

Communicating Sequential Processes

...

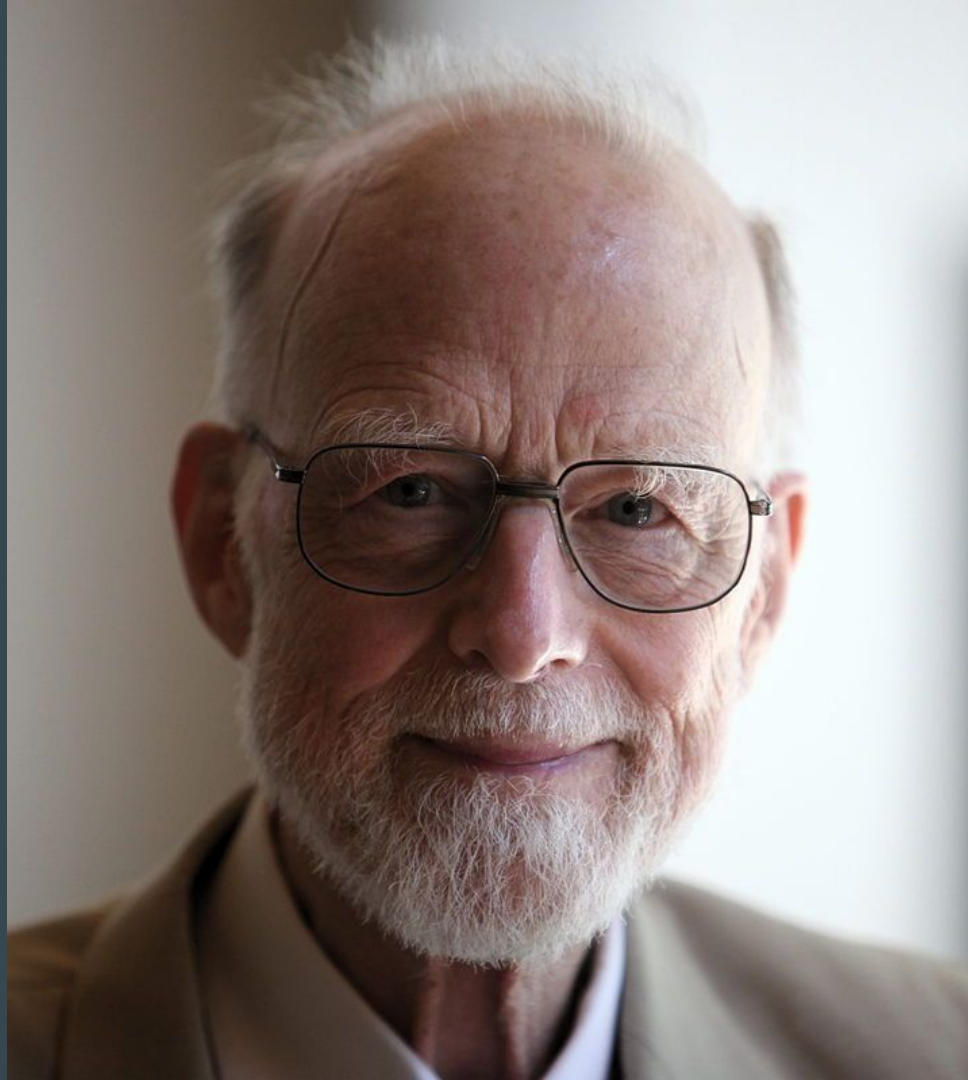
overview of a 1978 paper by C.A.R. Hoare

Sir Tony Hoare

Hoare logic

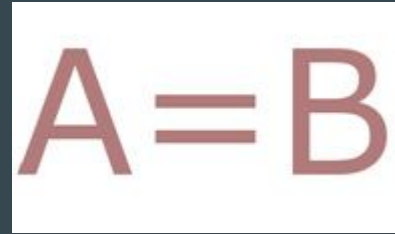
Quicksort

CSP



1. Introduction

- Assignment:
 - well understood
 - built into programming languages

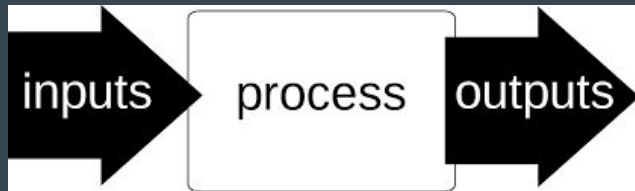


1. Introduction

- Assignment:
 - well understood
 - built into programming languages
- Input/output:
 - *not* well understood
 - often left to libraries



A diagram showing the assignment statement $A=B$ in a large, red, serif font. The text is centered within a white rectangular box.



