

Data Structures

Lecture 17 & 18: Heaps

- Binary heaps, ADT, operations, construction, Heap sort

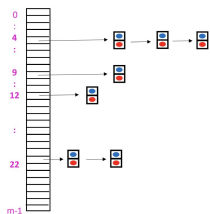
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Recap

Collision Resolution Techniques:

- Separate chaining
- Open addressing
 - Linear probing
 - Quadratic probing
 - Double hashing



```

hash ( 89, 10 ) = 9
hash ( 18, 10 ) = 8
hash ( 49, 10 ) = 9
hash ( 58, 10 ) = 8
hash ( 9, 10 ) = 9

```

After insert 89 After insert 18 After insert 49 After insert 58 After insert 9

0			49	49	49
1				58	58
2					9
3					
4					
5					
6					
7					
8		18	18	18	18
9	89	89	89	89	89

```

hash ( 89, 10 ) = 9
hash ( 18, 10 ) = 8
hash ( 49, 10 ) = 9
hash ( 58, 10 ) = 8
hash ( 9, 10 ) = 9

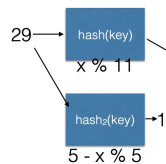
```

After insert 89 After insert 18 After insert 49 After insert 58 After insert 9

0			49	49	49
1				58	58
2					9
3					
4					
5					
6					
7					
8		18	18	18	18
9	89	89	89	89	89

$f(i) = i \cdot \text{hash}_2(x)$ Compute a second hash function to determine a linear offset for this key.

$f(1) = 1 \cdot \text{hash}_2(x) = 1$
 $f(2) = 2 \cdot \text{hash}_2(x) = 2$

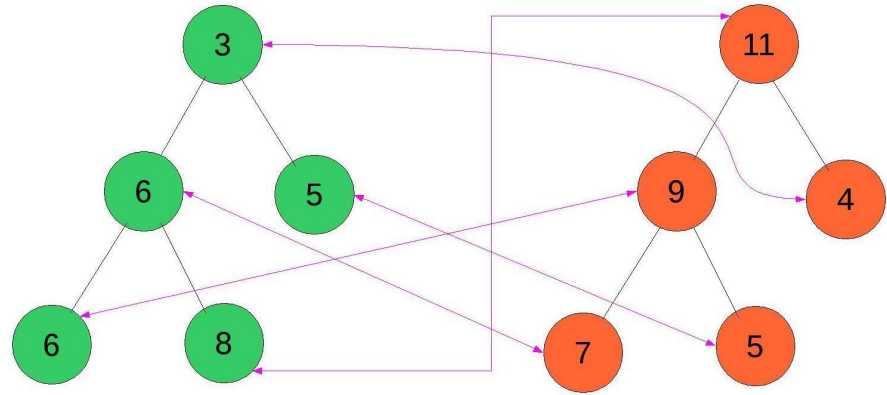


0	
1	
2	
3	
4	
5	
6	
7	40
8	84
9	29
10	62



Onto...

Heaps



Min Heap

Max Heap



Recap | Priority Queues

- **Min Priority Queue**
- **Max Priority Queue**



Recap | Operations in a Priority Queue

- **insert(DataType item, int pri)**
Add an item to the queue which has priority pri
- **DataType peekMax()** or **DataType peekMin()**
Peek at the item in the queue with the highest/lowest priority
- **DataType extractMax()** or **DataType extractMin()**
Remove and return the item in the queue with the highest priority



Heaps...

Is a binary tree satisfying 2 properties:

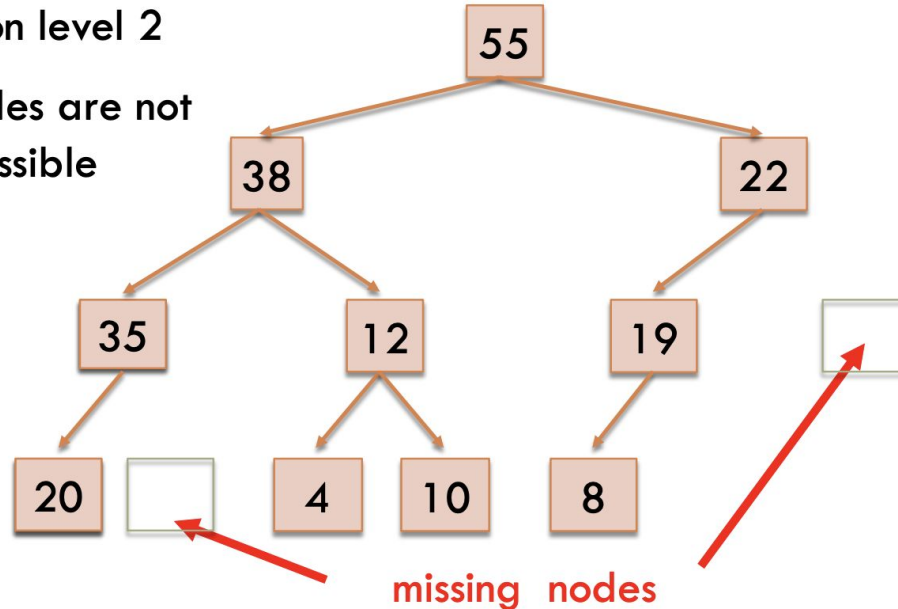
- 1) **Completeness.** Every level of the tree (except last) is completely filled, and on last level nodes are as far left as possible.

