

Syntax

Type Variables	X, Y, Z	
Term Variables	x, y, z	
Types	T, S, U	$::= X$ $ I$ $ T \oplus T$ $ T \otimes T$ $ T \multimap T$ $ \mu X. T$
Terms	t, u, v	$::= x$ $ ()$ $ \text{inl } t$ $ \text{inr } t$ $ t \times t$ $ t \mapsto t$ $ \text{fold } t$ $ t \text{ trace}_T$ $ t + t$ $ - t$ $ \emptyset$
Type Contexts	Θ	$::=$
Term Contexts		$ \Theta, X$
	Γ	$::=$
		$ \Gamma, t : T$
Type Judgements		$\Theta \vdash T$
Term Judgements		$\Gamma \vdash t : T$
Expressions	e, f, g	$::= t$ $ f @ t$ $ f^\dagger$

Type Formation rules

$$\frac{}{\Theta, X \vdash X} \quad \frac{}{\Theta \vdash I} \quad \frac{\Theta \vdash T_1 \quad \Theta \vdash T_2}{\Theta \vdash T_1 \oplus T_2} \quad \frac{\Theta \vdash T_1 \quad \Theta \vdash T_2}{\Theta \vdash T_1 \otimes T_2} \quad \frac{\Theta \vdash T_1 \quad \Theta \vdash T_2}{\Theta \vdash T_1 \multimap T_2} \quad \frac{\Theta, X \vdash T}{\Theta \vdash \mu X. T}$$

Substitution

$$\begin{aligned} X[X/S] &= S \\ Y[X/S] &= Y \\ I[X/S] &= I \\ T_1 \oplus T_2[X/S] &= T_1[X/S] \oplus T_2[X/S] \\ T_1 \otimes T_2[X/S] &= T_1[X/S] \otimes T_2[X/S] \\ T_1 \multimap T_2[X/S] &= T_1[X/S] \multimap T_2[X/S] \\ \mu Y. T[X/S] &= \mu Y. (T[X/S]) \end{aligned}$$

Term Typing rules

$$\begin{array}{c}
\text{Variable} \frac{}{x : T \vdash x : T} \quad \text{Exchange} \frac{\Gamma_2, \Gamma_1 \vdash t : T}{\Gamma_1, \Gamma_2 \vdash t : T} \\
\\
I_L \frac{\Gamma \vdash t : T}{\Gamma, () : I \vdash t : T} \quad \frac{}{\vdash () : I} I_R \\
\\
\oplus_{L_l} \frac{\Gamma, t_1 : T_1 \vdash t : T}{\Gamma, \text{inl } t_1 : T_1 \oplus T_2 \vdash t : T} \quad \frac{\Gamma \vdash t_1 : T_1}{\Gamma \vdash \text{inl } t_1 : T_1 \oplus T_2} \oplus_{R_l} \\
\\
\oplus_{L_r} \frac{\Gamma, t_2 : T_2 \vdash t : T}{\Gamma, \text{inr } t_2 : T_1 \oplus T_2 \vdash t : T} \quad \frac{\Gamma \vdash t_2 : T_2}{\Gamma \vdash \text{inr } t_2 : T_1 \oplus T_2} \oplus_{R_r} \\
\\
\otimes_L \frac{\Gamma, t_1 : T_1, t_2 : T_2 \vdash t : T}{\Gamma, t_1 \times t_2 : T_1 \otimes T_2 \vdash t : T} \quad \frac{\Gamma_1 \vdash t_1 : T_1 \quad \Gamma_2 \vdash t_2 : T_2}{\Gamma_1, \Gamma_2 \vdash t_1 \times t_2 : T_1 \otimes T_2} \otimes_R \\
\\
\multimap_L \frac{\Gamma_1 \vdash t_1 : T_1 \quad \Gamma_2, t_2 : T_2 \vdash t : T}{\Gamma_1, \Gamma_2, t_1 \mapsto t_2 : T_1 \multimap T_2 \vdash t : T} \quad \frac{\Gamma, t_1 : T_1 \vdash t_2 : T_2}{\Gamma \vdash t_1 \mapsto t_2 : T_1 \multimap T_2} \multimap_R \\
\\
\mu_L \frac{\Gamma, u : U[X/\mu X.U] \vdash t : T}{\Gamma, \text{fold } u : \mu X.U \vdash t : T} \quad \frac{\Gamma \vdash t : T[X/\mu X.T]}{\Gamma \vdash \text{fold } t : \mu X.T} \mu_R \\
\\
\text{Linearity}_L \frac{\Gamma, t_1 : U \vdash t : T \quad \Gamma, t_2 : U \vdash t : T}{\Gamma, t_1 + t_2 : U \vdash t : T} \quad \frac{\Gamma \vdash t_1 : T \quad \Gamma \vdash t_2 : T}{\Gamma \vdash t_1 + t_2 : T} \text{Linearity}_R \\
\\
\text{Opposition}_L \frac{\Gamma, u : U \vdash t : T}{\Gamma, -u : U \vdash t : T} \quad \frac{\Gamma \vdash t : T}{\Gamma \vdash -t : T} \text{Opposition}_R \\
\\
\text{Trace}_L \frac{\Gamma, u : S \oplus U_1 \multimap S \oplus U_2 \vdash t : T}{\Gamma, u \text{ trace}_S : U_1 \multimap U_2 \vdash t : T} \quad \frac{\Gamma \vdash t : U \oplus T_1 \multimap U \oplus T_2}{\Gamma \vdash t \text{ trace}_U : T_1 \multimap T_2} \text{Trace}_R
\end{array}$$