1. Introduction

- Programming languages usually agree on having:
 - repetitive constructs (loops)
 - alternative constructs (conditionals)
 - sequential composition

1. Introduction

- Programming languages usually agree on having:
 - repetitive constructs (loops)
 - alternative constructs (conditionals)
 - sequential composition
- but *not* on the design of **abstractions**:
 - subroutines (Fortran)
 - coroutines (Unix)
 - classes (Simula 67)
 - actors (*Actor model* by Carl Hewitt, 1973)

The image shows the 'Need for Speed' logo in a bright blue, stylized, italicized font. The logo is positioned over the front of a green car, with a circular headlight visible on the left. The background is a blurred night scene with warm, golden bokeh lights.

NEED FOR SPEED™

Traditional computer was designed for
“deterministic execution of a single sequential program”

CDC 6600

Used at CERN since 1965

Had 10 parallel functional units



Multiprocessor machine

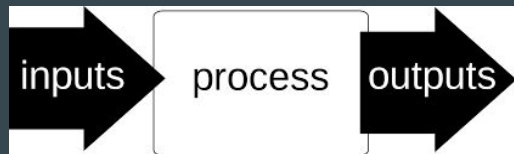
“... may become more powerful, capatious, reliable, and economical than a machine which is disguised as a monoprocessor.”

Need for synchronization

- Shared mutable state
- Monitors
- Semaphores
- Queues

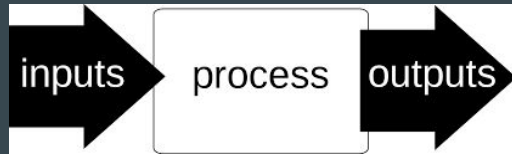


CSP's Conjecture

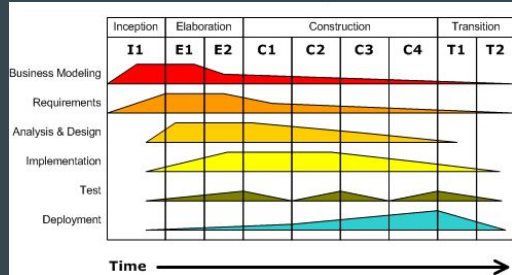


“Input and output
are basic primitives of programming”

CSP's Conjecture



“Input and output
are basic primitives of programming”



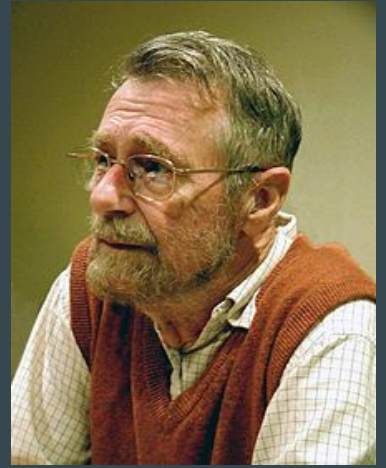
“Parallel composition of
communicating sequential processes
is a fundamental program structuring method”

CSP's Conjecture

- ❖ These concepts are “surprisingly versatile”
when combined with Dijkstra’s guarded commands

Guarded Command Language

*Edsger
Dijkstra*



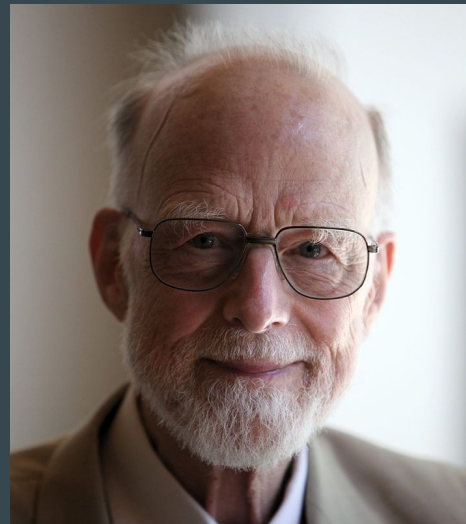
A simple pseudo-language that makes it easier to prove the correctness of programs.

