

# Bucket-Sort (§ 4.5.1)



- $\bullet$  Let be S be a sequence of n(key, element) items with keys in the range [0, N-1]
- Bucket-sort uses the keys as indices into an auxiliary array B of sequences (buckets)

Phase 1: Empty sequence S by moving each item (k, o) into its bucket B[k]

Phase 2: For i = 0, ..., N - 1, move the items of bucket B[i] to the end of sequence S

- Analysis:
  - Phase 1 takes O(n) time

• Phase 2 takes O(n + N) time Bucket-sort takes O(n + N) time

#### Algorithm *bucketSort(S, N)*

**Input** sequence *S* of (key, element) items with keys in the range [0, N-1]

**Output** sequence S sorted by increasing keys

 $B \leftarrow$  array of N empty sequences

while  $\neg S.isEmpty()$ 

 $f \leftarrow S.first()$ 

 $(k, o) \leftarrow S.remove(f)$ 

B[k].insertLast((k, o))

for  $i \leftarrow 0$  to N-1

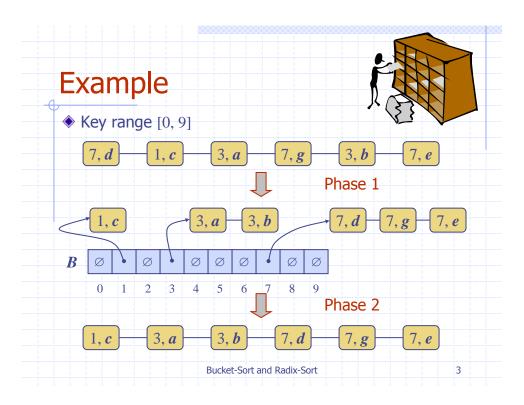
while  $\neg B[i].isEmpty()$ 

 $f \leftarrow B[i].first()$ 

 $(k, o) \leftarrow B[i].remove(f)$ 

S.insertLast((k, o))

**Bucket-Sort and Radix-Sort** 



# **Properties and Extensions**



- Key-type Property
  - The keys are used as indices into an array and cannot be arbitrary objects
  - No external comparator
- Stable Sort Property
  - The relative order of any two items with the same key is preserved after the execution of the algorithm

### **Extensions**

- Integer keys in the range [a, b]
  - Put item (k, o) into bucket B[k-a]
- String keys from a set D of possible strings, where D has constant size (e.g., names of the 50 U.S. states)
  - Sort D and compute the rank r(k) of each string k of D in the sorted sequence
  - Put item (k, o) into bucket
     B[r(k)]

Bucket-Sort and Radix-Sort

## Lexicographic Order



- lacktriangle A *d*-tuple is a sequence of *d* keys  $(k_1, k_2, ..., k_d)$ , where key  $k_i$  is said to be the *i*-th dimension of the tuple
- Example:
  - The Cartesian coordinates of a point in space are a 3-tuple
- The lexicographic order of two d-tuples is recursively defined as follows

$$(x_1, x_2, ..., x_d) < (y_1, y_2, ..., y_d)$$
 $\Leftrightarrow$ 
 $x_1 < y_1 \lor x_1 = y_1 \land (x_2, ..., x_d) < (y_2, ..., y_d)$ 

I.e., the tuples are compared by the first dimension, then by the second dimension, etc.

Bucket-Sort and Radix-Sort

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# Lexicographic-Sort

- Let C<sub>i</sub> be the comparator that compares two tuples by their i-th dimension
- Let stableSort(S, C) be a stable sorting algorithm that uses comparator C
- Lexicographic-sort sorts a sequence of d-tuples in lexicographic order by executing d times algorithm stableSort, one per dimension
- Lexicographic-sort runs in O(dT(n)) time, where T(n) is the running time of stableSort

Algorithm *lexicographicSort(S)* 

**Input** sequence *S* of *d*-tuples **Output** sequence *S* sorted in lexicographic order

for  $i \leftarrow d$  downto 1 stableSort(S, C.)

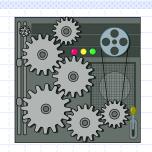
## Example:

(7,4,6) (5,1,5) (2,4,6) (2, 1, 4) (3, 2, 4) (2, 1, 4) (3, 2, 4) (5,1,5) (7,4,6) (2,4,6) (2, 1, 4) (5,1,5) (3, 2, 4) (7,4,6) (2,4,6) (2, 1, 4) (2,4,6) (3, 2, 4) (5,1,5) (7,4,6)

Bucket-Sort and Radix-Sort

## Radix-Sort (§ 4.5.2)

- Radix-sort is a specialization of lexicographic-sort that uses bucket-sort as the stable sorting algorithm in each dimension
- ◆ Radix-sort is applicable to tuples where the keys in each dimension i are integers in the range [0, N-1]
- Radix-sort runs in time
  O(d(n+N))



## Algorithm radixSort(S, N)

**Input** sequence S of d-tuples such that  $(0, ..., 0) \le (x_1, ..., x_d)$  and  $(x_1, ..., x_d) \le (N-1, ..., N-1)$  for each tuple  $(x_1, ..., x_d)$  in S **Output** sequence S sorted in lexicographic order

for  $i \leftarrow d$  downto 1 bucketSort(S, N)

Bucket-Sort and Radix-Sort

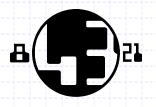
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# Radix-Sort for Binary Numbers

Consider a sequence of n
 b-bit integers

$$\boldsymbol{x} = \boldsymbol{x_{b-1}} \dots \boldsymbol{x_1} \boldsymbol{x_0}$$

- ◆ We represent each element as a b-tuple of integers in the range [0, 1] and apply radix-sort with N = 2
- This application of the radix-sort algorithm runs in O(bn) time
- For example, we can sort a sequence of 32-bit integers in linear time



#### Algorithm *binaryRadixSort(S)*

**Input** sequence *S* of *b*-bit integers

**Output** sequence S sorted replace each element x of S with the item (0, x)

for  $i \leftarrow 0$  to b-1

replace the key k of each item (k, x) of S with bit  $x_i$  of x

bucketSort(S, 2)

Bucket-Sort and Radix-Sort