Expression Typing rules

$$\begin{array}{c}
\text{Application} \frac{\vdash f : T_1 \multimap T_2 \quad \vdash t : T_2}{\vdash f t : T_1} \quad \frac{\vdash f : T_1 \multimap T_2 \quad \vdash g : T_2 \multimap T_3}{\vdash f \circ g : T_1 \multimap T_3} \text{Composition} \\
\\
\frac{\vdash f : T \multimap U}{\vdash f^\dagger : U \multimap T} \text{Dagger}
\end{array}$$

$$\begin{array}{c}
 \text{Match-var} \frac{x = t \hookrightarrow [x \rightarrow t]}{[x \rightarrow t] x \rightsquigarrow t} \quad \text{Subst-var} \\
 \\
 \text{Match-unit} \frac{() = () \hookrightarrow []}{[] () \rightsquigarrow ()} \quad \text{Subst-unit} \\
 \\
 \text{Match-inl} \frac{t = u \hookrightarrow [\Xi]}{\text{inl } t = \text{inl } u \hookrightarrow [\Xi]} \quad \frac{[\Xi] t \rightsquigarrow t'}{[\Xi] \text{inl } t \rightsquigarrow \text{inl } t'} \quad \text{Subst-inl} \\
 \\
 \text{Match-inr} \frac{t = u \hookrightarrow [\Xi]}{\text{inr } t = \text{inr } u \hookrightarrow [\Xi]} \quad \frac{[\Xi] t \rightsquigarrow t'}{[\Xi] \text{inr } t \rightsquigarrow \text{inr } t'} \quad \text{Subst-inr} \\
 \\
 \text{Match-tensor} \frac{t_1 = u_1 \hookrightarrow [\Xi_1] \quad t_2 = u_2 \hookrightarrow [\Xi_2]}{t_1 \times t_2 = u_1 \times u_2 \hookrightarrow [\Xi_1, \Xi_2]} \quad \frac{[\Xi_1] t_1 \rightsquigarrow t'_1 \quad [\Xi_2] t_2 \rightsquigarrow t'_2}{[\Xi_1, \Xi_2] t_1 \times t_2 \rightsquigarrow t'_1 \times t'_2} \quad \text{Subst-tensor} \\
 \\
 \text{Match-arrow} \frac{t_1 = u_1 \hookrightarrow [\Xi_1] \quad t_2 = u_2 \hookrightarrow [\Xi_2]}{t_1 \mapsto t_2 = u_1 \mapsto u_2 \hookrightarrow [\Xi_1, \Xi_2]} \quad \frac{[\Xi_1] t_1 \rightsquigarrow t'_1 \quad [\Xi_2] t_2 \rightsquigarrow t'_2}{[\Xi_1, \Xi_2] t_1 \mapsto t_2 \rightsquigarrow t'_1 \mapsto t'_2} \quad \text{Subst-arrow} \\
 \\
 \text{Match-fold} \frac{t = u \hookrightarrow [\Xi]}{\text{fold } t = \text{fold } u \hookrightarrow [\Xi]} \quad \frac{[\Xi] t \rightsquigarrow t'}{[\Xi] \text{fold } t \rightsquigarrow \text{fold } t'} \quad \text{Subst-fold} \\
 \\
 \text{Match-lin1} \frac{t = u_1 \hookrightarrow [\Xi] \quad t = u_2 \hookrightarrow [\Xi]}{t = u_1 + u_2 \hookrightarrow [\Xi]} \quad \frac{[\Xi] t_1 \rightsquigarrow t'_1 \quad [\Xi] t_2 \rightsquigarrow t'_2}{[\Xi] t_1 + t_2 \rightsquigarrow t'_1 + t'_2} \quad \text{Subst-lin} \\