	<u>Guards</u>	<u>Statements</u>
if	$a \geq b \rightarrow$	$\text{max} := a$
	$ \quad b \geq a \rightarrow$	$\text{max} := b$
fi		

	<u>Guards</u>	<u>Statements</u>
do	$a > b \rightarrow$	$a := a - b$
	$ \quad b > a \rightarrow$	$b := b - a$
od		

CSP Language

Based on Guarded Command Language



by Tony Hoare

2. Concepts and Notations

- Assignment Command

$$(1) \quad x \coloneqq x + 1$$

$$(2) \quad (x, y) \coloneqq (y, x)$$

$$(3) \quad x \coloneqq \text{cons}(\text{left}, \text{right})$$

$$(4) \quad \text{cons}(\text{left}, \text{right}) \coloneqq x$$

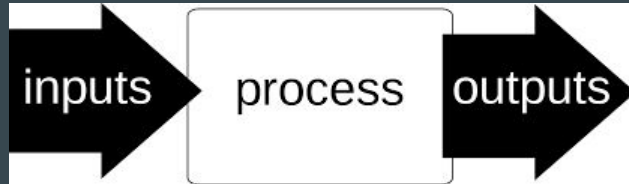
⋮

$$(5) \quad \text{insert}(n) \coloneqq \text{insert}(2 * x + 1)$$

$$(6) \quad c \coloneqq P()$$

2. Concepts and Notations

- Input Command
- Output Command



(1) **cardreader?cardimage**

(2) **lineprinter!lineimage**

(3) **$X?(x, y)$**

(4) **DIV!(3*a + b, 13)**

2. Concepts and Notations

- Alternative Command

MAX $\{x \geq y \rightarrow m \vdash x \parallel y \geq x \rightarrow m \vdash y\}$

- Repetitive Command

SEARCH $i \vdash 0; *[i < \text{size}; \text{content}(i) \neq n \rightarrow i \vdash i + 1]$

2. Concepts and Notations

- Parallel Command

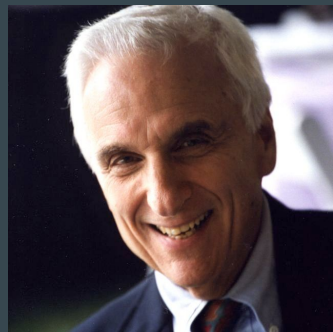
[cardreader?cardimage||lineprinter!lineimage]

$X(i:1..n) :: \text{CL}$

❖ "mutually disjoint"

3. Coroutines

First coined by *Melvin Conway* and published in his 1963 paper:

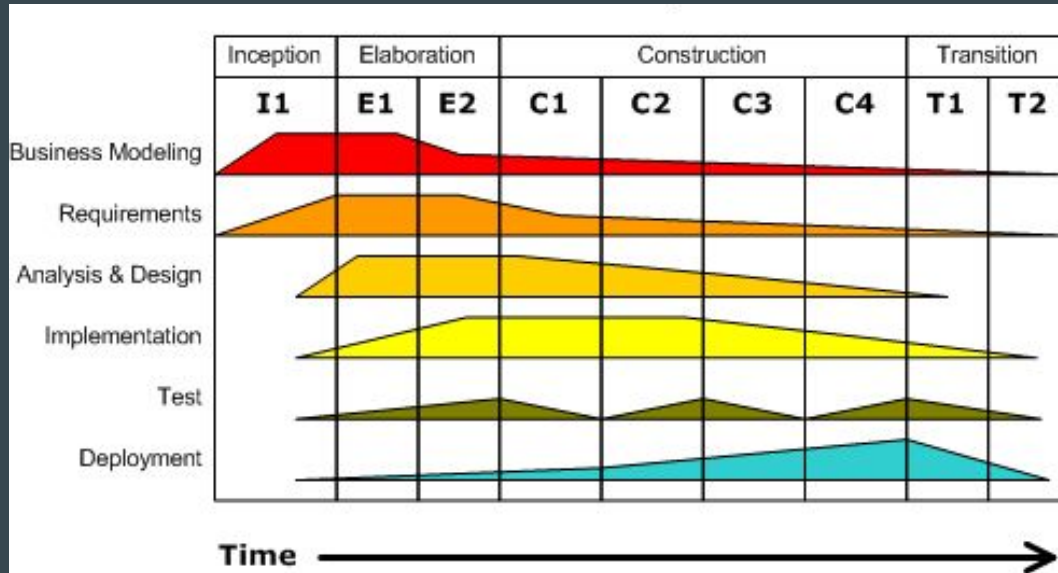


*Coroutine ... an autonomous program which communicates with adjacent modules as if they were *input* or *output* subroutines.*

*Thus, coroutines are subroutines *all at the same level*, each acting as if it were the master program when in fact *there is no master program*.*

from “Design of a Separable Transition-Diagram Compiler”

3. Coroutines



3. Coroutines



3. Coroutines

COPY

$X :: *[c:\text{character}; \text{west}?c \rightarrow \text{east}!c]$

SQUASH

$X :: *[c:\text{character}; \text{west}?c \rightarrow$
 $[c \neq \text{asterisk} \rightarrow \text{east}!c$
 $\parallel c = \text{asterisk} \rightarrow \text{west}?c;$
 $[c \neq \text{asterisk} \rightarrow \text{east}!\text{asterisk}; \text{east}!c$
 $\parallel c = \text{asterisk} \rightarrow \text{east}!\text{upward arrow}$
 $]\parallel]$

3. Coroutines

DISASSEMBLE

```
lineimage:(1..125)character;  
i:integer; i := 1;  
*[c:character; X?c →  
    lineimage(i) := c;  
    [i ≤ 124 → i := i + 1  
    []i = 125 → lineprinter!lineimage; i := 1  
]    ];  
[i = 1 → skip  
[]i > 1 → *[i ≤ 125 → lineimage(i) := space; i := i + 1];  
    lineprinter!lineimage  
]
```

ASSEMBLE

```
*[cardimage:(1..80)character; cardfile?cardimage →  
    i:integer; i := 1;  
    *[i ≤ 80 → X!cardimage(i); i := i + 1]  
    X!space  
]
```


3. Coroutines (parallel composition)

REFORMAT

```
[west::DISASSEMBLE||X::COPY||east::ASSEMBLE]
```

CONWAYS_PROBLEM

```
[west::DISASSEMBLE||X::SQUASH||east::ASSEMBLE]
```

4. Subroutines and Data Representations

Coroutines can be used to simulate:

- subroutines
- recursion
- ADTs / objects / actors

4. Subroutines and Data Representations

Subroutine

```
[DIV::*[x,y:integer; X?(x,y) →  
    quot,rem:integer; quot := 0; rem := x;  
    *[rem ≥ y → rem := rem - y; quot := quot + 1];  
    X!(quot,rem)  
    ]  
||X::USER  
]
```

4. Subroutines and Data Representations

Recursion

```
{fac(i:1..limit)::  
  *[n:integer; fac(i - 1)?n →  
    [n = 0 → fac(i - 1)!]  
    ||n > 0 → fac(i + 1)!n - 1;  
      r:integer; fac(i + 1)?r, fac(i - 1)!(n * r)  
    ]]  
||fac(0)::USER  
}
```

4. Subroutines and Data Representations

ADT

SET

```
content:(0..99)integer; size:integer; size := 0;
*[n:integer; X?has(n) → SEARCH; X!(i < size)
[]n:integer; X?insert(n) → SEARCH;
    { i < size → skip
    []i = size; size < 100 →
        content (size) := n; size := size + 1
    }
[]X?scan( ) → i:integer; i := 0;
    *[i < size → X!next(content(i)); i := i + 1];
    X!noneleft( )
]
```