Do not confuse with **heap memory** – different use of the word **heap**.

Heaps : Identifying a heap

Not a heap because:

- missing a node on level 2
- bottom level nodes are not as far left as possible





Heaps : Properties

Is a binary tree satisfying 2 properties:

1) Completeness. Every level of the tree (except last) is completely filled, and on last level nodes are as far left as possible.

2) Heap-order.

“max on top”

Max-Heap: every element in tree is \leq its parent

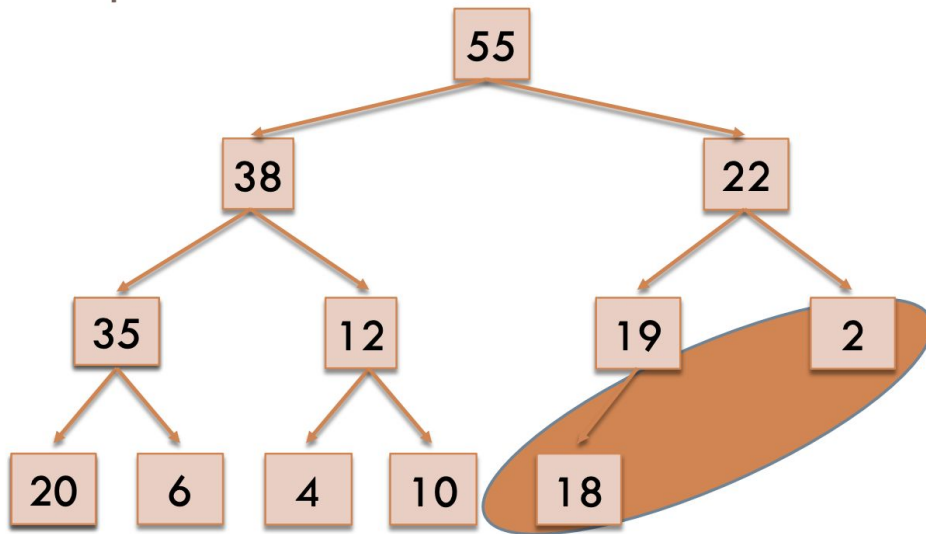
Min-Heap: every element in tree is \geq its parent

“min on top”



Heaps: Properties

Every element is \leq its parent



Note: Bigger elements
can be deeper in the tree!



Heaps: ADT implementation

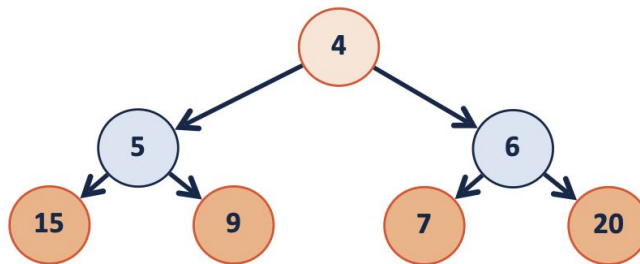
By storing as a complete tree, can avoid using pointers at all!

Can index from 0 or 1 (we will index from 1 in slides)

