# RHEINISCH-WESTFÄLISCHE TECHNISCHE HOCHSCHULE AACHEN KNOWLEDGE-BASED SYSTEMS GROUP PROF. GERHARD LAKEMEYER, Ph. D.

### Seminararbeit

# On CRIKEY - A Temporal Planner Looking at the Integration of Scheduling and Planning

Mostafa Gomaa

Matrikelnummer: 340822

Seminar Dynamics of Knowledge and Belief Sommersemester 2017

## **Contents**

1	Intro	oduction	1
	1.1	The Planning Problem	1
	1.2	Conceptual Model	2
	1.3	Classical Planning	2
		1.3.1 PDDL	2
		1.3.2 State Space Planning	2
	1.4	Non-classical Planning	2
		1.4.1 Heuristic Search Planning	2
		1.4.2 Graphplan	2
	1.5		2
		1.5.1 FF	2
		1.5.2 MetricFF	2
	1.6	The Scheduling Problem	2
		1.6.1 STN	2
2	Tem	nporal Planning	3
	2.1	PDDL2.1	4
		2.1.1 Temporal Constraints	4
3	CRI	KEY	8
	3.1	motivation	8
	3.2		9
		3.2.1 Action Splitting	0
	3.3	The Workings of CRIKEY	0
	3.4	Conclusion	3
Bi	bliogi	raphy 1	4

## 1 Introduction

Temporal planning is basically the same problem as classical planning but with time introduced. In Temporal planning classical planning moves closer to scheduling. The combination of ,tow already hard problems, only makes temporal planning harder. Many Approaches has been introduced to attempt of solving the temporal planning problem. CRIKEY a temporal planner that separates the logical and temporal reasoning, is main focus of this seminar paper. It attempts to solve the temporal planning problem by breaking it down into two smaller sub-problems. As with the case with most interacting engineering problems, some issues arise and the separation is not always easily achievable.

## 1.1 The Planning Problem

The planning problem is the problem of finding a sequence of actions that lead us from some initial state of the world to a state where some objectives are met. For humans, planning comes very natural. With a purpose of achieving a certain goal, choosing a specific sequence of actions from the space of all possible actions a human can perform, anticipating the outcome of each and ordering them as to meet or avoid some constraints, is a simple every day task for us. Automated Planning is area of Artificial Inelegance that is concerned with choosing and deliberating a sequence of actions, for their desired expected effects. It is concerned with questions like how to represent those actions, to how to model the world they interact with and how to reason about their effects, that would lead eventually to choosing and executing them. "Planning is the reasoning side of acting" (Ghallab et al., 2004)

1 Introduction 2

## 1.2 Conceptual Model

## 1.3 Classical Planning

- 1.3.1 PDDL
- 1.3.2 State Space Planning

**Forward Search** 

**Backward Search** 

- 1.4 Non-classical Planning
- 1.4.1 Heuristic Search Planning
- 1.4.2 Graphplan
- 1.5 Planners
- 1.5.1 FF
- 1.5.2 MetricFF
- 1.6 The Scheduling Problem
- 1.6.1 STN

## 2 Temporal Planning

Temporal planning is basically planning with time. In classical planning time has been ignored with restricted assumption A???. This allowed the study of a simpler problem, giving rise to great algorithms to tackle the abstracted classical planning problem. Yet in real world domains time could rarely be ignored. Temporal Planning considers the relaxation of the abstraction a step further. It introduces time to the problem. "Temporal Planning is the same problem as classical planning, but with time introduced" (Halsey et al., 2003). It seems when time is introduced, that the planning problem extends its focus from only reasoning about logical constraints, and finding set of actions that needs to achieve a goal, to also arranging this sequence in time as not to break any temporal constraints. This sounds a lot like what the study on scheduling is about as discussed briefly in section 1.6. Finding an order for executing a set of actions that need to use limited set of resources, is what the study of Scheduling is concerned with (Ghallab et al., 2004). Scheduling, has been an area of research for decades and it has many well developed schedulers written to solve the problem.