5. Monitors and Scheduling

Monitor

```
*[(i:1..100)X(i)?(value parameters) → ... ; X(i)!(results)]
```

5. Monitors and Scheduling

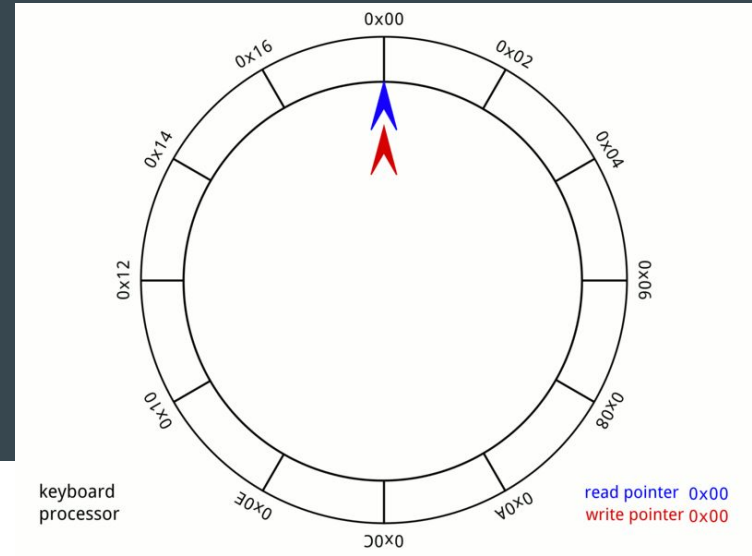
Semaphore

```
S::val:integer; val := 0;  
  *[(i:1..100)X(i)?V( ) → val := val + 1  
  ][(i:1..100)val > 0; X(i)?P( ) → val := val - 1  
  ]
```

5. Monitors and Scheduling

Bounded Buffer

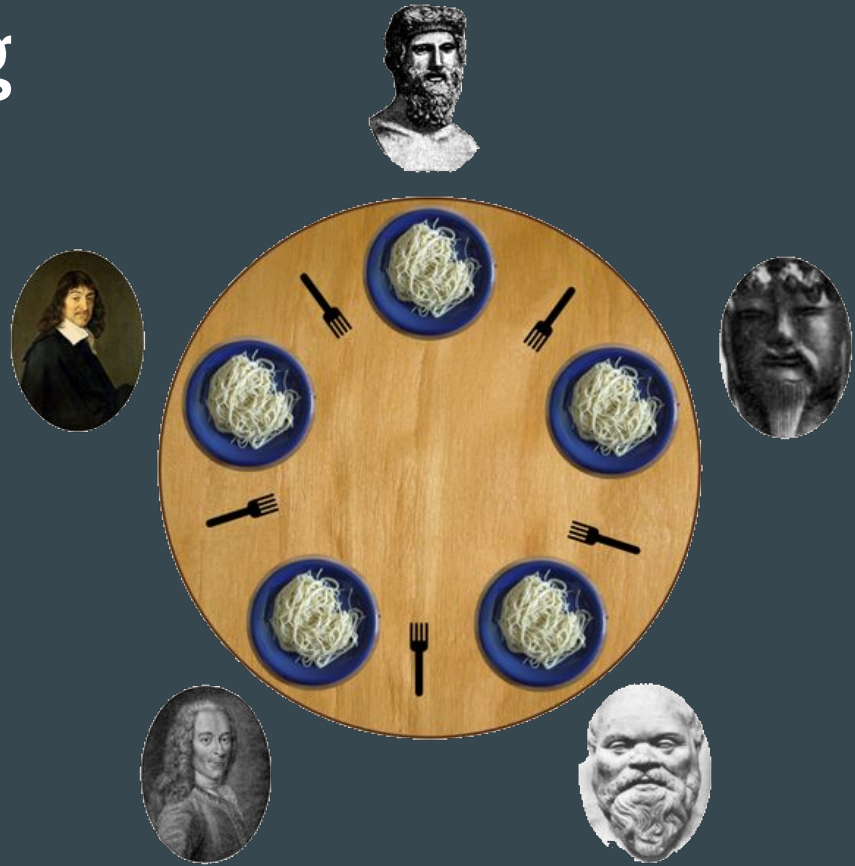
```
buffer:(0..9) portion;  
in,out:integer; in := 0; out := 0;  
comment  $0 \leq out \leq in \leq out + 10$ ;  
  * $[in < out + 10$ ; producer? $buffer(in \bmod 10) \rightarrow in := in + 1$   
  ] $out < in$ ; consumer? $more( ) \rightarrow consumer!buffer(out \bmod 10)$ ;  
    out := out + 1  
  ]
```



