (IPC)
  - basic scheduling
- All other OS services are processes in user space

### Monolithic versus Microkernel

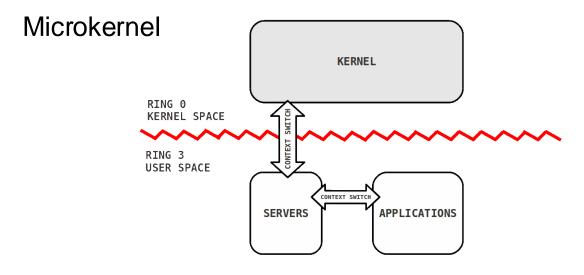


#### Pro:

- less context switches
- no expensive communication

#### Contra:

complications when exchanging functionality



#### Pro:

- strict interfaces
- less complexity, clear structure

#### Contra:

- speed
- synchronization

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