Looking from this perspective, What if we could solve the each of those problems independently. What if we could find away to separate time from the problem to be able to use the already written classical planners on the smaller problem and then care about fitting them into the time constraints later on. As it has been with the engineering approach for years, to break down the complicated problems into its smaller constituting problems and solve each alone, it has also been clear that the problem separation is not easily achievable. There is usually a margin that the problems are interdependent, and cannot be resolved one after the other. The problems interact too much or too coupled, that solution of one of the problems highly dependent on the solution of the other.

There has been already a great deal of research oriented to how to solve the temporal planning problem by separating it into classical planning and scheduling when possible and solve each independently, detecting when the problems are too tightly coupled to separate and finding clever ways to go about these cases.

CRIKEY, the system discussed in this seminar paper does exactly that. Details on how it attempts to separate, detect and solve the temporal planning problem will be detailed in Chapter 3. In this Chapter, the focus will be on the more general and representational aspects of the temporal planning, as well as, a broader look into the problems arise when attempting to isolate the planning and scheduling. How CRIKEY tackles those problems is deferred Chapter 3.

#### 2.1 PDDL2.1

PDDL2.1 is domain and problem description language, that extends the modelling power of PDDL () to meet some of challenges in modelling real world problems. One of those challenges is to have some notion of time. There have been many models tackling this problem, yet the most common is for actions to have a duration. PDDL2.1 does exactly that. It introduced away for formalizing durative actions /citep. Moreover, since actions are no longer instantaneous as in PDDL, it is no longer sufficient to assume that preconditions or effects would also apply instantaneously. In the real world scenario logical changes (i.e. conditions required or effect applied) can occur at point of time during the action's execution. Ranging from, only occurring at certain points of an action's execution, to holding over the whole period of the action, it is clean that more elaborate approach for representing when logical changes occur is needed. PDDL2.1 allow for actions to have effects or preconditions apply either at the start or at the end of an action. Furthermore, it allows to have conditions hold over the whole period of an action, called invariant.

It might still seem that some domains might require a more complex model of time, yet a simple compilation of the domain and the use of dummy actions and predicates can give great power. Modelling of, Logical changes happening at a given point during the actions (that is neither the start nor the end), or more complex temporal constraints (like enforcing actions to partly overlap), becomes possible.

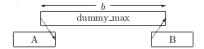
#### 2.1.1 Temporal Constraints

This section introduces how temporal constrains can be represented in PDDL2.1. It attempts to summarize a section, that carries the same name, in the work of Halsey et al. (2003). All the figures and arguments here are credited to (Halsey et al., 2003).

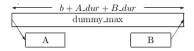
In PDDL2.1, there is no explicit way to represent temporal constraint. It is only possible implicitly through the use of durative actions, which is the only way to get temporal information into the plan. Whether those durative actions are actually meant to account for a real desirable action taking place, or just used as a dummy action to ensure a certain temporal constrain, is secondary in this discussion. There is mostly two important temporal constraint to consider here, the maximum and minimum time by which an action x follows an action y. All other constraint that do not use disjunction are specialised cases of these. Each of those constraints can be represented in two different ways. namely "clipping" or "enveloping".

Figure 2.1 shows how to represent the maximum difference constraint x - y <= b (i.e, the maximum time allowed for x to follow y). Notice that when expressing a maximum constraint, regardless of the representation, the direction of the arrows is from the start of the durative (potentially dummy) action to the first action, and from the second action to the durative (potentially dummy) action. The ordering of the actions is enforced here

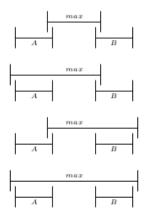
using logical constraints (i.e. By manipulating the preconditions, invariants and effects of actions to enforce the order). Note that, with the maximum constraint, there is no minimum ensured. Nothing would stop the second action from being ordered before the first.



(a) Expressing Maximum Elapsed time with Clipping



(b) Expressing Maximum Elapsed time with Enveloping



(c) Compact notation of possible combinations

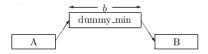
Figure 2.1

Figure 2.2 shows how to represent the minimum difference constraint x - y >= b (i.e, the minimum time allowed for x to follow y). Notice that when expressing a minimum constraint, regardless of the representation, the direction of the arrows is from the first action to the start of the durative (potentially dummy) action, and from the end of durative (potentially dummy) action to the second action (i.e. By manipulating the preconditions, invariants and effects of actions to enforce the order). Note that, with the minimum constraint, there is no maximum ensured. The second action can occur at any point after the first once it passed the minimum (or never).