 \\
 \text{Match-lin2} \frac{t_1 = u \hookrightarrow [\Xi] \quad t_2 = u \hookrightarrow [\Xi]}{t_1 + t_2 = u \hookrightarrow [\Xi]} \\
 \\
 \text{Match-opp-c} \frac{v = u \hookrightarrow [\Xi]}{-v = -u \hookrightarrow [\Xi]} \quad \frac{[\Xi] v \rightsquigarrow v'}{[\Xi] -v \rightsquigarrow -v'} \quad \text{Subst-opp} \\
 \\
 \text{Match-opp-l} \frac{v = -u \hookrightarrow [\Xi]}{-v = u \hookrightarrow [\Xi]} \\
 \\
 \text{Match-opp-r} \frac{v = u \hookrightarrow -[\Xi]}{v = -u \hookrightarrow [\Xi]} \\
 \\
 \text{Match-trace} \frac{t = u \hookrightarrow [\Xi]}{t \text{ trace}_T = u \text{ trace}_T \hookrightarrow [\Xi]} \quad \frac{[\Xi] t \rightsquigarrow t'}{[\Xi] t \text{ trace}_T \rightsquigarrow t' \text{ trace}_T} \quad \text{Subst-trace} \\
 \\
 \text{App-arrow} \frac{t_1 = t \hookrightarrow [\Xi] \quad [\Xi] t_2 \rightsquigarrow t'}{(t_1 \mapsto t_2) @ t \Rightarrow t'} \quad \frac{t_1 @ t \Rightarrow t'_1 \quad t_2 @ t \Rightarrow t'_2}{(t_1 + t_2) @ t \Rightarrow t'_1 + t'_2} \quad \text{App-lin} \\
 \\
 \frac{f^\dagger \rightsquigarrow f' \quad f' @ t \Rightarrow t'}{f^\dagger @ t \Rightarrow t'} \quad \text{App-dagger} \\
 \\
 \text{App-trace-nix} \frac{t @ (\text{inr } u) \Rightarrow \text{inr } u' \quad (t \text{ trace}_X) @ u \Rightarrow u'}{t @ (\text{inl } u) \Rightarrow \text{inl } u' \quad (t \text{ trace}_X) @ \text{inl } u \Rightarrow u''} \quad \text{App-trace-loop} \\
 \\
 \text{App-trace-start} \frac{t @ (\text{inr } u) \Rightarrow \text{inl } u' \quad (t \text{ trace}_X) @ (\text{inl } u') \Rightarrow u''}{(t \text{ trace}_X) @ u \Rightarrow u''} \quad \frac{t @ (\text{inl } u) \Rightarrow \text{inr } u' \quad (t \text{ trace}_X) @ \text{inl } u \Rightarrow u'}{(t \text{ trace}_X) @ u \Rightarrow u''} \quad \text{App-trace-end} \\
 \\
 \text{Flip-arrow} \frac{}{(t_1 \mapsto t_2)^\dagger \rightsquigarrow t_2 \mapsto t_1} \quad \frac{t_1^\dagger \rightsquigarrow t'_1 \quad t_2^\dagger \rightsquigarrow t'_2}{(t_1 + t_2)^\dagger \rightsquigarrow t'_1 + t'_2} \quad \text{Flip-lin} \\
 \\
 \text{Flip-app} \frac{f @ t \Rightarrow t' \quad t'^\dagger \rightsquigarrow t''}{(f @ t)^\dagger \rightsquigarrow t''} \quad \frac{t^\dagger \rightsquigarrow t'}{(t \text{ trace}_X)^\dagger \rightsquigarrow t' \text{ trace}_X} \quad \text{Flip-trace}
 \end{array}$$

