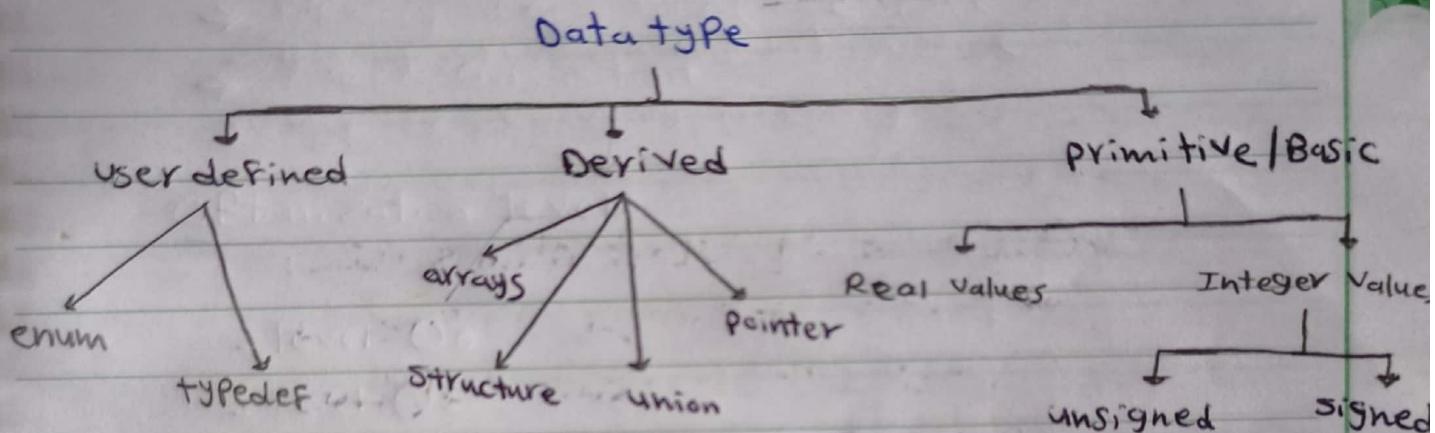


typeDef Command :

typedef Keyword in C used to create **an alias-name** to user defined data type to **primitive** and **derived datatype**.

**Syntax**

typedef <primitive or derived datatype> <new-name>

Some Notes to make your code more readable

1- Make **NEW_NAME** Capital letter
or like that **New_Name**

2- Put **t** after Name
ex **Name-t**

3- Put **S** Capital if **structure** and **U**
if **union** at beginning → **Sname-t**
if it is a **Enum** put **e** → **Name-e**

Header protection :Syntax

#ifndef T_H

#define T_H

#endif

usage

يمكننا استخدام #include في حالة إذا عدنا لفайл file أكثر من مرة فـ #include يتحقق
 لدينا واحد نعم إذا استخدمنا الملفات في ملف واحد فهو #include
 #include في definition لـ Text replacement
 #include في compilation error line
 #include في text replacement(يعني
 مرة واحدة بس والكور ميعرفش
 مدعانا ايرور.

optimization 1:-

is a series of actions taken by the compiler on your program's code generation process to reduce.

- number of instructions (Code space optimization)
- Memory access time (time space optimization)
- Power consumption

