

Big O Notation - Detailed Notes

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Key Concepts

Common Big O Notations

1. $O(1)$ - Constant Time

- The algorithm's runtime does not change with the size of the input.
- Example: Accessing a specific element in an array by index.

2. $O(N)$ - Linear Time

- The runtime grows linearly with the size of the input.
- Example: Iterating through an array of size N .

3. $O(N^2)$ - Quadratic Time

- The runtime is proportional to the square of the input size. Often occurs with algorithms involving nested loops.
- Example: A double loop over an array.

4. $O(\log N)$ - Logarithmic Time

- The runtime grows logarithmically as the input size increases. Common in algorithms that divide the problem in half at each step (e.g., binary search).
- Example: Binary search in a sorted array.

5. $O(N \log N)$ - Log-Linear Time

- The runtime is slightly worse than linear but better than quadratic. Common in efficient sorting algorithms.
- Example: Merge sort and quicksort.

6. $O(2^N)$ - Exponential Time

- The runtime doubles with each additional element in the input. Extremely inefficient for large inputs.
- Example: Recursive algorithms that solve subproblems multiple times.

7. $O(N!)$ - Factorial Time

- The runtime is proportional to the factorial of the input size. Very poor performance, typically seen in algorithms that generate all possible permutations.
- Example: Brute-force solutions to the traveling salesman problem.

How to Calculate Big O Notation

1. Identify the Input Size (N):

- Determine what the input is and define its size as **N**. The input size could be the length of an array, the number of elements in a list, or the number of nodes in a tree.

2. Analyze the Loops:

- **Single Loops:** If an algorithm has a loop that iterates **N** times, the time complexity is typically $O(N)$.
- **Nested Loops:** If there are nested loops, multiply the complexities of the loops. For example, a loop inside another loop that both run **N** times results in $O(N^2)$.
- **Logarithmic Loops:** If the loop decreases the problem size by half each time (e.g., $i = i * 2$), the time complexity is $O(\log N)$.

3. Analyze Conditional Statements:

- Conditional statements like **if-else** do not usually affect the time complexity unless the conditions contain loops or recursive calls. In such cases, analyze the code within those conditions.

4. Focus on the Most Significant Term:

- When combining terms, keep only the most significant one (i.e., the one that grows the fastest). For example, if an algorithm has steps that take $O(N)$ and $O(N^2)$ time, the overall time complexity is $O(N^2)$.

5. Drop Constants and Lower-Order Terms:

- Big O notation describes the upper bound, so drop any constants or less significant terms. For instance, $O(3N)$ simplifies to $O(N)$, and $O(N^2 + N)$ simplifies to $O(N^2)$.

Examples of Big O Calculation

1. Example 1: Single Loop

```
def example_function(arr):  
    for i in range(len(arr)):  
        print(arr[i])
```

- **Analysis:** The loop runs **N** times (where **N** is the size of the array).
- **Big O Notation:** $O(N)$.

2. Example 2: Nested Loop

```
def example_function(arr):  
    for i in range(len(arr)):  
        for j in range(len(arr)):  
            print(arr[i], arr[j])
```

- **Analysis:** The outer loop runs N times, and for each iteration of the outer loop, the inner loop also runs N times.
- **Big O Notation:** $O(N) * O(N) = O(N^2)$.

3. Example 3: Logarithmic Loop

```
def example_function(n):
    while n > 1:
        n = n // 2
```

- **Analysis:** The loop decreases n by half each time, resulting in a logarithmic number of iterations.
- **Big O Notation:** $O(\log N)$.

4. Example 4: Multiple Terms

```
def example_function(arr):
    for i in range(len(arr)): # O(N)
        print(arr[i])

    for j in range(len(arr)): # O(N)
        for k in range(len(arr)): # O(N)
            print(arr[j], arr[k])
```

- **Analysis:** The first loop runs in $O(N)$, and the nested loop runs in $O(N^2)$.
- **Big O Notation:** $O(N) + O(N^2)$ simplifies to $O(N^2)$ since the quadratic term dominates.

Big O Notation Comparison

- **Hierarchy of Common Complexities:**

- $O(1) < O(\log N) < O(N) < O(N \log N) < O(N^2) < O(2^N) < O(N!)$

Common Algorithms and their Big O Notations

1. Bubble Sort - $O(n^2)$

```
def bubble_sort(arr):
    n = len(arr)
    # O(n^2) - Overall Big O
    for i in range(n):
        for j in range(0, n-i-1):
            if arr[j] > arr[j+1]:
                arr[j], arr[j+1] = arr[j+1], arr[j]
```

which is approximately n times.
operation.

$O(n)$ - Outer loop runs n times.
$O(n)$ - Inner loop runs $(n-i-1)$ times,
$O(1)$ - Comparison is a constant time
$O(1)$ - Swapping elements is

```
a constant time operation.
return arr
```

2. Merge Sort - $O(n \log n)$

```
def merge_sort(arr):
    #  $O(n \log n)$  - Overall Big O
    if len(arr) > 1:
        mid = len(arr) // 2
        L = arr[:mid]
        R = arr[mid:]

        merge_sort(L) #  $O(\log n)$  - Recursive call on the left half.
        merge_sort(R) #  $O(\log n)$  - Recursive call on the right half.

        i = j = k = 0

        while i < len(L) and j < len(R): #  $O(n)$  - Merging the sorted halves.
            if L[i] < R[j]: #  $O(1)$  - Comparison is a constant time
operation.
                arr[k] = L[i]
                i += 1
            else:
                arr[k] = R[j]
                j += 1
            k += 1

        while i < len(L): #  $O(n)$  - Merging remaining elements.
            arr[k] = L[i]
            i += 1
            k += 1

        while j < len(R): #  $O(n)$  - Merging remaining elements.
            arr[k] = R[j]
            j += 1
            k += 1
```

Quick Sort - $O(n \log n)$ on average, $O(n^2)$ in the worst case

```
def quick_sort(arr):
    #  $O(n \log n)$  on average,  $O(n^2)$  in the worst case - Overall Big O
    if len(arr) <= 1:
        return arr
    pivot = arr[len(arr) // 2]
    left = [x for x in arr if x < pivot]
    middle = [x for x in arr if x == pivot]
```

```
right = [x for x in arr if x > pivot]
return quick_sort(left) + middle + quick_sort(right)
```

Linear Search - $O(n)$

```
def linear_search(arr, x):
    #  $O(n)$  - Overall Big O
    for i in range(len(arr)):          #  $O(n)$  - Looping through n elements in the
array.                                #  $O(1)$  - Comparison is a constant time
        if arr[i] == x:
operation.
            return i
    return -1
```

Binary Search - $O(\log n)$

```
def binary_search(arr, x):
    #  $O(\log n)$  - Overall Big O
    left, right = 0, len(arr) - 1
    while left <= right:                #  $O(\log n)$  - The loop runs  $\log(n)$  times,
halving the search space each time.    #  $O(1)$  - Calculating the midpoint is a
        mid = (left + right) // 2      #  $O(1)$  - Comparison is a constant time
operation.                             #  $O(1)$  - Comparison and reassignment are
        if arr[mid] == x:              #  $O(1)$  - Reassignment is a constant time
operation.                             #  $O(1)$  - Reassignment is a constant time
            return mid
        elif arr[mid] < x:
operation.                             #  $O(1)$  - Reassignment is a constant time
            left = mid + 1
        else:
            right = mid - 1
operation.
    return -1
```