Equivalence relation

$$\begin{array}{c}
\frac{}{t_1 + t_2 \equiv t_2 + t_1} \quad \frac{}{(t_1 + t_2) + (t_3 + t_4) \equiv (t_1 + t_3) + (t_2 + t_4)} \\
\frac{t_1 \equiv t_2}{t_1 + t_2 \equiv t_1} \quad \frac{t_1 \equiv t_2}{t_1 + -t_2 \equiv \emptyset} \quad \frac{}{t + \emptyset \equiv t} \\
\frac{}{() \equiv ()} \quad \frac{}{x \equiv x} \\
\frac{t \equiv t'}{\text{inl } t \equiv \text{inl } t'} \quad \frac{t \equiv t'}{\text{inr } t \equiv \text{inr } t'} \\
\frac{t_1 \equiv t'_1 \quad t_2 \equiv t'_2}{t_1 \times t_2 \equiv t'_1 \times t'_2} \quad \frac{t_1 \equiv t'_1 \quad t_2 \equiv t'_2}{t_1 \mapsto t_2 \equiv t'_1 \mapsto t'_2} \\
\frac{t \equiv t'}{\text{fold } t \equiv \text{fold } t'} \quad \frac{t \equiv t'}{t \text{ trace}_T \equiv t' \text{ trace}_T} \\
\frac{t \equiv t'}{-t \equiv -t'} \\
\frac{}{\text{inl } (t_1 + t_2) \equiv \text{inl } t_1 + \text{inl } t_2} \quad \frac{}{\text{inr } (t_1 + t_2) \equiv \text{inr } t_1 + \text{inr } t_2} \\
\frac{}{(t_1 + t_2) \times t_3 \equiv (t_1 \times t_3) + (t_2 \times t_3)} \quad \frac{}{t_1 \times (t_2 + t_3) \equiv (t_1 \times t_2) + (t_1 \times t_3)} \\
\frac{}{(t_1 + t_2) \mapsto t_3 \equiv (t_1 \mapsto t_3) + (t_2 \mapsto t_3)} \quad \frac{}{t_1 \mapsto (t_2 + t_3) \equiv (t_1 \mapsto t_2) + (t_1 \mapsto t_3)} \\
\frac{}{\text{fold } (t_1 + t_2) \equiv \text{fold } t_1 + \text{fold } t_2} \quad \frac{}{(t_1 + t_2) \text{ trace}_T \equiv (t_1 \text{ trace}_T) + (t_2 \text{ trace}_T)} \\
\frac{}{-(t_1 + t_2) \equiv -t_1 + -t_2} \\
\frac{}{-\text{inl } t \equiv \text{inl } (-t)} \quad \frac{}{-\text{inr } t \equiv \text{inr } (-t)} \\
\frac{}{-(t_1 \times t_2) \equiv (-t_1) \times (-t_2)} \quad \frac{}{-(t_1 \mapsto t_2) \equiv (-t_1) \mapsto (-t_2)} \\
\frac{}{-\text{fold } t \equiv \text{fold } (-t)} \quad \frac{}{-(t \text{ trace}_T) \equiv (-t) \text{ trace}_T} \\
\frac{}{-(-t) \equiv t}
\end{array}$$

partial order

$$x < () < \text{inl } t < \text{inr } t < t_1 \times t_2 < t_1 \mapsto t_2 < \text{fold } t < t \text{ trace}_T < t_1 + t_2 < -t$$

Type Interpretation

\mathbf{V} is Compact Closed Category with Finite Biproduct.