`leftChild(i) : 2i`

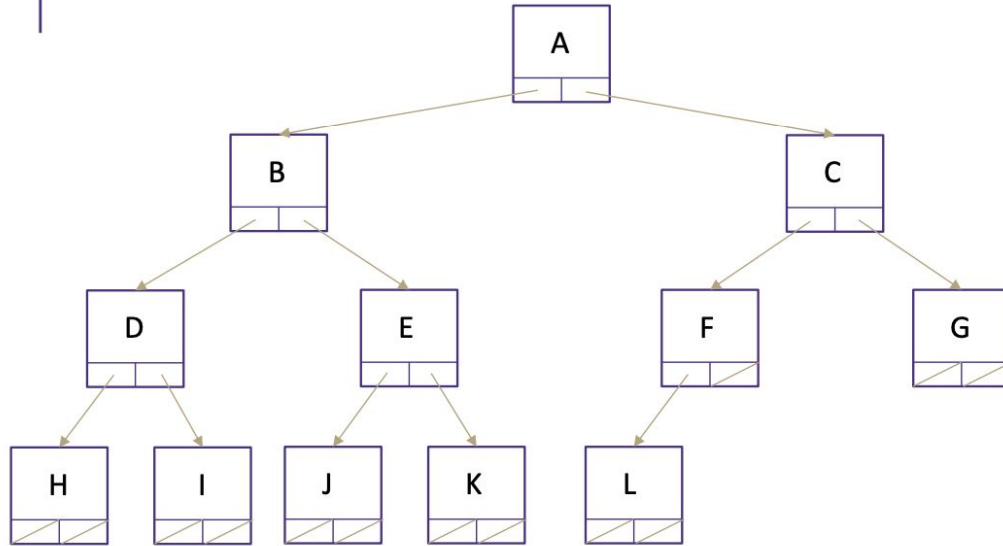
`rightChild(i) : 2i+1`

`parent(i) : floor(i/2)`



	4	5	6	15	9	7	20
0	1	2	3	4	5	6	7

Heaps: Array implementation



Fill array in **level-order** from left to right

0	1	2	3	4	5	6	7	8	9	10	11	12	13
A	B	C	D	E	F	G	H	I	J	K	L		

How do we find the minimum node?

$$\text{peekMin}() = \text{arr}[0]$$

How do we find the last node?

$$\text{lastNode}() = \text{arr}[\text{size} - 1]$$

How do we find the next open space?

$$\text{openSpace}() = \text{arr}[\text{size}]$$

How do we find a node's left child?

$$\text{leftChild}(i) = 2i + 1$$

How do we find a node's right child?

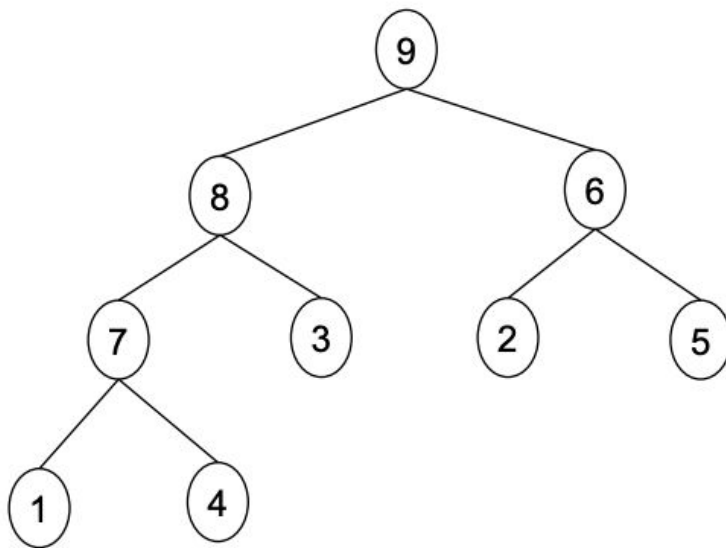
$$\text{rightChild}(i) = 2i + 2$$

How do we find a node's parent?

$$\text{parent}(i) = \frac{(i - 1)}{2}$$



Max Heap



Max Priority Queue ADT

State

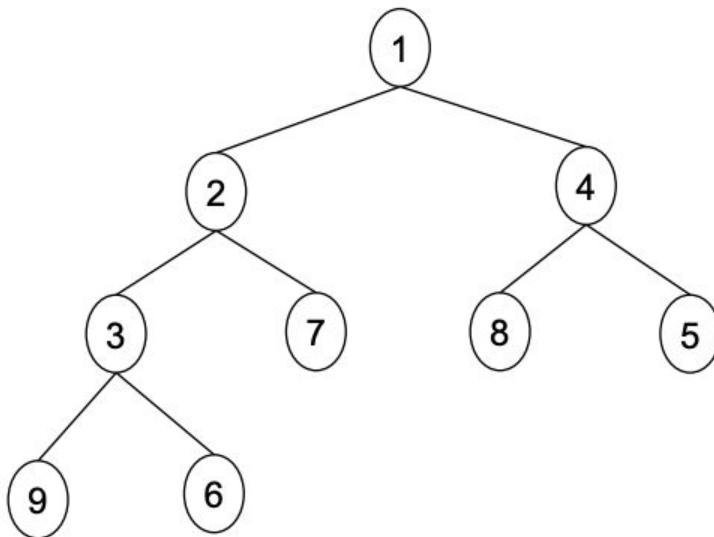
Set of comparable values ordered based on 'priority'