Optimization levels 2:-

- level O0 zero
- level O1 one
- level O2 two
- level O3 three

Compiler make optimization level by default is → level O0

نحوه ١٥٣ ملء الحالات التي تؤدي

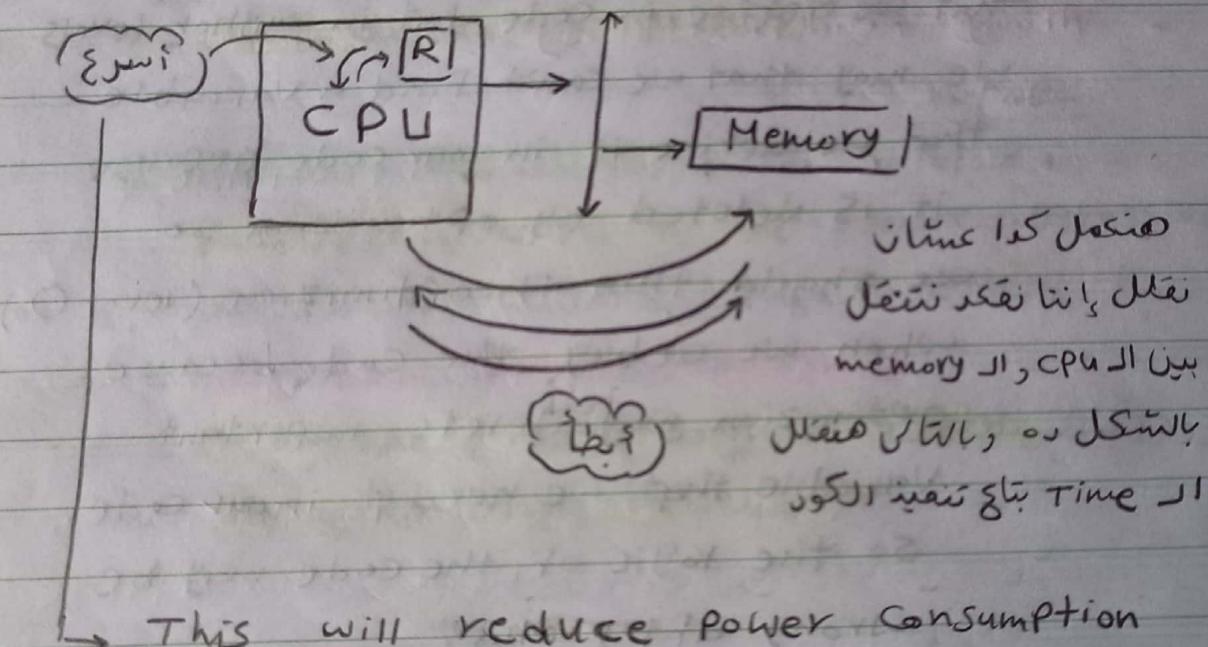
• reduce number of instructions

This will reduce .text section that has all assembly instructions so the code size will decrease so the binary file will take less memory in Flash memory and reduce the instructions will make the processor execute the code fast and less the execution time of the program.

Memory access time

If we have a variable and this variable used in number of equations and after every equation we save the result in the memory so this operation will take time so we can save it in a register in CPU to take less time and switch to memory in each operation.

in general
purpose
register



This will reduce power consumption

الكلام ده كله يعني دى اللي زيارة الـ زىارة الـ

زيارت الكود و تقليل الـ

والـ Memory

Optimization level O0

- No optimization
- Not recommended for production

if you have limited

Code

Ram Space

- Has fastest compilation time
- This is debugging friendly why??
because in debugging we find all variables that we define in code but in other levels we may that we can't find a variable that we define it in our code because it is deleted by optimization so we should turn off optimization (level O0) when we debug the code because optimization may delete important variable that we need it in our code. So the logic of the code may be wrong after this.

- A code which work with O0
may not work with Oot
optimization levels.

How to open the debug on Keil ??

بعد ما نكتب الكود ناصل build ونعمل stm32f103c8t6 في جو project رباعي الـ .elf من debug ملأه كوبس ونرجع لأن options for Target في Project settings في keil وبدر كوبس وندر كوبس بعد كذا نروح على الـ output المفتوح اللي أخذه كوبس يعني اختيار him debug himas executable ونعمل كذا كده المهمية البسيطة 😊

Optimization level O1 ..Optimization level O1 ..

Moderate optimization to reduce

- Memory access
- Code space

Optimization level O2 ..

- Full optimization.

- Slow compilation time.

- Not debugging friendly.

Optimization level O3 ..

- Full optimization of O2 + more aggressive steps will be taken by the compiler.

- slowest compilation time.

- May cause bugs in the program.

- Not debugging friendly.

Note

when using debug like level

level O2 or optimization II

→ less function

→ symbols known

→ pure assembly (only)

→ easier to debug

Volatile TYPE Qualifier

- Volatile keyword is a qualifier that is applied to a variable when it is declared

- it tells the compiler that the value of the variable may change at any time. "without any action being taken by the code" it will start
So it can't optimize.

