

# SRS Documentation

## Notown Record Management System

Prepared by,

DEBOPRIYO GHOSH	Roll: 35
-----------------	----------

SAIKAT JANA	Roll: 50
-------------	----------

MOUSUMI MONDAL	Roll: 40
----------------	----------

SARADINDU RANA	Roll: 28
----------------	----------

**B.Tech., Sem - V, Group- 1**

**Dated: 07/04/2021**

# Contents

---

1. Data Requirement
2. Assumptions
3. Entity - Relationship Diagram
4. Relationships
5. Relational Schema
6. Normalization of Relational Schema

**Draw the ER diagram; Design the relational schema with minimum redundancy.**

Notown Records has decided to store information about musicians who perform on its albums (as well as other company data) in a database. The company has wisely chosen to hire you as a database designer (at your usual consulting fee of \$2,500/day).

Each musician that records at Notown has an SSN, a name, an address, and a phone number. Poorly paid musicians often share the same address, and no address has more than one phone.

Each instrument that is used in songs recorded at Notown has a name (e.g., guitar, synthesizer, flute) and a musical key (e.g., C, B-flat, E-flat).

Each album that is recorded on the Notown label has a title, a copyright date, a format (e.g., CD or MC), and an album identifier.

Each song recorded at Notown has a title and an author.

Each musician may play several instruments, and a given instrument may be played by several musicians.

Each album has a number of songs on it, but no song may appear on more than one album.

Each song is performed by one or more musicians, and a musician may perform a number of songs.

Each album has exactly one musician who acts as its producer. A musician may produce several albums, of course.

## **1.Data Requirement**

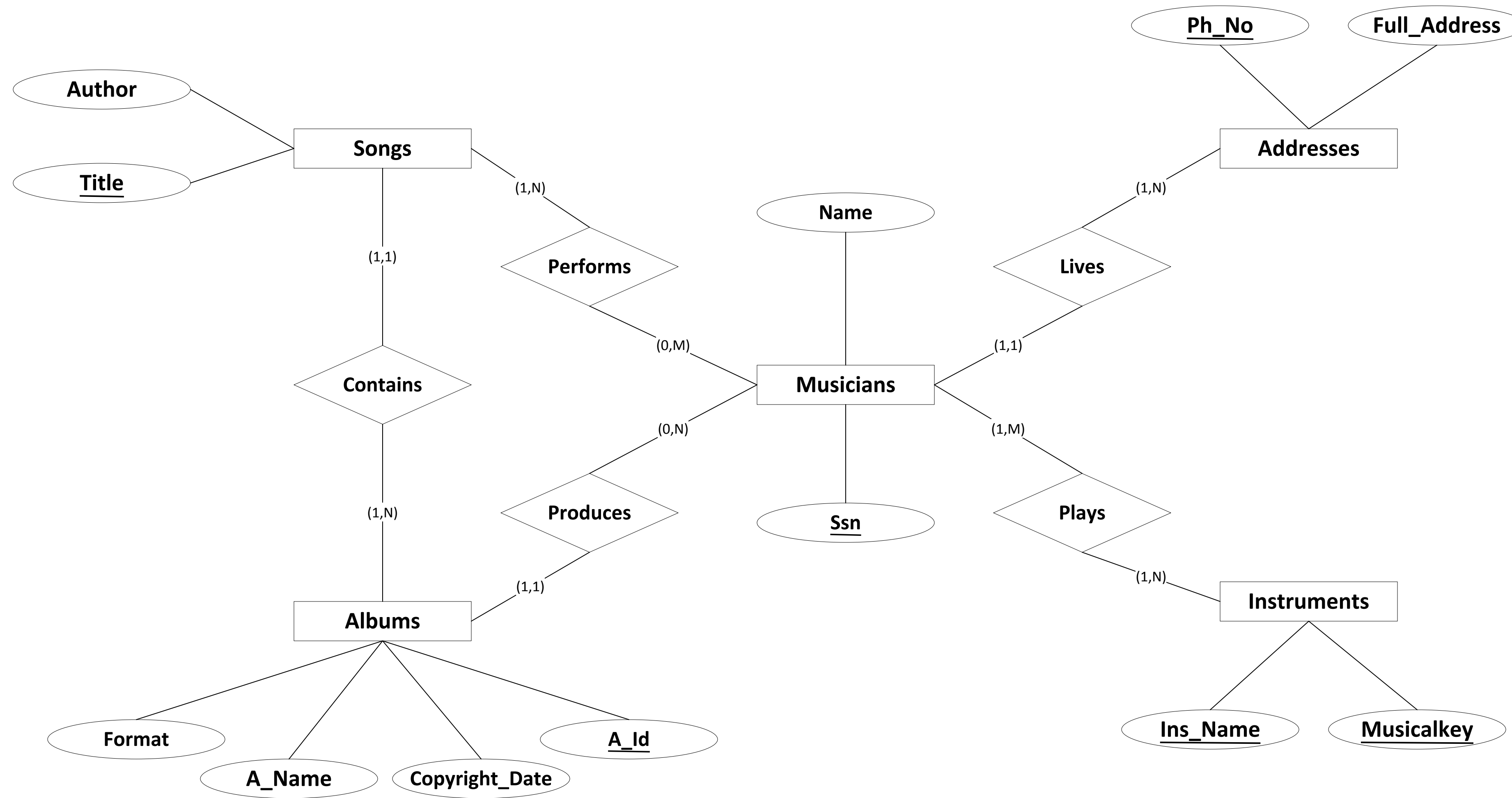
- **Songs:** Every song has a unique title. Each song is written by specific author.
- **Albums:** Every album has an unique id and respective album name. Eachone has specific format (e.g. CD or MC) and date of copyright.
- **Musicians:** Every musician has an unique Ssn number and their respective name .
- **Instruments:** Each instrument is identified by an unique name (e.g. guiter, synthesizer, flute ) and its musical key (e.g. C,B-flat ,E-flat ).
- **Addresses:** Each address of musicians identified by a contact phone number. Each one has full address.