Behaviour

- **extractMax()** - returns the element with the largest priority, removes it from the collection
- **peekMax()** - find, but do not remove the element with the largest priority
- **insert(value)** - add a new element to the collection



Min heap



Min Priority Queue ADT

State

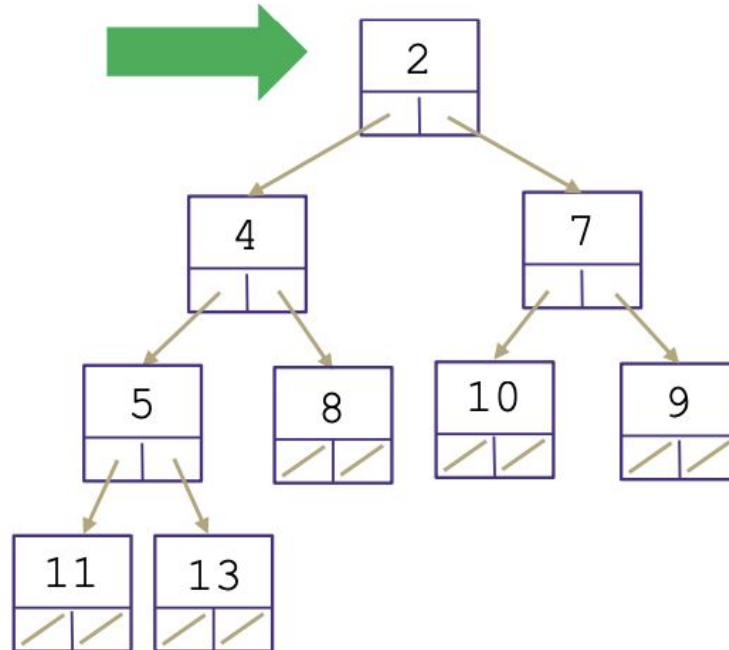
Set of comparable values ordered based on 'priority'

Behaviour

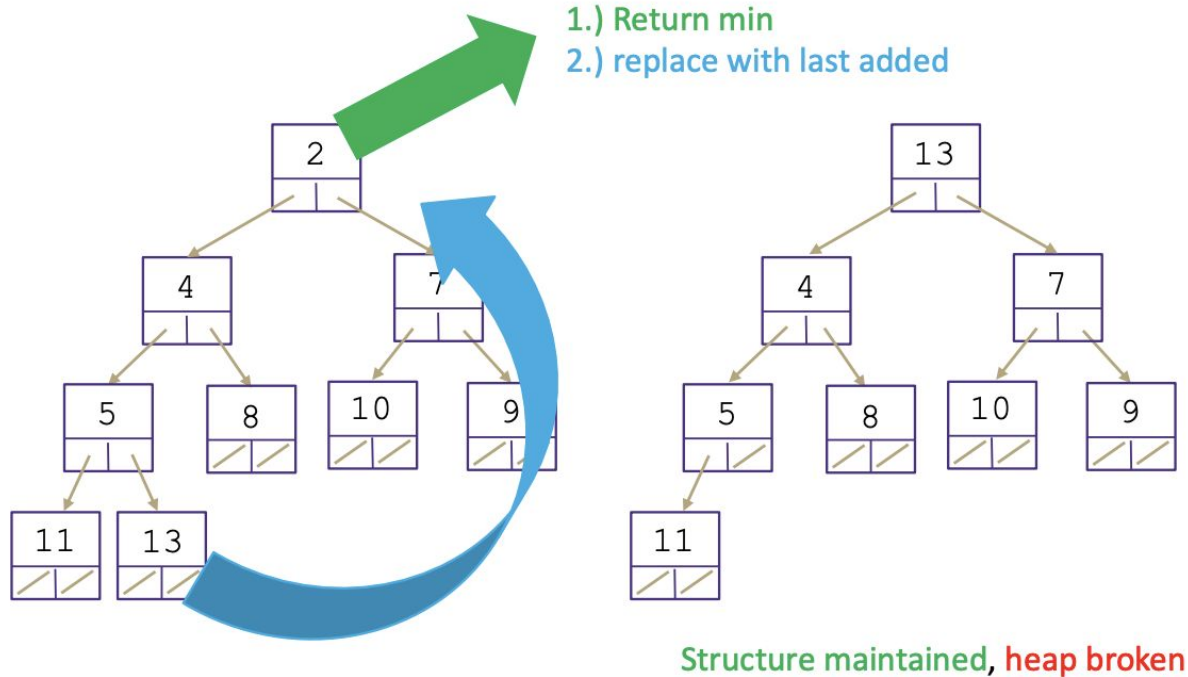
- **extractMin()** - returns the element with the smallest priority, removes it from the collection
- **peekMin()** - find, but do not remove the element with the smallest priority
- **insert(value)** - add a new element to the collection



