

# On interior point methods for LCPs

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# Why complementarity is so important?

- ① Linear and Nonlinear programming.
- ② Optimization Problems with Equilibrium Constraints.
- ③ Relaxation of binary variables.
- ④ ...

For given  $q \in \mathbb{R}^n$ ,  $M \in \mathbb{R}^{n \times n}$ , the  $LCP(M, q)$  is the following problem

$$\text{find } z \in \mathbb{R}^n : \quad 0 \leq z \perp w := Mz + q \geq 0$$

- $LCP(M, q)$  is an NP-hard problem for general matrices  $M$
- Fortunately, for different large classes of matrices  $M$ , there exist polynomial or at least efficient numerical methods.

# Classes of matrices

There are some usefull classes of matrices when studying LCPs.

- SS (Skew symmetric) matrices:

$$\forall x \in \mathbb{R}^n, \quad x^T M x = 0.$$

- SDP matrices:

$$M = {}^t M : \forall x \in \mathbb{R}^n, \quad x^T M x \geq 0.$$

- P matrices :

$$\forall x \in \mathbb{R}^n, \exists i \in \{1, \dots, n\} : \quad x_i (Mx)_i > 0.$$

- $P_0$  matrices :

$$\forall x \in \mathbb{R}^n, \exists i \in \{1, \dots, n\} : \quad x_i (Mx)_i \geq 0.$$

- SU : CSU (or RSU) :

$$\forall x \in \mathbb{R}^n, \forall i \in \{1, \dots, n\} : \quad x_i (Mx)_i \leq 0 \Rightarrow x_i (Mx)_i = 0.$$

(resp.)

$$\forall x \in \mathbb{R}^n, \forall i \in \{1, \dots, n\} : \quad x_i (M^T x)_i \leq 0 \Rightarrow x_i (M^T x)_i = 0.$$

## Other classes of matrices

- $P_*(\kappa)$  matrices ( $\kappa > 0$ ) such that

$$\forall x \in \mathbb{R}^n, (1 + 4\kappa) \sum_{i \in I^+(x)} x_i (Mx)_i + \sum_{i \in I^-(x)} x_i (Mx)_i \geq 0.$$

where  $I^+(x) := \{i : x_i (Mx)_i > 0\}$  et  $I^-(x) := \{i : x_i (Mx)_i < 0\}$

- $P_*$  matrices : Union of  $P_*(\kappa)$  matrices ( $\kappa > 0$ ).

We have in fact

$$SS \subset SDP \subset P_* = SU \subset P_0,$$

$$P \subset P_* \text{ and } P \cap SS = \emptyset.$$

# Principle of IPMs

We suppose  $M$  to be symmetric and positive definite.

Construct a sequence  $(z(\varepsilon), w(\varepsilon))$  satisfying

$$w(\varepsilon) \geq 0, z(\varepsilon) \geq 0, \|w(\varepsilon) - Mz(\varepsilon) - q\| \leq \varepsilon \quad \text{et} \quad z(\varepsilon)^T w(\varepsilon) \leq \varepsilon$$

To produce such sequence, perturb the equations  $LCP(M, q)$ , in the following scheme

$$\begin{cases} w = Mz + q \\ w \cdot z = \mu e \end{cases} \quad (2.1)$$

for a sequence of positive parameters  $\mu$  going to 0

Since  $M$  is positive definite, the perturbed system of equations has a unique solution for each  $\mu$  if  $F_0 \neq \emptyset$ .

## Principle of IPMs (2)

Remark: In general, we don't have (for free) any solution  $(\mu_0, z(\mu_0), w(\mu_0))$  of the perturbed system. If it is the case, the algorithm is said to be feasible.

- consider  $\mu_0 > 0$  and an initial point  $(z^0, w^0) > 0$  such that  $w^0 \cdot z^0 = \mu_0 e$ .
- Compute  $r^0$  such that  $\mu_0 r^0$  is the initial feasibility error:

$$\mu_0 r^0 = w^0 - Mz^0 - q$$

# The infeasible central path

- $(\mu, z(\mu), w(\mu)) > 0$  belongs to the Central path if

$$\begin{cases} w(\mu) - Mz(\mu) - q = \mu r^0 \\ w(\mu).z(\mu) = \mu e \end{cases} \quad (2.2)$$

Remark: we can use another parameter ( $\neq \mu$ ) to model the infeasibility.

- To measure the distance or proximity to the central path, one can use

$$\delta(\mu, z, w) = \frac{1}{2} \|v - v^{-1}\| \quad \text{where } v = \sqrt{\frac{z.w}{\mu}}.$$

Remark:

$$\delta(\mu, z, w) = 0 \Leftrightarrow v = e \Leftrightarrow (\mu, z, w) \in TC$$



# Generic Algorithm

- 1 **Initialization:** Choose a proximity tolerance  $\tau$ , a global precision  $\varepsilon$  and a reduction parameter  $\theta$ . Choose an initial point  $(\mu_0, z^0, w^0)$  and let  $(\mu, z, w) := (\mu_0, z^0, w^0)$ .
- 2 **Stopping criterion** While  $\max(\|w - Mz - q\|, z^T w) > \varepsilon$ , repeat

- **Prediction step:**

Compute a direction  $(\Delta z, \Delta w)$  such that

$$\begin{cases} M\Delta z - \Delta w = \theta \mu r^0 \\ w.\Delta z + z\Delta w = (1 - \theta)\mu e - z.w \end{cases} \quad (2.3)$$

$$z = z + \Delta z, \quad w = w + \Delta w, \text{ and } \nu := \sqrt{\frac{z.w}{\mu}}$$

Compute  $\delta(\nu)$ .

Reduce the central path parameter  $\mu = (1 - \theta)\mu$

- **Correction step:** While  $\delta(\nu) > \tau$ , repeat

Compute a direction  $(\Delta z, \Delta w)$  using

$$\begin{cases} M\Delta z - \Delta w = 0 \\ w.\Delta z + z\Delta w = \mu e - z.w \end{cases} \quad (2.4)$$

$$z = z + \Delta z, \quad w = w + \Delta w, \text{ and } \nu := \sqrt{\frac{z.w}{\mu}}$$

Compute  $\delta(\nu)$ .

# Convergence

I present here a simple convergence result without proof when  $\theta$  and  $\tau$  are constant and using complete steps.

## Theorem

*If  $\theta = \frac{1}{12n}$  and  $\tau = \frac{1}{4}$  then the generic algorithm needs at most  $12n \log \frac{33z_0^T w_0}{32\varepsilon}$  iterations to get an  $\varepsilon$  solution for the LCP( $M, q$ ) ( $O(n \log \frac{n}{\varepsilon})$  iterations).*

The best known complexity for IPM is  $O(\sqrt{n} \log \frac{n}{\varepsilon})$

- Write a matlab function *IPMLCP* with the following calling sequence

$$[z, w] = \text{IPMLCP}(M, q, \varepsilon, \theta, \tau);$$

- Generate a reduction parameter, a precision and a proximity tolerance

$$\varepsilon = \dots; \theta = \dots; \tau = \dots$$

- Generate an LCP with symmetric definite matrix  $M$  and some vector  $q$
- Solve the corresponding LCP using your matlab function

$$[z, w] = \text{IPMLCP}(M, q, \varepsilon, \theta, \tau);$$

- Compute the feasibility and the complementarity errors
- Make the size of the problem higher and higher, compute the outer iterations number and comment
- What can happen if  $M$  is not symmetric definite ?