- "Always Read From Memory" No optimization possible

Syntax

- Volatile int x_j
 - int volatile x_j

لهمانا Variable مدين وقيمة ميئ بتغير باستخاده ال Code حال optimization المفروضه هى ذيده فما هنا عاييز ال variable ده ميئ فش لأن قيمة بتغير بس ميئ الكور بس بتغير عن طريق اى حاجة موجودة او Hardware وبالتالي ده ميئ فش

Proper use of C's volatile keyword

- ### - Memory-mapped peripheral registers

و register هي registersips في الموديل ليس عالي

و registerN pointer درسته فاچتا values ایو یعنی

بس مکان قیمة ار pointer تغییر کی وئی فاختنا

عازفون نهرأ التعديل ده و تعرف بيه فاحنا هنخال

Java Pointer JI is Volatile as Pointer JI

موجود و میداد خش بینی اد optimization و کمان

نقطة أو تعدل بمحمل في قيمة الـ pointer

نقرة على زر تديل بحث عن قيمة الـ pointer

- Global variables modified by an interrupt service routine (ISR)
- Global variables accessed by multiple tasks within a multi-threaded application.

note

Volatile uint8-t * PReg;

uint8-t volatile * PReg;

Pointer to volatile data

Volatile unsigned 8-bit integer

int * volatile P;

Volatile Pointer to non volatile data

int volatile * volatile P;

Volatile Pointer to volatile data

notes

يجب أن يكون الـ register هي
نقطة إدخال متحركة إلى volatile data لـ
الـ optimization أو لـ memory屏障
فهي تكون ملحوظة في
الـ program.

How To access the register by absolute address ??

ex

write 0xFFFFFFF on SIU register which have absolute address 0x3061000

Solutions

1st method

Volatile int *p;

p = (Volatile int *) 0x3061000;

*p = 0xFFFFFFF;

2nd method

*((Volatile int *) (0x3061000)) = 0xFFFFFFF;

3rd method

#define SIU_Register

*((Volatile int *) (0x3061000))

SIU_Register = 0xFFFF FFFF;

2nd method = 3rd method

4th method

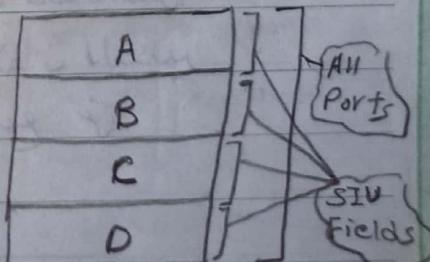
في الطريقة الرابعة يوجد union primitive list

تحت إنتا لوكاينز نستخدم ار

مع 4 Ports (32 bit) عادي ولو

كابيزين نستخدم كل Port

فبر فتو عادي (8 bit)



```

typedef union {
    uint32_t All_Ports;
    struct {
        uint32_t PortA:8;
        uint32_t PortB:8;
        uint32_t PortC:8;
        uint32_t PortD:8;
    } SIU_Fields;
} SIU_R;

```

Volatile SIU_R * Ports = (Volatile SIU_R *) 0x3061000;

- Ports → All_Ports = 0xFFFFFFF;

يمثل جميع Ports الـ 4 هنا بـ 16 بت

- Ports → SIU_Fields.PortA = 0xFF;

يمثل Port A الـ 8 بت هنا

Usage of const and volatile together

`uint32_t volatile * const Preg1 = (uint32_t *) 0xFFFF0000;`

Preg1 is a constant pointer to volatile data
of type `uint32-t`.

Const is used to guard the pointer from unexpected changes from the programmer.

العنوان الذي يكتب user ليس pointer الذي هو const للعنوان الذي يكتب user ليس memory الذي هو const

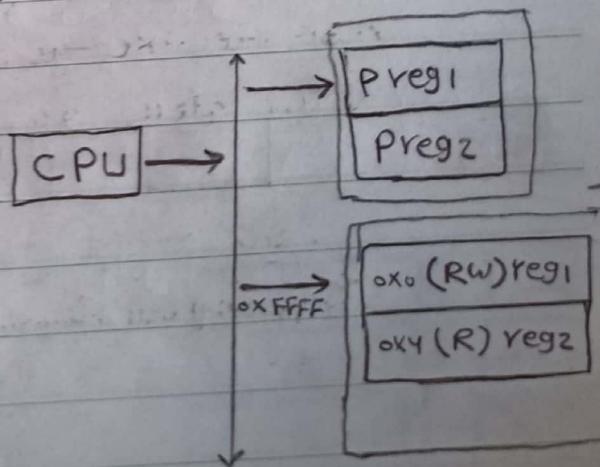
`Preg1++;` → This will make compilation error.