Heap Operation: PeekMin



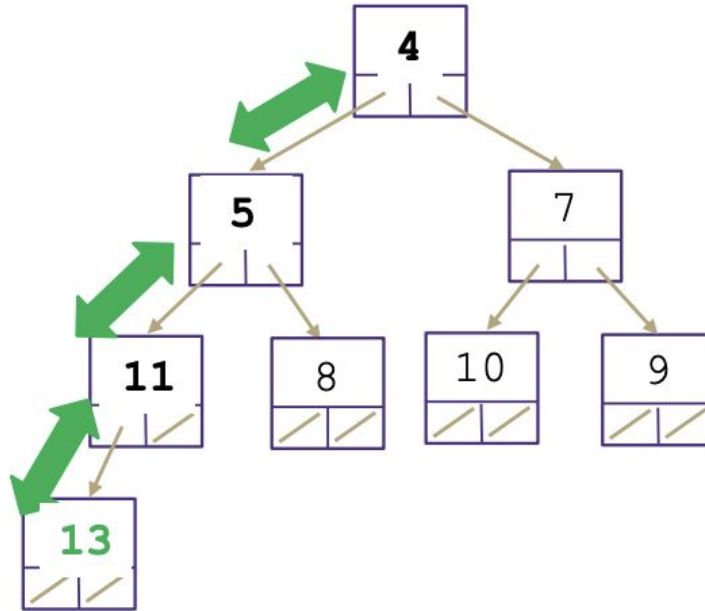
Heap Operation: ExtractMin



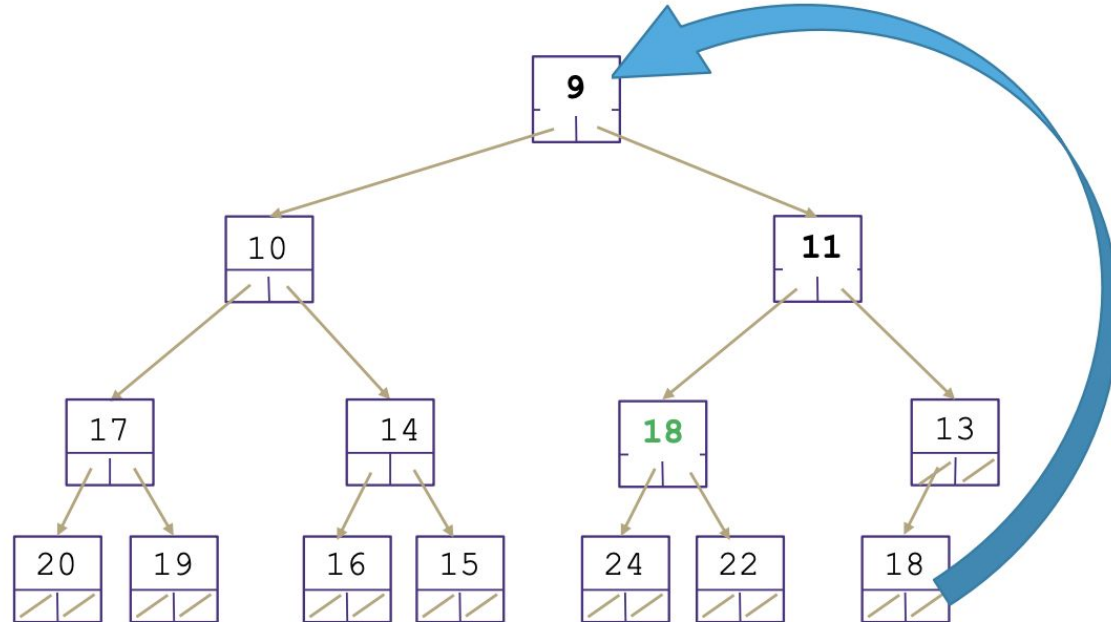
Heap Operation: ExtractMin

3.) percolateDown()