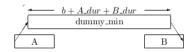
The more interesting cases would be when both constraints needs to be modelled at the same time Figure 2.3.

#### **Analysis of Representations**

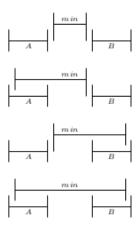
The two representation "clipping" vs "enveloping", though syntactically equivalent, they are not always semantically equivalent. Some domains would correspond in its encoding



(a) Expressing Minimum Elapsed time with Clipping

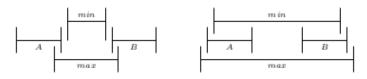


(b) Expressing Minimum Elapsed time with Enveloping



(c) Compact notation of possible combinations

Figure 2.2



**Figure 2.3:** Expressing both Minimum and Maximum Time Between Actions in PDDL2.1

to one representation but not the other, even though, it would seem equivalent in some other domains. Further more, In one representation, enveloping, the duration of the durative (potentially dummy) action, seems to depend on the duration of the actions it envelops (namely, the content actions). This results in a tight coupling between the content actions and the constraint they are fulfilling. Whilst in clipping, this is not the case. In clipping, planning and scheduling problems do not interact in a complex manner, resulting for both problems to be loosely coupled. To push it even further, what if those durative (potentially dummy) actions are not really dummy at all. What if they are actions with a specific duration, conditions and effects, that the content actions can only be executed during them. In this case, the problem just looks like a maximum constraint problem, but in fact it in not a maximum problem at all. In this case we have

very little control over the duration of the enveloping action. The specific arrangement of those actions, envelope and content, need to be ensured by the planner as not to break any temporal constraint.

In summery, translating the problem from one representation to the other, highly depends on the semantics of the domain, and not always possible.

On the bright side, we have just made clear when exactly, the planning and scheduling problems in temporal planning, are too tightly coupled to separate. Where the logical structure and the temporal constraint interact and can not be solved separately. This gives CRIKEY, and similar systems, the inspiration for when to would it make sense to separate the problems and solve each alone, and when would the separation effect the problem description in essence, and it would be required to find more clever ways to solve both together.

This chapter focuses on how CRIKEY, a planner that separates the logical and temporal reasoning, tackles the separation of the sub-problems

Section 3.1 provides some insight on a predecessor system, problems with it and what motivates CRIKEY. Section 3.2 provide some technical specification on the planning and scheduling technologies used at the heart. Section 3.3 overview of the system architecture, explanation of how and when it separates the problems when possible, how it detects the coupling of the problems and how it deals with it.

#### 3.1 motivation

A system was developed in "(CITE)" for separating out the temporal and logical reasoning requires in temporal planning. Figure ?? summarize the main idea of the system. Given a temporal domain and problem description in PDDL2.1. A translator, separates out the temporal aspects translating it into a PDDL classical planning problem and domain that preserves the key temporal relations. The temporal information about actions duration are stored in a separate file. The translator used was the same as in LPGP "(CITE)". It splits a durative actions into three instantaneous actions, modelling the at start conditions and effects, invariants and at end conditions and effects respectively. Later, the translated PDDL problem was solved with a classical planner producing a totally ordered plan. From the totally ordered plan, a partially ordered plan is lifted and feed, together with the duration file, separated at the begging, into a Simple Temporal Network. This schedules the plan by calculating relative and actual timing of actions. The product is a Temporal Plan.

This system always separates the planning and scheduling problem, and solves each sub problem independently, as proposed in chapter 2. This turns out to be a good idea for a variety of domains (including all problems in International Planning Completion 2002) "CITE". In these domains the solution plan could always be sequentialised (i.e. executed one action after the other). Concurrency was an optimization choice of the planner in case actions don't interfere. Yet, in those domains, there was no concurrency enforced by the domains itself. In those domains, the temporal and logical constraints are loosely coupled and the problems do not deeply interact.