Volatile used to tell the compiler that data pointed could receive unexpected changes so the compiler has not to apply any optimization on the data.

pointer يكتب مني يعني reg1 أو عيّن المؤلف

pointer يكتب مني يعني reg1 أو عيّن المؤلف external I/P

∴ reg1 → (RW) ✓✓



`uint32_t const volatile * const Preg2 = (uint32_t *) 0xFFFFd004;`

Preg2 is a constant pointer pointing to volatile constant data of type `uint32t`.

What that means ?? 😊

This Qualifier (`const + volatile`) on data.

it is fixed (☺) it can be changed at any time (☹)

Const means that the programmer can't change the data by the code and if he try to do that the compiler will give error !!

Volatile means that the data can be changed by the hardware "external input" not from the programmer.

البيانات المخزنة في الذاكرة يمكنها أن تغير قيمتها من دون إرادة المبرمج

it is useful for (`R (read only access)`) register like the status register which is changed from Hardware not from the programmer

بيانات الوضعية هي البيانات التي لا ي-Controlها المبرمج
إلا بتغييرها من قبل الكوادل الخارجية مثل
إذا أخذنا القيمة بـ `in reg1` ولو حاولنا تغيير قيمته
بـ `out reg1, value` فـ `value` لا يغير قيمة `reg1`

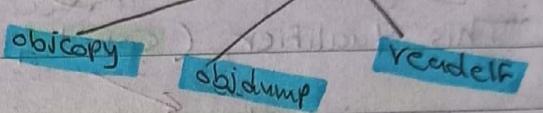
Tool Chain

consists of ⇒ **Compiler**, **Assembler**,

Linker, **Debugger**,

Libraries,

[**Binary utilities**] ← **number of helper programs**



the host PC (Ubuntu, redhat, ...) by default
is native toolchain.

ToolchainCROSS

in this tool chain we
setup it in our PC and
we write code on it and
it will run on PC but
generates code for another
architecture.

جهاز الآخرين (PC) يكتب
الكود بحسب arachitecture
bin file يجيء من run على
آخر arch (جهاز آخرين)

ex: Atmel studio

STM32 Cube

Native

in this tool chain we
set up it in our PC and
we write code on it and
it will run on PC and
generates code for the
same PC.

جهازنا (PC) يكتب
الكود بحسب arachitecture
bin file يجيء من run على
نفس arachitecture

ex: mingw

True or False

The executable file gets from cross tool chain
it can run on the same tool chain.

(False)

Why

لأن الـ code يكتب على cross toolchain وليس على executable file في وبالتالي لا يمكن تشغيله على another architecture

Illustration

Host machine

هي التي نكتب بها الكود

target machine

هي التي نريد تشغيل الكود عليها

Static library

هو مكتبة binary code لمجموعة من الـ functions

وتحتاج إلى استخدام أي حاجة من الـ functions

بس بدون ما نكرر implementation فنأخذ

هذا lib static lib نحملها الدوارده ونذا بالع

إن اسمها يبدأ بـ lib وبينها بـ .

وارد يحصل لهم static libraries

مع الـ object file جوا الـ linker

يخرج علينا الـ executable file

الـ main function لما يجيء نستخدم printf

عارفين الـ implement بـ include

لينك بين الـ stdio.h

بس اخذناه مع الـ static libraries

بسوندي مثلا كل

note

- library object files should not contain
 - a main function
- static library names should start with "lib" and end with ".a" \Rightarrow ex: libmath.a

important note

لأن linker يأخذ كل الأجزاء بالجزء المكتوب في file name .a

يأخذ كل الأجزاء التي تكتب في file name .a

يعني لو نكتب "printf" في برنامج فالـ linker يأخذ جزء file name .a من "stdio.h" ومن "math.h" لأن printf