Recursively swap parent with smallest child



Practice: ExtractMin

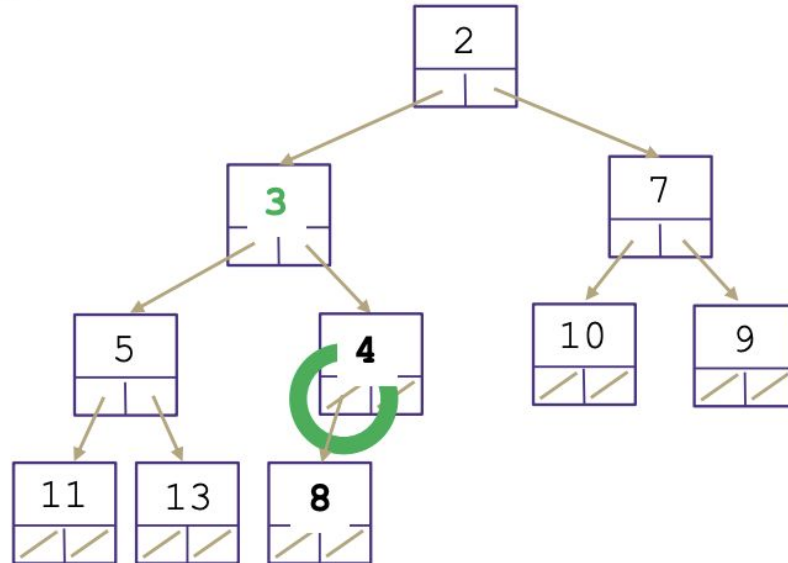




Heap Operation: Insert

Algorithm:

- Insert a node to ensure no gaps
- Fix heap invariant
- percolate **UP**

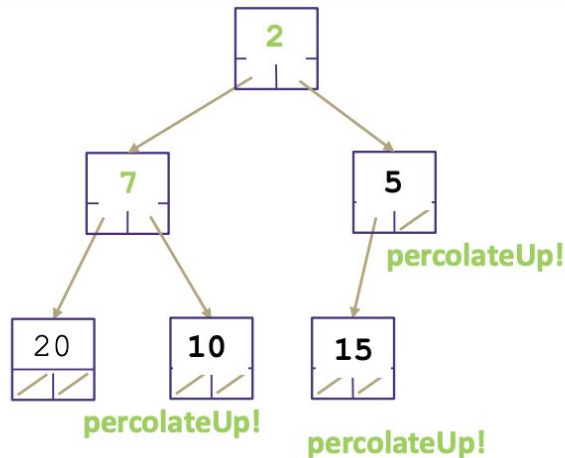


Heap Operation: Practice

5, 10, 15, 20, 7, 2

Min Binary Heap Invariants

1. **Binary Tree** – each node has at most 2 children
2. **Min Heap** – each node's children are larger than itself
3. **Level Complete** - new nodes are added from left to right completely filling each level before creating a new one





Heap construction: Process

1. Ensure that the heap retains the property of a max/min heap as the heap is built.
2. Build the heap and then transform the heap into a max/min heap (“heapify” the heap).



Heap construction: Method 1

Method 1: Add The First Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20					

- The corresponding tree





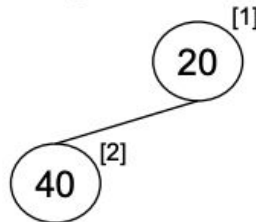
Heap construction: Method 1

Method 1: Add The Second Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40				

- The corresponding tree





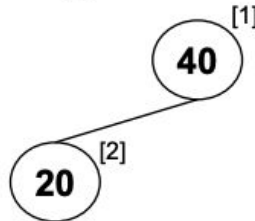
Heap construction: Method 1

Method 1: Swap The First And Second Elements

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	40	20				

- The corresponding tree





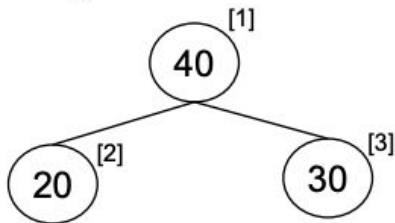
Heap construction: Method 1

Method 1: Add The Third Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	40	20	30			

- The corresponding tree





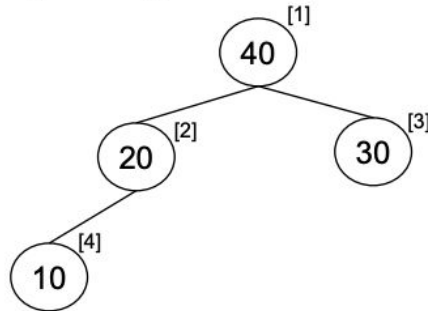
Heap construction: Method 1

Method 1: Add The Fourth Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	40	20	30	10		

