## Week 8:

Write a program to color the nodes in a given graph such that no two adjacent can have the same color using backtracking.

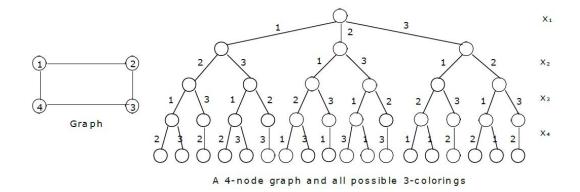
## **Graph coloring problem:**

- ✓ Let G be a graph and m be a given positive integer.
- ✓ Find whether the nodes of G can be colored in such a way that no two adjacent nodes have the same color, yet only m colors are used.
- ✓ This is termed the m-colorabiltiy decision problem.
- ✓ The m-colorability optimization problem asks for the smallest integer m for which the graph G can be colored.
- ✓ The function m-coloring will begin by first assigning the graph to its adjacency matrix, setting the array x [] to zero.
- The colors are represented by the integers 1, 2, ..., m and the solutions are given by the n-tuple  $(x_1, x_2, ..., x_n)$ , where  $x_i$  is the color of node i.
- ✓ A recursive backtracking algorithm for graph coloring is carried out by invoking the statement mcoloring(1);

## Algorithm:

```
Algorithm mColoring(k)
    This algorithm was formed using the recursive backtracking
// This algorithm was formed using the recursive backtracking
// schema. The graph is represented by its boolean adjacency
// matrix G[1:n,1:n]. All assignments of 1,2,\ldots,m to the
// vertices of the graph such that adjacent vertices are
// assigned distinct integers are printed. k is the index
// of the next vertex to color.
     repeat
     \{//\text{ Generate all legal assignments for } x[k].
          NextValue(k); // Assign to x[k] a legal color. if (x[k] = 0) then return; // No new color possible
                                  // At most m colors have been // used to color the n vertices.
          if (k=n) then
               write (x[1:n]);
          else mColoring(k+1);
     } until (false);
}
Algorithm NextValue(k)
//x[1], \ldots, x[k-1] have been assigned integer values in
// the range [1, m] such that adjacent vertices have distinct
// integers. A value for x[k] is determined in the range
//[0,m]. x[k] is assigned the next highest numbered color
// while maintaining distinctness from the adjacent vertices
 // of vertex k. If no such color exists, then x[k] is 0.
     repeat
     {
          x[k] := (x[k] + 1) \mod (m+1); // \text{Next highest color.}
          if (x[k] = 0) then return; // All colors have been used.
          for j := 1 to n do
               // Check if this color is
               // distinct from adjacent colors. if ((G[k,j] \neq 0) and (x[k] = x[j]))
               // If (k, j) is and edge and if adj.
               // vertices have the same color.
                   then break;
          if (j = n + 1) then return; // New color found
     } until (false); // Otherwise try to find another color.
}
```

Example:



## **Sample Implementation:**

```
#include<stdio.h>
#include<conio.h>
int m,n;
int c=0;
int sol=0;
int g[10][10];
int x[10];
void nextvalue(int k);
void graphcoloring(int k);
void main()
{
  int i,j,t;
  printf("\n enter the number of vertices in a graph: " );
  scanf("%d", &n);
  printf("\n enter graph edges\n");
    for(i=1;i \le n;i++)
      for(j=1;j<=n;j++)
         scanf("\%d",\&g[i][j]);
  printf("\n All possible solutions are\n");
     for(m=1;m<=n;m++)
     {
       if(c==1)
         { break; }
       graphcoloring(1);
     printf("\n chromatic number is %d",m-1);
     printf("\n total number of solutions is %d",sol);
getch();
void graphcoloring(int k)
int i;
    while(1)
       nextvalue(k);
         if(x[k]==0)
```

```
{ return; }
          if(k==n)
            {
                  c=1;
                  for(i{=}1;i{<}{=}n;i{+}{+})
                     { printf("%d ",x[i]); }
                        sol++;
                       printf("\n");
                  }
                  else
        graphcoloring(k+1);
}
void nextvalue(int k)
int j;
   while(1)
       x[k]=(x[k]+1)\%(m+1);
            if(x[k]==0)
                  return;
            for(j=1;j \le n;j++)
                  if(g[k][j]==1\&\&x[k]==x[j])
                        break;
            if(j==(n+1))
                  return;
Sample Output:
enter the number of vertices in a graph: 4
enter graph edges
\begin{smallmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{smallmatrix}
All possible solutions are
\begin{smallmatrix}1&2&1&2\\2&1&2&1\end{smallmatrix}
 chromatic number is 2 total number of solutions is 2
```