Π_i is projection functor.

K_I is constant I functor.

$[-, -]$ is internal hom functor.

$(-)^*$ is contravariant anafunctor. Using Axiom of Choice, we can define it on strict functor (see <https://ncatlab.org/nlab/show/rigid+monoidal+category#remarks>).

Id is identity functor.

$$\begin{aligned}
\llbracket \Theta \vdash T \rrbracket &: \mathbf{V}^{|\Theta|} \rightarrow \mathbf{V} \\
\llbracket \Theta \vdash X_i \rrbracket &= \Pi_i \\
\llbracket \Theta \vdash I \rrbracket &= K_I \\
\llbracket \Theta \vdash T_1 \oplus T_2 \rrbracket &= \oplus \circ \langle \llbracket \Theta \vdash T_1 \rrbracket, \llbracket \Theta \vdash T_2 \rrbracket \rangle \\
\llbracket \Theta \vdash T_1 \otimes T_2 \rrbracket &= \otimes \circ \langle \llbracket \Theta \vdash T_1 \rrbracket, \llbracket \Theta \vdash T_2 \rrbracket \rangle \\
\llbracket \Theta \vdash T_1 \multimap T_2 \rrbracket &= \otimes \circ [(-)^*, Id] \circ \langle \llbracket \Theta \vdash T_1 \rrbracket, \llbracket \Theta \vdash T_2 \rrbracket \rangle \\
\llbracket \Theta \vdash \mu X. T \rrbracket &= \llbracket \Theta, X \vdash T \rrbracket^\sharp
\end{aligned}$$

$\mathbf{V}(A, B)$ is morphism which Domain and Codomain are $A \in \text{Obj}(\mathbf{V})$ and $B \in \text{Obj}(\mathbf{V})$, respectively.

id_T is identity morphism of Object T .

$0_{A,B}$ is zero morphism of $\mathbf{V}(A, B)$.

ι_i is injection morphism.

π_i is projection morphism.

$$\pi_i = \iota_i^{-1}$$

$$\pi_i \circ \iota_j = id \text{ if } i = j$$

$$\pi_i \circ \iota_j = 0 \text{ otherwise}$$

$$unfold_{\mu X.T} = fold_{\mu X.T}^{-1}$$

$$f \circ (g \oplus h) = (f \circ g) \oplus (f \circ h)$$

$$(f \oplus g) \circ h = (f \circ h) \oplus (g \circ h)$$

$$id_A \otimes id_B = id_{A \otimes B}$$

$$(f \otimes g) \circ (h \otimes k) = (f \circ h) \otimes (g \circ k)$$

$$f^{**} = f$$

$$(f \otimes g)^* = f^* \otimes g^*$$

$$\Delta_T : T \rightarrow T \oplus T$$

$$\nabla_T : T \oplus T \rightarrow T$$

$$\llbracket t : T \rrbracket \in \mathbf{V}(I, T)$$

$$\llbracket x : T \rrbracket := x_T$$

$$\llbracket () : T \rrbracket := id_I$$

$$\llbracket \text{inl } t_1 : T_1 \oplus T_2 \rrbracket := \iota_1 \circ \llbracket t_1 : T_1 \rrbracket$$

$$\llbracket \text{inr } t_2 : T_1 \oplus T_2 \rrbracket := \iota_2 \circ \llbracket t_2 : T_2 \rrbracket$$

$$\llbracket t_1 \times t_2 : T_1 \otimes T_2 \rrbracket := \llbracket t_1 : T_1 \rrbracket \otimes \llbracket t_2 : T_2 \rrbracket$$

$$\llbracket t_1 \mapsto t_2 : T_1 \multimap T_2 \rrbracket := \llbracket t_1 : T_1 \rrbracket^* \otimes \llbracket t_2 : T_2 \rrbracket$$

$$\llbracket \text{fold } t : \mu X.T \rrbracket := fold_{\mu X.T} \circ \llbracket t : T[X/\mu X.T] \rrbracket$$

$$\llbracket t_1 + t_2 : T \rrbracket := \nabla_T \circ \llbracket t_1 : T \rrbracket \oplus \llbracket t_2 : T \rrbracket \circ \Delta_T$$