- The corresponding tree





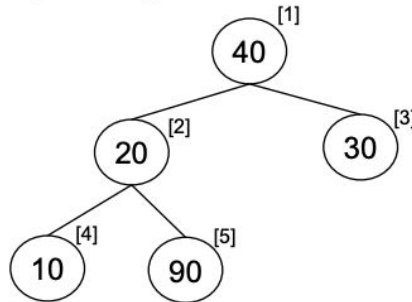
Heap construction: Method 1

Method 1: Add The Fifth Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	40	20	30	10	90	

- The corresponding tree





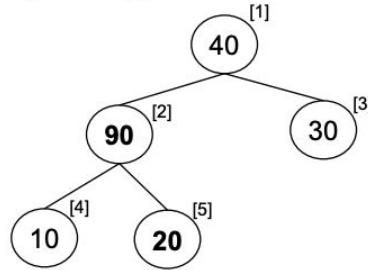
Heap construction: Method 1

Method 1: Swap the Second And Fifth Elements

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	40	90	30	10	20	

- The corresponding tree





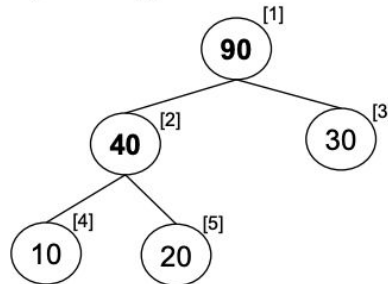
Heap construction: Method 1

Method 1: Swap the First And Second Elements

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	40	30	10	20	

- The corresponding tree





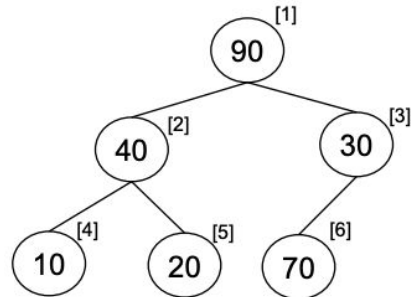
Heap construction: Method 1

Method 1: Add The Sixth Element

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	40	30	10	20	70

- The corresponding tree





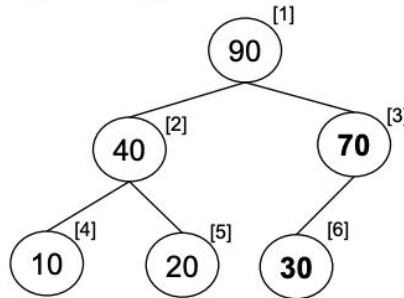
Heap construction: Method 1

Method 1: Swap The Third And Sixth Elements

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	40	70	10	20	30

- The corresponding tree





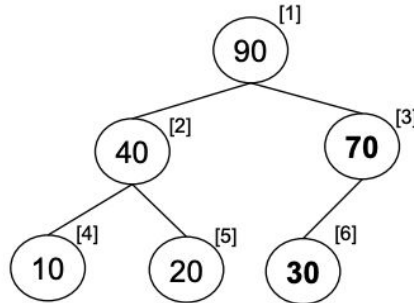
Heap construction: Method 1

Method 1: The Final State Of The Heap

- Array representation

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	40	70	10	20	30

- The corresponding tree



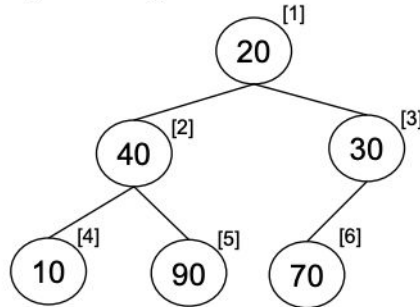
Heap construction: Method 2

Method 2 For Building A Heap

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40	30	10	90	70

- The corresponding tree



Also called
Floyd's method

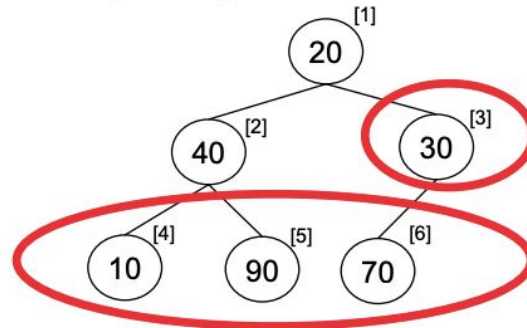
Heap construction: Method 2

Method 2 For Building A Heap: Where To Start

- Start with node $\lfloor \text{NoNodes}/2 \rfloor$, examine if the heap is a maxheap

