## **Definition: Conditional Mutual Information**

Applying the **definition of mutual information** to the conditional distribution  $P_{XY|\mathcal{A}}$  naturally defines  $I(X;Y|\mathcal{A})$ , the mutual information of X and Y conditioned on the event  $\mathcal{A}$ :

## **Definition: Conditional mutual information**

Let X,Y,Z be random variables. Then the conditional mutual information of X and Y given Z is defined as

$$I(X;Y|Z) = \sum_{z} P_{Z}(z) I(X;Y|Z\!=\!z) \, ,$$

with the convention that the corresponding argument in the summation is 0 for z with  $P_Z(z)=0$ .

Conditional mutual information has properties similar to the ones we saw for mutual information:

$$\begin{split} I(X;Y|Z) &= I(Y;X|Z)\\ I(X;Y|Z) &\geq 0\\ I(X;Y|Z) &= 0 \text{ iff } X \text{ and } Y \text{ are independent given } Z. \end{split}$$

Furthermore, the previous bounds  $H(X) \geq 0$ ,  $H(X|Y) \geq 0$ , and  $I(X;Y) \geq 0$ , can all be seen as special cases of  $I(X;Y|Z) \geq 0$ . These bounds, and any bound they imply, are called **Shannon inequalities**. It is important to realize that I(X;Y|Z) may be larger or smaller than (or equal to) I(X;Y).

The following is sometimes used as definition of I(X;Y|Z): verify it for yourself using the definition above.

## **Alternative definition**

Let X,Y,Z be random variables. Then

$$I(X;Y|Z) = H(X|Z) - H(X|YZ).$$

Proof

$$\begin{split} I(X;Y|Z) &= \sum_{z \in \mathcal{Z}} P_Z(z) I(X;Y|Z=z) \\ &= \sum_{z \in \mathcal{Z}} P_Z(z) (H(X|Z=z) - H(X|Y,Z=z)) \\ &= \sum_{z \in \mathcal{Z}} P_Z(z) H(X|Z=z) - \sum_{z \in \mathcal{Z}} P_Z(z) H(X|Y,Z=z) \\ &= H(X|Z) - H(X|YZ). \end{split}$$
 (by definition of mutual information of the property of the prop

created: 2018-12-12