To create a static library file

ar rcs lib-name.a File1.o File2.o

To add another object file to the library

ar r lib-name.a File3.o

To remove an object file from the library

ar d lib-name.a File3.o

To View the object files in the library

ar t lib-name.a

To extract object files from the library

ar x lib-name.a

note

- arm-none-eabi
- arm-linux-gnueabi

ابن الفرقانين (٦)**arm-none-eabi**

هـ يستخدم مع الـ Bare metal applications
بـ يـ نـ كـ تـ بـ

direct ← arm v/s run alien code

arm-linux-gnueabi

هـ يستخدم في إنـ نـ كـ تـ بـ مـ شـ ظـ ظـ عـ لـ application

ولـ يـ كـ لـ الـ linux وـ يـ كـ لـ direct ← arm

الـ microcontroller

who is used to write baremetal application?

on Arm Board??

a) arm-none-eabi

✓

b) arm-linux-gnueabi

طـ وـ جـ دـ حـ وـ جـ دـ حـ

main.c , lib.c , lib.h ← line 3Files ١- نـ كـ لـ

lib.h لـ prototype جـ وـ اـ لـ lib.c رـ نـ طـ اـ لـ Func بـ تـ بـ

main.c دـ يـ جـ وـ اـ لـ Func N call ٢-

٣- نـ ظـ اـ لـ lib.c وـ main.c مـ نـ اـ لـ main.o مـ نـ اـ لـ

٤- نـ كـ لـ lib.c object file جـ وـ اـ لـ static lib