5. Monitors and Scheduling

Dining Philosophers

```
[room::ROOM  
  ||fork(i:0..4)::FORK  
  ||phil(i:0..4)::PHIL]
```



5. Monitors and Scheduling

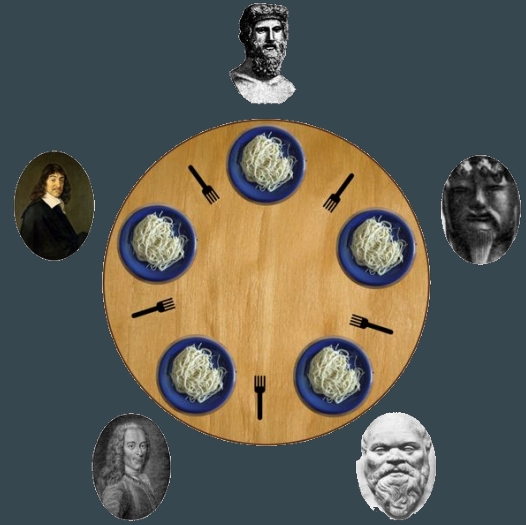
Dining Philosophers

```
PHIL = * [... during ith lifetime ... →  
    THINK;  
    room!enter( );  
    fork(i)!pickup( ); fork((i + 1) mod 5)!pickup( );  
    EAT;  
    fork(i)!putdown( ); fork((i + 1) mod 5)!putdown( );  
    room!exit( )  
    ]
```



5. Monitors and Scheduling

Dining Philosophers



FORK =

```
*[phil(i)?pickup( ) → phil(i)?putdown( )  
  ||phil((i - 1) mod 5)?pickup( ) → phil((i - 1) mod 5)?putdown( )  
  ]
```

5. Monitors and Scheduling

Dining Philosophers



```
ROOM = occupancy:integer; occupancy := 0;  
*[(i:0..4)phil(i)?enter( ) → occupancy := occupancy + 1  
  [(i:0..4)phil(i)?exit( ) → occupancy := occupancy - 1  
  ]
```

6. Miscellaneous

7. Discussion

- APL-like brevity (7.1)
- Input/output notation is not assignment (7.1)

7. Discussion

- Explicit source/destination makes processes hard to reuse (e.g. in libraries) (7.2)

7. Discussion

- ❖ *Channels* (7.3) to connect ports exposed by processes
 - *option: multichannels*

7. Discussion

- ❖ *Buffers (7.4)* to hold output messages
 - *option*: automatic buffering (actor-like “mailboxes”)

7. Discussion

- Possibly introduce unbounded process arrays (7.5)

7. Discussion

- ❖ *Fairness* (7.6) is hard to ensure
 - should be the programmer's responsibility
 - e.g. Go lang's `select` is unfair

```
*[continue; X?stop( ) → continue := false  
  ||continue → n := n + 1  
  ]
```

7. Discussion

- CSP's imperative approach is more machine-oriented (7.7)
 - vs. Gilles Kahn's "Language for parallel programming" which is strictly deterministic (e.g. no alternative commands)

7. Discussion

- ❖ *Output guards (7.8)* is an option
 - e.g. Go lang's `select`
 - e.g. Clojure's `alt!`

```
Z :: [true → X!2; Y!3][true → Y!3; X!2]
```

7. Discussion

- Automatic termination of a repetitive command on termination of the sources of all its input guards (7.9) is
 - convenient VS. implementable?

8. Conclusion

- CSP is not a practical language
- CSP lacks more useful abstractions
- but CSP primitives can be used to construct:
 - objects
 - actors
 - monitors
 - streams
 - etc.

Questions

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