However, Some domains exists where this is not the case. Where it is necessary to perform some actions in parallel, and where the choice of actions depends on both their

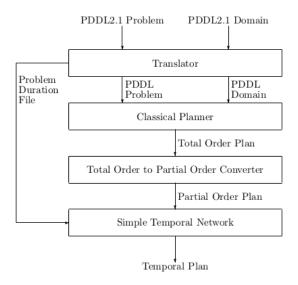


Figure 3.1: An Architecture for Separating Planning and Scheduling

logical and temporal merits. In such domains the planning and scheduling problems are tightly coupled in the manner demonstrated in chapter 2. In these domains, the planner of predecessor of CRIKEY, would produce plans that the scheduler can not schedule. The un-schedulable plans comes from temporal constraint in the problem that are not met. The planner here lacks any temporal knowledge, and only makes its choices based on the logical constraints.

This is where CRIKEY comes. CRIKEY basically inherits the same basic architecture for separating and solving the planning and scheduling problems independently when possible. Moreover, it always tries to detect cases where the problems are tightly coupled and solves them together. Details on how CRIKEY detects those cases and what does it do when it finds them are in the main focus of this chapter.

## 3.2 Technical Specification

CRIKEY planner was implemented in java. The planning algorithm implemented in cirkey is forward chaining heuristic planner. The implementation is based very closely on MetricFF "CITE". As in MetricFF, it uses a relaxed plan length as a heuristic, helpful actions (which appear in the first layer of the relaxed plan) to aid its action selection and performs Enforced Hill-Climbing. Unlike MetricFF, upon finding a plateau in the search space, it performs Best First Search to exit the plateau. This means expanding the same states as FF at the start, but if no exit found, will try those states that move away from the goal.

#### 3.2.1 Action Splitting

CRIKEY, separates the durative actions (when necessary) into two instantaneous actions. At start and at end instantaneous actions. Invariants of the durative actions, become preconditions of the start and end spontaneous actions. Unless, if that invariant was achieved by the at start effects of the durative actions, it only become a condition of the end spontaneous action. For each state, the planner keeps a list of invariants that holds in that state. Actions that delete an invariant are considered inapplicable in this state. When an action is chosen, its invariants are either added to (if it was a start action) or removed from (if it was an end action) the invariants list.

Another list keeps track of semi-complete actions. Those that the start action is chosen, yet the end action is not. A final state could only be achieved if this list is empty and the goal is reached. The start actions are given a cost equal to its duration and end actions are given a small tolerance value (to motivate their selection in case they were not selected for some useful effect).

This list also comes in handy during the extraction of the heuristic value from the relaxed plan. If during graph plan generation, a start action was choose for its desired effect, but not its end action, only the start action will show in the relaxed plan. However if later, during plan relaxed plan extraction this start action was selected, then its end action needs to also be selected, increasing the size of the relaxed plan by one.

CRIKEY does not allow negative conditions in a durative action.

## 3.3 The Workings of CRIKEY

As discussed in chapter 2, when a sequence of content actions needs to be completed before another set of enveloping actions do, is exactly when caution needs to be taken in their selection and ordering. If there not enough time to execute content actions during the enveloping actions we risk producing an unscheduled plan. CRIKEY, avoids this risk by looking for "potential envelopes". Here we are looking for actions that allow other actions to only execute during their duration.

The detection happens in two different place in the algorithm with very similar reasoning. Once during the construction of the relaxed plan graph and extraction of the relaxed plan (during calculation of the heuristic). And another, during the actions selection. When a durative action is identified as a "potential envelops". The durative action is then split into a start and end instantaneous actions (as described in the previous section). When an durative actions is clearly not an envelope, the durative action is condensed into one instantaneous action (with effects combined and the conditions are the at start conditions combined with any invariants unsatisfied by the at start effects). If the "potential envelope" action is later selected by the FF planner, on adding the start action, a mini-scheduler is associated with it to ensure there is enough time for any

potential content actions to execute during the enveloping action. A content action is only selected if there is enough time for it to execute with the consideration of any other content actions already selected within the envelope. Once the envelope end action is selected the mini-scheduler is discarded.