## **2. Assumption**

- 1) An album may contain multiple songs. But the recorded albums has atleast one song.
- 2) A musician may or may not perform songs . Musician may produce one or more albums.
- 3) A musician can play one or more instruments.
- 4) A musician lives only in one address identified by it's phone number.
- 5) In one address there may be more than one poorly paid musicians.
- 6) An instrument can be played by one or more than one customer.
- 7) Every album has only one producer.
- 8) A song can be performed by one or more than one musician.

## **3. Entity Relationship Diagram**

### **A. Diagram**



## B. Description

- a) In this diagram the entities are Albums, Musicians , Songs , Instruments , Addresses.
- b) Albums contains songs so they are connected by the relationship 'Contains'.
- c) Musicians performs songs. They are connected by the relationship 'Performs'.
- d) Musicians produces album, so they are connected by the relationship 'Produces'.
- e) A musician lives in an addresses. So musician and addresses are connected by the relationship 'Lives'.
- f) Musician plays instruments, so musician and instruments is connected by the relationship 'Plays'.

## C. Attributes

Albums: {A\_Id, A\_Name, Format, Copyright\_Date}

Songs: {Title , Author}

Musicians: {Ssn, Name }

Instruments: {Ins\_Name , MusicalKey}

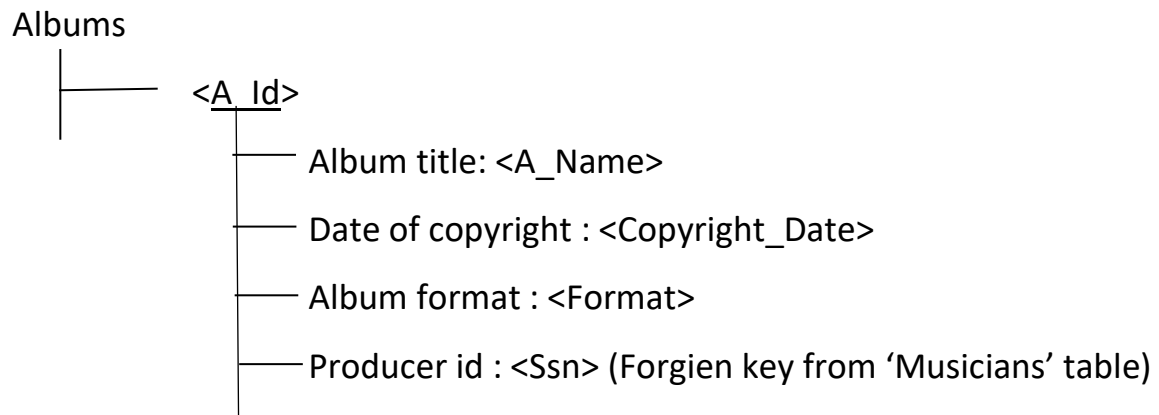
Addresses: {Ph\_no, FullAddress}

## 4.Relationships

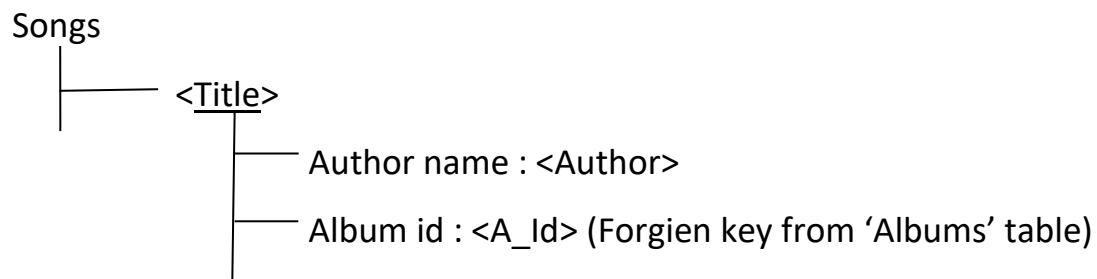
- a. Musicians-Addresses (N:1) comprises.
- b. Musicians-Instruments (N:M) comprises.
- c. Musicians-Albums (1:N) comprises.
- d. Musicians-Songs (N:M) comprises.
- e. Albums-Songs (1:N) holds.