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40	30	10	90	70

- The corresponding tree



Start
examining the
first non-leaf
node

Nodes after
node 3 will be
leaves



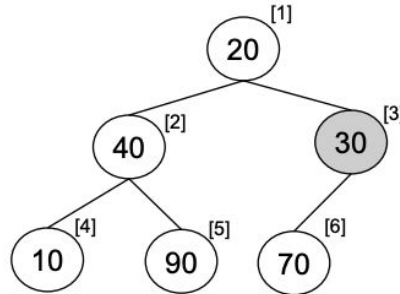
Heap construction: Method 2

Method 2 For Building A Heap: Examine Element [3]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40	30	10	90	70

- The corresponding tree



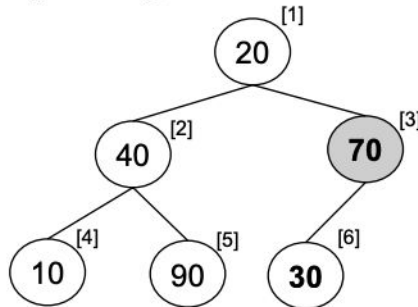
Heap construction: Method 2

Method 2 For Building A Heap: Reheap Related To Element [3]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40	70	10	90	30

- The corresponding tree





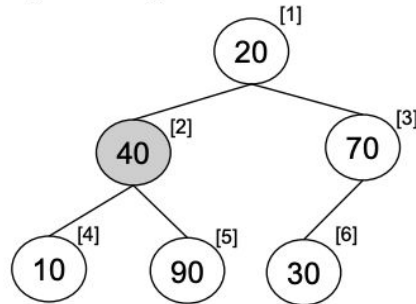
Heap construction: Method 2

Method 2 For Building A Heap: Examine Element [2]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	40	70	10	90	30

- The corresponding tree



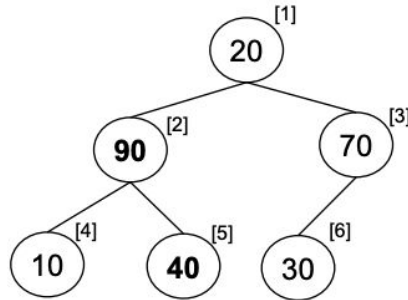
Heap construction: Method 2

Method 2 For Building A Heap: Reheap Related To Element [2]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	90	70	10	40	30

- The corresponding tree



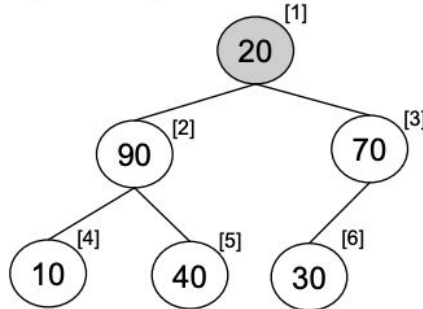
Heap construction: Method 2

Method 2 For Building A Heap: Examine Element [1]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	20	90	70	10	40	30

- The corresponding tree



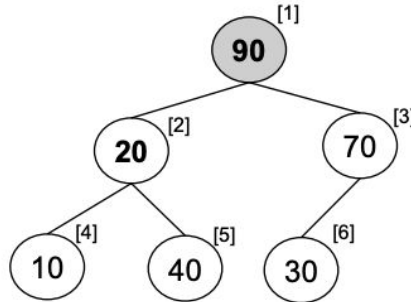
Heap construction: Method 2

Method 2 For Building A Heap: First Reheap Related To Element [1]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	20	70	10	40	30

- The corresponding tree





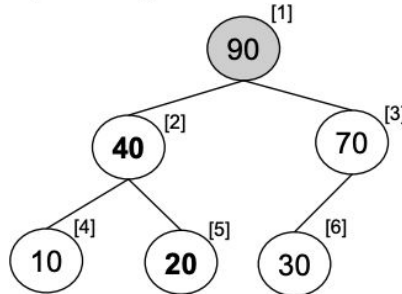
Heap construction: Method 2

Method 2 For Building A Heap: Second Reheap Related To Element [1]

- The information is read into an array

[0]	[1]	[2]	[3]	[4]	[5]	[6]
	90	40	70	10	20	30

- The corresponding tree





Heap construction: Method 2

Min-heap example

(UW slides: 18 - 23)



Heap Sort

MIT | 6-006 | Slide (21-27)



Recap

Heaps

- Properties
- Array implementation
- ADT
 - Max-heap
 - Min-heap
- Operations:
 - PeekMin()
 - ExtractMin()
 - Insert(val)
- Two methods for heap construction
- HeapSort