٥- نـ كـ لـ link بين اـ لـ lib وـ main.o وـ static lib وـ نـ ظـ اـ لـ

٦- نـ ظـ اـ لـ lib.exe وـ main.exe وـ static lib

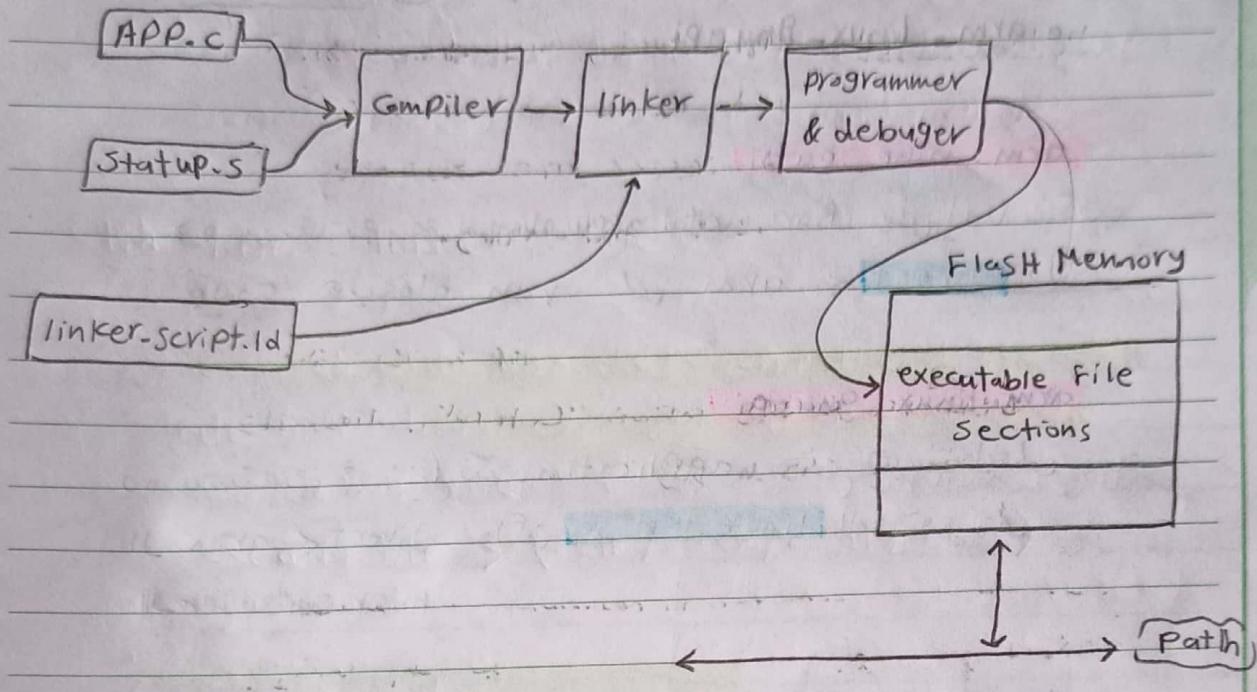
gcc -c -I . main.c -o main.o

ar rcs lib-trystatic.a lib.o

gcc main.o lib-trystatic.a -o lib.exe

lib.exe ✓

Compilation Process :-



1st stage Pre-processing stage

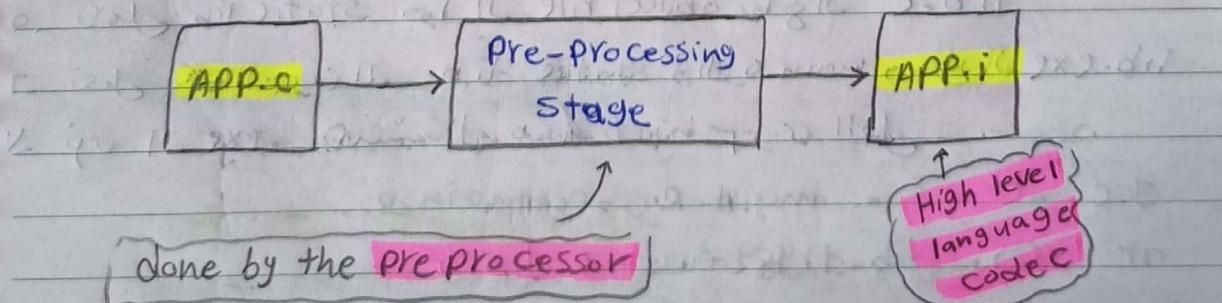
in this stage \Rightarrow all pre processing directives

will be resolved

1. #include
2. Macros \rightarrow #include \rightarrow #ifdef

3. intermediate file \leftarrow (File.i) lines

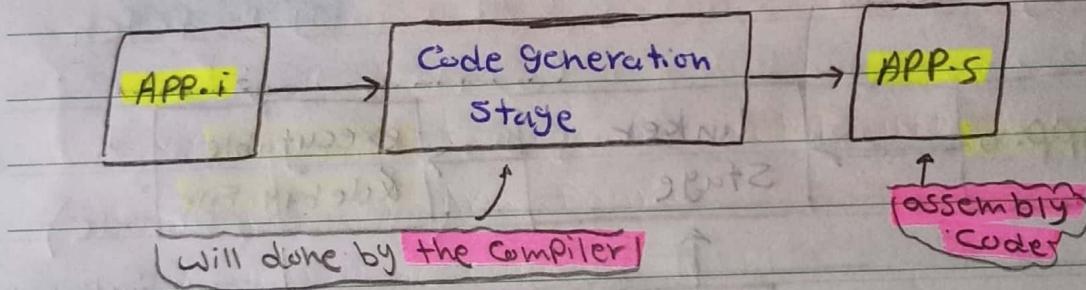
4. #pragma \leftarrow #define \leftarrow #include



2nd stage Code generation Stage

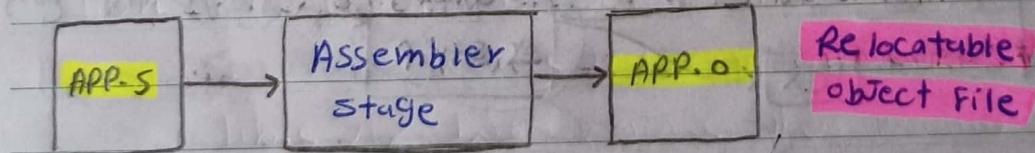
in this stage \Rightarrow High level language code will be converted into processor architecture level assembly.

This stage dependent on Architecture
why because every architecture has its own assembly.



3rd stage Assembler Stage

Assembly will be converted into opcode & operands (binary)