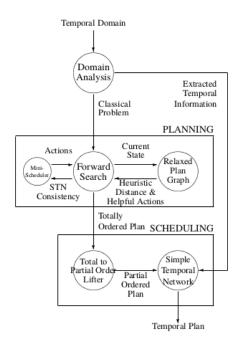


Figure 3.2: Architecture of CRIKEY

#### **Detecting Envelops**

During the construction and extraction of the relaxed plan, delete lists are ignored. A durative action is considered safe and condensed into a single action, if all of their at end conditions are satisfied by at start the add effect, Figure 3.3. Otherwise, when there is an at end cond unsatisfied by the at start effects, it is considered an envelop, and split into start and end actions.

$$\begin{aligned} & ((D_{cond_{snd}} \setminus D_{add_{start}}) = \emptyset) \Rightarrow \\ A_{cond} &= D_{cond_{start}} \cup (D_{cond_{invs}} \setminus D_{add_{start}}) \\ A_{add} &= D_{add_{start}} \cup D_{add_{snd}} \end{aligned}$$

Figure 3.3: A Durative Action D is condensed into an Instantaneous Action A

During the action selection, delete lists needs to be taken into account. Similar reasoning is used to identify none envelope cases, Figure 3.4.

```
((cond_{end} \setminus add_{start}) = \emptyset)

\land ((add_{start} \cup del_{end}) = \emptyset)

\land ((del_{start} \cap cond_{end}) = \emptyset)
```

Figure 3.4: When a durative action is considered none envelop

Follows is the reasoning presented at "CITE". "s1,s2 and s3 are three states immediately before a start of the action, immediately after and immediately after the end of the action. An action applicable in s2 but not s1 must have been achieved by the at start add effect (since there is no negative conditions it could not have been achieved by an at start delete effect). Taking it further, there are no actions the could be applied in s2 and not s3 which could not have been applied in s1, apart from those achieved by the at start add effects and then deleted by the at end delete effects. Alternatively, an action could be achieved by the at start effect, and the effects of this action needed to achieve the end conditions. They are called potential envelopes since there is no effort (at this moment) to find out if there are any content actions that must go in these envelops"

To recap, actions which allow other actions to happen only during their duration is what we are looking for. This formalization of these enveloping actions is described in figure.

```
(end_{cond} \setminus start_{add} \neq \emptyset \land start_{add} \setminus end_{cond} \neq \emptyset)

\lor start_{del} \cap end_{cond} \neq \emptyset

\lor add_{start} \cap del_{end} \neq \emptyset
```

Figure 3.5: Envelops Formalization

#### The Mini-scheduler

Each mini-scheduler associated with an envelope action, consists of a Simple Temporal Network, a list of content actions that need to fit inside the envelope and their ordering. The mini scheduler and the main scheduler in CRIKEY both use the same algorithm. After an enveloping start action is selected, the mini-scheduler makes sure to only select a content action if the STN is consistent. Meaning there will be enough time to execute the content action within the envelope. Content actions that are deemed inapplicable by the STN, are pruned. Figure 3.6 shows the pseudo code describing the mini-scheduler.

#### The Scheduler

After the partial order lifter, a modified version of "CITE", produces a partial order plan. It is passed through a Simple Temporal Network that uses Floyd-Warshall's algorithm

- 1. Check  $A_{cond}$  are satisfied.
- 2. Check  $A_{del}$  do not delete any currently open invariants.
- 3. If A is a start of an envelope
- (a) Create a new envelope for A and add to list of envelopes.
- 4. Else If A is an end of an envelope
- (a) Remove A's envelope from the list of envelopes.
- 5. For Each envelope E in the list
- (a) Get orderings for A in E.
- (b) If no orderings, return true.
- (c) Add orderings to the STN.
- (d) return the consistency of the STN.

Figure 3.6: Algorithm to decide whether an action A is applicable

to compute the actual timing of actions, producing a temporal plan.

#### 3.4 Conclusion

# **Bibliography**

- M. Ghallab, D. Nau, and P. Traverso. *Automated Planning: theory and practice*. Morgan Kaufmann, San Francisco, CA, 2004.
- K. Halsey, D. Long, and M. Fox. Isolating Where Planning and Scheduling Interact. In Proceedings from the 22nd UK Planning and Scheduling Special Interest Group PlanSIG'03, 14:104–114, 2003.