## 5.Relational Schema

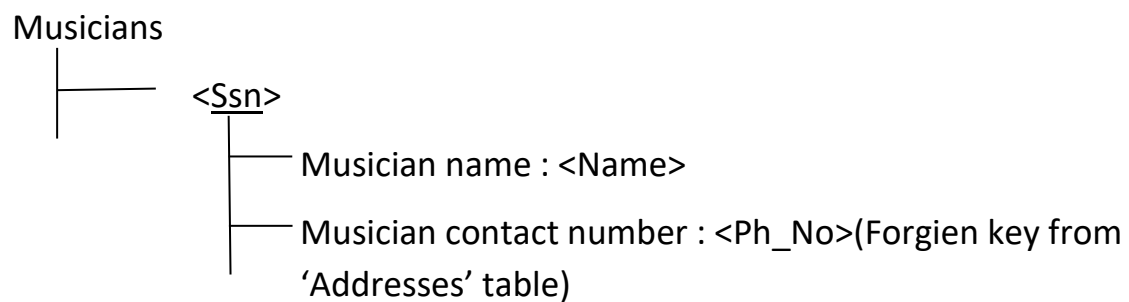
### Albums:



### Songs:

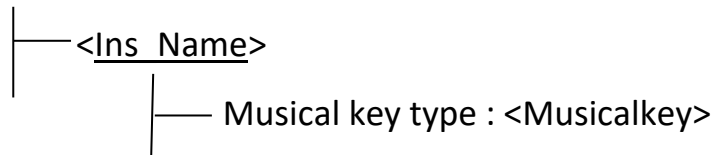


### Musicians:



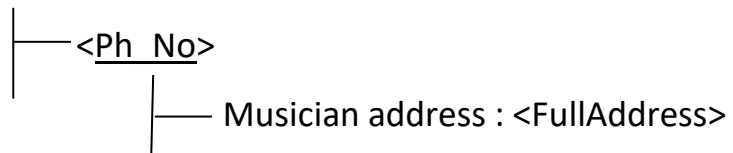
### **Instruments:**

Instruments



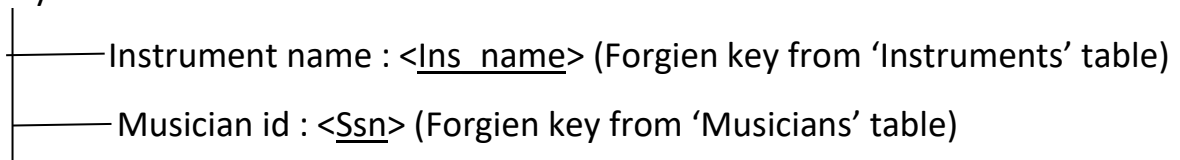
### **Addresses:**

Addresses



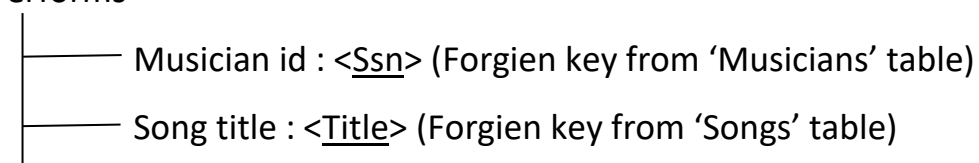
### **Plays:**

Plays



### **Performs:**

Performs





## 6. Normalization of Relational Schema

### **I) Musicians {Ssn, Name, Ph\_No }**

Ssn  $\rightarrow$  { Name, Ph\_No }

Ssn is the candidate key.

There is no multi valued attribute so the table is in **1NF**.

Ssn is the primary key and there is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

### **II) Songs { Title, Author, A\_Id }**

Title  $\rightarrow$  { Author, A\_Id }

Title is the candidate key.

There is no multi valued attribute so the table is in **1NF**.

Title is the primary key and there is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

### **III) Albums {A\_id, A\_Name, Format, Copyright\_Date, Ssn}**

A\_Id  $\rightarrow$  { A\_name, Format, Copyright\_Date, Ssn }

A\_Id is the candidate key.

There is no multi valued attribute so the table is in **1NF**.

A\_Id is the primary key and there is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

#### **IV)Instruments { Ins\_Name, Musicalkey}**

(Ins\_Name, MusicalKey} is the composite primary key.

There is no multi valued attribute so the table is in **1NF**.

There is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

#### **V)Addresses {Ph\_No, FullAddress}**

Ph\_No -> { FullAddress }

Ph\_No is the candidate key.

There is no multi valued attribute so the table is in **1NF**.

Ph\_No is the primary key and there is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

#### **VI)Performs{Ssn,Title}**

{Ssn,title} is a composite primary key.

There is no multi valued attribute so the table is in **1NF**.

There is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.

## **VII)Plays{Ssn,Ins\_Name}**

{Ssn, Ins\_Name} is a composite primary key.

There is no multi valued attribute so the table is in **1NF**.

There is no partial dependency, so the table is in **2NF**.

The table is in **3NF** as there is no transitive dependency.