مادى عن ال Relocatable object File

binary code أو object code هو باردة وتحتاج إلى ملء ببيانات

Virtual addresses على ال needs بجزئي

memory و physical addresses على ال

- Processor architecture specific machine

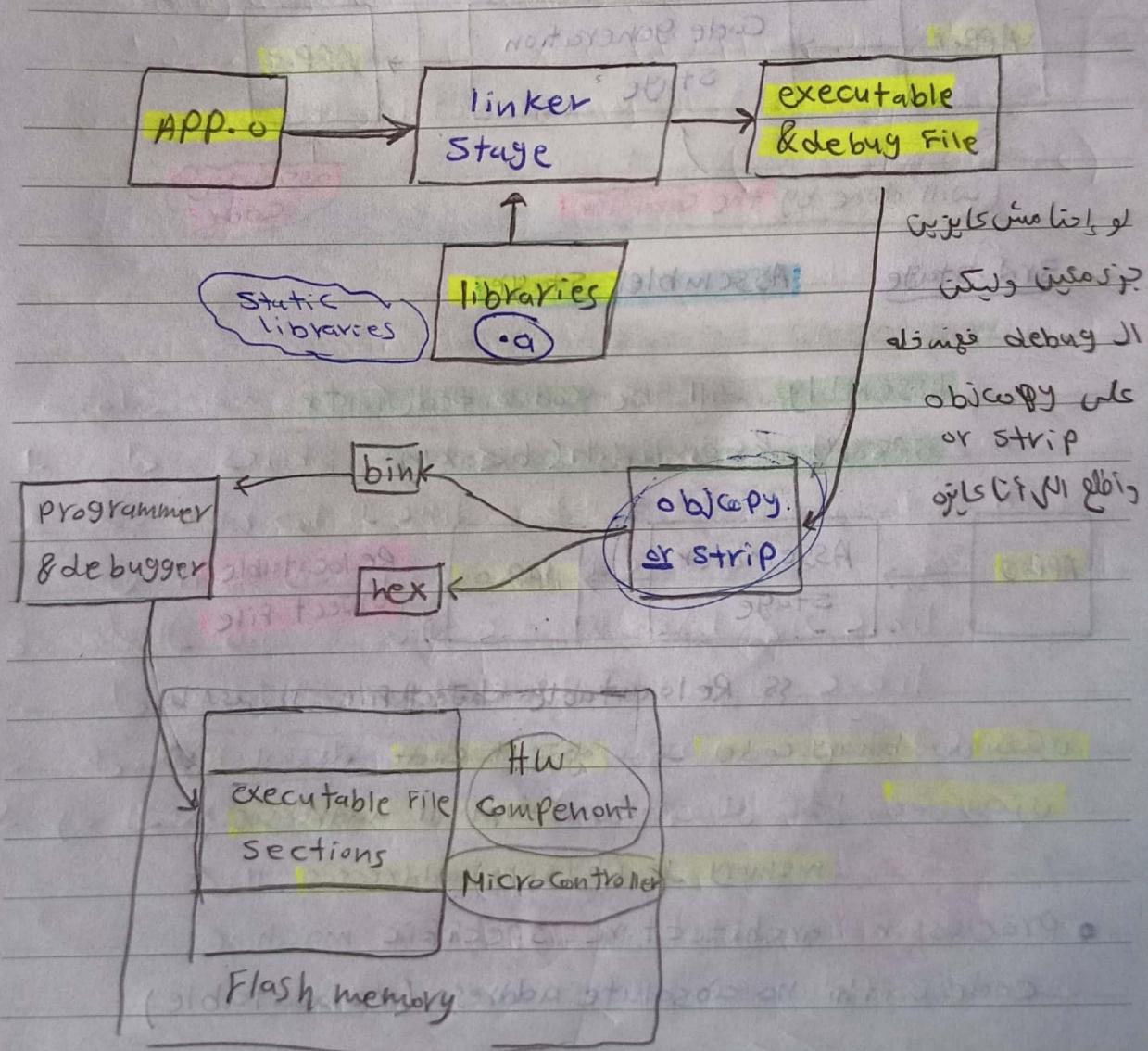
code with no absolute address (relocatable)

المسؤول عن ترتيب السكن بين المنشآت assembly
وي 自动生成 addresses واحد ويسيطر على microcontroller addresses مع المنشآت هو linker script

4th stage linker stage

the linker will link between the object file and libraries and out the executable and debug file for us.

executable → (.elf / .axf / .out / .exe)



- map File

it is a file we can get it ~~from~~ after linking stage.

it is includes the actual length of the sections and the names and locations of the public symbols found in the relocatable program.

ex symbols → global variables, Functions,

A processor executes what kind of code??

- a) assembly code
- b) Machine code
- c) C code
- d) Morse code

Linker file

- provides details on the memory size and location.
- maps compiled sections to physical memory.
- can check if the memory was over located.
- can provide the entry point to the program.

if a program is linked against a static library then

Machine code of the used function is copied in the executable.