$$\emptyset \vdash t : T \triangleright \emptyset$$

$$\text{Variable} \frac{\Gamma(x) = T}{\Gamma \vdash x : T \triangleright \Gamma \setminus \{x\}}$$

$$I_L \frac{\Gamma \vdash t : T \triangleright \Gamma'}{\Gamma, () : I \vdash t : T \triangleright \Gamma'} \quad \frac{}{\emptyset \vdash () : I \triangleright \emptyset} I_R$$

$$\oplus_{L_l} \frac{\Gamma, t_1 : T_1 \vdash t : T \triangleright \Gamma'}{\Gamma, \text{inl } t_1 : T_1 \oplus T_2 \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t_1 : T_1 \triangleright \Gamma'}{\Gamma \vdash \text{inl } t_1 : T_1 \oplus T_2 \triangleright \Gamma'} \oplus_{R_l}$$

$$\oplus_{L_r} \frac{\Gamma, t_2 : T_2 \vdash t : T \triangleright \Gamma'}{\Gamma, \text{inr } t_2 : T_1 \oplus T_2 \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t_2 : T_2 \triangleright \Gamma'}{\Gamma \vdash \text{inr } t_2 : T_1 \oplus T_2 \triangleright \Gamma'} \oplus_{R_r}$$

$$\otimes_L \frac{\Gamma, t_1 : T_1, t_2 : T_2 \vdash t : T \triangleright \Gamma'}{\Gamma, t_1 \times t_2 : T_1 \otimes T_2 \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t_1 : T_1 \triangleright \Gamma' \quad \Gamma' \vdash t_2 : T_2 \triangleright \Gamma''}{\Gamma \vdash t_1 \times t_2 : T_1 \otimes T_2 \triangleright \Gamma''} \otimes_R$$

$$\multimap_L \frac{\Gamma \vdash t_1 : T_1 \triangleright \Gamma' \quad \Gamma', t_2 : T_2 \vdash t : T \triangleright \Gamma''}{\Gamma, t_1 \mapsto t_2 : T_1 \multimap T_2 \vdash t : T \triangleright \Gamma''} \quad \frac{\Gamma, t_1 : T_1 \vdash t_2 : T_2 \triangleright \Gamma'}{\Gamma \vdash t_1 \mapsto t_2 : T_1 \multimap T_2 \triangleright \Gamma'} \multimap_R$$

$$\mu_L \frac{\Gamma, u : U[X/\mu X.U] \vdash t : T \triangleright \Gamma'}{\Gamma, \text{fold } u : \mu X.U \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t : T[X/\mu X.T] \triangleright \Gamma'}{\Gamma \vdash \text{fold } t : \mu X.T \triangleright \Gamma'} \mu_R$$

$$\text{Linearity}_L \frac{\Gamma, t_1 : U \vdash t : T \triangleright \Gamma' \quad \Gamma, t_2 : U \vdash t : T \triangleright \Gamma'}{\Gamma, t_1 + t_2 : U \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t_1 : T \triangleright \Gamma' \quad \Gamma \vdash t_2 : T \triangleright \Gamma'}{\Gamma \vdash t_1 + t_2 : T \triangleright \Gamma'} \text{Linearity}_R$$

$$\text{Opposition}_L \frac{\Gamma, u : U \vdash t : T \triangleright \Gamma'}{\Gamma, -u : U \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t : T \triangleright \Gamma'}{\Gamma \vdash -t : T \triangleright \Gamma'} \text{Opposition}_R$$

$$\text{Trace}_L \frac{\Gamma, u : S \oplus U_1 \multimap S \oplus U_2 \vdash t : T \triangleright \Gamma'}{\Gamma, u \text{ trace}_S : U_1 \multimap U_2 \vdash t : T \triangleright \Gamma'} \quad \frac{\Gamma \vdash t : U \oplus T_1 \multimap U \oplus T_2 \triangleright \Gamma'}{\Gamma \vdash t \text{ trace}_U : T_1 \multimap T_2 \triangleright \Gamma'} \text{Trace}_R$$